

Q) what is a distributed system?

→ a system where I can't get my work as somewhere any computer has crashed which I never heard of (Joke)

by father of DS

→ It's about failure.

Martin Kleppmann's definition author of Data Intensive Applications

→ running on several nodes connected by a network

→ characterized by partial failure.

↳ It means → some parts of the system might be broken while other parts of it are okay

and furthermore it's often impossible (at least impossible for the software) to know what exactly is broken.

"an instance of some software running on a computer in order to accomplish some task"

Partial failures can be:

- Either a computer failed
- or Connection btw computers.  
broke / failed

This is different from the kind of  
failure that we are often  
used to thinking

or a single machine if part of OS fails.  
then all fails. and there's no way to  
operate.

but if you have many machines then the  
point of keep operating. When an  
failure is very imp

(restarting all machines is  
not a good option)

"The more machines you have the  
more inevitable that some  
systems will fail"



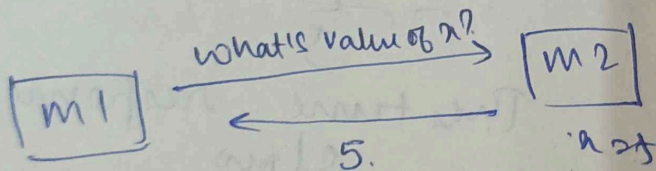
# Cloud Computing v/s HPC Computing

→ working around partial failures.

→ treats partial failures as total failures.

→ checkpoints involved for long run.

Case 1. We have two machines.



Is simple req-response model mai kya kya dikkate/failures aa sakte h?

(i) Req from  $M_1$  gets lost

(ii) Req from  $M_1$  is slow ( $M_2$  khamcha ki chiz wait hozi kar sakti h)

Reason can be "congestion" or "maybe a misbehaving server"

~~(i)  $M_1$~~

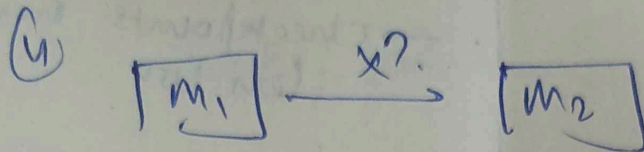
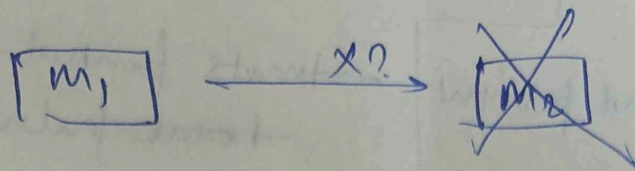
$M_1$

$x$

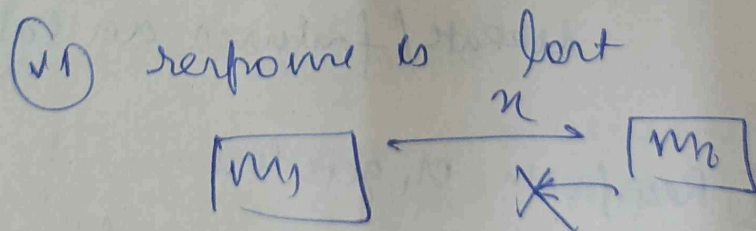
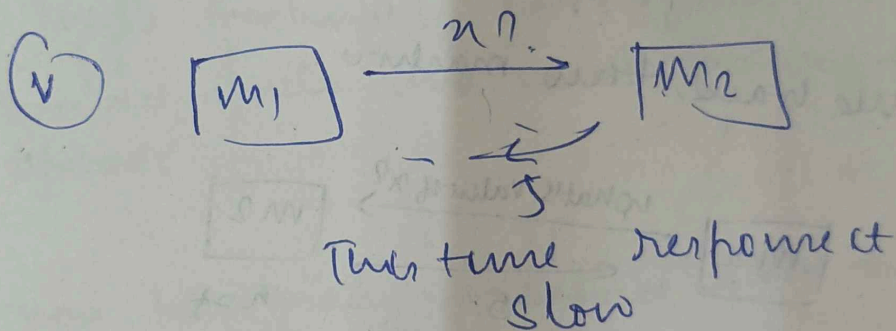
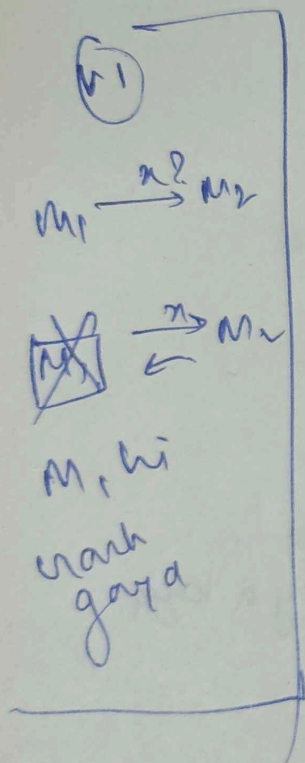
→

