

Networking and Interworking

Network issues

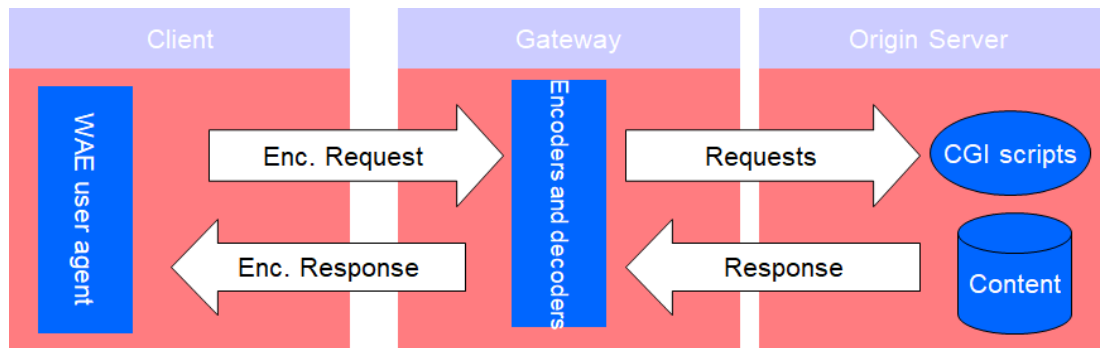
- Performance
 - Latency
 - Data transfer rate
 - $\text{Data transfer rate} = \text{amount of Data} / \text{time}$
 - $\text{Message transfer rate time} = \text{latency} + \text{message length} / \text{data transfer rate}$
 - Total system bandwidth of network
 - Throughput (吞吐量) in the end systems
 - Total volume of traffic can be transferred across network in a time
- Scalability No designable to cope with size and load about network growing
- Reliability Recoverable from communication failures
- Security Protecting network and computers
- Mobility
 - Portability of computer and handled digital devices using wireless network
 - Location and identification are depicted with each other
 - No designable to cope with size and load about network growing
- QoS (Quality of Service)

Guarantee for requirements of computer and network to meet deadline, bandwidth, bounded latency
- Multicasting One-to-many communication

Types of Networks

- LANs (Local Area Networks) Technology suitable for small area, wire/fiber
 - High bandwidth
 - Low latency
 - ATM, better QoS, but more expensive
- WANs (Wide Area Networks) Large distances, inter-city/country/continental
 - Low bandwidth, high latency
 - Satellite/wire/cable
 - Routers introduce delays
- MANs (Metropolitan Area Networks) Intra-city, cable based, multimedia
- Wireless networks WLANs, WPANs (wireless personal area network)
 - WLANs: Wireless Local Area Networks to replace wired LANs

- WaveLAN technology
 - WIFI: Wireless Internet Platform for Interoperability
 - IEEE 802.11
- WMAN: Wireless Metropolitan Area Networks
 - WiMAX: Worldwide Interoperability for Microwave Access
 - IEEE 802.16
- WPANs: Wireless Personal Area Networks
 - Infrared (infra-red) links
 - Bluetooth: IEEE 802.15.1 (standard no longer maintained)
 - Mobile phone network (Wireless WAN)
 - European GSM: Global system for Mobile communication
 - US: analogue AMPS: Advanced Mobile Phone System Cellular radio network, Cellular digital packet data Replaced by GSM
 - WAP: Wireless Applications Protocol For use on wireless potable devices



Network Principles

Mode of Transmission

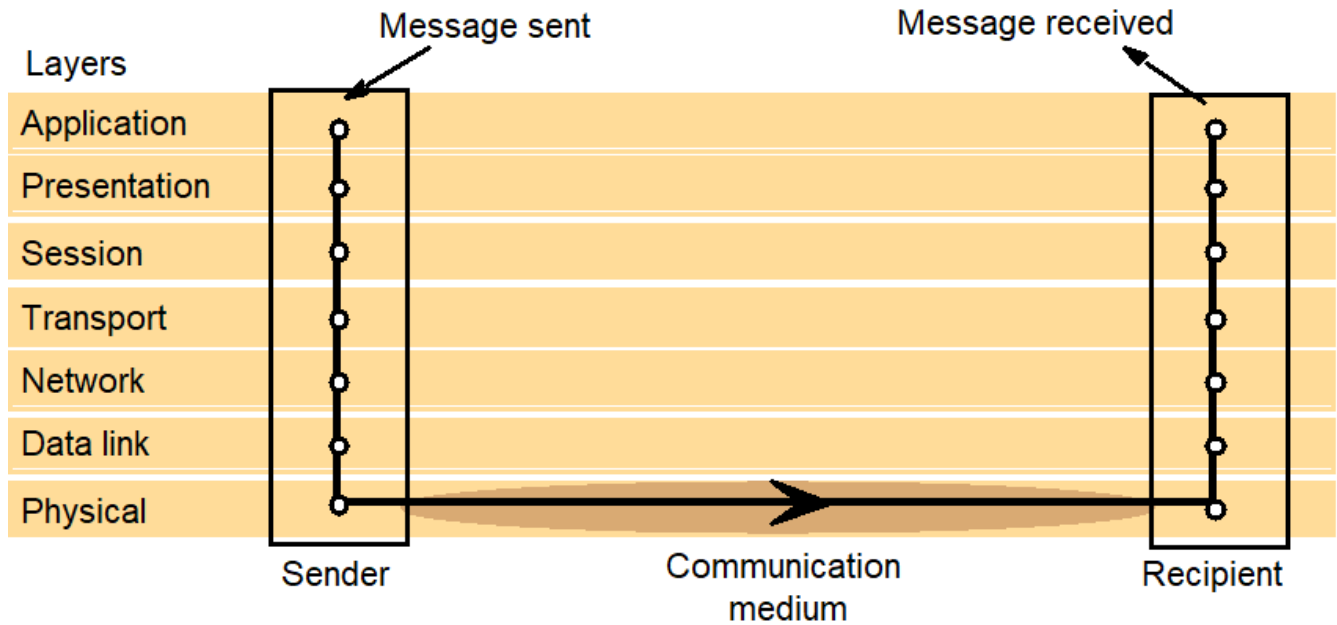
- Packets
 - 1. Messages divided into packets (on Transport Layer)
 - 2. Packets **queued in buffers** before sent onto link
 - QoS not guaranteed
- Data streaming
 - Links **guarantee QoS** (rate of delivery)
 - For multimedia traffic
 - Need higher bandwidth

