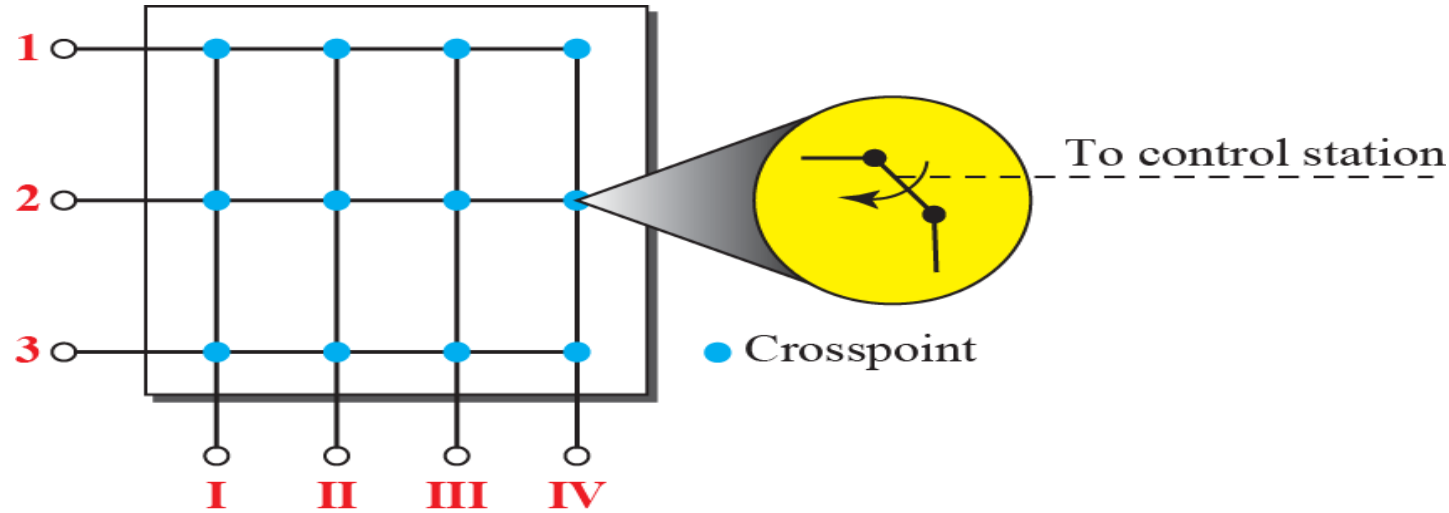


# Structure of Circuit Switches

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Circuit switching can use either of two technologies: the **space-division switch** or the **time-division switch**.

