

# UNIT-1

## Introduction to DATA COMMUNICATION and NETWORK MODEL

### Data Communication:-

Data communication systems is effectiveness depends on four fundamental characteristics.

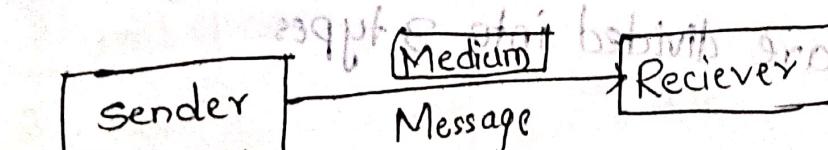
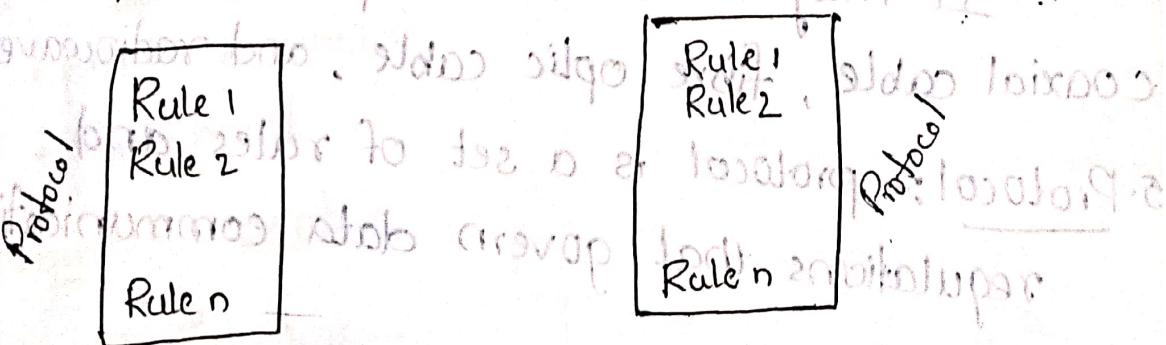
1. Delivery.

2. Accuracy

3. Timeliness

4. Jitter.

### Components:-



1. Sender
2. Reciever
3. Message
4. Medium
5. Protocol

A data communications system has five components.

1. Message :- Text, Numbers, Audio, Video, images etc..

2. Sender :- It may be a computer, work station, telephone handset, video camera etc..

3. Receiver :- Computer, work station, telephone etc..

4. Medium :- It is the physical part to which a message is transmitted from the sender to the receiver.

It may be a twisted-pair wire, coaxial cable, fibre optic cable, and radiowaves.

5. Protocol :- protocol is a set of rules and regulations that govern data communication.

## II DATA FLOW :-

Dataflow are divided into 3 types.

1. Simplex

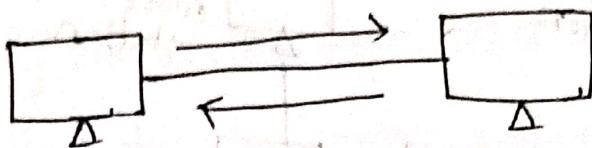
2. Half-Duplex.

