

Bonus-2 Battle Over Cities - Hard Version

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1. Introduction

1.1 Description

With the information of all the cities and the highways connected to them (destroyed or remaining), we're supposed to find the city that we'll take the largest cost to restore the connectivity of the graph if that city is conquered.

Actually, we are asked to find the largest **minimum spanning tree** for that graph with each time one vertex deleted.

2.1 Algorithm Specification

In this algorithm, I used the **Kruska's Algorithm** to find the minimum spanning tree for each condition. Then general implementation is as follows.

After finishing reading in the map, I firstly used `qsort` to sort the edges according the decreasing status. That is, after this sort, those highways (edges) which are in use will be put prior to those destroyed. The reason for this sort is that we will add the existing edges to the *minimum spanning tree* first (since we don't need to spend any cost to repair it). Then, I separately sort the edges in use and edges that are destroyed according to the ascending cost. So after these three `qsort`, we just need to add edges to the minimum spanning tree according to the ascending index sequence.

The `MST` function is where the *Kruska's Algorithm* takes place. Each time I omit one vertex and all the edges incident to it, and finding the minimum spanning tree by add the edges according to the previously sorted edge sequence. Note that if the total number of edges in the final minimum spanning tree is less than $N - 2$ (N is the number of vertices in the graph), then the minimum spanning tree is not connected (The minimum number of edges needed to connect $N - 1$ vertices is $N - 2$). So that deleted vertex must be the answer (one of the answer).

Pseudocode is as follows. Note that N is the number of vertices while M is the number of edges, `result[]` is the array to store the minimum spanning trees and `S[]` is the array of the minimum spanning tree union. For example, `result[i]` is the minimum spanning tree after deleting the vertex i and the edges incident to it, and `S[i]` is the previous vertex in the union of vertex i .

```
1  read in the Graph
2
3  sort all the edges by putting the destroyed edges last
4  sort the used edges in ascending cost
5  sort the destroyed edges in ascending cost
6
7  for each i from 1 to N:
8      result[i] <- 0
9      count <- 0
10     for each j from 1 to M:
11         S[j] <- (-1)
12     for each j from 0 to M-1:
13         if one of the end points of edge[j] is i:
14             continue
15         root_source <- the root of the end_point_1 in the union
16         root_dest <- the root of the end_point_2 in the union
17         if root_source != root_dest:           // no cycle produced
18             S[root_source] <- root_dest
```

```

19         count++
20         if edge.status == 0:           // add to the MST if it's destroyed
21             result[i] <- result[i] + edge[j].cost
22     if count < N-2:
23         result[i] <- INF
24     ans <- MAX(ans, result[i])    // update the ans
25
26 for each element in result[]:
27     if element == ans
28         output index

```

3. Testing Results

- **Test point 1: MAX N, connected**

Input :

See the file "MAX_N_connected.txt"

Output :

0

- **Test point 2: MAX N, output all**

Input :

See the file "MAX_N_outputAll.txt"

Output :

2 3 4 ... 499

4. Analysis and Comments

4.1 Time Complexity

The average time complexity for this algorithm is $O(E \log V)$, because of the Kruskal's Algorithm with union find operation. Other functions like `initGraph` and `printResult` are all $O(N)$.

4.2 Space Complexity

The space complexity for this algorithm is $O(E + 2V)$, since in the function `initGraph`, I used the a structure that has an array of size E , in the function `MST`, there're two utility arrays of size E and size V respectively. So the total space complexity is $O(E + 2V)$.

5. Appendix

Source Code in C

```

1  #include<stdio.h>
2  #include<stdlib.h>
3
4  #define MAX (501)
5  #define FINDMAX(x,y) ((x)>(y)?(x):(y))
6  #define ROOT (-1)           // to indicate the root if a union
7  #define INF (0x7FFFFFFF)    // to mark that if remove this vertex, the
                                graph cannot be connected
8
9  /*****
10     Data Structure Used in This Code
11
12     * GNode: Used to store the information the graph.
13       That is the all the cities(vertices) and
14       the highways(edges) connected to them.
15
16     * Edge: Used to store the information of one edge
17       (highway), including the end points, cost
18       and the status
19
20     *****/
21
22  typedef struct GNode* Graph;
23  typedef struct Edge* PtrToEdge;
24
25  struct GNode {
26      int NumOfVertices;
27      int NumOfEdges;
28      PtrToEdge edge;
29  };
30
31  struct Edge {
32      int source;
33      int dest;
34      int cost;
35      int status;
36  };
37
38  /* S[i] represents the previous ponint in the union of the minimum
39     spanning tree */
40  /* result[i] represents the cost we need to take if city[i] is being
41     conquered*/
42  int S[MAX], result[MAX];
43
44  /*****
45     Functions
46
47     * InitGraph: Read in all the information of the graph
48
49     * myCompare_cost: Compare function used in qsort.Determine
50       the sort according to the ascending cost of rebuild
51       -ing the highway.
52
53     * myCompare_status: Compare function used in qsort.Determine
54       the sort according to the status in decreasing order.
55       That is, put the highways that are in use prior to
56       those destroyed.

