



May 24 to July 28

Optimizing the IPL Schedule

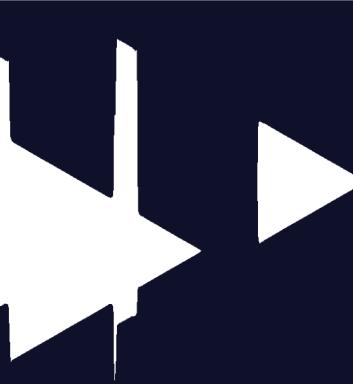




Travel and Accommodation Cost Optimization for the Indian Premier League (IPL) Cricket Tournament

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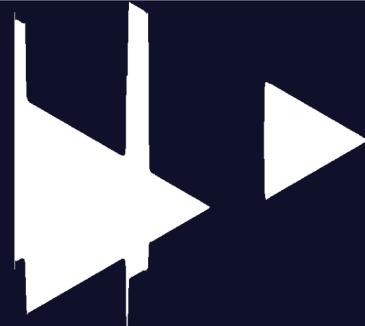




Let's Formulate!!!!

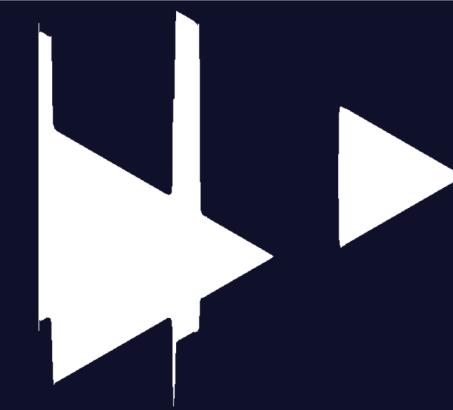
Formulating the Optimization Model

- The Indian Premier League (IPL): A Global T20 Cricket Powerhouse
- Scheduling Challenges: Teams, Venues, Broadcast Windows
- The Optimization Goal: Minimizing Travel & Accommodation Costs
- Our Approach: The Travelling Tournament Problem (TTP) Framework



Assumptions

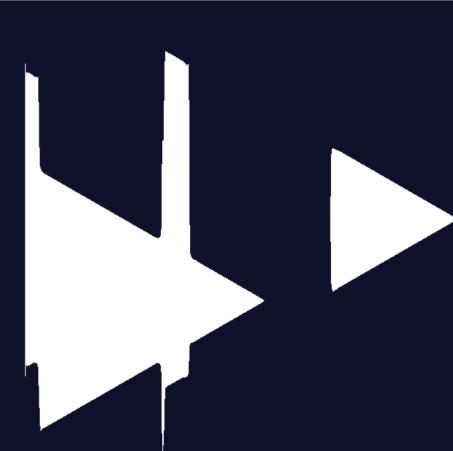
- For a route from location A to location B, we always choose the airline who provides us the minimum fare.
- For any location we always choose the hotel which provides us with the minimum tariff.
- The cost of flights and hotels are constant for weekdays and then increase by a factor of ‘o’ during weekends
- For weekends, we can randomly assign any timing slot to the two matches being played on the same weekend day



Problem Formulation

Defining the Building Blocks

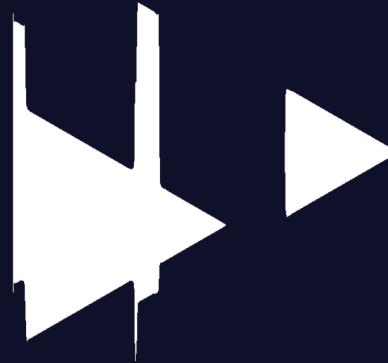
- I: Teams (MI, CSK, KKR, etc) -> 8 Teams
- V: Venues (Wankhede Stadium, Eden Gardens, etc) -> 9 Venues
- D: Days of the Season (from May 24th - July 28th) -> Represented by natural numbers
- O: Broadcast Slots (Weekend, Weekday) -> 2 slots
- J: Airline Partners (e.g., Air India, Vistara, Indigo)
- H: Hotel Partner (e.g., Taj, Marriott, ITC)
- P: Playoff Matches (e.g., Semi-finals, Final)
- Parameters:
- c_vwijo: Travel Cost (Cost depends on origin, destination, team, airline, and broadcast slot)
- a_voh: Accommodation Cost (Cost depends on venue, hotel and broadcast slot)



Decision Variables & Objective Function

Decision Variables:

