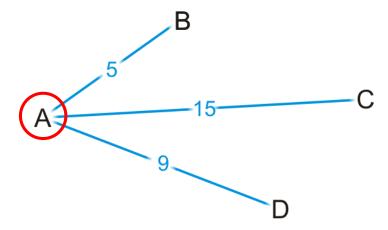
Dijkstra's algorithm

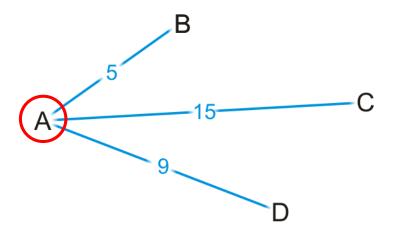
Suppose you are at vertex A

- You are aware of all vertices adjacent to it
- This information is either in an adjacency list or adjacency matrix

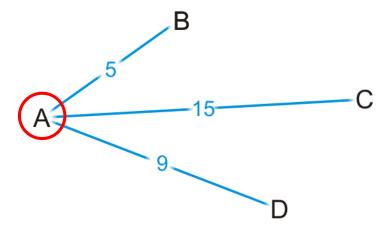


Is 5 the shortest distance to B via the edge (A, B)?

• Why or why not?

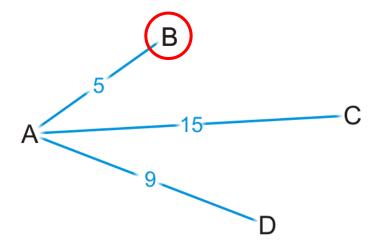


Are you guaranteed that the shortest path to C is (A, C), or that (A, D) is the shortest path to vertex D?



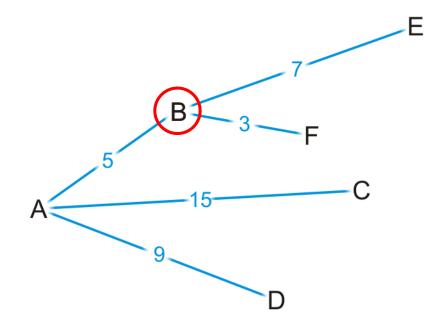
We accept that (A, B) is the shortest path to vertex B from A

Let's see where we can go from B



By some simple arithmetic, we can determine that

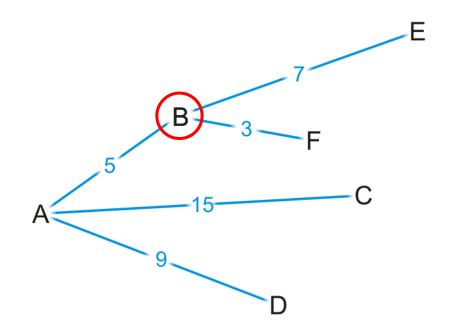
- There is a path (A, B, E) of length 5 + 7 = 12
- There is a path (A, B, F) of length 5 + 3 = 8



Is (A, B, F) is the shortest path from vertex A to F?

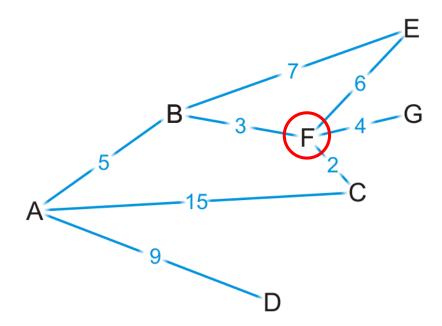
• Why or why not?

Are we guaranteed that any other path we are currently aware of is also going to be the shortest path?



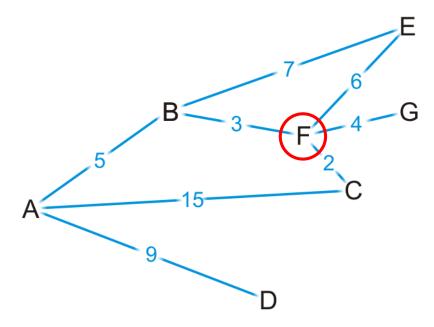
Okay, let's visit vertex F

• We know the shortest path is (A, B, F) and it's of length 8



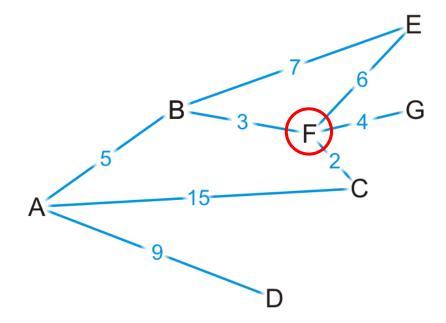
There are three edges exiting vertex F, so we have paths:

- (A, B, F, E) of length 8 + 6 = 14
- (A, B, F, G) of length 8 + 4 = 12
- (A, B, F, C) of length 8 + 2 = 10