3. Full-Duplex.

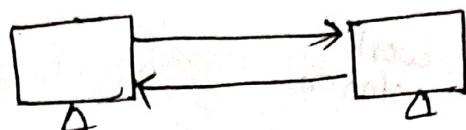
1. Simplex :-



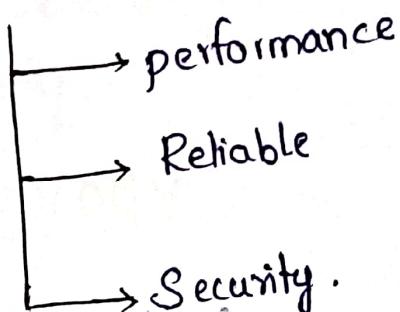
## 2. Half-Duplex:-



## 3. Full-Duplex:-



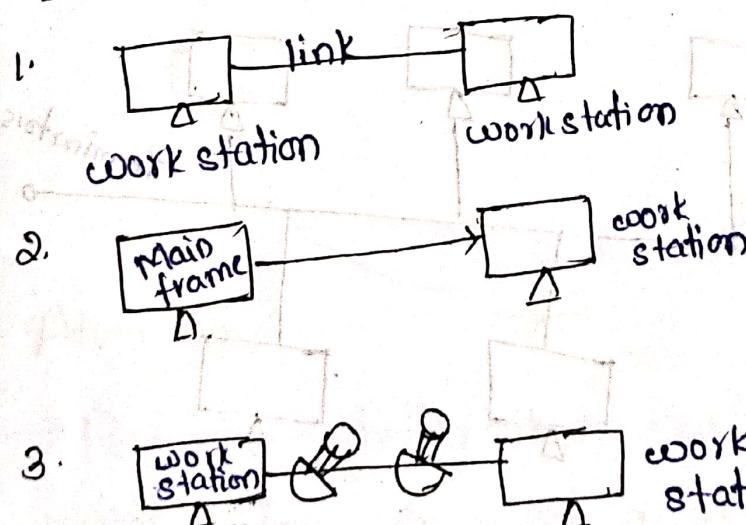
## III NETWORK CRITERIA :-



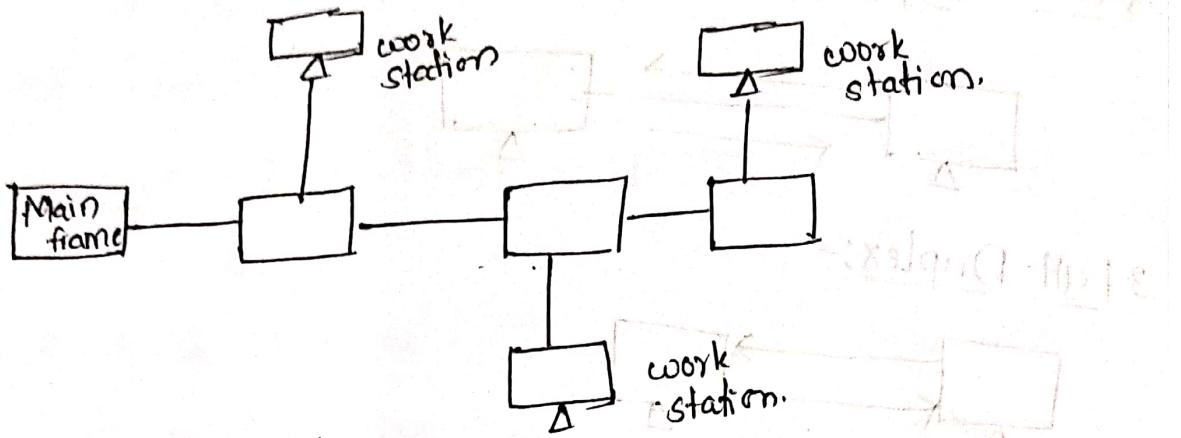
## IV Types of connections:-

- point-to-point
- Multipoint.