# Crossbar Switch with 3 Inputs and 4 Outputs

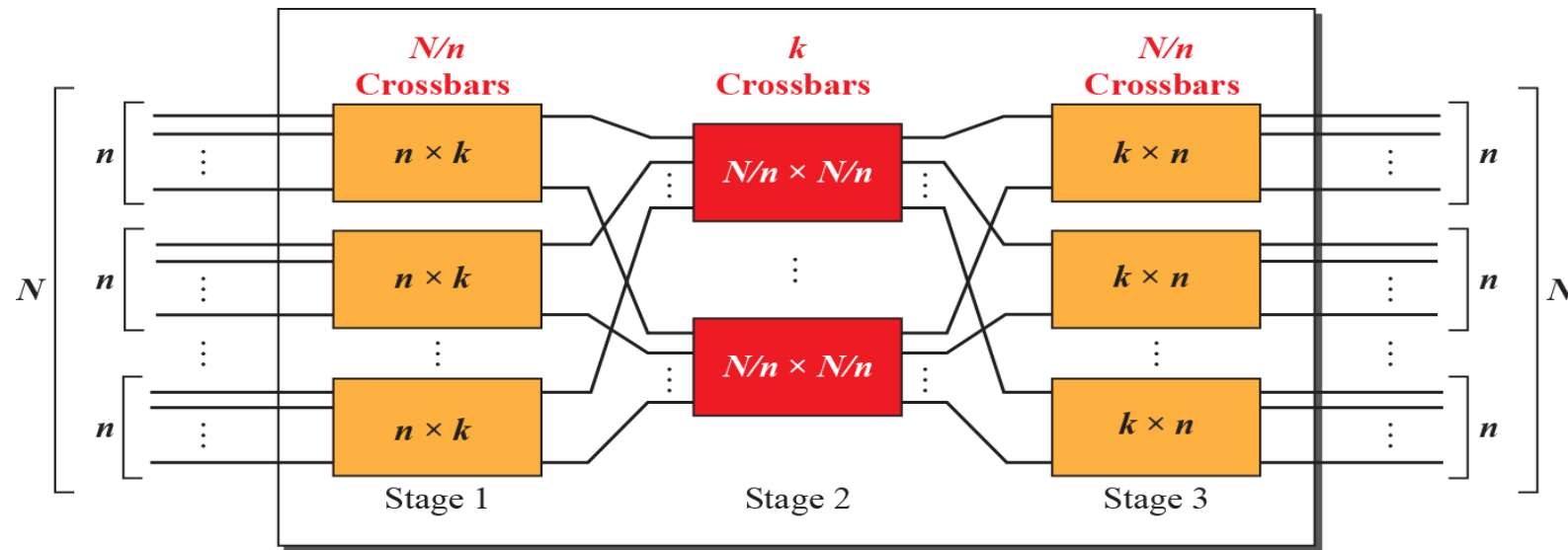


$N^2$  crosspoints

Most of the crosspoints may not be in use at any given time

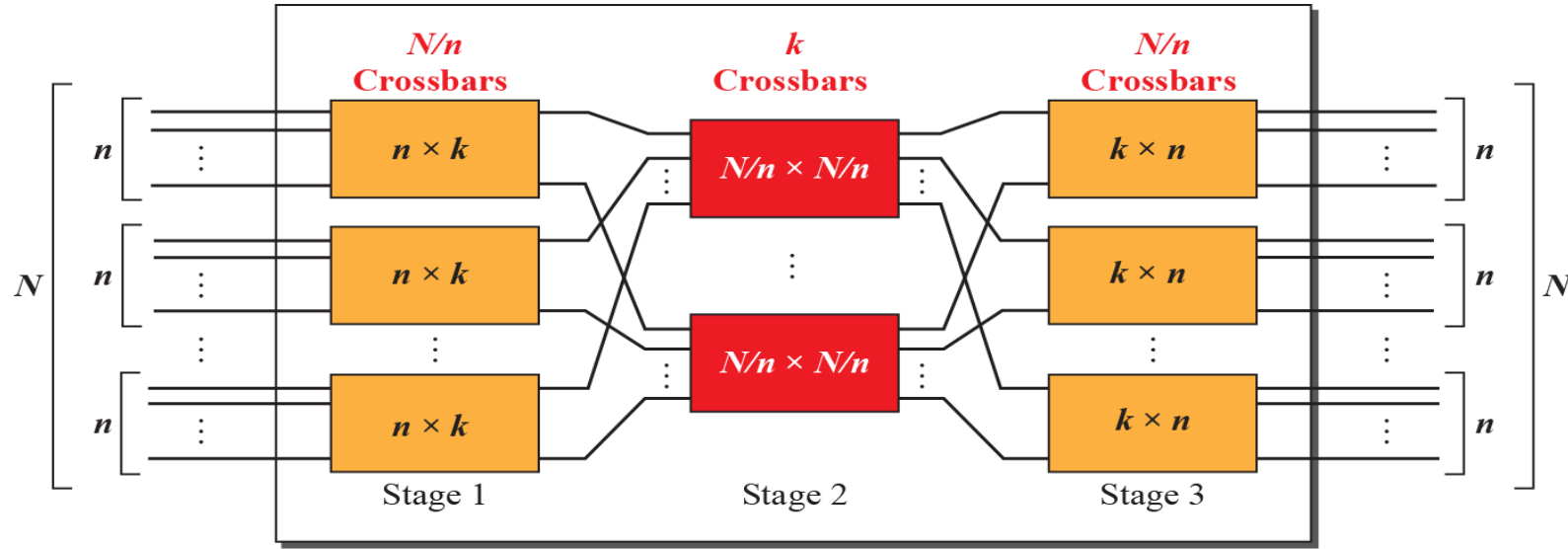
*Non-blocking*; connection never denied due to lack of crosspoints

# Multistage Switch



- Three stages of smaller space-division switches
- Each set of  $n$  input lines shares  $k$  paths to any one of the switches in the last stage

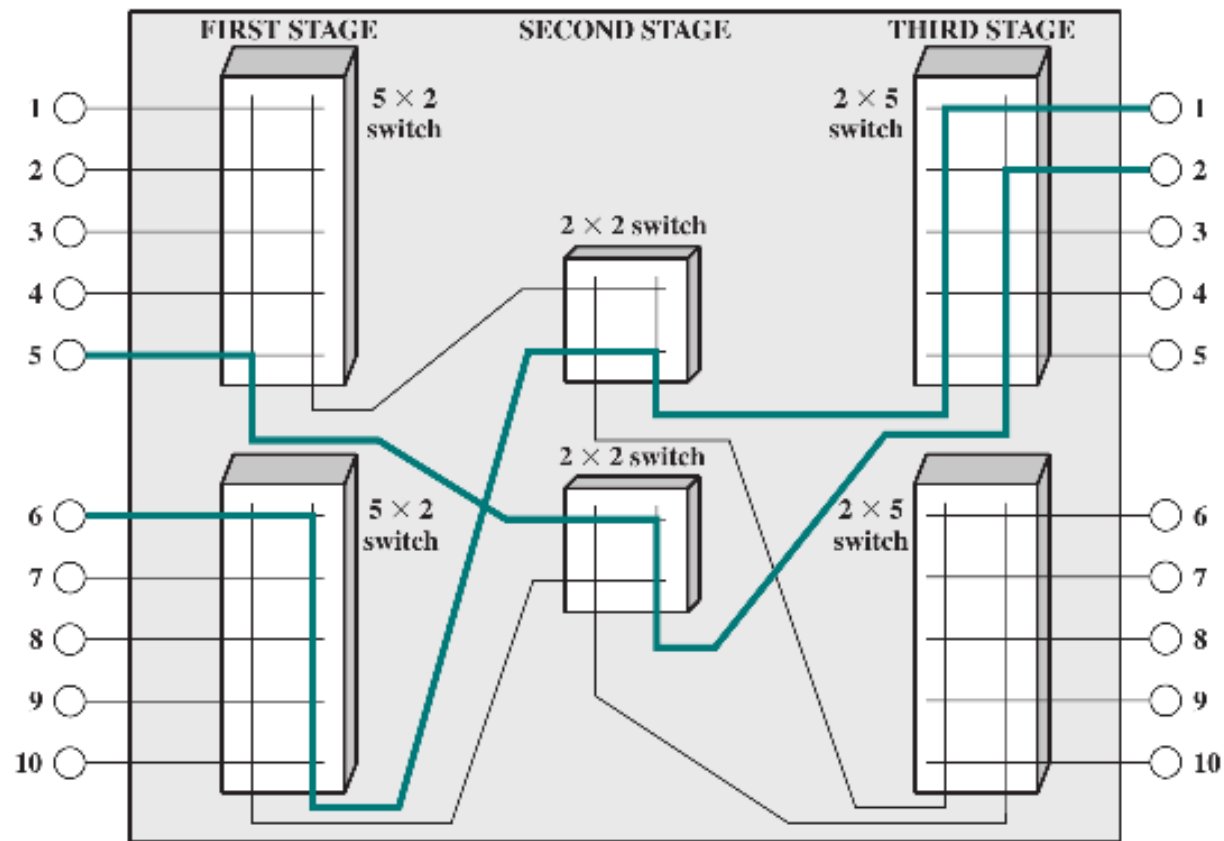
# Multistage Switch



Total number of cross points:  $\frac{N}{n} (n \times k) + k \left( \frac{N}{n} \times \frac{N}{n} \right) + \frac{N}{n} (k \times n) = 2kN + k \left( \frac{N}{n} \right)^2$

# Multistage Switch

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## Example

Design a three-stage,  $200 \times 200$  switch ( $N = 200$ ) with  $k = 4$  and  $n = 20$ .