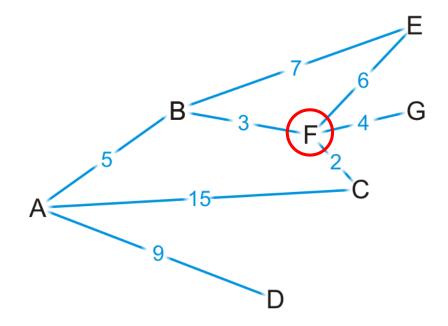
By observation:

- The path (A, B, F, E) is longer than (A, B, E)
- The path (A, B, F, C) is shorter than the path (A, C)



At this point, we've discovered the shortest paths to:

- Vertex B: (A, B) of length 5
- Vertex F: (A, B, F) of length 8



Dijkstra's algorithm

Dijkstra's algorithm solves the single-source shortest path problem

- It is very similar to Prim's algorithm
- Assumption: all the weights are positive

Like Prim's algorithm,

We initially don't know the distance to any vertex except the initial vertex We require an array of distances, all initialized to infinity except for the source vertex, which is initialized to 0

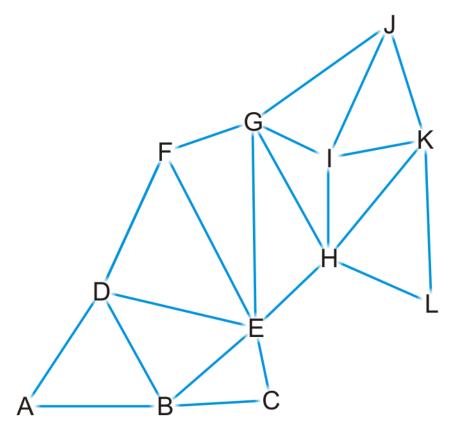
Each time we visit a vertex, we will examine all adjacent vertices We need to track visited vertices—a Boolean table of size |V| We need an array of previous vertices, all initialized to null

Dijkstra's algorithm

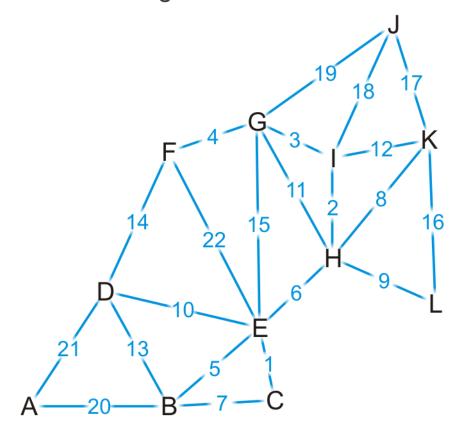
Thus, we will iterate |V| times:

- Find that unvisited vertex v that has a minimum distance to it
- Mark it as having been visited
- Consider every adjacent vertex w that is unvisited:
 - Is the distance to v plus the weight of the edge (v, w) less than our currently known shortest distance to w
 - \circ If so, update the shortest distance to w and record v as the previous pointer
- Continue iterating until all vertices are visited or all remaining vertices have a distance to them of infinity

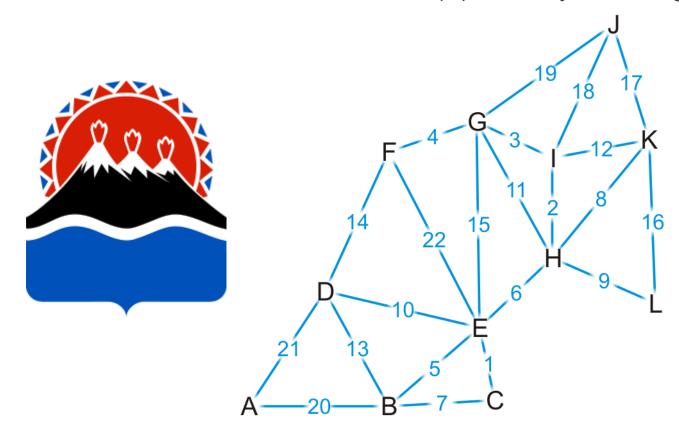
Here is our abstract representation



Let us give a weight to each of the edges

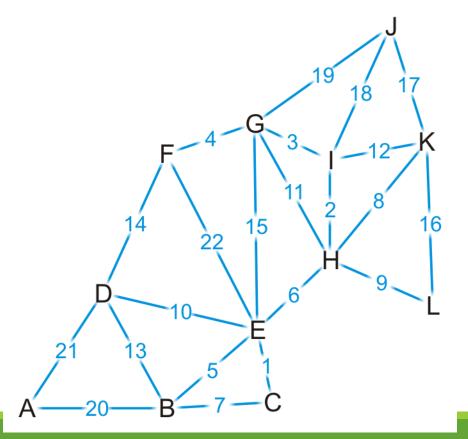


Find the shortest distance from Kamchatka (K) to every other region



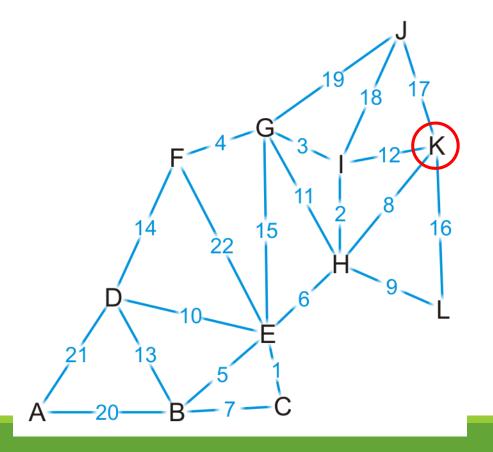
We set up our table

• Which unvisited vertex has the minimum distance to it?