Switching Schemes

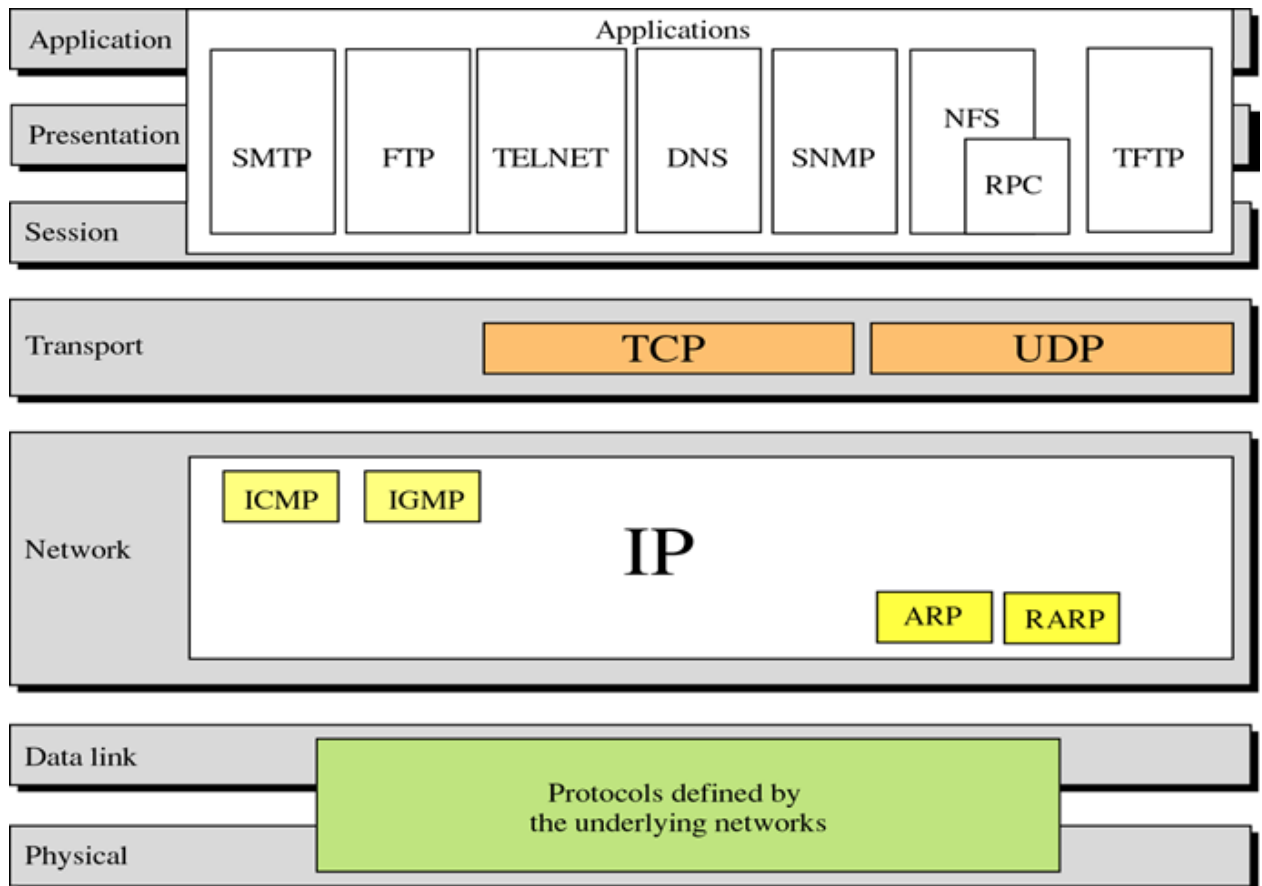
- Broadcasts (Ethernet, wireless)
 - Send messages to all nodes
 - Nodes listen for own messages (carrier sensing)
- Circuit switching (phone networks)

- Packet switching (TCP/IP)
 - Store-and-forward
 - Unpredictable delays
- Frame/Cell relay (ATM)
 - Bandwidth & latency guaranteed (virtual path)
 - Small, fixed size packets (padded if necessary) 53 bytes= header 5 + body 48
 - Avoids error checking at nodes (use reliable links)