### Point-to-point:-



## Multipoint:-



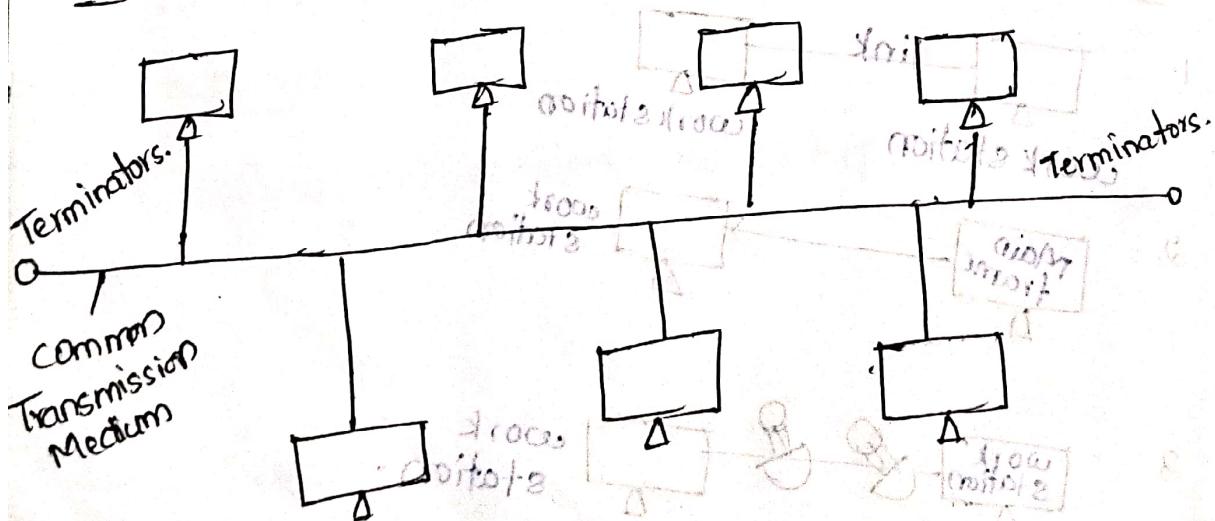
## Physical Topology:-

- 1) Bus Topology
- 2) Ring Topology
- 3) Star Topology
- 4) Mesh Topology
- 5) Hybrid Topology.

→ The Relationship among the two nodes is off two nodes.

- a) peer -to- peer
- b) primary - secondary.