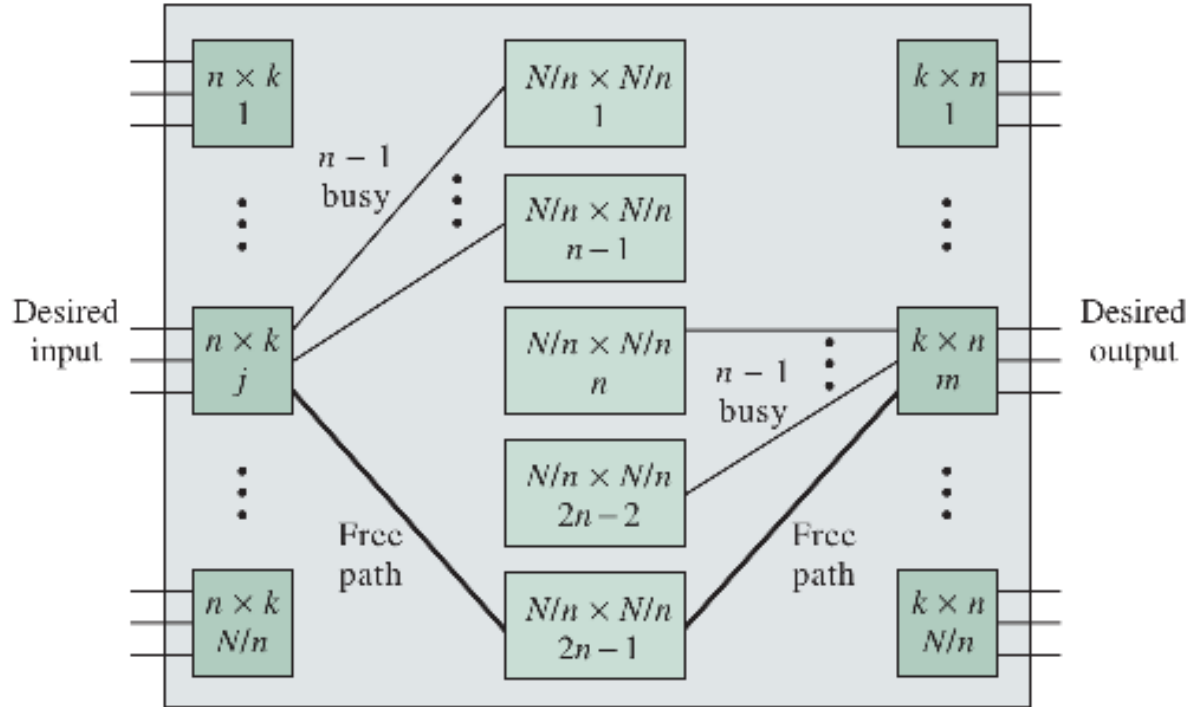
### Solution

In the first stage we have  $N/n$  or 10 crossbars, each of size  $20 \times 4$ . In the second stage, we have 4 crossbars, each of size  $10 \times 10$ . In the third stage, we have 10 crossbars, each of size  $4 \times 20$ . The total number of crosspoints is

$$2kN + k(N/n)^2 = 2000 \text{ crosspoints.}$$

This is 5 percent of the number of crosspoints in a single-stage switch ( $200 \times 200 = 40,000$ ).

# Clos Criterion for Non-blocking Switch



To avoid blocking:  $k \geq 2n-1$

# Clos Criterion for Non-blocking Switch

---

Total number of cross points:  $2Nk + k(N/n)^2$

With non-blocking criterion:  $k = 2n - 1 \implies 2N(2n - 1) + (2n - 1)(N/n)^2$

Min. number of  $n$  under non-blocking criterion:  $n \approx (N/2)^{\frac{1}{2}}$