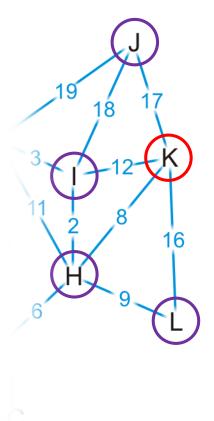
Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	F	∞	Ø
F	F	∞	Ø
G	F	∞	Ø
Н	F	∞	Ø
I	F	∞	Ø
J	F	∞	Ø
K	F	0	Ø
L	F	∞	Ø

We visit vertex K



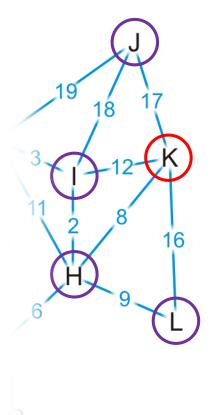
Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	F	∞	Ø
F	F	∞	Ø
G	F	∞	Ø
Н	F	∞	Ø
I	F	∞	Ø
J	F	∞	Ø
K	Т	0	Ø
L	F	∞	Ø

Vertex K has four neighbors: H, I, J and L



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	00	Ø
F	F	00	Ø
G	F	00	Ø
Н	F	∞	Ø
I	F	∞	Ø
J	F	∞	Ø
K	T	0	Ø
L	F	∞	Ø

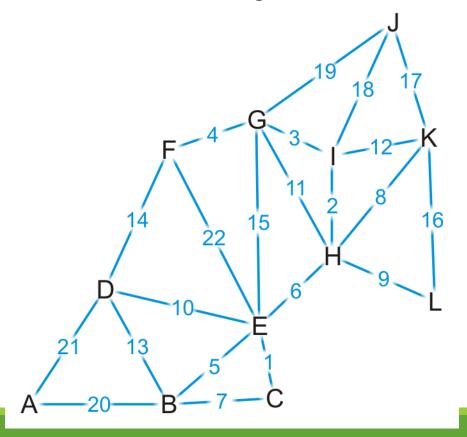
We have now found at least one path to each of these vertices



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	00	Ø
F	F	00	Ø
G	F	00	Ø
Н	F	8	K
I	F	12	K
J	F	17	K
K	T	0	Ø
L	F	16	K

We're finished with vertex K

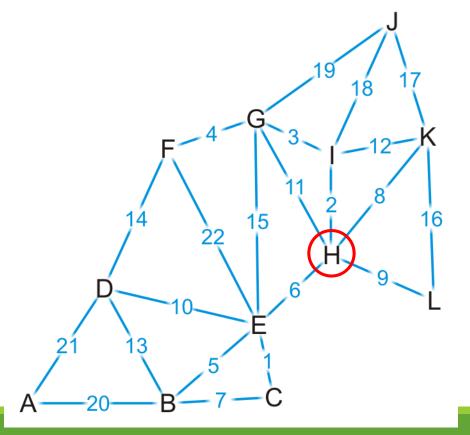
• To which vertex are we now guaranteed we have the shortest path?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	F	∞	Ø
F	F	∞	Ø
G	F	∞	Ø
Н	F	8	K
I	F	12	K
J	F	17	K
K	Т	0	Ø
L	F	16	K

We visit vertex H: the shortest path is (K, H) of length 8

Vertex H has four unvisited neighbors: E, G, I, L

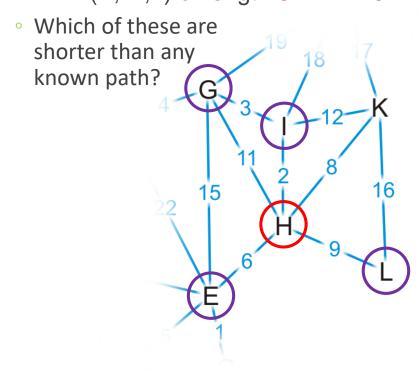


Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	F	∞	Ø
F	F	∞	Ø
G	F	∞	Ø
Н	T	8	K
I	F	12	K
J	F	17	K
K	Т	0	Ø
L	F	16	K

Consider these paths:

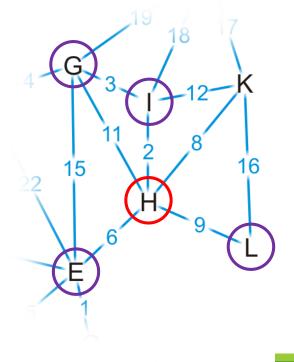
(K, H, E) of length 8 + 6 = 14 (K, H, I) of length 8 + 2 = 10

(K, H, G) of length 8 + 11 = 19(K, H, L) of length 8 + 9 = 17



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	∞	Ø
F	F	00	Ø
G	F	∞	Ø
Н	T	8	K
I	F	12	K
J	F	17	K
K	T	0	Ø
L	F	16	K

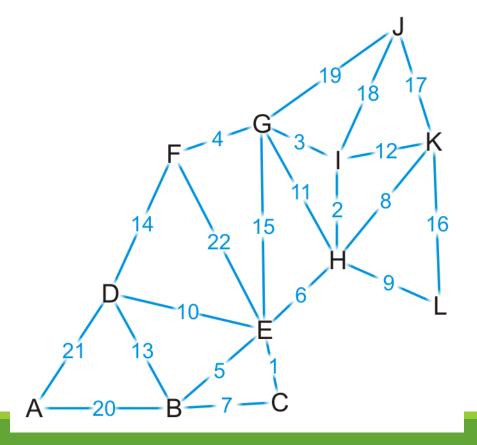
We already have a shorter path (K, L), but we update the other three



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	14	Н
F	F	00	Ø
G	F	19	Н
Н	T	8	K
I	F	10	Н
J	F	17	K
K	T	0	Ø
L	F	16	K

We are finished with vertex H

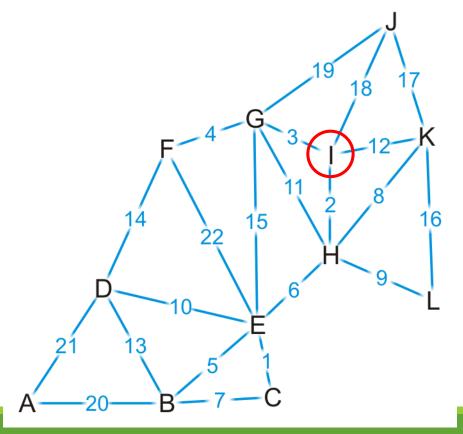
• Which vertex do we visit next?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	F	14	Н
F	F	∞	Ø
G	F	19	Н
Н	Т	8	K
I	F	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

The path (K, H, I) is the shortest path from K to I of length 10

Vertex I has two unvisited neighbors: G and J

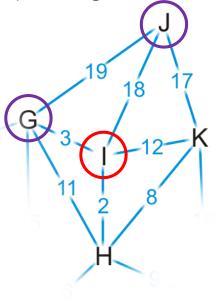


Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	F	14	Н
F	F	∞	Ø
G	F	19	Н
Н	Т	8	K
I	T	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

Consider these paths:

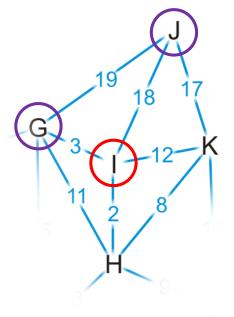
(K, H, I, G) of length 10 + 3 = 13

(K, H, I, J) of length 10 + 18 = 28



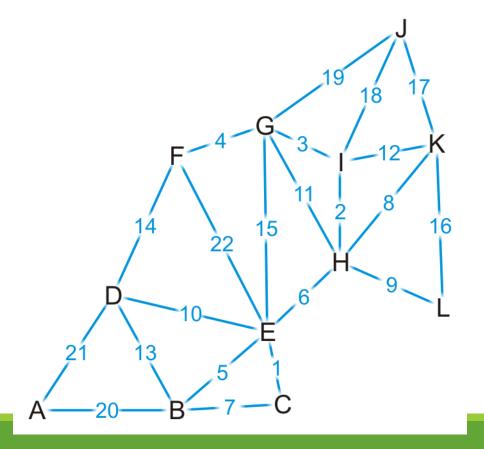
Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	14	Н
F	F	00	Ø
G	F	19	Н
Н	Т	8	K
I	T	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

We have discovered a shorter path to vertex G, but (K, J) is still the shortest known path to vertex J



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	14	Н
F	F	00	Ø
G	F	13	
Н	Т	8	K
1	T	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

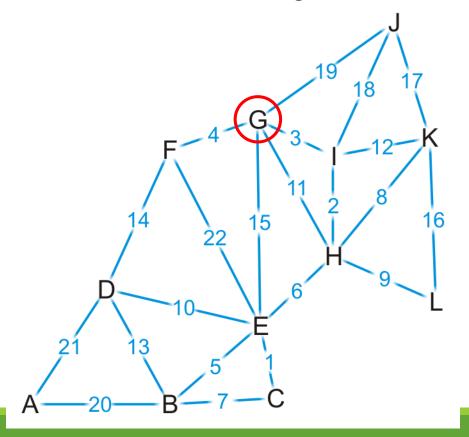
Which vertex can we visit next?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	F	14	Н
F	F	∞	Ø
G	F	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