## Bus Topology:-



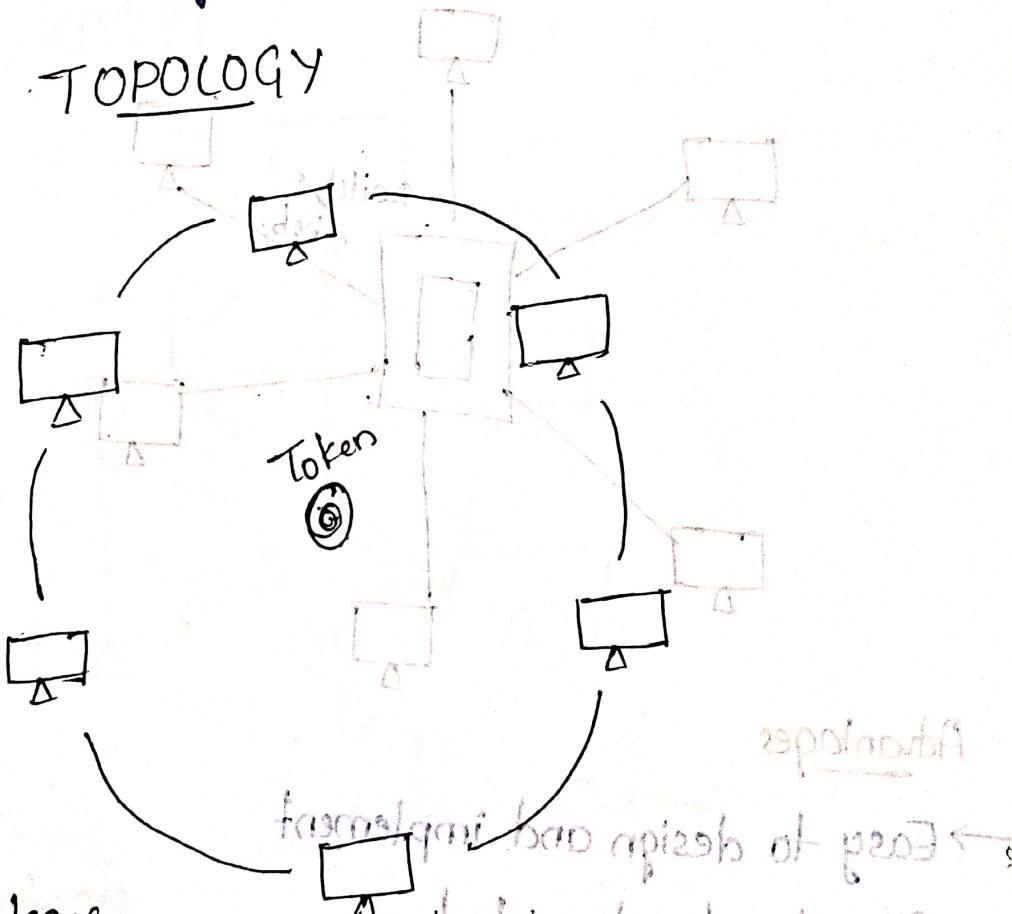
## Advantages of Bus Topology:-

- Only one wire - less expensive.
- Suited for temporary network.
- Node failures does not effect others.

## Disadvantages of Bus Topology:-

- Not fault tolerant
- Limited cable length
- No security.

## RING TOPOLOGY



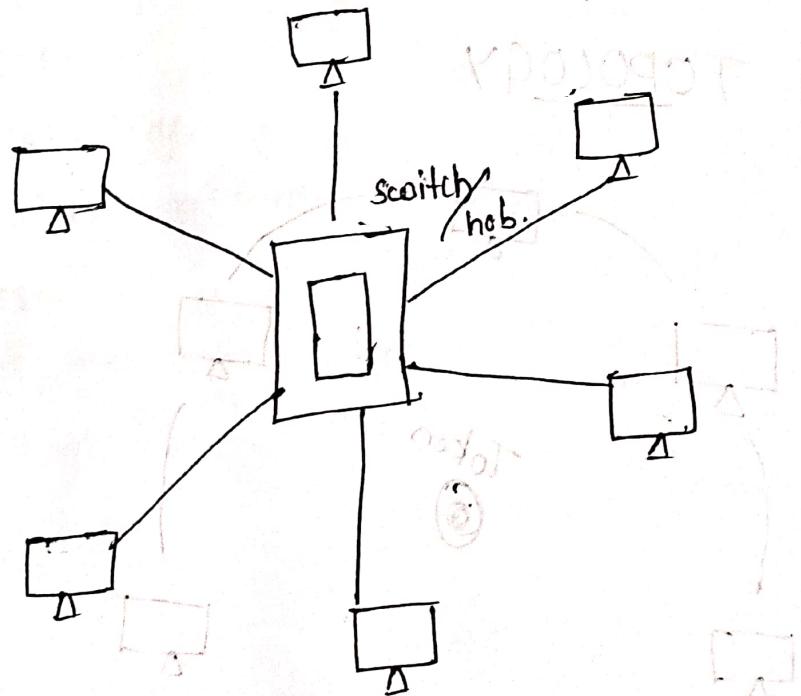
## Advantages

- Performance better than bus topology.
- All nodes with equal access.

## Disadvantages

- Uni directional - single point of failure will effect the whole network.
- Heavy load - low performance.
- No security.