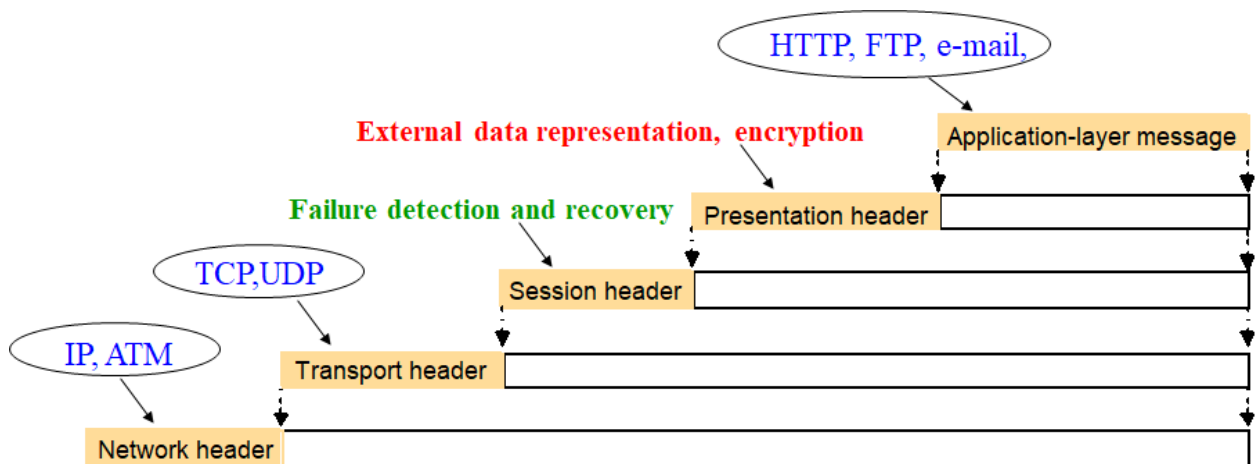
Protocols



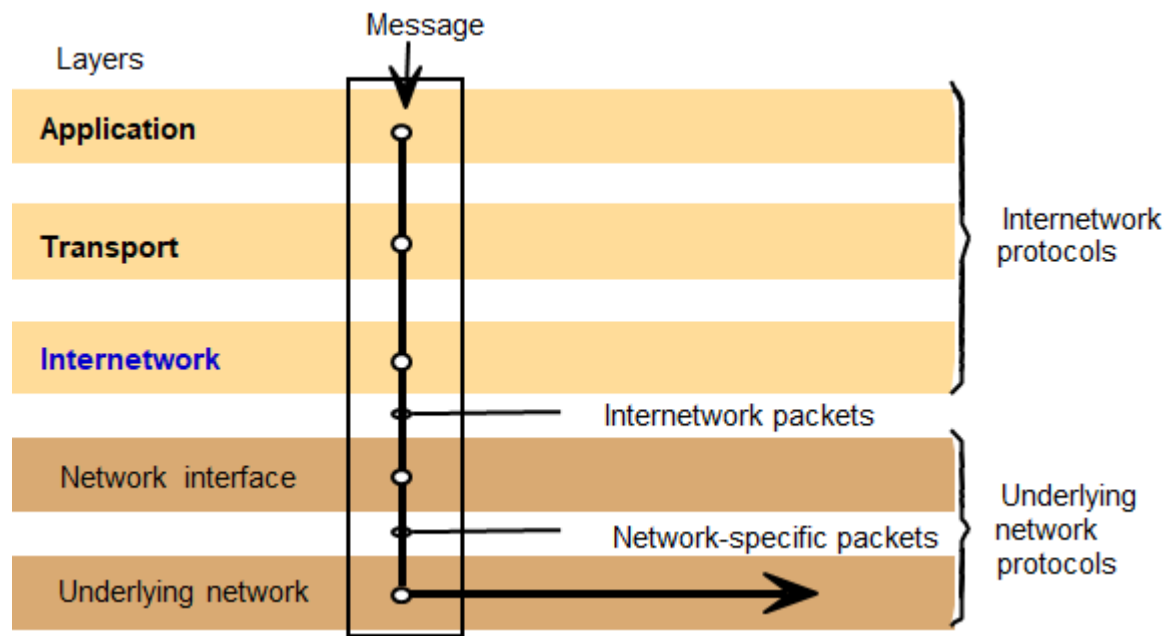
- OSI Model



- Message encapsulation



- Internetwork protocol
 - Internetwork layer(=Virtual network layer) internet packet → destination (by datagram protocol)
 - Network interfaces layer internetwork packets → suitable packets → underlying layer
 - Underlying network layer