The path (K, H, I, G) is the shortest path from K to G of length 13

Vertex G has three unvisited neighbors: E, F and J



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	F	14	Н
F	F	∞	Ø
G	T	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

Consider these paths:

(K, H, I, G, E) of length
$$13 + 15 = 28$$

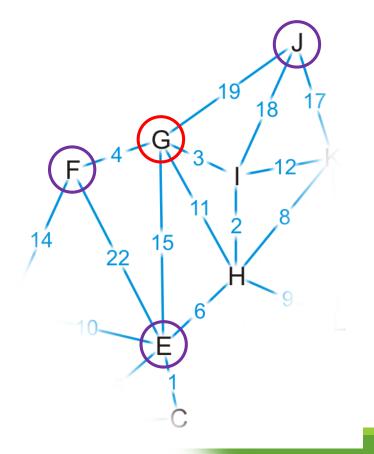
(K, H, I, G, J) of length $13 + 19 = 32$

(K, H, I, G, F) of length 13 + 4 = 17

• Which do we update?	19 18 17
F)-4	3 12-12-11 8
14 22	15 H 9
10	
	-c

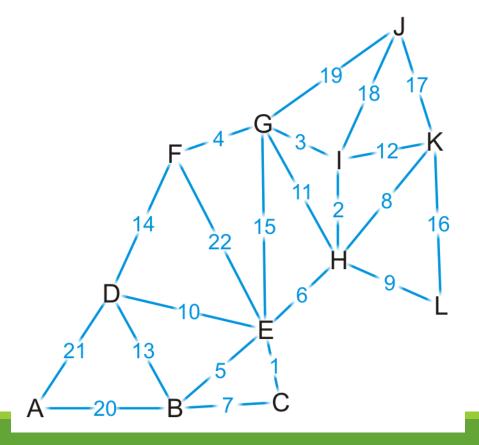
Vertex	Visited	Distance	Previous
Α	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	14	Н
F	F	∞	Ø
G	T	13	I
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

We have now found a path to vertex F



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	00	Ø
С	F	00	Ø
D	F	00	Ø
Е	F	14	Н
F	F	17	G
G	T	13	I
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

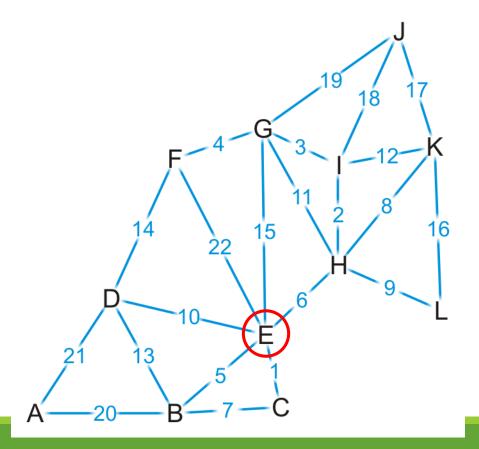
Where do we visit next?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	F	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

The path (K, H, E) is the shortest path from K to E of length 14

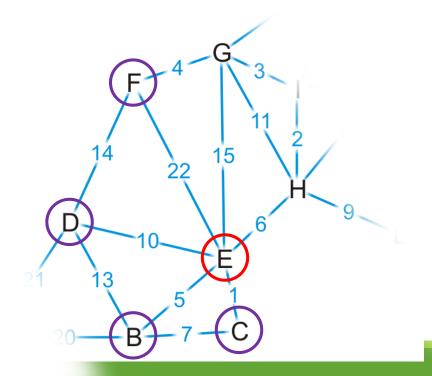
Vertex G has four unvisited neighbors: B, C, D and F



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
Е	T	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

The path (K, H, E) is the shortest path from K to E of length 14

Vertex G has four unvisited neighbors: B, C, D and F



Vertex	Visited	Distance	Previous
Α	F	00	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	T	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

Consider these paths:

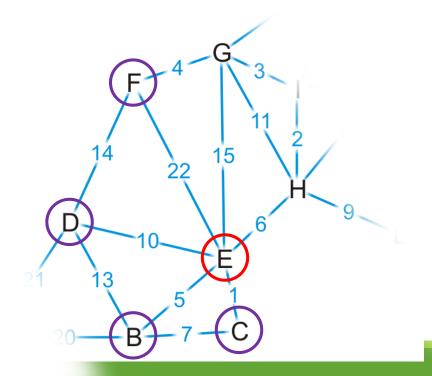
$$(K, H, E, B)$$
 of length $14 + 5 = 19$
 (K, H, E, D) of length $14 + 10 = 24$