Min. number of non-blocking crosspoints:  $4N([2N]^{\frac{1}{2}} - 1)$



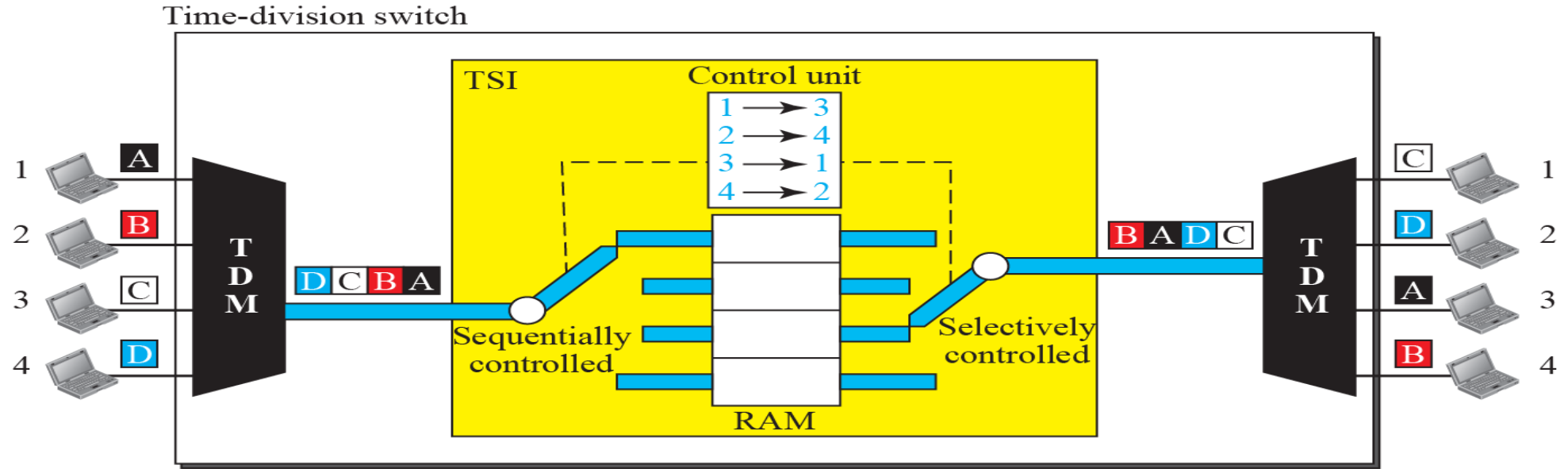
## Example

Redesign the previous three-stage,  $200 \times 200$  switch, using the Clos criteria with a minimum number of crosspoints.

### Solution

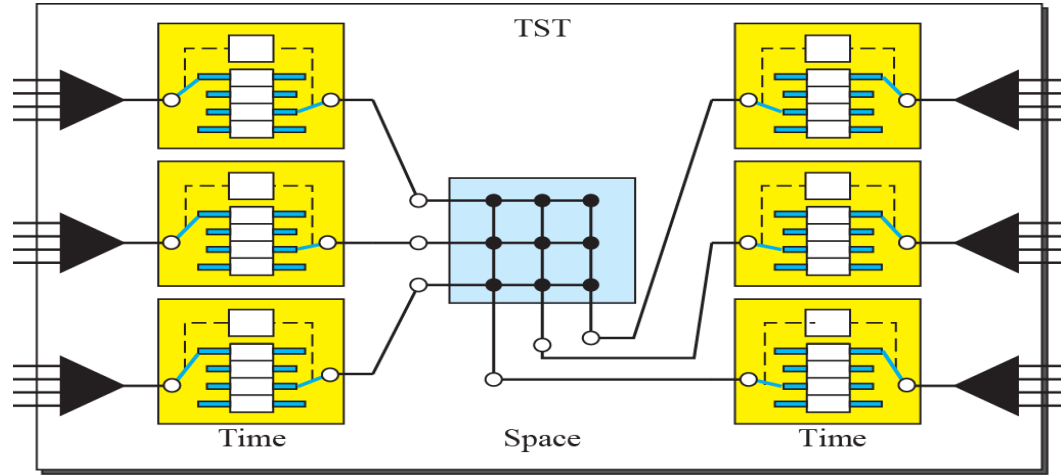
We let  $n = (200/2)^{1/2}$ , or  $n = 10$ . We calculate  $k = 2n - 1 = 19$ . In the first stage, we have  $200/10$ , or 20, crossbars, each with  $10 \times 19$  crosspoints. In the second stage, we have 19 crossbars, each with  $10 \times 10$  crosspoints. In the third stage, we have 20 crossbars each with  $19 \times 10$  crosspoints. The total number of crosspoints is  $20(10 \times 19) + 19(10 \times 10) + 20(19 \times 10) = 9500$ . If we use a single-stage switch, we need  $200 \times 200 = 40,000$  crosspoints. The number of crosspoints in this switch is 24 percent that of a single-stage switch.

# Time-slot Interchange

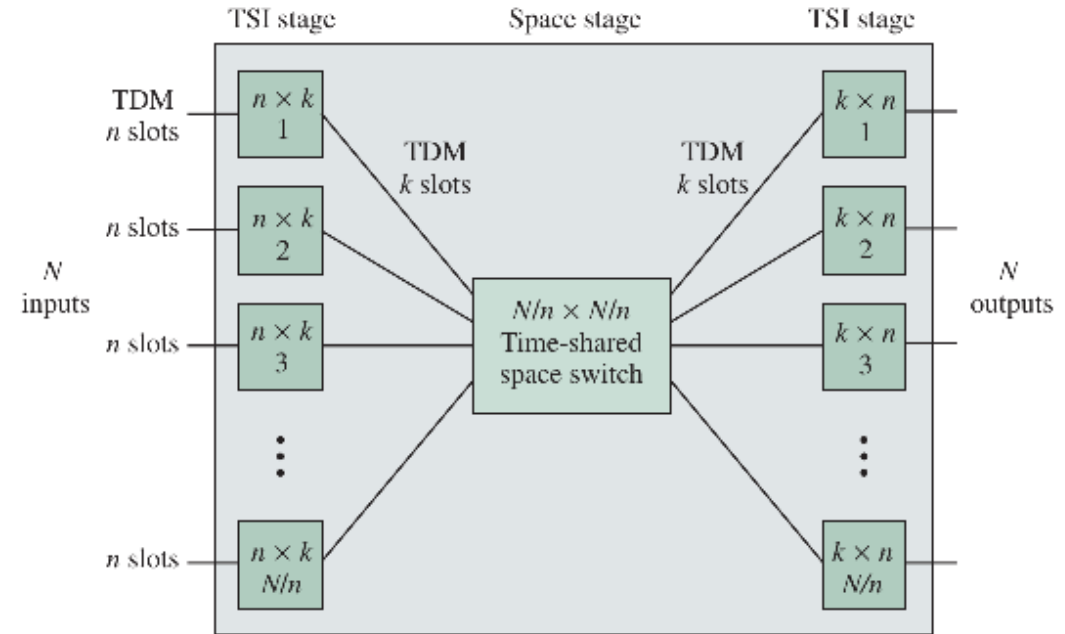
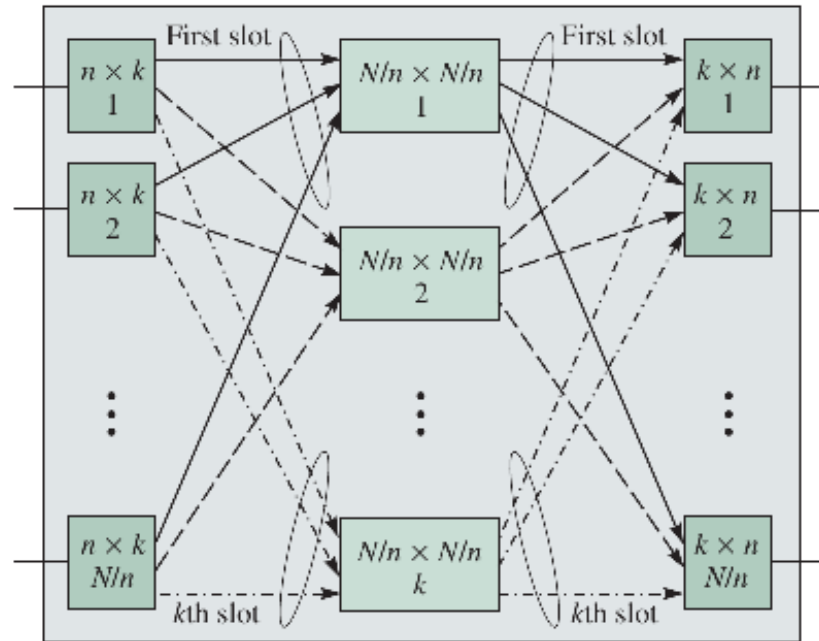