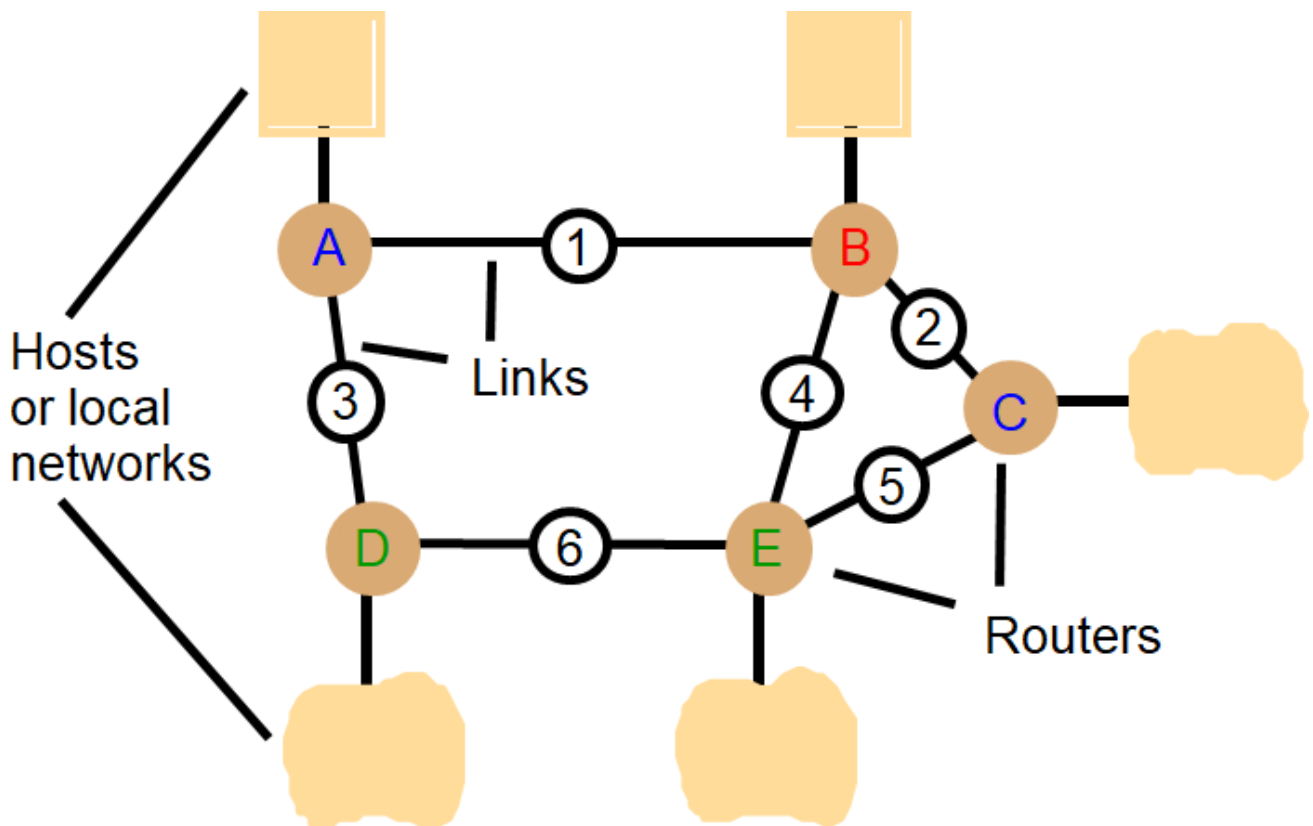


- Port & Addressing
 - Port
 - Network-independent message transport service between networks ports
 - Software-definable destination points for communications
 - Addressing
 - Delivering messages to destination with transport addresses
 - Transport address Network address + port number
- Packet delivery
 - In network layer
 - Datagram packet delivery (IP in Ethernet, most wired and wireless LAN technologies)
 - Virtual circuit packet delivery (ATM)
 - In transport layer
 - Connection-oriented transmission(TCP) Reliable communication with static routing table (ISO, X.25) Ex) remote login(Telnet), FTP, HTTP(big-sized file), stream data
 - Connectionless transmission (UDP) Unreliable communication with pre-defined routing table Ex) rcp, rwho, RPC, HTTP(small-sized file), FTP(non-error bulk file)

Routing

- Necessary in non-broadcast networks (cf Internet) : Hop by Hop
- Distance-vector algorithm for each node
 - Stores table of state & cost information of links, cost infinity for faulty links
 - Determines route taken by packet (the next hop)
 - Periodically updates the table and sends to neighbors
 - May converge slowly [Bellman-Ford]
- RIP-1 (Router Information Protocol) for Internet

- Local router table changes
- Use default routes, plus multicast and authentication
- Better convergence (routes better route to an existing destination)



- RIP routing algorithm

Variables: T_l local table, T_r remote table received.

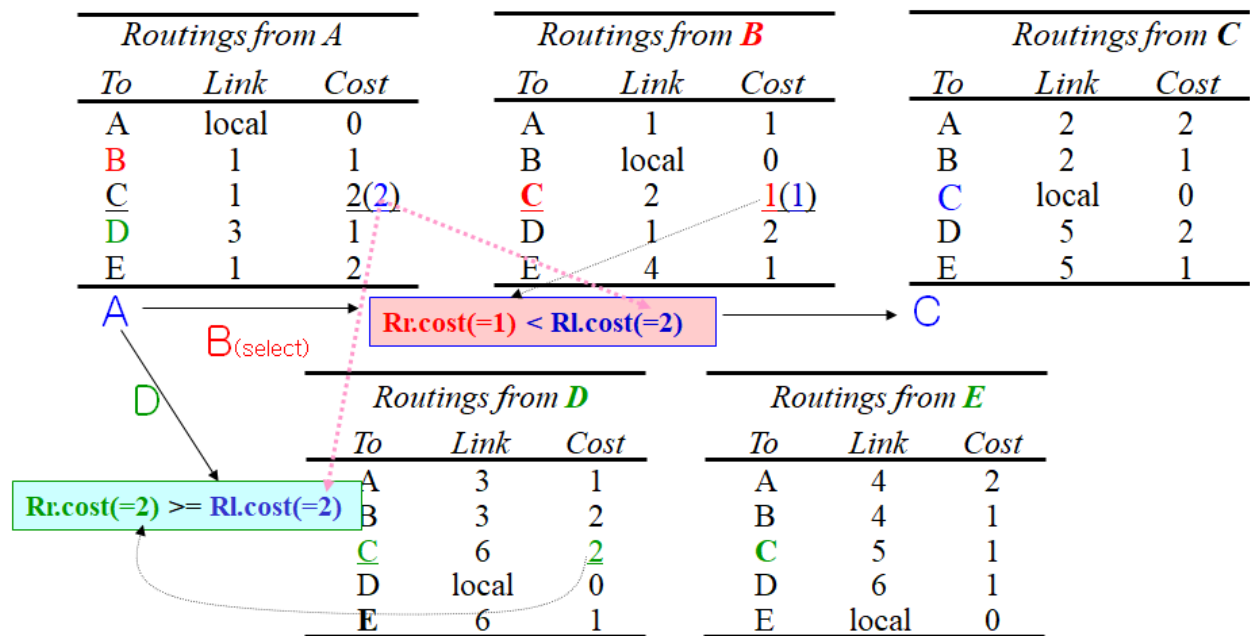
- R_r : remote, R_l : local

Send: Each t seconds or when T_l changes, send T_l on each non-faulty outgoing link.