$M_2$

③  $m_2$  crashed



$m_1$  sends, but  $m_2$  is slower than  $m_1$  expects, so  $m_1$  assumes I would not get a response.



All these conditions are different but they are same from  $m_1$  perspective.

$m_1$  ko kabhi pata hi chalega ki response byu ni rahi ya nahi  
 hiye saare situations same hi h.



"So, if you send a req to another node and you don't receive a response it is impossible to know why"  
(without global knowledge of the system)

So, how do real systems deal with this type of situations?

If one machine wants to send a message to another machine and it doesn't know for sure if it's going to get a response.

So what one should do when sending a req and expecting a response.

In reality  $M_1$  should have some sort of timeout.

Ye kya h?

↓  
after sometime when  $M_1$  has send req and if there is no response then assume failure.

It is a one of the best approach in reality.

P.T.O

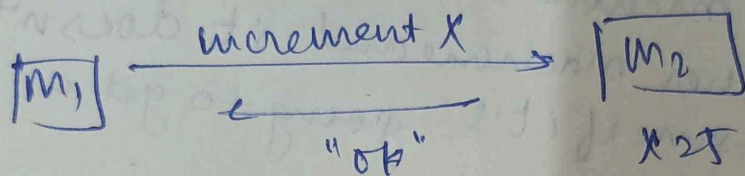
But why it might be a mistake to assume failure

↳ Abhi to dekh sakte

M1 n bya h whi puch raha tha,

par aise situation ka kya jaha.

M1 ka req M2 ko ~~effect~~ affect kar raha



x = 25

x++

x = 26

Agar M2 ok ni bhy paya, to ya timeout se phle w naya

ye mana kr n++ hua

galat aur nasa mana bhi risky h.

or inevitable

and this sort of uncertainty is fundamental characteristic of distributed system.



$$\text{Distributed System} = \text{partial failure} + \text{unbounded latency}$$

So in Distributed System we have to deal with partial failures and unbounded latency and that is what makes Distributed systems so hard.

By why you like a system like this that is clearly very terrible, hard to debug, to build, why would you want it?

→ data too big to fit on one machine.

→ You want things to be faster. So you use more computers.

→ for redundant ~~for~~ copies to fail independently.

# Time and clocks

what are clocks for?

① mark points in time.

"This class starts at 9:20 am"

② duration or intervals of time.

"This class is 65 minutes long".

## time-of-day clocks

- Sync'd with  
NTP (Network  
time protocol)

- bad for  
measuring  
duration &  
intervals.

because here  
time can jump  
forward &  
backward.

- Don't use for  
implementing  
timeouts

- ok for timestamps  
but not great

## monotonic clocks

- only go forward

- bad for timestamping

- good for duration,  
intervals &  
timeouts



lec 08 1: CAPS

4e do no physical clocks h. (Peechle page wall)

Physical clocks.	logical clocks.
Time of day, monotonic clocks.	only measure <u>ordering of events</u>  which event happened before another.

Suppose A happened before B

↳  $A \rightarrow B$

Inference / Causality -

- 'A' could have caused 'B'

- 'B' could not have caused 'A'

"My style of teaching might be  
influenced by Sir teaching but  
Sir teaching was not influenced  
by my teaching" not sure  
afterwards