```

```

56  * MST: calculate the minimum spanning tree after deleting
57      the vertex and the edges incident with it using
58      Kruska's Algorithm.
59
60  * find: Union find operation. Used path compression.
61
62  * printResult: Print the cities that we must pay most attention
63      to in increasing order.
64  *****/
65
66  Graph initGraph(int* destroy_cnt);
67  int myCompare_cost(const void* a, const void* b);
68  int myCompare_status(const void* a, const void* b);
69  int MST(Graph G);
70  int find(int vertex);
71  void printResult(Graph G, int* result, int ans);
72
73  int main()
74  {
75      int i, ans, destroy_cnt = 0;    // ans is the final answer,
76                                     // destroy_cnt is the number of destroyed highways
77
78      Graph G = initGraph(&destroy_cnt); // read in the graph
79
80      qsort(G->edge, G->NumOfEdges, sizeof(G->edge[0]), myCompare_status);
81      // sort the edges according to the status in decreasing order
82      qsort(G->edge, G->NumOfEdges - destroy_cnt, sizeof(G->edge[0]),
83      myCompare_cost); // sort the edges in use according to the ascending cost
84      // order
85      qsort(G->edge + (G->NumOfEdges - destroy_cnt), destroy_cnt, sizeof(G->
86      edge[0]), myCompare_cost); // sort the destroyed edges according to
87      // ascending cost order
88
89      /*calculate the minimum spanning tree after deleting each vertex and the
90      edges incident with it using Kruska's Algorithm. Store the result*/
91      ans = MST(G);
92
93      printResult(G, result, ans);    // output the answer
94
95      return 0;
96  }
97
98  Graph initGraph(int* destroy_cnt)
99  {
100      Graph G = (Graph)malloc(sizeof(struct GNode));
101      scanf("%d %d", &G->NumOfVertices, &G->NumOfEdges); // read in the
102      // number of cities and highways
103      G->edge = (PtrToEdge)malloc(sizeof(struct Edge) * G->NumOfEdges);
104
105      for (int i = 0; i < G->NumOfEdges; i++) {
106          scanf("%d %d %d %d", &(G->edge[i]).source, &(G->edge[i]).dest, &
107          (G->edge[i]).cost, &(G->edge[i]).status); //read in the information for
108          // each highway
109          if (!(G->edge[i]).status) // count the number of destroyed
110          // highways
111              (*destroy_cnt)++;
112      }
113      return G;

```

```

103 }
104
105 int MST(Graph G)
106 {
107     int i, j, ans, root_source, root_dest, count;
108
109     /*****
110         Variables Used
111
112     * i, j: Temporary counter.
113
114     * ans: The maximum effort we need to take to restore the connectivity
115     of
116         the graph. (final answer)
117
118     * root_source: The root of the source vertex of a highway in the MST
119     union.
120
121     * count: A counter to count the number of edges in the final minimum
122     spanning
123         tree. If it's less than N-2 then it imply that the minimum
124     spanning
125         tree is not conneced. Then that vertex must be the answer.
126
127     *****/
128     ans = 0;
129
130     for (i = 1; i <= G->NumOfVertices; i++) { // loop for each vertex,
131         calculate the MST without that vertex
132         result[i] = 0; // initialize the result[] array and the
133         counter
134         count = 0;
135         for (j = 1; j <= G->NumOfVertices; j++) // initialize the
136         union array
137             S[j] = ROOT;
138         for (j = 0; j < G->NumOfEdges; j++) {
139             if ((G->edge[j]).source == i || (G->edge[j]).dest == i) //
140             if the highway is from/toward that deleting city, continue
141                 continue;
142             root_source = find((G->edge[j]).source); // find the root
143             of the end points
144             root_dest = find((G->edge[j]).dest);
145             // if the root of the end points is identical ,then adding
146             this edge will produce a cycle
147             if (root_source != root_dest) {
148                 S[root_source] = root_dest; // connect the edge into
149                 the union
150                 count++;
151                 if (!(G->edge[j]).status) // only the
152                 destroyed highway needs to be repaired
153                     result[i] += (G->edge[j]).cost;
154             }
155         }
156     }
157 }

```

```

145         if (count < G->NumOfVertices - 2)           // if the edge in the
final MST is less than N-2, then the MST is not connected
146             result[i] = INF;
147             ans = FINDMAX(ans, result[i]); // update the ans
148     }
149     return ans;
150 }
151
152 int find(int vertex)           // union find operation with path compression
153 {
154     if (S[vertex] == ROOT)
155         return vertex;
156     else
157         return S[vertex] = find(S[vertex]);
158 }
159
160
161 void printResult(Graph G, int* result, int ans)      // print the results
162 {
163     int i, flag;
164
165     if (!ans)
166         printf("0\n");
167     else {
168         for (i = 1, flag = 0; i <= G->NumOfVertices; i++)
169             if (result[i] == ans) {
170                 if (flag)
171                     printf(" %d", i);
172                 else {
173                     printf("%d", i);
174                     flag = 1;
175                 }
176             }
177         printf("\n");
178     }
179 }
180
181 int myCompare_cost(const void* a, const void* b) // sort according to the
cost
182 {
183     PtrToEdge a1 = (PtrToEdge)a;
184     PtrToEdge b1 = (PtrToEdge)b;
185     return a1->cost - b1->cost;
186 }
187
188 int myCompare_status(const void* a, const void* b) // sort according to
the status
189 {
190     PtrToEdge a1 = (PtrToEdge)a;
191     PtrToEdge b1 = (PtrToEdge)b;
192     return b1->status - a1->status;
193 }

```