Receive: Whenever a routing table T_r is received on link n :

```
for all rows  $R_r$  in  $T_r$  {
    if ( $R_r.link \neq n$ ) {
         $R_r.cost = R_r.cost + 1$ ; // hop
         $R_r.link = n$ ;
        if ( $R_r.destination$  is not in  $T_l$ ) add  $R_r$  to  $T_l$ ;
        // add new destination to  $T_l$ 
        else for all rows  $R_l$  in  $T_l$  {
            if ( $R_r.destination = R_l.destination$  and ( $R_r.cost < R_l.cost$  or
 $R_l.link = n$ ))
                 $R_l = R_r$ ;
            //  $R_r.cost < R_l.cost$  : remote node has better route
            //  $R_l.link = n$  : remote node is more authoritative
        }
    }
}
```

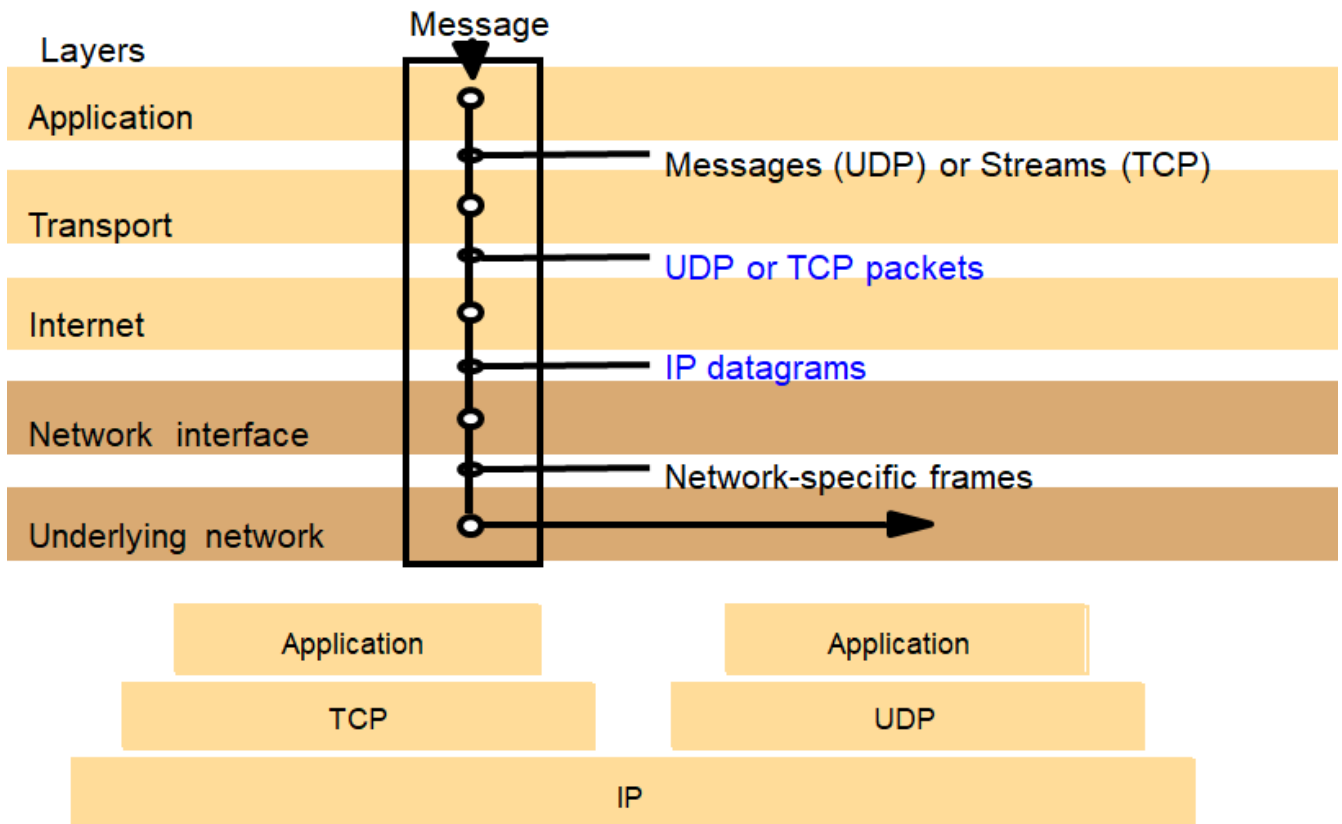
- Routing tables (A \rightarrow C)



Congestion Control

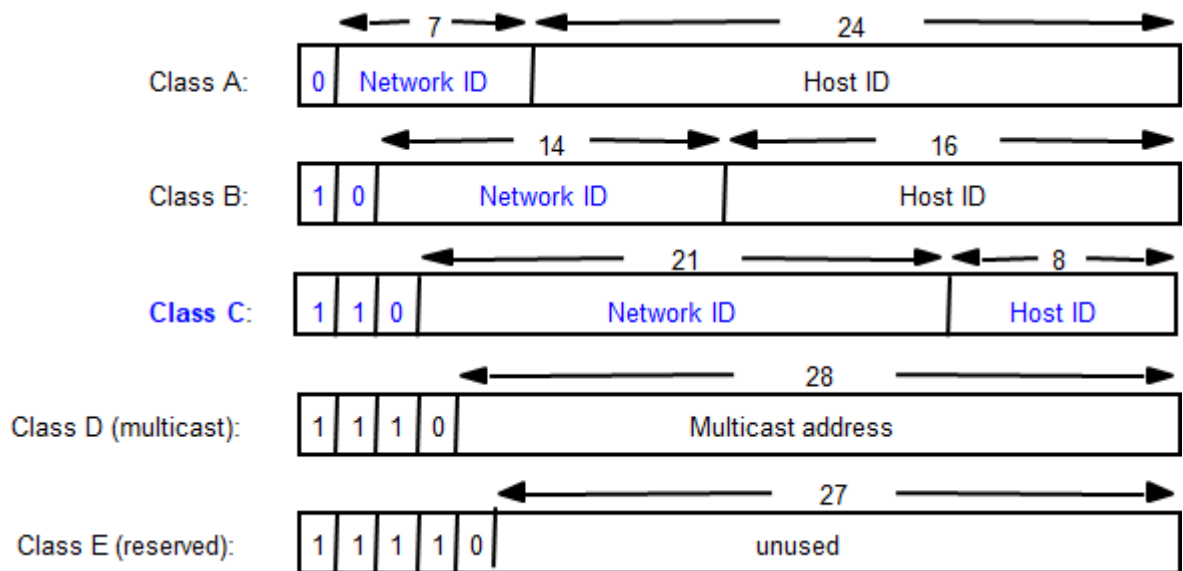
- When load on network exceeds 80% of its capacity packet queues long, links blocked
- Solutions (in datagram-based network layers)
 - packet dropping reliable of delivery at higher levels
 - reduce rate of transmission
 - nodes send choke packets (Ethernet) special message requesting a reduction in transmission rate
 - transmission control (TCP)
 - transmit congestion information to each node QoS guarantees (ATM)