## STAR TOPOLOGY



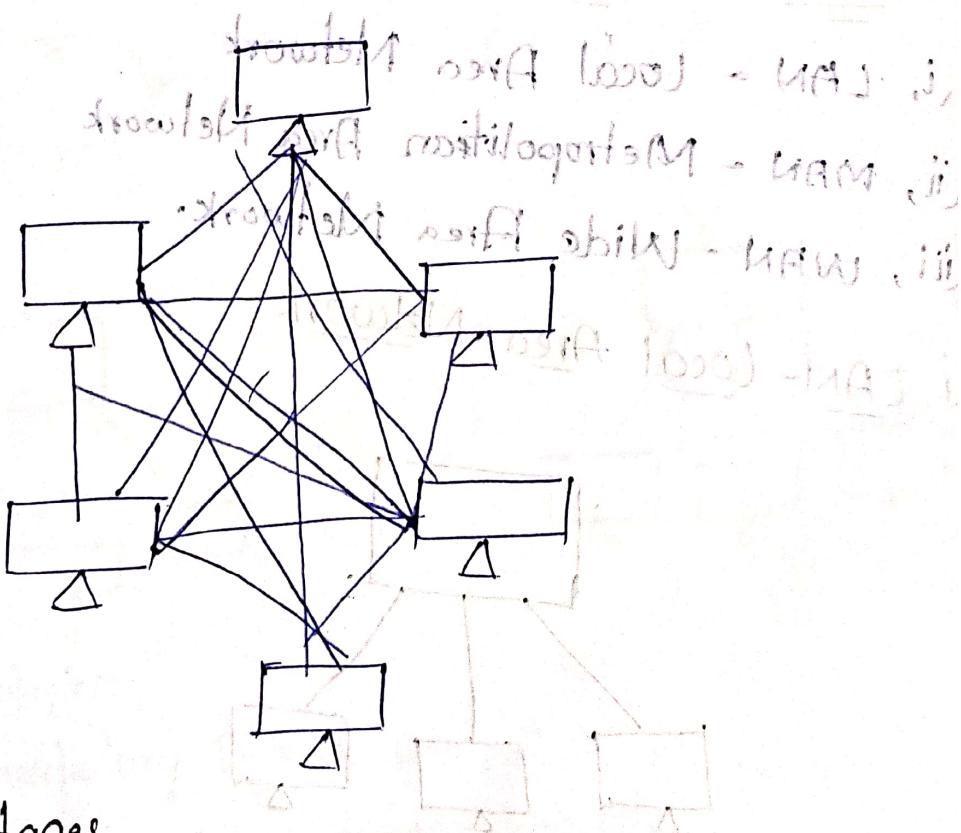
## Advantages

- Easy to design and implement
- Centralized administration
- Scalable

## Disadvantages

- Single point of failure effects the whole network.
  - Communication of performance will be effected when there is overloaded switch/hub.
  - Increased cost due to switch/hub.
- Note:- Extended star topology using repeater

## Mesh Topology



## Advantages

- Fault tolerant
- Reliable

(transmits) besides  
(it-is) also nice

## Disadvantages

- Issues with broadcasting messages.
- Expensive and impractical for large network.

Note:- Broadcast,

↓  
one sender and all receivers

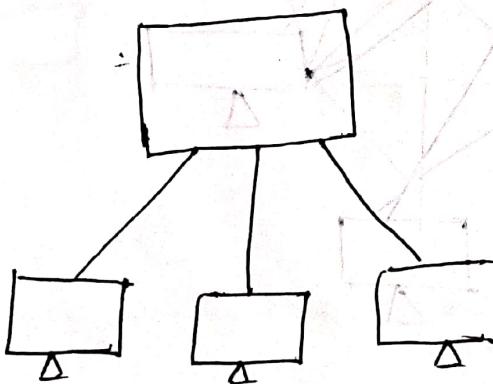
## HYBRID TOPOLOGY

Combination of different topologies

## NETWORK MODEL

- i. LAN - Local Area Network
- ii. MAN - Metropolitan Area Network
- iii. WAN - Wide Area Network.

### i. LAN - Local Area Network



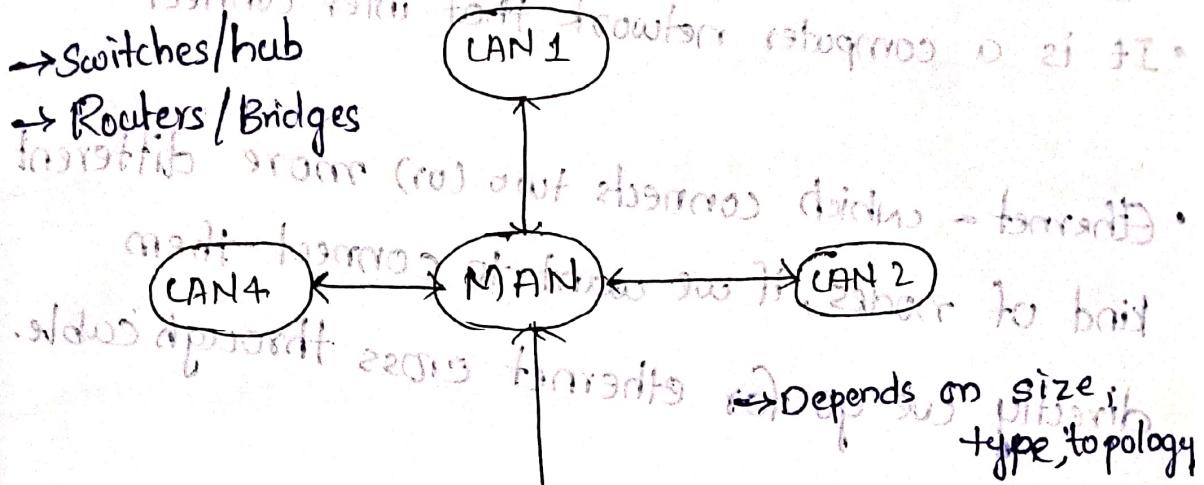
wired (ethernet)

