PTC (Fall 2021) – Assignment 2

Name: _____ Dept: _____

Grade: _____ ID: ____

Due: Nov. 30, 2021

Problem 1

Give context free grammars that generate the following languages, and give a brief description of the functionality of each variable in your grammars (in natural language).

- a. $\{w \in \{0,1\}^* \mid w = w^R \text{ and } |w| \text{ is even } \}$
- b. $\{w \in \{a,b\}^* | in \ w$, the number of $b = 1 + the number of a\}$.
- c. $\{a^i b^j c^k \mid i, j, k \ge 0, \text{ and } i = j \text{ or } i = k\}$

Let the language L consist of all the regular expressions over the alphabet $\{0, 1\}$. Thus, L is a language over the alphabet $\{0, 1, \epsilon, \emptyset, +, \circ, *, (,)\}$ and includes, for example, the string $(1 \circ (0 + 0 \circ 0))^*$. (For simplicity let us assume that concatenation is always written out explicitly using the symbol \circ) Please construct two context-free grammars for L, one ambiguous and one unambiguous. For the ambiguous one, please give an example of a string with two different parse trees, and the corresponding left derivations.

Consider the following context free grammar: $G = (\{S, B\}, \{0, 1\}, P, S)$, where P consists of:

$$S \to 0 BB$$

$$B \to 0 S|1S|0$$

- a. For the string 010000, give its parse tree and rightmost derivation according to G.
- b. Provide a nondeterministic PDA P that accepts the language L(G) by empty stack.

Begin with the grammer:

$$\begin{split} S &\to aAa \mid bBb \mid \epsilon \\ A &\to C \mid a \\ B &\to C \mid b \\ C &\to CD \mid \epsilon \\ D &\to A \mid B \mid ab \end{split}$$

- 1. Eliminate ϵ -productions.
- 2. Eliminate any unit productions in the resulting grammar of (1.).
- 3. Eliminate any useless symbols in the resulting grammar of (2.).
- 4. Put the resulting grammar of (3.) into Chomsky normal form.

Given grammer G:

$$S \rightarrow AB \mid BC$$

$$A \rightarrow BA \mid a$$

$$B \rightarrow CC \mid b$$

 $C \to AB \mid a$

Please use CYK algorithm to decide whether string aaaaba belongs to $\mathcal{L}(G)$.

Use the CFL pumping lemma to show each of these languages are not context free.

- a. $L = \{ 0^p \mid p \text{ is a prime } \}$
- b. $L = \{ 0^p 1^q 2^p 3^{(p+q)} \mid p, \ q \ge 0 \}$
- c. L = { $\omega \in \{0,\ 1,\ 2\}^* \mid \omega$ has equal number of 0's, 1's and 2's. }

For a language L, let $\operatorname{prefix}(L)$ be the language $\{w \mid wv \in L, \text{ for some } v \in \Sigma^*\}$

- a. Show that if L is context-free, so is $\mathbf{prefix}(L)$.
- b. Show that if $\mathbf{prefix}(L) \subseteq L$, L is infinite and L is context-free, then there is an infinite regular subset of L.

Proof.

We define an operation \bowtie for language A and B to be

$$A\bowtie B=\{w\mid w=a_1b_1a_2b_2\cdots a_nb_n, \text{ where } a_1a_2\cdots a_n\in A \text{ and } b_1b_2\cdots b_n\in B, \text{ each } a_i,b_i\in \Sigma\}$$

- a. Show that if A is context-free and B is regular, then $A \bowtie B$ is context-free.
- b. Show that the class of CFL is not closed under ⋈ operation.

Proof.