# Time-Space-Time (Hybrid) Switch

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# Time Sharing of Crossbar Switch in TST

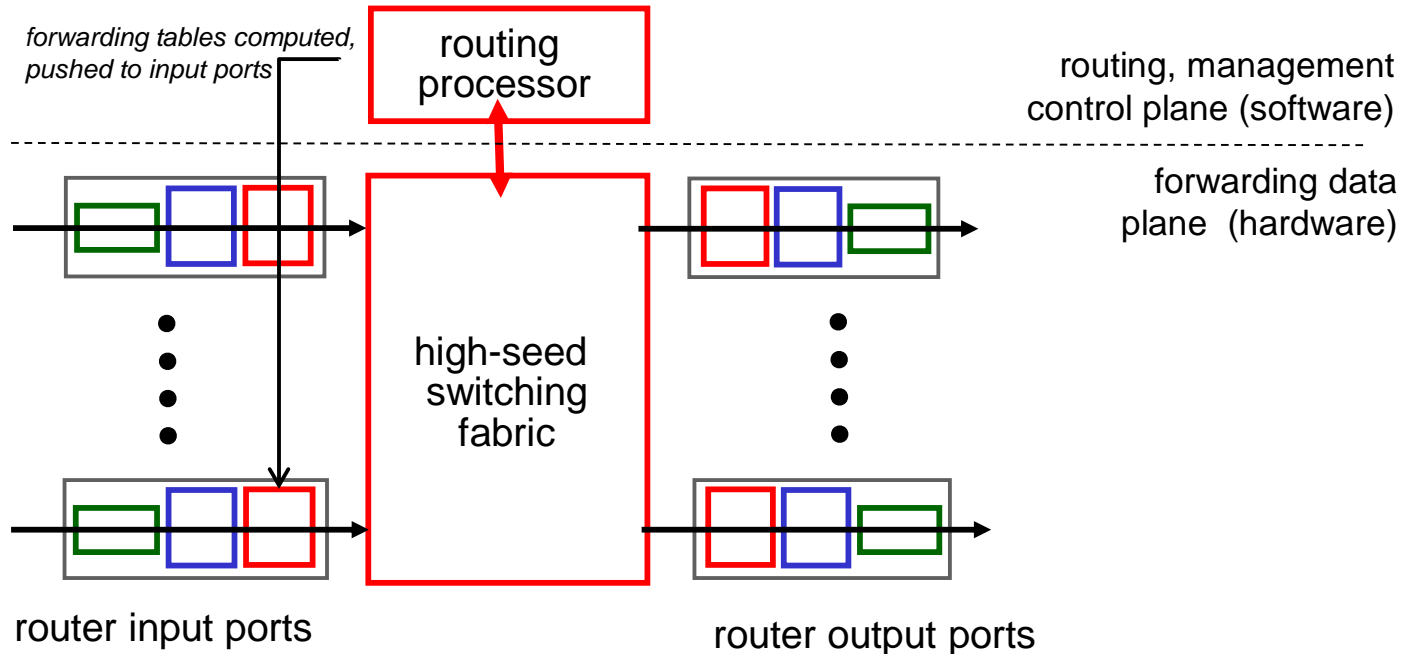


At any time-slot, only one intermediate switch is active

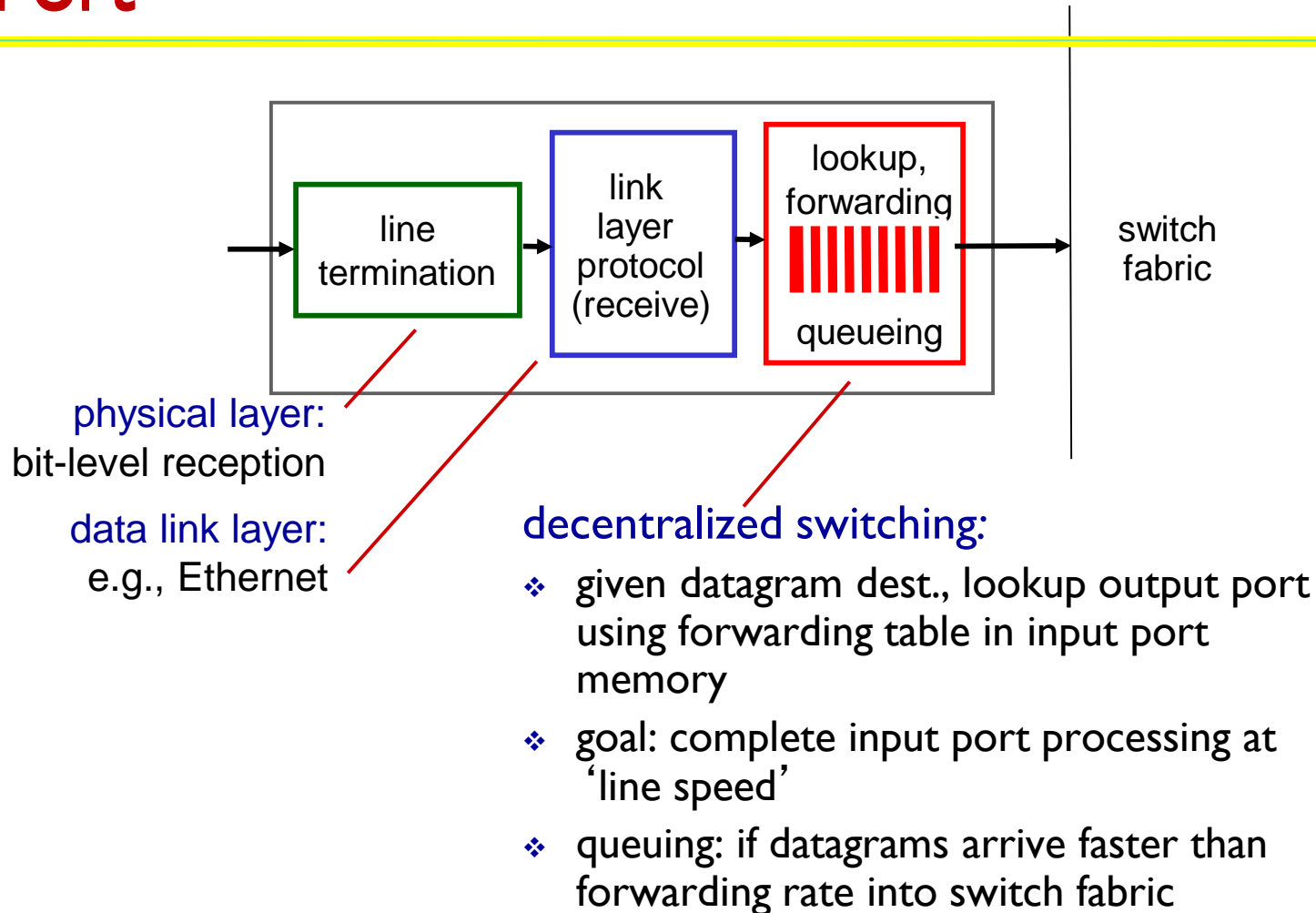
# Structure of Packet Switches

four components: input ports, output ports, the routing processor, and the switching fabric.

- ❖ run routing algorithms/protocol
