

# DAA C2 Assignment-1

## Group Members –

IIT2019108 - Harsha Vardhan Madasi

IIT2019109 - Rahul kumar

IIT2019110 - Sumit Katiyar

21 March 2021

## **Question -**

Given an array of  $2n$  elements in the following format  $a_1, a_2, a_3, a_4, \dots, a_n, b_1, b_2, b_3, b_4, \dots, b_n$  .

The task is shuffle the array to  $a_1, b_1, a_2, b_2, a_3, b_3, \dots, a_n, b_n$  without using extra space.

# Example -

INPUT :

8

1 3 5 7 2 4 6 8

OUTPUT:

1st iteration: 1 3 2 4 5 7 6 8

2nd iteration: 1 2 3 4

3rd iteration: 5 6 7 8

Final Shuffled Array : 1 2 3 4 5 6 7 8

# Some Definitions:-

**Divide and Conquer Algorithm:-** A divide-and-conquer algorithm recursively breaks down a problem into two or more sub-problems of the same or related type, until these become simple enough to be solved directly. The solutions to the sub-problems are then combined to give a solution to the original problem.

# Algorithm Description:-

Here we are using Divide and Conquer Technique.

Firstly we will divide the given array into half and swap second half element of arr1[] with first half element of arr2[]. Recursively do this for both the arrays arr1 and arr2.

## Algorithmic Steps:-

- Let us consider an array be {a1, a2, a3, a4, b1, b2, b3, b4} .
- Now split the array into two halves i.e., {a1, a2, a3, a4} ; {b1, b2, b3, b4}.
- Now we will swap elements which are present around the center i.e., we will swap a3, a4 with b1, b2 correspondingly.
- After swapping we will get: a1, a2, b1, b2, a3, a4, b3, b4.
- Now we will recursively split the above subarrays and swap the elements around the center for each subarray.
- Finally we will get a1, b1, a2, b2 and a3, b3, a4, b4.
- Now final shuffled array is {a1, b1, a2, b2, a3, b3, a4, b4}

## Pseudo Code:-

```
//Function to shuffle an array  
void shufflearray(int arr[],int start , int end)  
// base condition  
if(end < start )then  
return ;  
//if only two elements present in the subarray  
if(end = start + 1)then  
return ;
```

## Pseudo Code(contd.):

//Finding middle of the array to divide the array

int mid  $\leftarrow$  (start + end) / 2;

//using temp in order to swap first half of second array

int temp  $\leftarrow$  mid + 1;

//using first in order to swap second half of first array

int firstmid  $\leftarrow$  (start + mid) / 2;

//Swapping the center elements

for(i = firstmid + 1 to i  $\leq$  mid )



## Pseudo Code(contd.):-

```
swap (arr[i],arr[temp++]);
```

```
i ← i+1;
```

```
// recursively calling the function for first and second subarrays
```

```
shuflearray(arr, start, mid);
```

```
shuflearray(arr, mid + 1, end);
```

```
//MAIN FUNCTION
```

```
int main()
```

```
    int n =1+(rand())%10;
```

```
    cout <<n<<"\n";
```

## Psuedo Code(contd.):

```
int length=pow(2,n);
int arr[length];
for(int i=0;i<length;i++)
    arr[i]=(rand())%100;
for (int i = 0; i <length; i++)
    cout << arr[i] << " ";
cout <<"\n";
shufleArray(arr, 0, length - 1);

for (int i = 0; i < length; i++)
    cout << arr[i] << " ";

return 0;
```