(K, H, E, C) of length 14 + 1 = 15(K, H, E, F) of length 14 + 22 = 36

0	Which do we update?	G
	F	4 3
	14	15 22 15
	(D)	H 9
	21 13	
	20—(B)	7-C

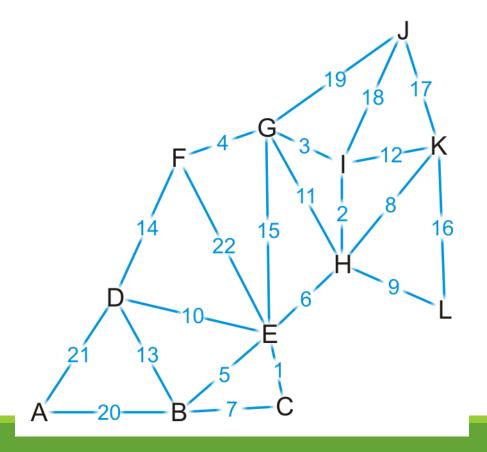
Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	∞	Ø
С	F	∞	Ø
D	F	∞	Ø
E	T	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

We've discovered paths to vertices B, C, D



Vertex	Visited	Distance	Previous
A	F	00	Ø
В	F	19	Е
С	F	15	E
D	F	24	E
E	T	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

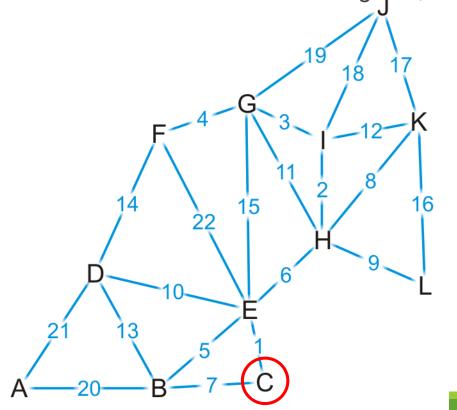
Which vertex is next?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	F	15	E
D	F	24	E
Е	Т	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

We've found that the path (K, H, E, C) of length 15 is the shortest path from K to C

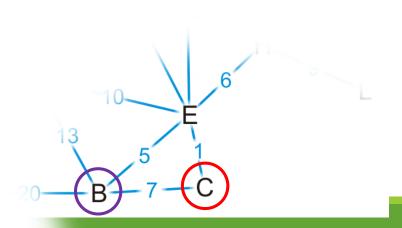
Vertex C has one unvisited neighbor, B



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	Е
C	T	15	E
D	F	24	Е
Е	Т	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

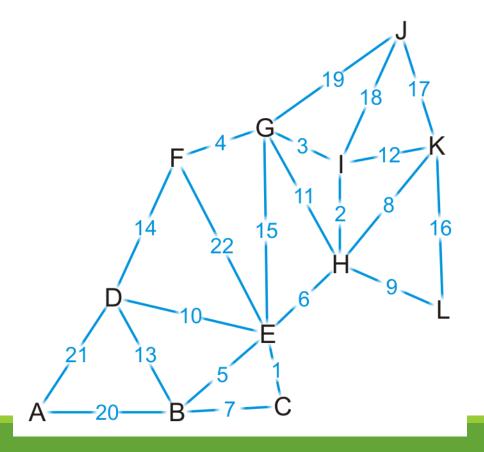
The path (K, H, E, C, B) is of length 15 + 7 = 22

We have already discovered a shorter path through vertex E



Vertex	Visited	Distance	Previous
Α	F	00	Ø
В	F	19	Е
C	T	15	E
D	F	24	E
Е	Т	14	Н
F/	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

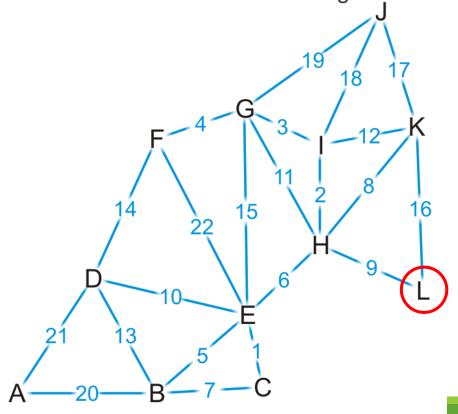
Where to next?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	Е
Е	Т	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	F	16	K

We now know that (K, L) is the shortest path between these two points

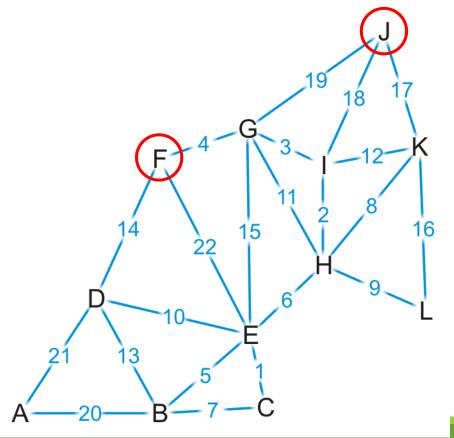
Vertex L has no unvisited neighbors



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	E
Е	Т	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	T	16	K

Where to next?