Transport level protocols

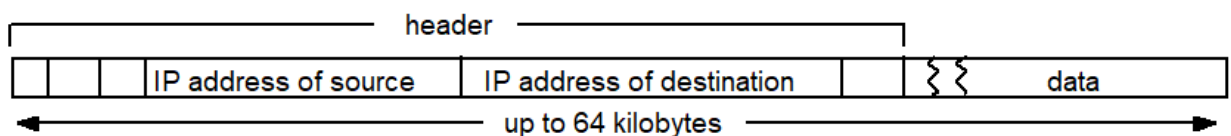


IP(TCP/IP) Addressing

- IP Structure(Universal)(IPv4 : 4bytes = 32bits)



- IP Packet layout



IPv6

	octet 1	octet 2	octet 3	octet 4	Range of addresses
	Network ID		Host ID		
Class A:	1 to 127	0 to 255	0 to 255	0 to 255	1.0.0.0 to 127.255.255.255
	Network ID		Host ID		
Class B:	128 to 191	0 to 255	0 to 255	0 to 255	128.0.0.0 to 191.255.255.255
	Network ID		Host ID		
Class C:	192 to 223	0 to 255	0 to 255	1 to 254	192.0.0.0 to 223.255.255.255
	Multicast address				
Class D (multicast):	224 to 239	0 to 255	0 to 255	1 to 254	224.0.0.0 to 239.255.255.255
Class E (reserved):	240 to 255	0 to 255	0 to 255	1 to 254	240.0.0.0 to 255.255.255.255

- Large Address Space – 128 bit addresses Every toaster can have its own IP address
- Aggregation-based address hierarchy Efficient backbone routing
- Efficient and Extensible IP datagram
 - No fragmentation by routers
 - 64 bits field alignment
 - Simpler basic header
- Auto-configuration
- Security
- IP Renumbering part of the protocol

TLA	NLA(s)	SLA	Interface ID
16 bits	32 bits	16 bits	64 bits

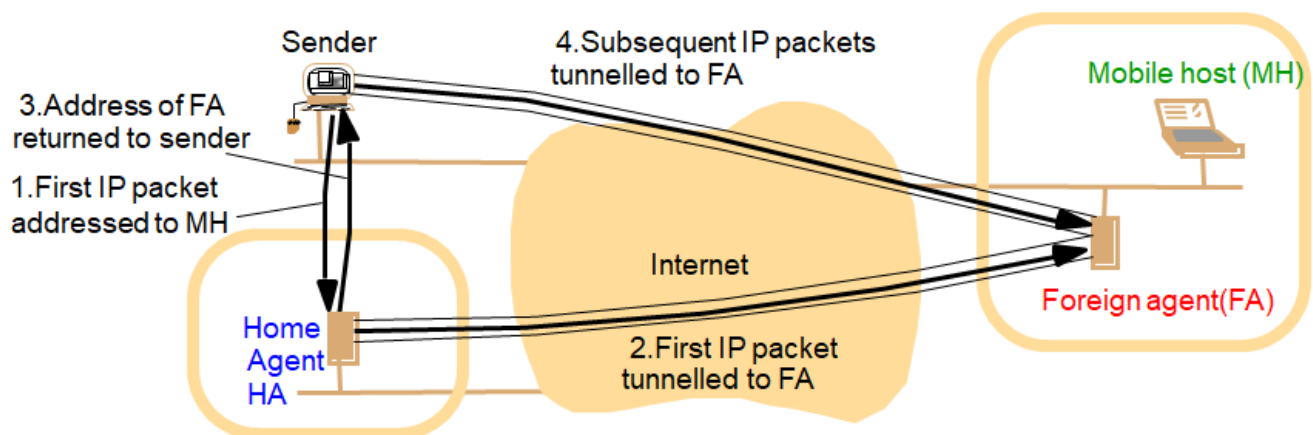
- TLA – top level aggregator Primary providers
- NLA: Next Level Aggregator Can have multiple NLA as sub-NLA
- SLA: Site Level Aggregator Your site (16 bits)
- Addresses are allocated from your provider If you change provider, your prefix changes But renumbering (of hosts, routers and sites) has been included in the IPv6 protocol
- IPv6 Header layout (16bytes: 128bits)

Version (4 bits)	Traffic class (8 bits)	Flow label (20 bits)
Payload length (16 bits)	Next header (8 bits)	Hop limit (8 bits)
Source address (128 bits)		
Destination address (128 bits)		

- IPv6's main advances (Adapted by IETF in 1994)
 - Address space ($2^{128} = 3 \times 10^{38}$ IPs), routing speed up
 - Real-time and other special services
 - Future evolution
 - Multicasting & anycasting
 - Security