- x_{ijvd} : Binary variable (1 if team i plays team j on day d at venue v, 0 otherwise)
- z_{ijvdw} : Binary variable (1 if team i travels directly from venue v to venue w for their next match after playing team j on day d, 0 otherwise)
- p_{mndo} : Binary variable (1 if playoff match m is played in timeslot o of day d in broadcast category n, 0 otherwise)



Objective Function

Minimize: $\sum (\sum c_{vwijo} * x_{ijvd} * o) + \sum (\sum c_{uwjio} * z_{ijvdw} * o)$

Travel Cost

i l, j l, v V, w V, d D, o O
i ≠ j, v ≠ w, u ≠ w

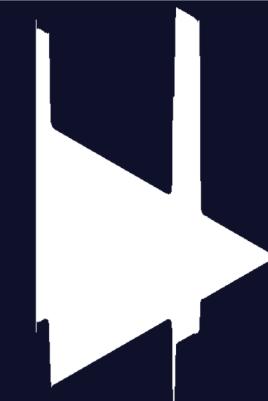
+ $\sum (a_{voh} * (\sum x_{ijvd} + \sum y_{ijvd}) * o)$

Accommodation Cost

i l, j l, v V, d D, o O, h H

+ $\sum (a_{vh} * \sum p_{mndo}) // \text{Accommodation (Playoff + Finals)}$

m P, v V, d D, n N, h H

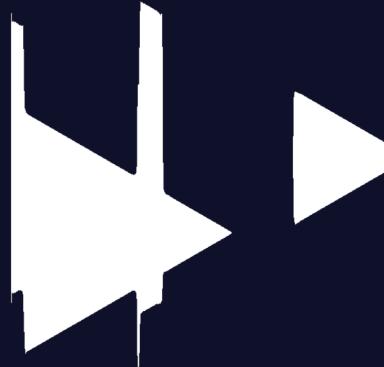


Objective Function Breakdown:

Minimize: We want to find a schedule that minimizes the total cost for both travel and accommodation.

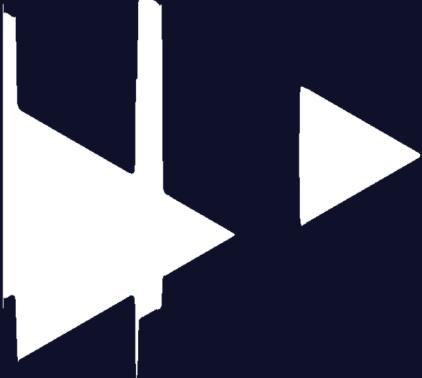
Travel Cost:

- $\Sigma (\Sigma c_{vwijo} * x_{ijvd} * o)$: This part calculates the outbound travel cost for all teams throughout the season.
 - Σ (summation): Loops through all possible combinations of teams (i), venues (v), days (d), and broadcast slots (o).
 - c_{vwijo} : Travel cost for team i from venue v to venue w using airline partner j for a match in slot o.
 - x_{ijvd} : Binary variable (1 if team i plays team j on day d at venue v, 0 otherwise).
 - o : Broadcast slot for the match (e.g., 1 - Weekend, 2 - Weekday).
- $\Sigma (\Sigma c_{uwjio} * z_{ijvdw} * o)$: This part calculates the return travel cost for teams (if they travel directly back after an away match).
 - Similar logic to the outbound cost, but uses z_{ijvdw} which indicates if a team travels directly from venue v (played against j) to venue w for their next match (d).



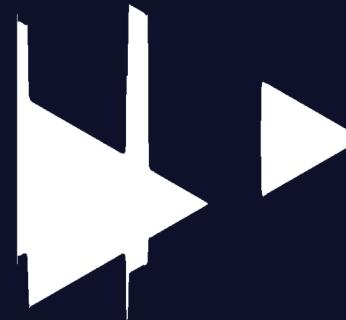
Accommodation Cost (Regular Season):

- $\sum (a_{voh} * (\sum x_{ijvd} + \sum y_{ijvd}) * o)$: This part calculates the total accommodation cost for teams during the regular season, considering both home and away matches.
 - a_{voh} : Accommodation cost per day for a team at venue v during a match in slot o for hotel h .
 - $\sum x_{ijvd}$: Summation of x_{ijvd} for all teams (i) playing at venue v on day d . This captures the accommodation cost for away teams.
 - $\sum y_{ijvd}$ (introduced earlier): Binary variable (1 if team i plays a return match against team j at their home venue (v) on day d , 0 otherwise). This captures additional accommodation for visiting teams during return matches.