## Time Complexity:-

In the above algorithm we get the recurrence

$$T(n)=2T(n/2)+O(n)$$

for the time complexity, which results in

$$T(n)=O(n\log n)$$

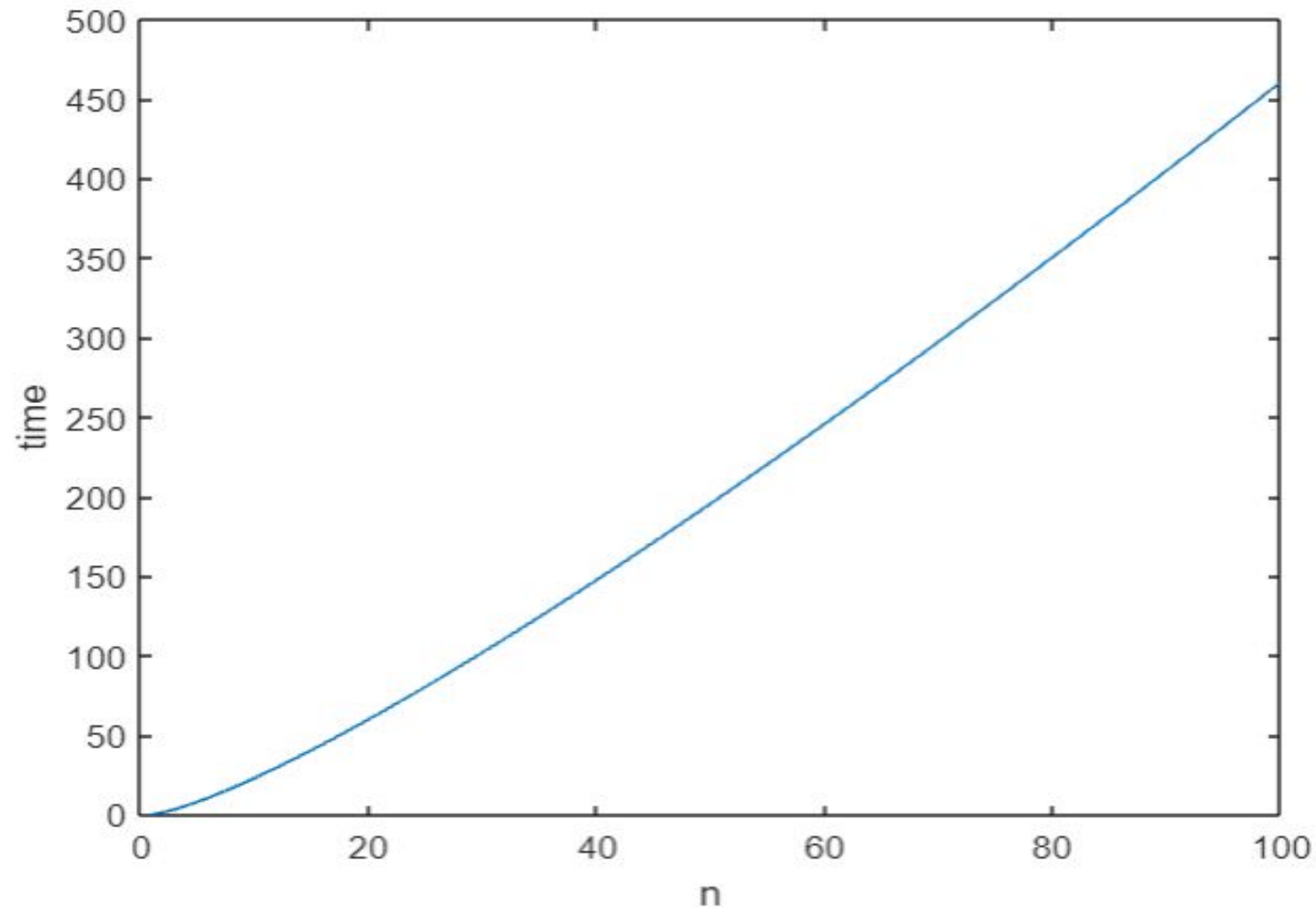
So the final computational time to shuffle the given array is  $O(n\log n)$ .

## **Space Complexity:-**

To be precise , as we are swapping the elements in the given array itself there is no extra space required.

So the space complexity will be  $O(1)$ .

# Time Complexity Graph( $O(n \log n)$ ):-



## **Conclusion:-**

Hence, from the above graph and analysis, we come to the conclusion that this problem can be solved using the time complexity of  $O(n \log n)$  and space complexity of  $O(1)$ .

## **References:-**

- 1) Introduction to Algorithms by Cormen, Charles, Rivest and Stein.
- 2) Geeks for Geeks.

***Thank You***