wireless (wi-fi)

→ Depends on size, type and topology.

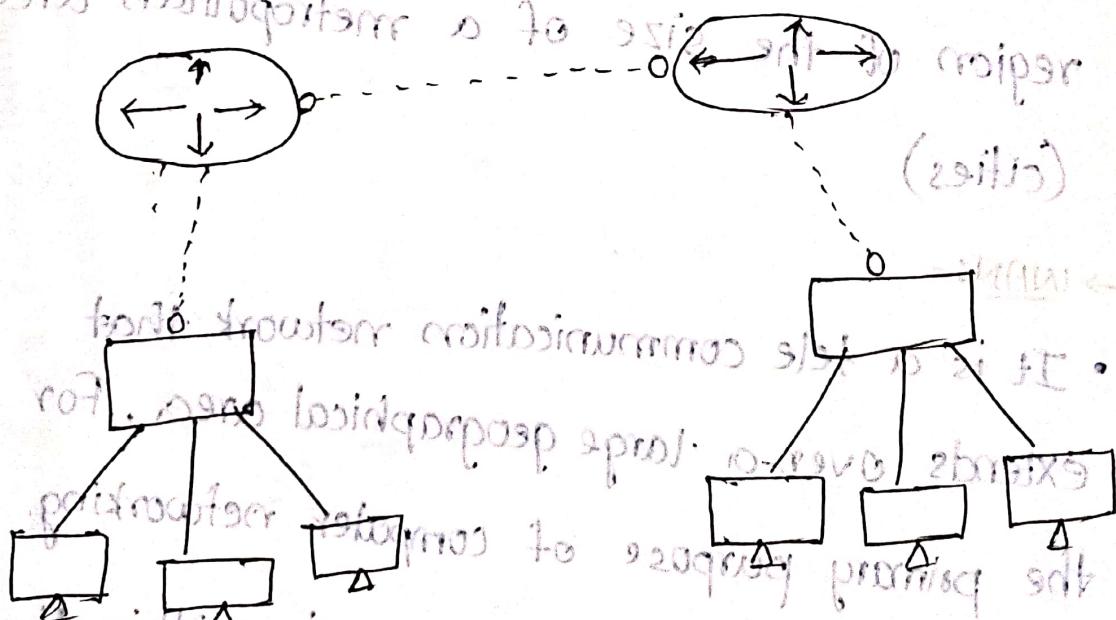
### iii, MAN - Metropolitan Area Network

- Switches/hub
- Routers/Bridges



### iii, WAN - Wide Area Network

- Antilogarithmic to size



→ End devices.

→ Intermediary Devices

→ Depends on sizes, types, topology.

## → CAN

- It is a computer network that inter connects nodes.
- Ethernet - which connects two (or) more different kind of nodes, if we want to connect them directly we go for ethernet cross through cable.

## → MAN

- It is a computer network that inter connects users with computer resources in a geographical region of the size of a metropolitan area (cities).