- ❖ *forwarding* datagrams from incoming to outgoing link

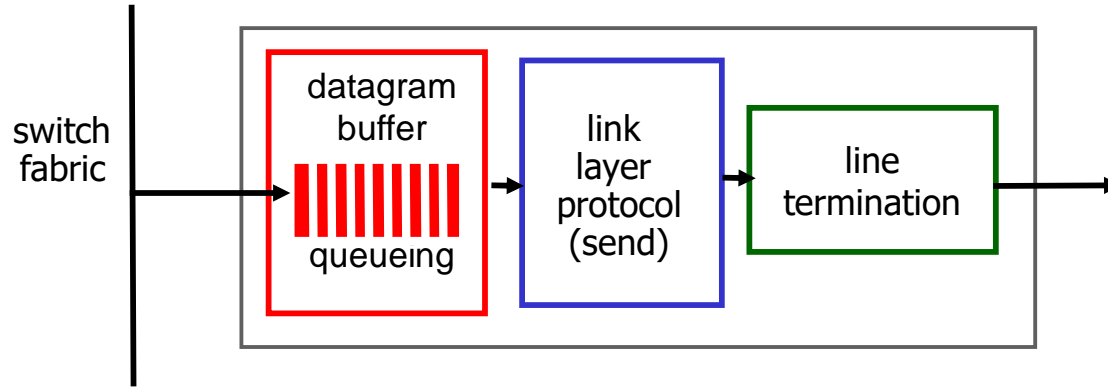


# Input Port



# Output Port

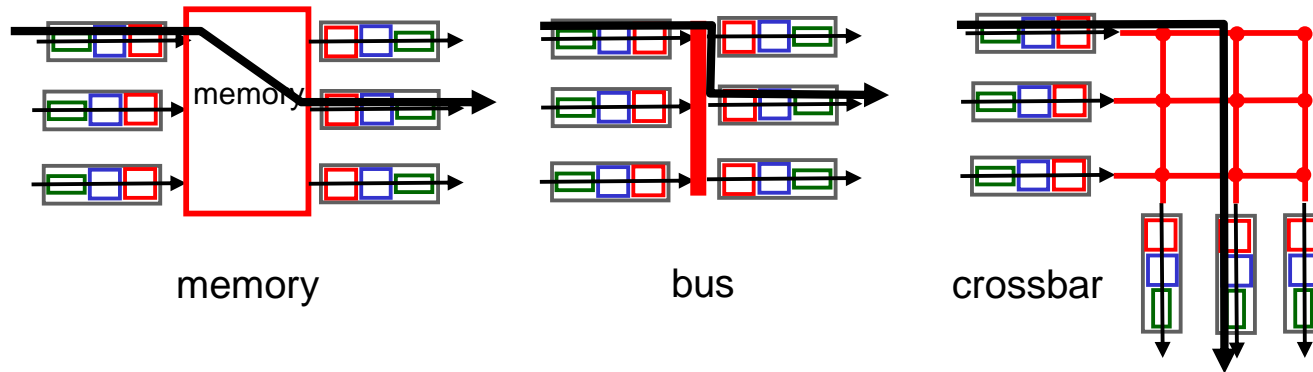
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- ❖ *buffering* required from fabric faster rate  
Datagram (packets) can be lost due to congestion, lack of buffers
- ❖ *scheduling* datagrams  
Priority scheduling – who gets best performance, network neutrality