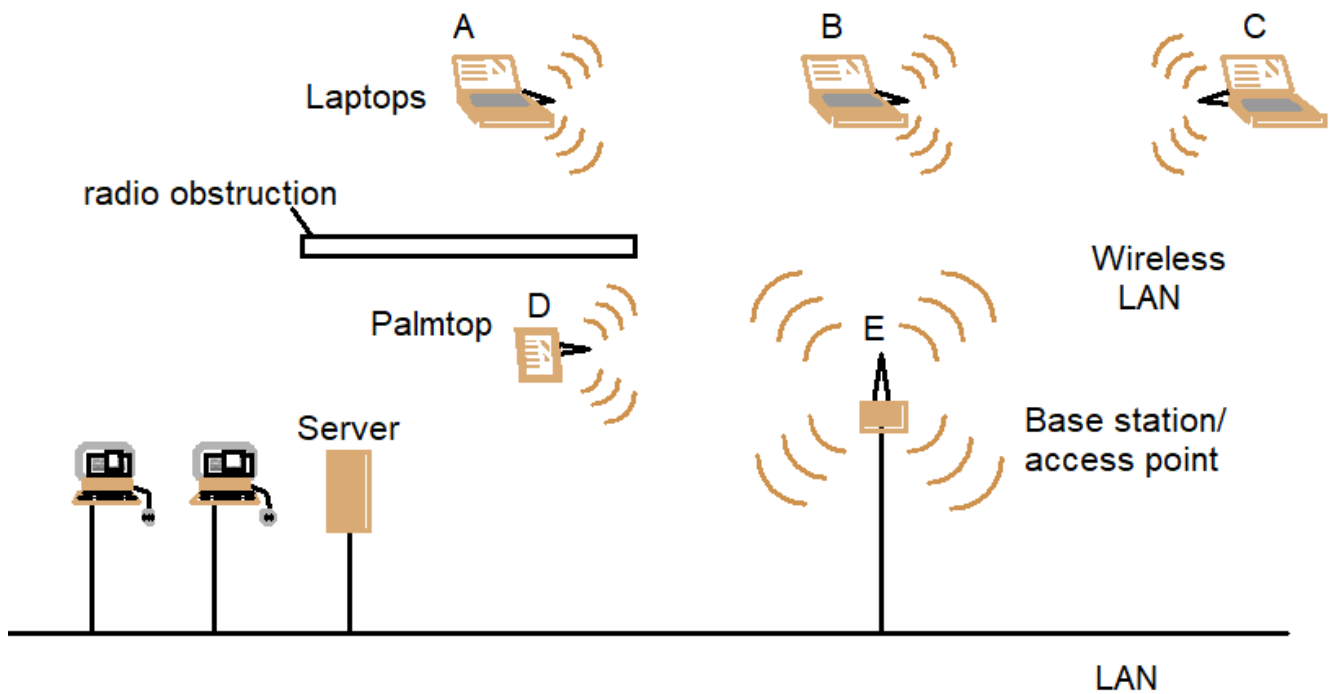
MobileIP

- At home normal, when elsewhere mobile host: notifies HA(Home Agent) before leaving informs FA(Foreign Agent), who allocates temporary care-of IP address & tells HA
- Packets for mobile host (MH): first packet routed to HA, encapsulated in MobileIP packet and sent to FA (tunneling) FA unpacks MobileIP packet and sends to mobile host sender notified of the care-of address for future communications which can be direct via FA
- Problems efficiency low, need to notify HA



Wireless LAN: IEEE 802.11

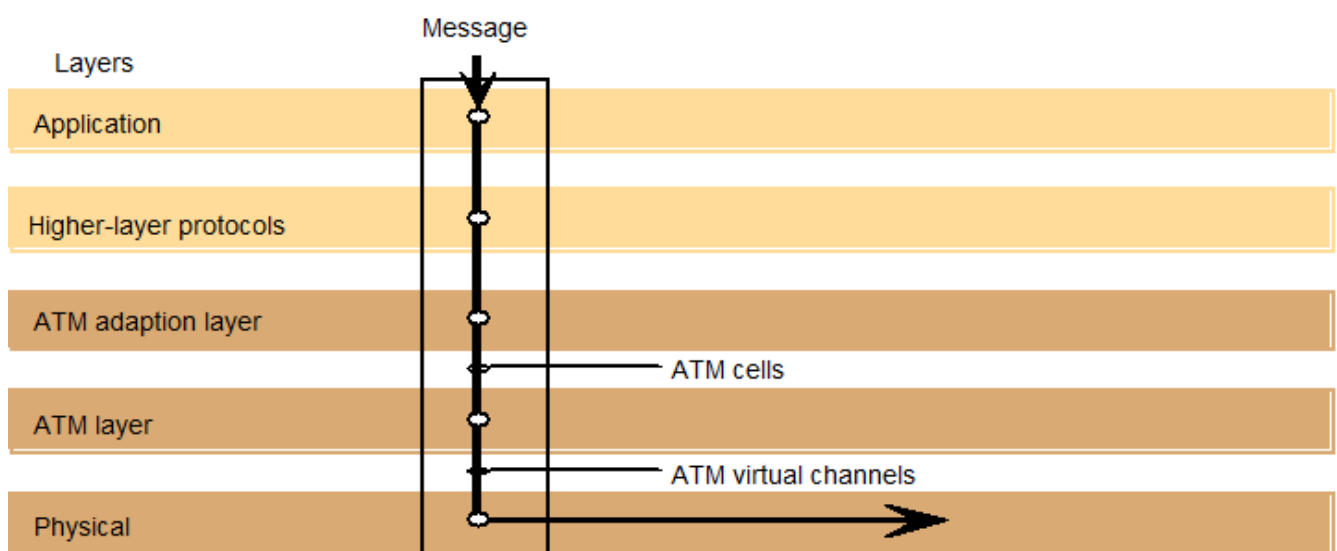
- Radio broadcast (fading strength, obstruction)
- Collision avoidance by Slot reservation mechanism by Request to Send (RTS) and Clear to Send (CTS) Stations in range pick up RTS/CTS and avoid transmission at the reserved times Collisions less likely than Ethernet, since RTS/CTS short Random back off period
- Problems Security (eaves dropping), use shared-key authentication