## → WAN:-

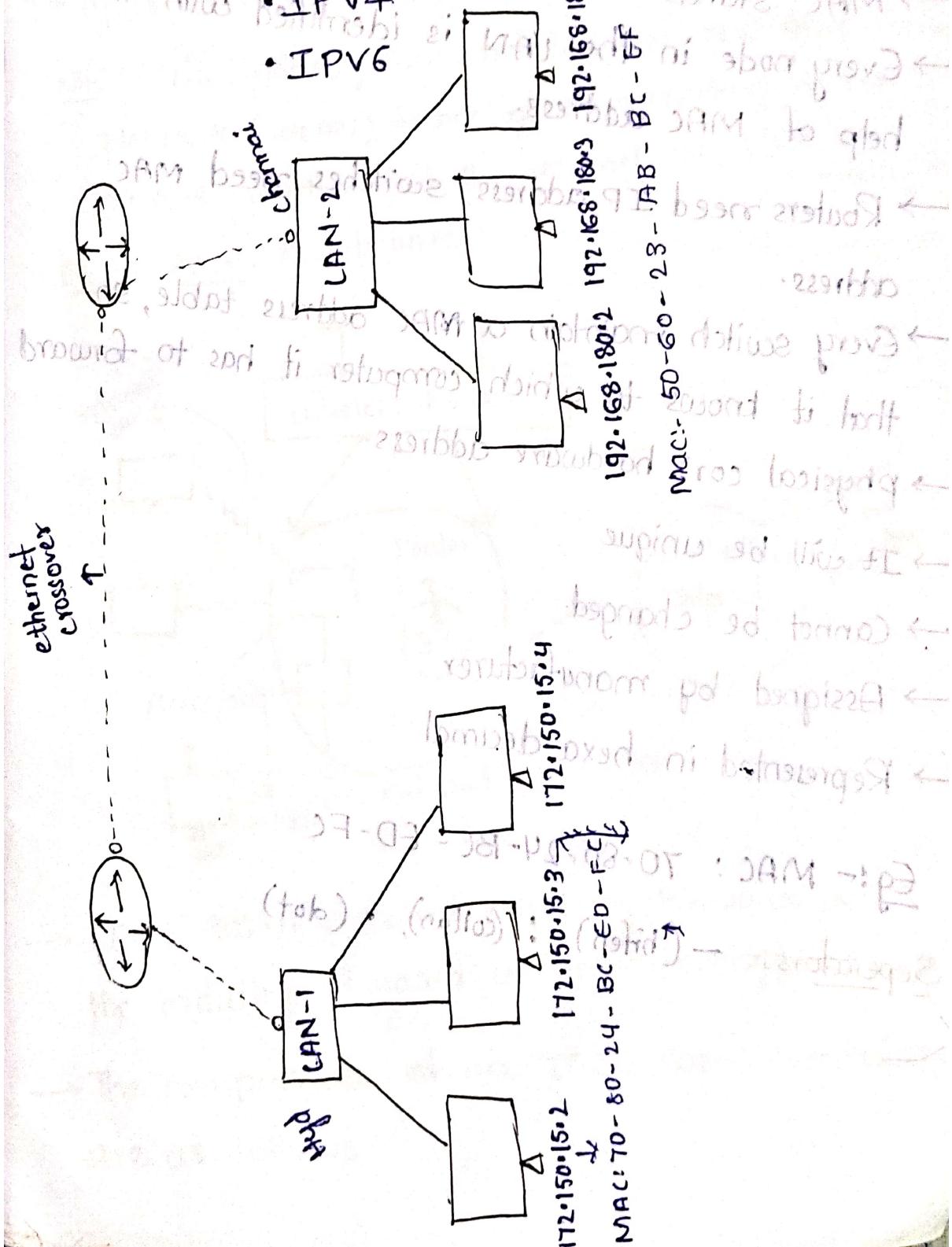
- It is a telecommunication network that extends over a large geographical area, for the primary purpose of computer networking.
- In WAN two (or) more LAN's within the country can communicate each other.