# Switching Fabric

- transfer packet from input buffer to appropriate output buffer
- switching rate: rate at which packets can be transfer from inputs to outputs
  - often measured as multiple of input/output line rate
  - N inputs: switching rate N times line rate desirable
- three types of switching fabrics

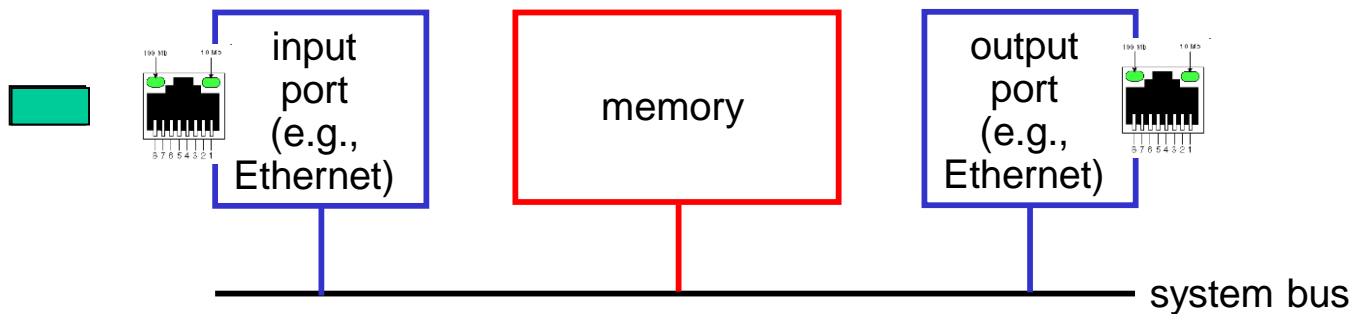




# Switching via Memory

## *first generation routers:*

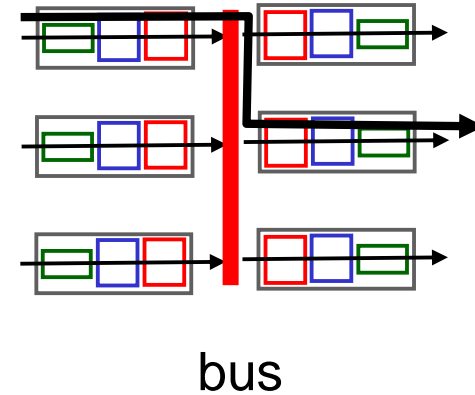
- traditional computers with switching under direct control of CPU
- packet copied to system's memory
- speed limited by memory bandwidth (2 bus crossings per datagram)



# Switching via a Bus

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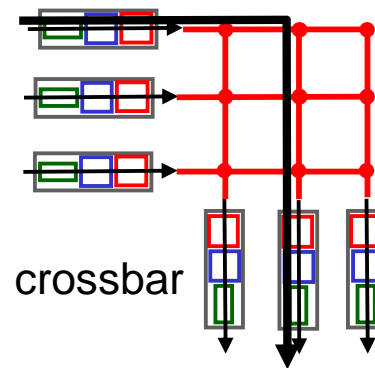
- datagram from input port memory to output port memory via a shared bus
- *bus contention*: switching speed limited by bus bandwidth
- 32 Gbps bus, Cisco 5600: sufficient speed for access and enterprise routers



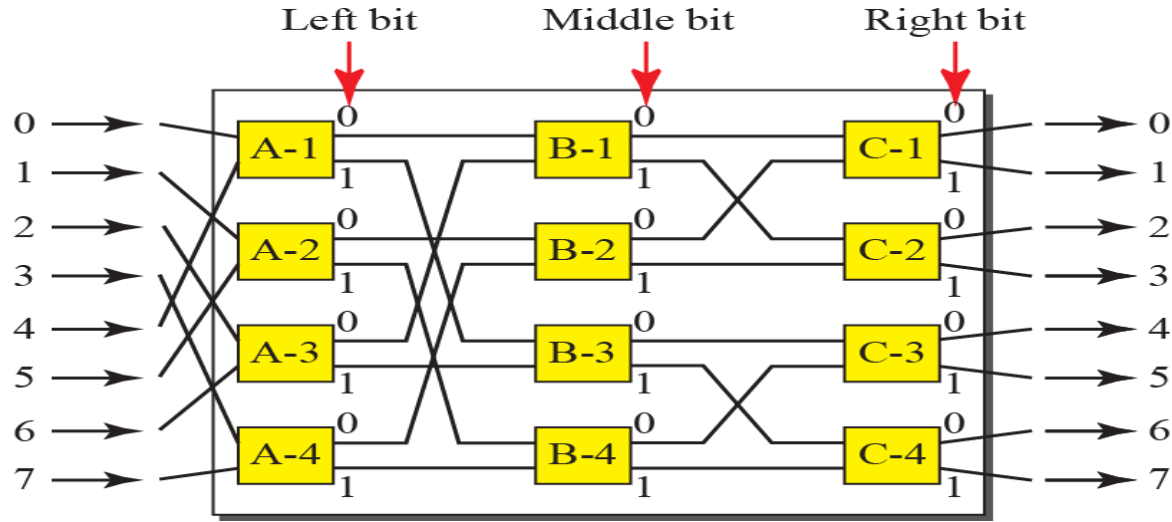
# Switching via Interconnection Network

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- overcome bus bandwidth limitations
- banyan networks, crossbar, other interconnection nets initially developed to connect processors in multiprocessor
- advanced design: fragmenting datagram into fixed length cells, switch cells through the fabric.
- Cisco I2000: switches 60 Gbps through the interconnection network