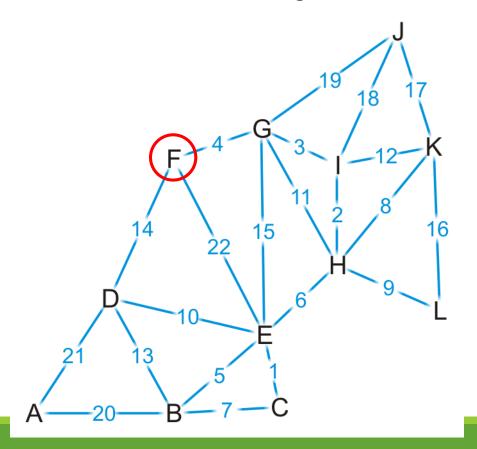
Does it matter if we visit vertex F first or vertex J first?



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	E
Е	Т	14	Н
F	F	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	Т	16	K

Let's visit vertex F first

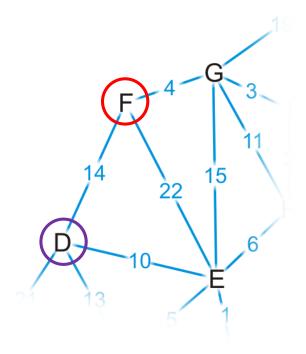
It has one unvisited neighbor, vertex D



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	E
Е	Т	14	Н
F	T	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	Т	16	K

The path (K, H, I, G, F, D) is of length 17 + 14 = 31

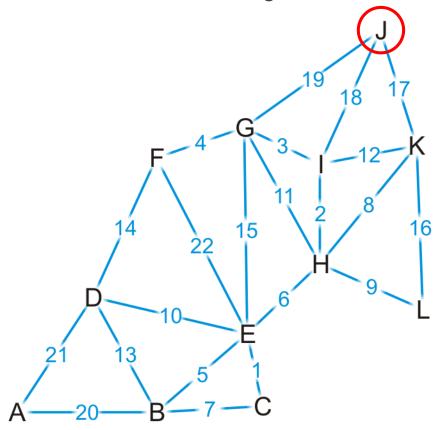
This is longer than the path we've already discovered



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	Е
Е	Т	14	Н
F	T	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	F	17	K
K	Т	0	Ø
L	T	16	K

Now we visit vertex J

It has no unvisited neighbors



Vertex	Visited	Distance	Previous
Α	F	∞	Ø
В	F	19	E
С	Т	15	Е
D	F	24	E
Е	Т	14	Н
F	Т	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	T	17	K
K	Т	0	Ø
L	T	16	K

Next we visit vertex B, which has two unvisited neighbors:

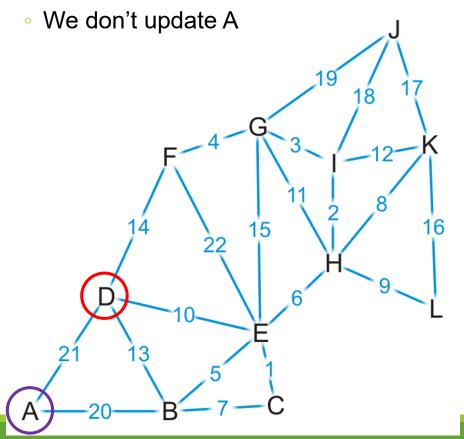
(K, H, E, B, A) of length 19 + 20 = 39 (K, H, E, B, D) of length 19 + 13 = 32

 We update the path length to A 15

Vertex	Visited	Distance	Previous
Α	F	39	В
В	T	19	E
С	Т	15	Е
D	F	24	Е
Е	Т	14	Н
F	Т	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	Т	17	K
K	Т	0	Ø
L	Т	16	K

Next we visit vertex D

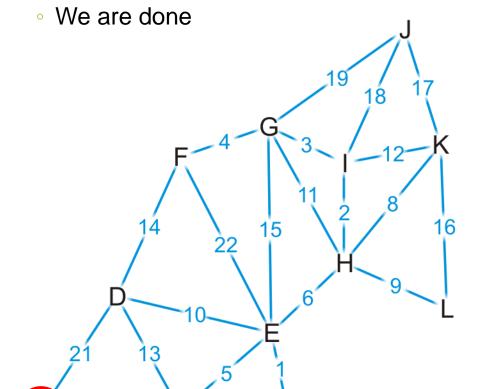
• The path (K, H, E, D, A) is of length 24 + 21 = 45



Vertex	Visited	Distance	Previous
Α	F	39	В
В	Т	19	Е
С	Т	15	Е
D	T	24	E
Е	Т	14	Н
F	Т	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	Т	17	K
K	Т	0	Ø
L	Т	16	K

