

COMPUTER NETWORKS LAB (CS315)

Assignment-13

TLS

Date: 28 March 2024



Indian Institute of Technology Dharwad
भारतीय प्रौद्योगिकी संस्थान धारवाड़

Transport-layer security (TLS)

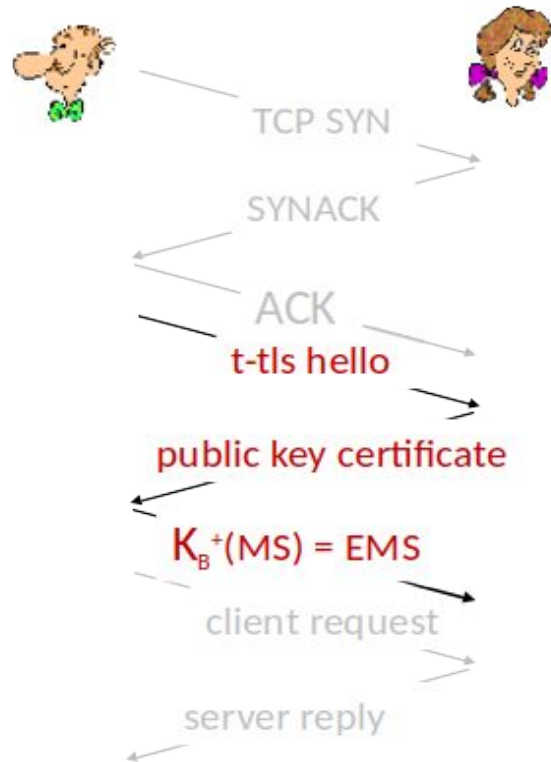
- widely deployed security protocol above the transport layer
 - supported by almost all browsers, web servers: https (port 443)
- provides:
 - **confidentiality**: via *symmetric encryption*
 - **integrity**: via *cryptographic hashing*
 - **authentication**: via *public key cryptography*

} all techniques we have studied!
- history:
 - early research, implementation: secure network programming, secure sockets
 - secure socket layer (SSL) deprecated [2015]
 - TLS 1.3: RFC 8846 [2018]

Transport-layer security: what's needed?

- let's *build* a toy TLS protocol, *t-tls*, to see what's needed!
- we've seen the “pieces” already:
 - **handshake**: Alice, Bob use their certificates, private keys to authenticate each other, exchange or create shared secret
 - **key derivation**: Alice, Bob use shared secret to derive set of keys
 - **data transfer**: stream data transfer: data as a series of records
 - not just one-time transactions
 - **connection closure**: special messages to securely close connection

t-tls: initial handshake



t-tls handshake phase:

- Bob establishes TCP connection with Alice
- Bob verifies that Alice is really Alice
- Bob sends Alice a master secret key (MS), used to generate all other keys for TLS session
- potential issues:
 - 3 RTT before client can start receiving data (including TCP handshake)

t-tls: cryptographic keys

- considered bad to use same key for more than one cryptographic function
 - different keys for message authentication code (MAC) and encryption
- four keys:
 - 🔑 K_c : encryption key for data sent from client to server
 - 🔑 M_c : MAC key for data sent from client to server
 - 🔑 K_s : encryption key for data sent from server to client
 - 🔑 M_s : MAC key for data sent from server to client
- keys derived from key derivation function (KDF)
 - takes master secret and (possibly) some additional random data to create new keys

t-tls: encrypting data

- recall: TCP provides data *byte stream* abstraction
- Q: can we encrypt data in-stream as written into TCP socket?
 - A: where would MAC go? If at end, no message integrity until all data received and connection closed!
 - solution: break stream in series of “records”
 - each client-to-server record carries a MAC, created using M_c
 - receiver can act on each record as it arrives
- t-tls record encrypted using symmetric key, K_c , passed to TCP:

$$K_c \left(\begin{array}{|c|c|c|} \hline \text{length} & \text{data} & \text{MAC} \\ \hline \end{array} \right)$$

t-tls: encrypting data (more)

- possible attacks on data stream?
 - *re-ordering*: man-in middle intercepts TCP segments and reorders (manipulating sequence #s in unencrypted TCP header)
 - *replay*
- solutions:
 - use TLS sequence numbers (data, TLS-seq-# incorporated into MAC)
 - use nonce