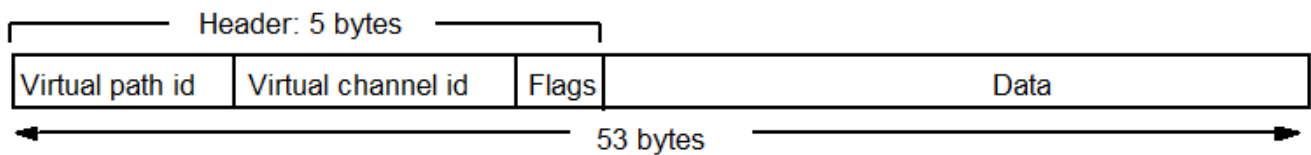
Asynchronous Transfer Mode (ATM)

- Multimedia data(voice and video), distributed system services are available
- Packet switching network based on Cell-relay (a method of packet routing)
- Avoiding flow-control and error checking at the intermediate nodes
- Small and fixed length unit of data transmitted (53bytes= header 5 + body 48)
reduction of buffer size, complexity, queuing delay at intermediate nodes
- B-ISDN (CCITT I.150 standard)
- Optical fiber transmission medium (155 - 622 megabits/sec)
- ATM protocol layers (next page)

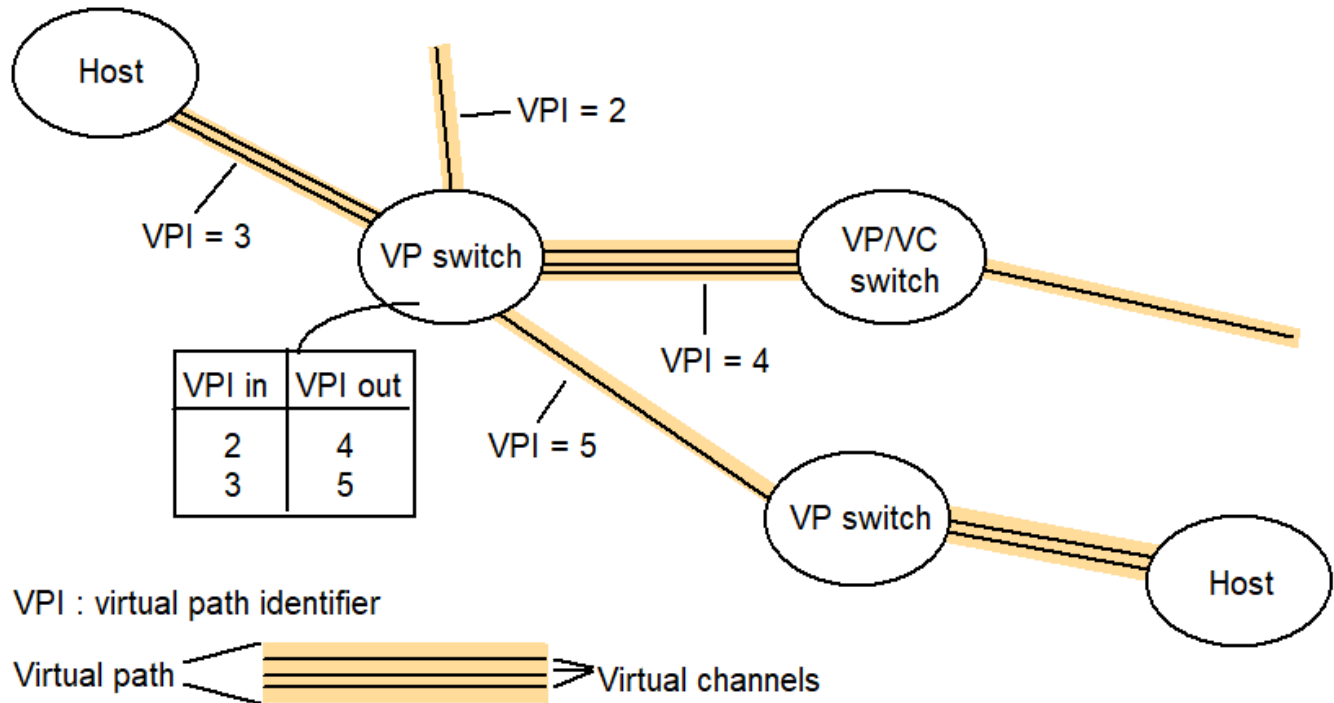
ATM Protocol Layer



ATM Cell Layout



Switching virtual in an ATM network



• From 202.30.11.187 to 202.30.11.44

```

202.30.11.187    1 1 0 0 1 0 1 0 0 0 0 1 1 1 1 0 0 0 0 0 1 0 1 1 1 0 1 1 1 0 1 1
                  AND
255.255.255.0    1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0
                  ||
202.30.11.0      1 1 0 0 1 0 1 0 0 0 0 1 1 1 1 0 0 0 0 0 1 0 1 1 0 0 0 0 0 0 0 0

202.30.11.44     1 1 0 0 1 0 1 0 0 0 0 1 1 1 1 0 0 0 0 0 1 0 1 1 0 0 1 0 1 1 0 0
                  AND
255.255.255.0    1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0 0
                  ||
202.30.11.0      1 1 0 0 1 0 1 0 0 0 0 1 1 1 1 0 0 0 0 0 1 0 1 1 0 0 0 0 0 0 0 0
  
```