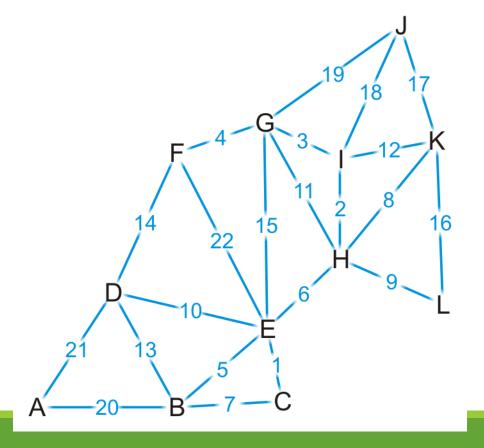
Finally, we visit vertex A

• It has no unvisited neighbors and there are no unvisited vertices left



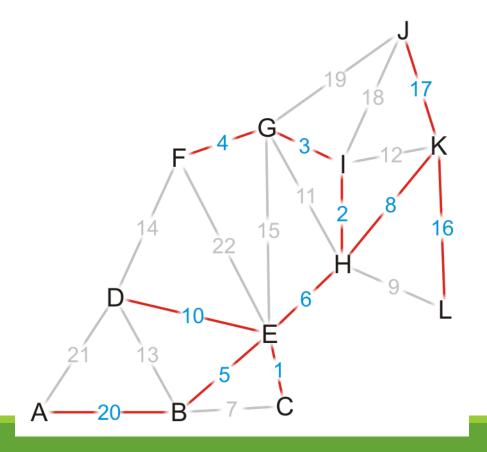
Vertex	Visited	Distance	Previous
A	T	39	В
В	Т	19	Е
С	Т	15	Е
D	Т	24	Е
Е	Т	14	Н
F	Т	17	G
G	Т	13	
Н	Т	8	K
	Т	10	Н
J	Т	17	K
K	Т	0	Ø
L	T	16	K

Thus, we have found the shortest path from vertex K to each of the other vertices



Vertex	Visited	Distance	Previous
Α	Т	39	В
В	Т	19	E
С	Т	15	Е
D	Т	24	Е
Е	T	14	Н
F	T	17	G
G	T	13	
Н	T	8	K
I	Т	10	Н
J	Т	17	K
K	Т	0	Ø
L	Т	16	K

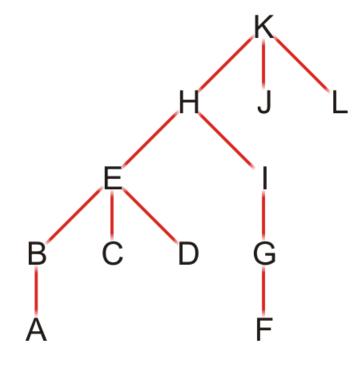
Using the *previous* pointers, we can reconstruct the paths



Vertex	Visited	Distance	Previous
Α	Т	39	В
В	Т	19	E
С	Т	15	Е
D	Т	24	Е
Е	Т	14	Н
F	Т	17	G
G	Т	13	
Н	Т	8	K
I	Т	10	Н
J	Т	17	K
K	Т	0	Ø
L	T	16	K

Note that this table defines a rooted parental tree

- The source vertex K is at the root
- The previous pointer is the parent of the vertex in the tree



Vertex	Previous
Α	В
В	E
С	E
D	E
Е	Н
F	G
G	l
Н	K
I	Н
J	K
K	Ø
L	K

Comments on Dijkstra's algorithm

Questions:

- What if at some point, all unvisited vertices have a distance ∞?
- What if we just want to find the shortest path between vertices v_i and v_k ?
- Does the algorithm change if we have a directed graph?

Implementation and analysis

The initialization requires $\Theta(|V|)$ memory and run time

We iterate |V| - 1 times, each time finding next closest vertex to the source

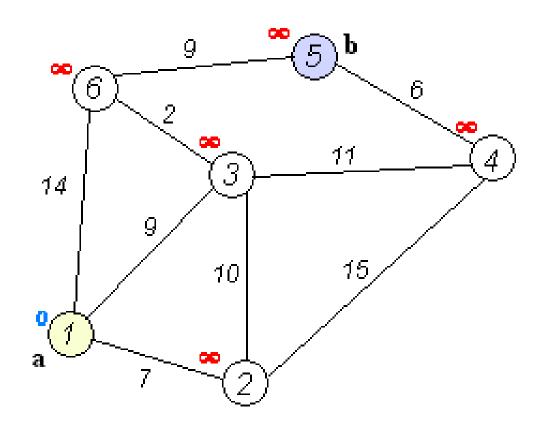
- Iterating through the table requires is $\Theta(|V|)$ time
- Each time we find a vertex, we must check all of its neighbors
- With an adjacency matrix, the run time is $\Theta(|V|(|V|+|V|)) = \Theta(|V|^2)$
- With an adjacency list, the run time is $\Theta(|V|^2 + |E|) = \Theta(|V|^2)$ as $|E| = O(|V|^2)$

Can we do better?

- Recall, we only need the closest vertex
- How about a priority queue?
 - Assume we are using a binary heap
 - We will have to update the heap structure—this requires additional work
- the total run time is $O(|V| \ln(|V|) + |E| \ln(|V|)) = O(|E| \ln(|V|))$



Another Example





Practice Example