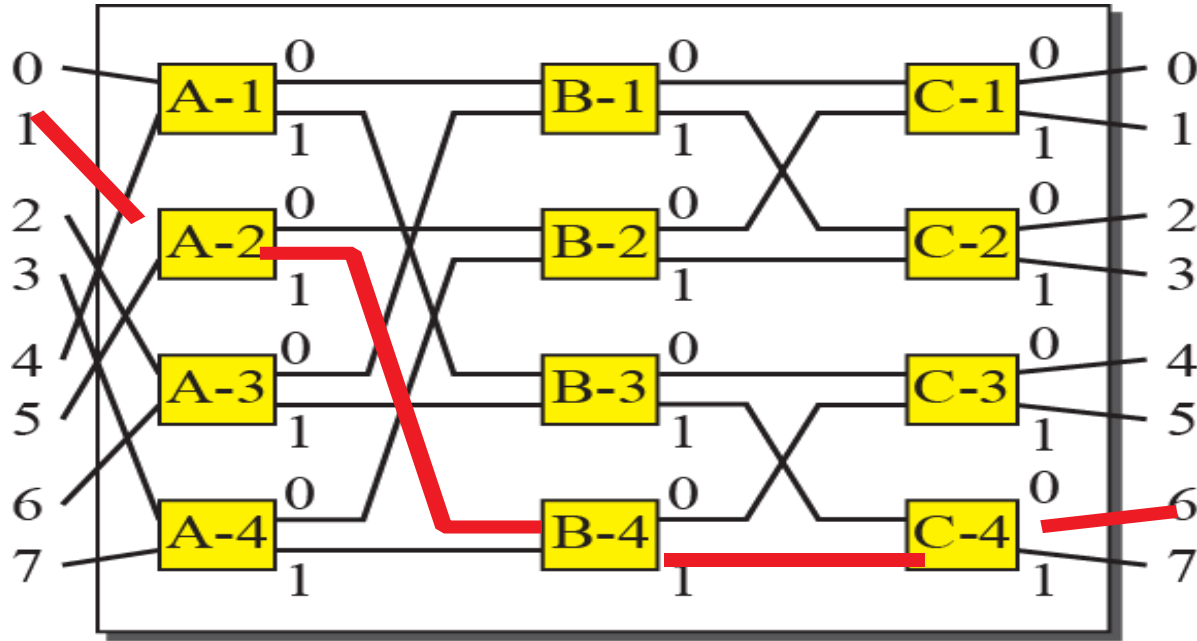


# A Banyan Switch



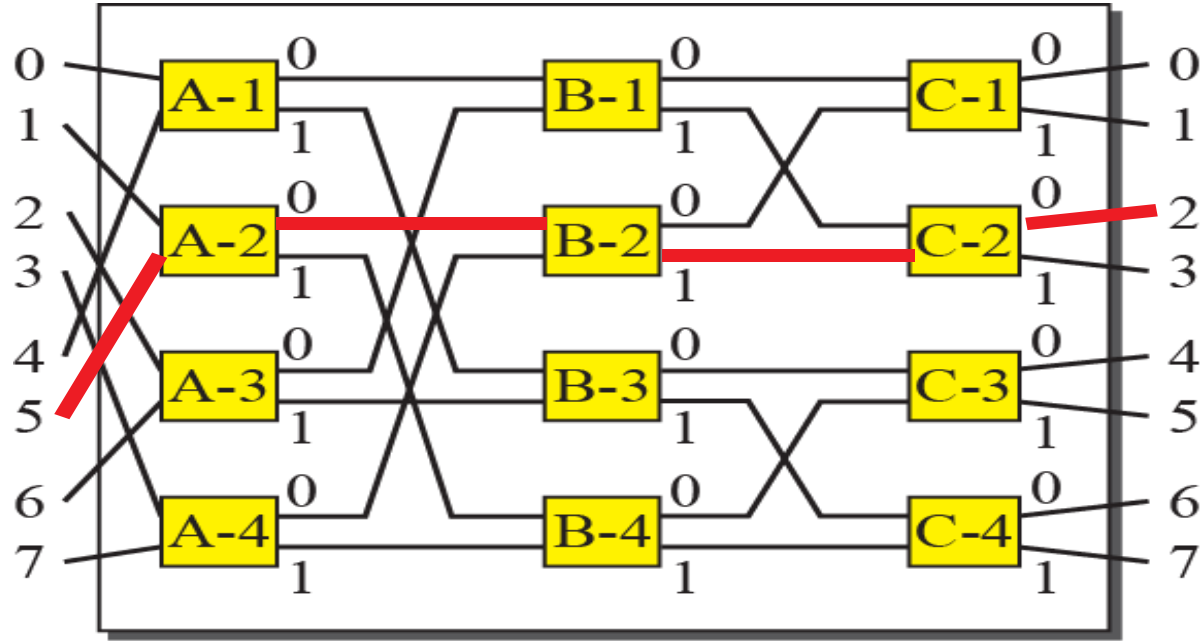
- A banyan switch is a multistage switch with microswitches at each stage that route the packets based on the output port represented as a binary string.
- For  $n$  inputs and  $n$  outputs, we have  $\log_2(n)$  stages with  $n/2$  microswitches at each stage.

# Example of Routing in a Banyan Switch



a. Input 1 sending a cell to output 6 (110)

# Example of Routing in a Banyan Switch



b. Input 5 sending a cell to output 2 (010)

# *Batcher-Banyan Switch*

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