Accommodation Cost (Playoffs):

- $\Sigma (a_{vh} * \Sigma p_{mndo})$: This part calculates the accommodation cost for playoff and final matches.
 - **a_vh:** Accommodation cost per day at venue v at hotel h. (Assumes same cost for all teams during playoffs).
 - **Σp_{mndo} :** Summation of p_{mndo} for all playoff matches (m) played at venue v on day d in broadcast slot o.



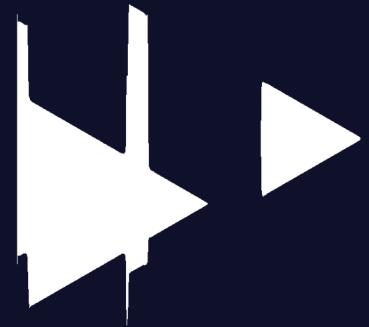
Constraints

- Match Constraint: Total 14 matches and 2 between each pair

$$\sum_{j \in d} x_{egd} = 2^*(n-1) \quad \forall i \in I, n = |I|$$

- One match per day (for weekdays):

$$\sum_{ij} x_{egd} \leq 1 \quad \forall i \in I$$

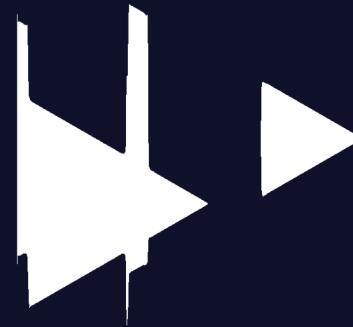


Constraints

- No match within same team

$$\chi_{ii} = 0 \quad \forall i \in I$$

- Venue constraints: Limit options of venues for a pair to 3 {H,O,N}



Constraints

- Rest days: Regular & Playoffs

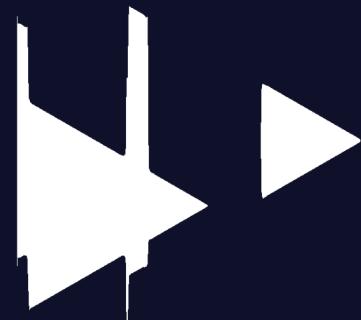
4) Rest days :

a) Regular season -

$$x_{ejd} + x_{ej(d+1)} + x_{ej(d+2)} \leq 1 \quad \forall j \in I \quad d \in D$$

b) Playoffs & Finals -

$$x_{ejd} + x_{ej(d+1)} \leq 1 \quad \forall j \in I, d \in D$$



Constraints

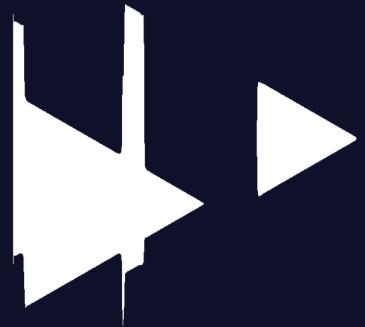
- Broadcast slot constraint: Weekdays and weekends

a) Weekdays -

$$\sum_{i,j} x_{i,j,o} \leq 4 \quad \forall o \in \text{weekdays}$$

b) Weekends -

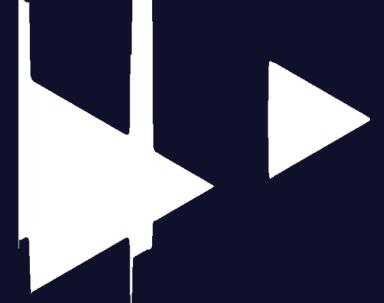
$$\sum_{i,j} x_{i,j,o} \leq 6 \quad \forall o \in \text{weekends}$$



Constraints

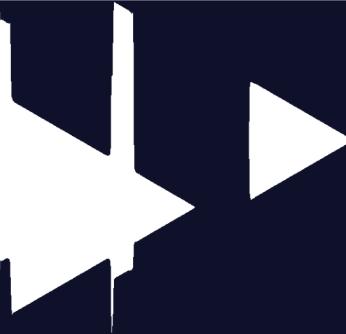
- Final Match Constraint

$$\sum_{ij} \chi_{ij}(d=28\text{ July}) = 1$$



Solution Approaches

- Mixed-Integer Linear Programming (MILP): Given the binary decision variables (x_{ijvd} , z_{ijvdw} , p_{mndo}), and linear constraints representing the scheduling rules, MILP is a powerful and widely used optimization technique for solving this problem. It allows us to incorporate both integer and continuous variables within the model.
- Solvers: Gurobi, CBC
- Metaheuristic: Due to the large computational complexity of the problem, we can go for this non exact algorithm. It will provide us with a near optimal solution. Eg: Genetic algorithm, Tabu search, Ant Colony Optimization



Thank You