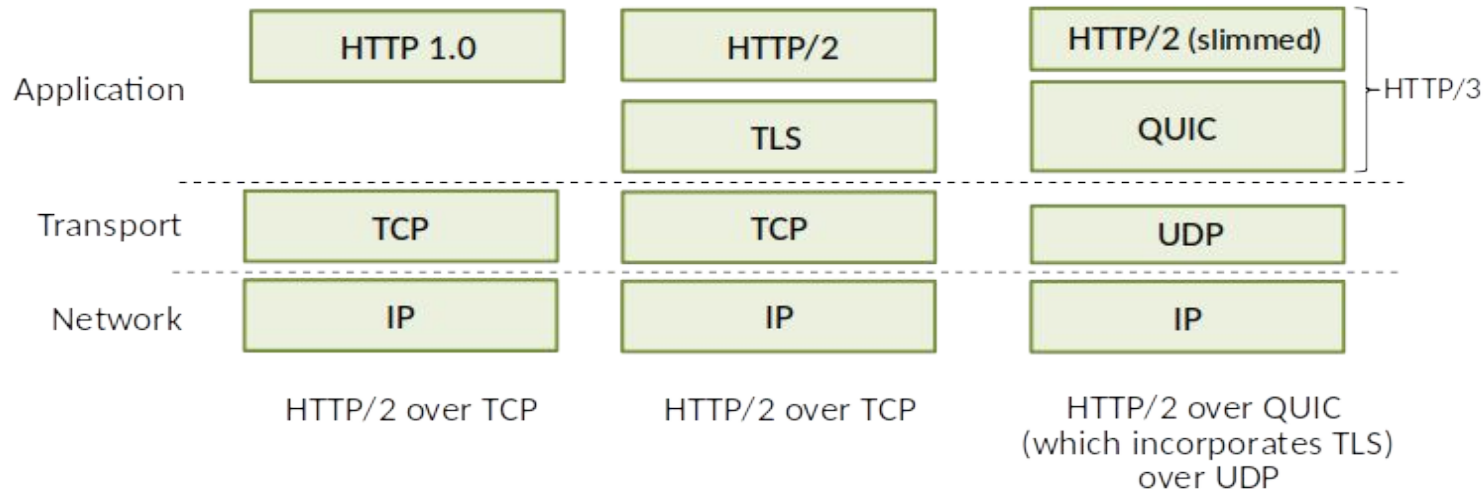
t-tls: connection close

- truncation attack:
 - attacker forges TCP connection close segment
 - one or both sides thinks there is less data than there actually is
- **solution:** record types, with one type for closure
 - type 0 for data; type 1 for close
- MAC now computed using data, type, sequence #

$$K_c \left(\begin{array}{|c|c|c|c|} \hline \text{length} & \text{type} & \text{data} & \text{MAC} \\ \hline \end{array} \right)$$

Transport-layer security (TLS)

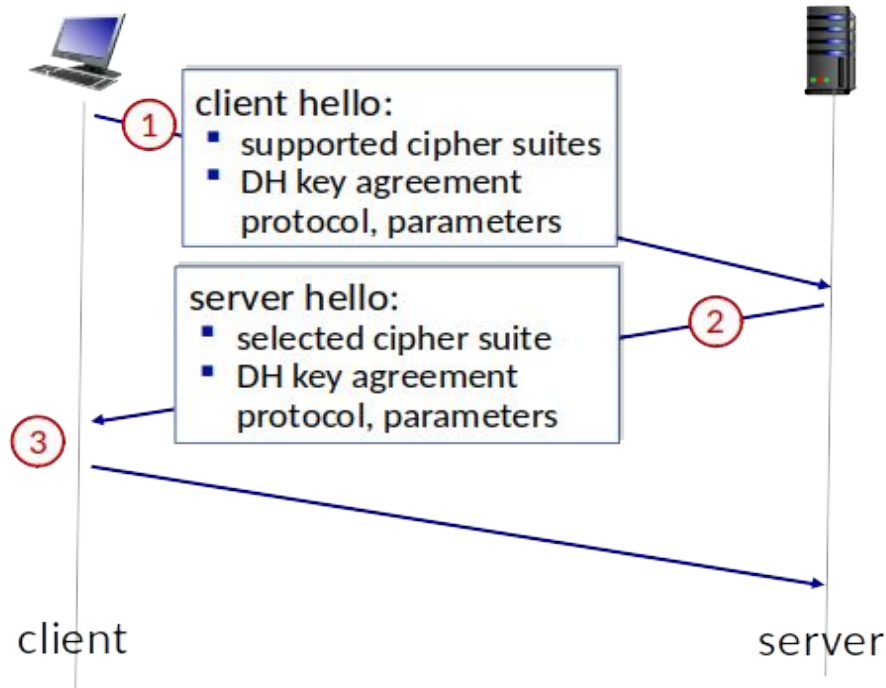
- TLS provides an API that *any* application can use
- an HTTP view of TLS:



TLS: 1.3 cipher suite

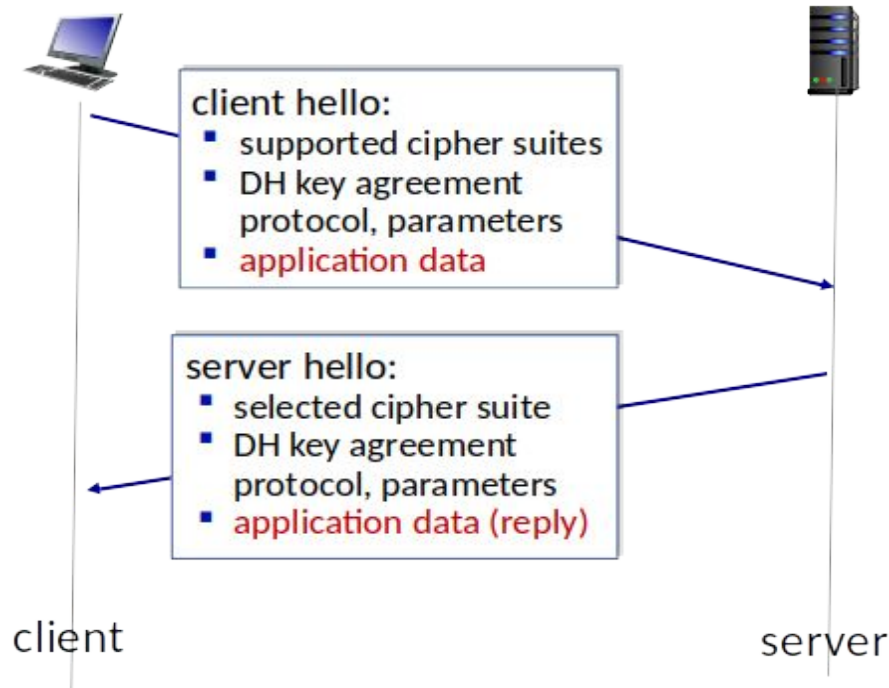
- “cipher suite”: algorithms that can be used for key generation, encryption, MAC, digital signature
- TLS: 1.3 (2018): more limited cipher suite choice than TLS 1.2 (2008)
 - only 5 choices, rather than 37 choices
 - *requires* Diffie-Hellman (DH) for key exchange, rather than DH or RSA
 - combined encryption and authentication algorithm (“authenticated encryption”) for data rather than serial encryption, authentication
 - 4 based on AES
 - HMAC uses SHA (256 or 284) cryptographic hash function

TLS 1.3 handshake: 1 RTT



- 1** client TLS hello msg:
 - *guesses* key agreement protocol, parameters
 - indicates cipher suites it supports
- 2** server TLS hello msg chooses
 - key agreement protocol, parameters
 - cipher suite
 - server-signed certificate
- 3** client:
 - checks server certificate
 - generates key
 - can now make application request (e.g., HTTPS GET)

TLS 1.3 handshake: 0 RTT



- initial hello message contains encrypted application data!
 - “resuming” earlier connection between client and server
 - application data encrypted using “resumption master secret” from earlier connection
- vulnerable to replay attacks!
 - maybe OK for get HTTP GET or client requests not modifying server state

Thank you