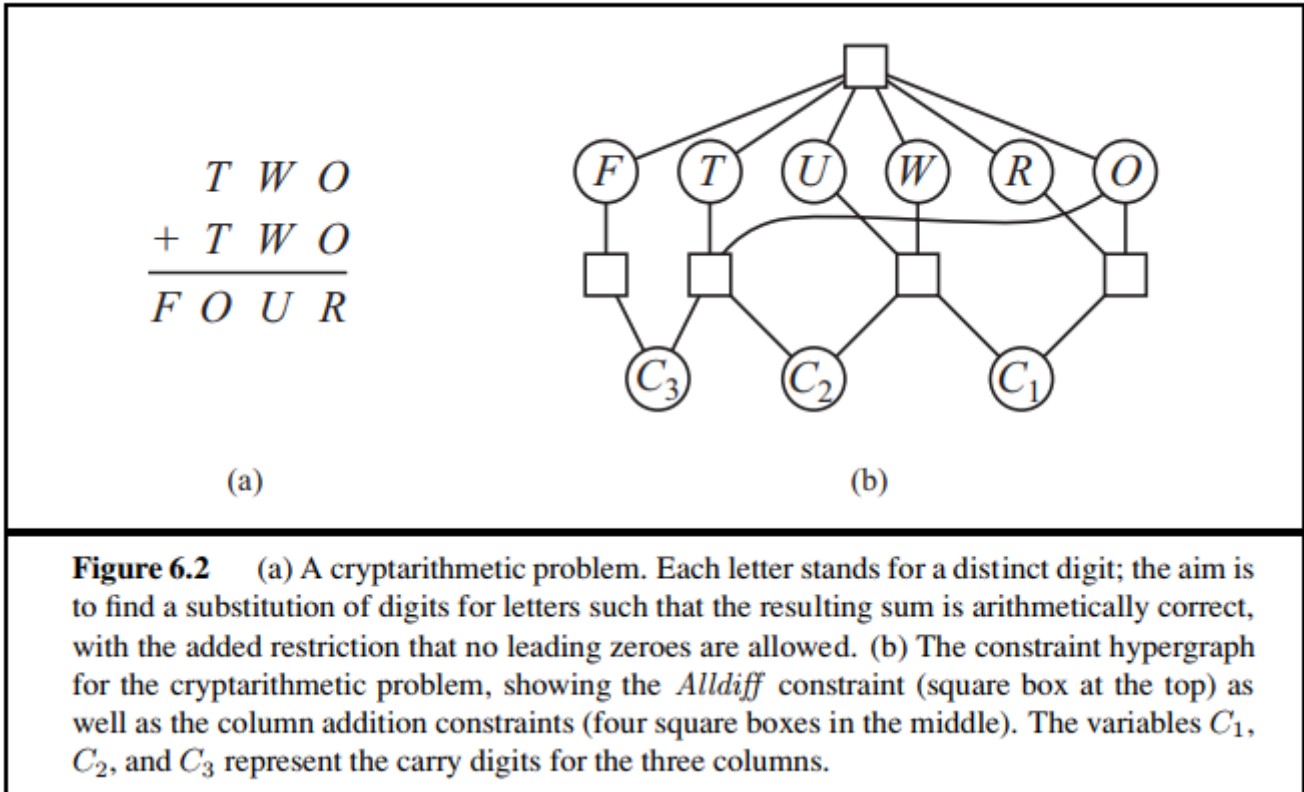


第三次作业答案

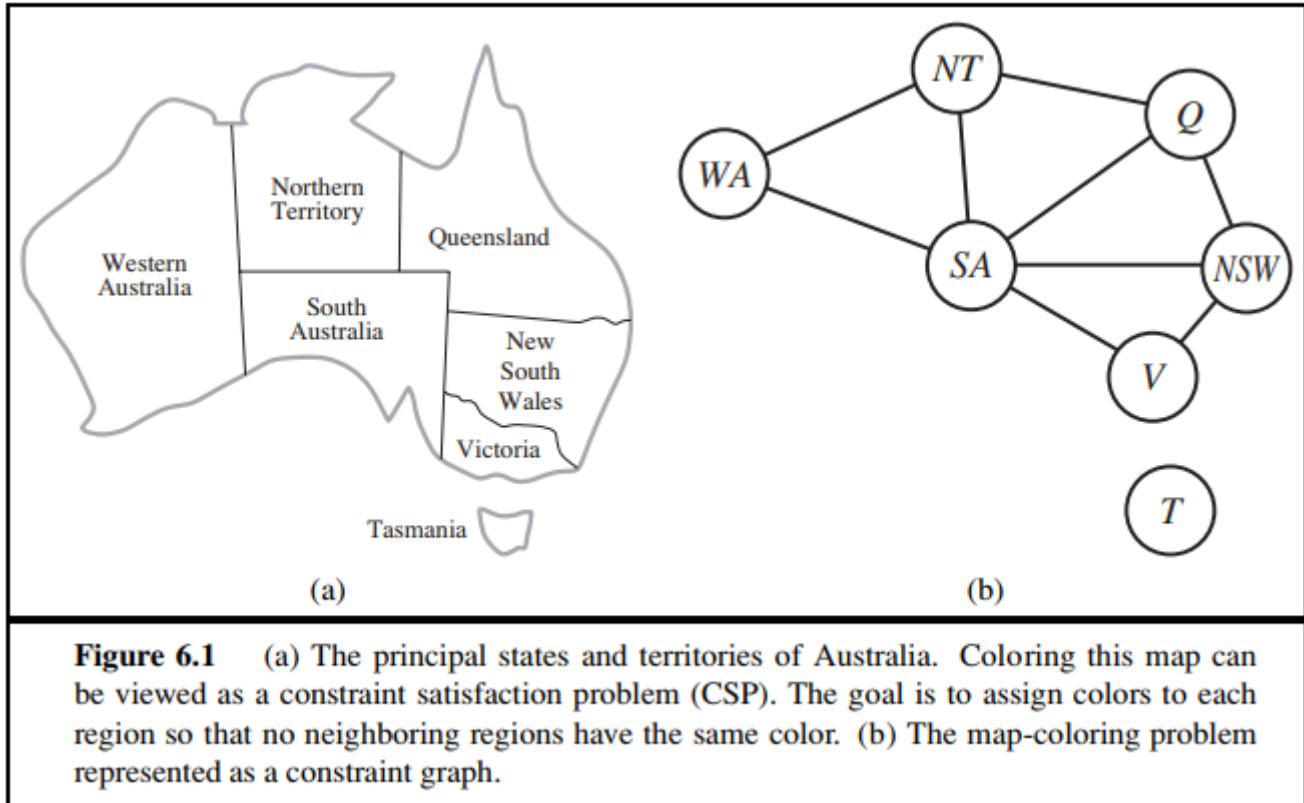
6.5 分别用带有前向检验、MRV和最少约束值启发式的回溯算法手工求解图6.2中的密码算数问题



- Choose the X_3 variable. Its domain is $\{0, 1\}$.
- Choose the value 1 for X_3 . (We can't choose 0; it wouldn't survive forward checking, because it would force F to be 0, and the leading digit of the sum must be non-zero.)
- Choose F , because it has only one remaining value.
- Choose the value 1 for F .
- Now X_2 and X_1 are tied for minimum remaining values at 2; let's choose X_2 .
- Either value survives forward checking, let's choose 0 for X_2 .
- Now X_1 has the minimum remaining values.
- Again, arbitrarily choose 0 for the value of X_1 .
- The variable O must be an even number (because it is the sum of $T + T$ less than 5 (because $O + O = R + 10 \times 0$). That makes it most constrained.
- Arbitrarily choose 4 as the value of O .
- R now has only 1 remaining value.
- Choose the value 8 for R .
- T now has only 1 remaining value.

- n. Choose the value 7 for T .
- o. U must be an even number less than 9; choose U .
- p. The only value for U that survives forward checking is 6.
- q. The only variable left is W .
- r. The only value left for W is 3.
- s. This is a solution.

6.11 用AC-3算法说明弧相容对图6.1中问题能够检测出部分赋值, $WA = green, V = red$ 的不相容。



- a. Remove $SA - WA$, delete G from SA .
- b. Remove $SA - V$, delete R from SA , leaving only B .
- c. Remove $NT - WA$, delete G from NT .
- d. Remove $NT - SA$, delete B from NT , leaving only R .
- e. Remove $NSW - SA$, delete B from NSW .
- f. Remove $NSW - V$, delete R from NSW , leaving only G .
- g. Remove $Q - NT$, delete R from Q .
- h. Remove $Q - SA$, delete B from Q .
- i. remove $Q - NSW$, delete G from Q , leaving no domain for Q .

6.12 用**AC-3**算法求解树结构**CSP**在最坏情况下的复杂度是多少？

- 采用逆拓扑序检验，保证每条弧只需要检验一次
- 假设每个顶点的值域最多有 d 个取值，则每条弧检验复杂度为 $O(d^2)$
- 假设有 n 个顶点，则用**AC-3**算法在最坏情况下的复杂度是 $O(nd^2)$