## BASICS OF IP ADDRESSING

- IP stands for Internet Protocol and it has a range of 0.0.0.0 to 255.255.255.255
- Every node in the computer Network is identified with the help of IP address
- IP addresses are of two types:

• IPV4

• IPV6



- IPV4 addresses are represented in decimals  
and it has four octets (x.x.x.x)  
Each of x values takes from 0.0.0.0 to 255.

Basics of MAC light out to Access Control Unit.  
→ MAC stands for Medium Access Control  
→ Every node in the LAN is identified with the help of MAC address.

- Routers need IP address, switches need MAC address.  
→ Every switch maintains a MAC address table, so that it knows to which computer it has to forward physical or hardware address.  
→ It will be unique.  
→ Cannot be changed.  
→ Assigned by manufacturer  
→ Represented in hexa-decimal

Eg:- MAC : 70-80-24-BC = ED-FC

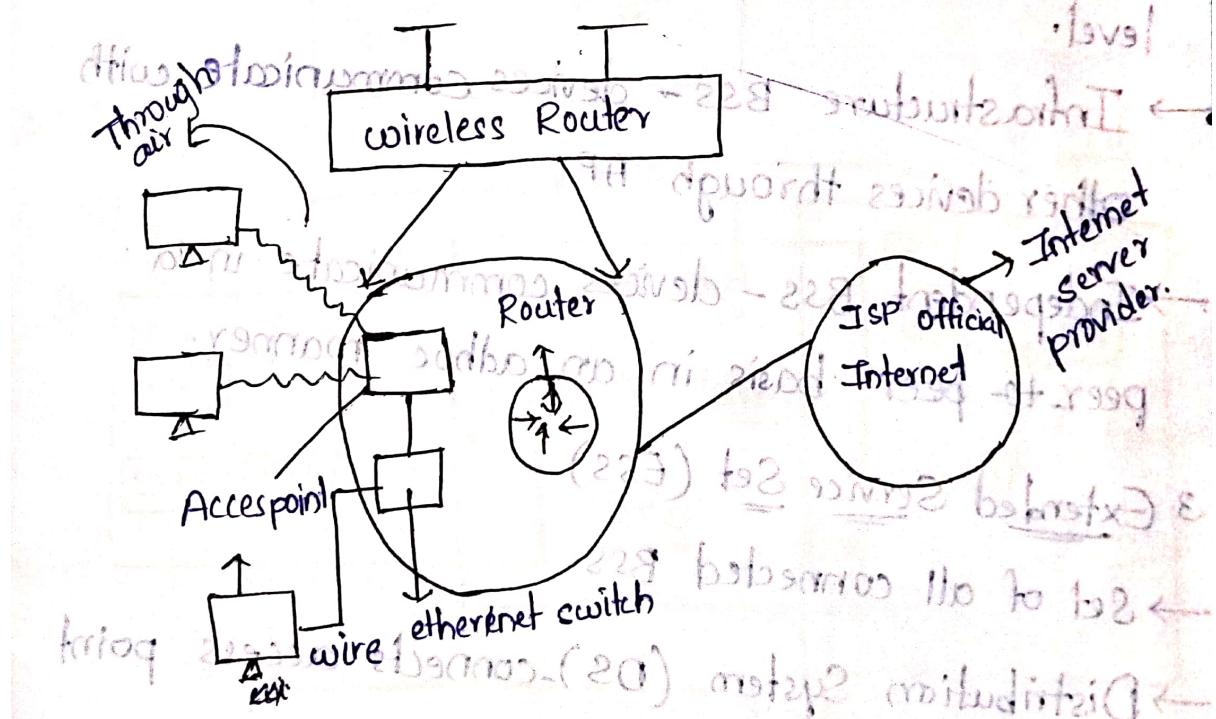
Separators:- - (hifen) : (colon), • (dot)

## Wireless Network

- Wireless is the connecting devices without wire is made up of electromagnetic wave (or) infrate waves.
- The wireless devices contains antennas (or) sensors
- Types of Networks:

Wifi:- The industry standard wireless Local Area Network (WLAN) technology for connecting devices to one another and the internet wi-fi (wireless fidelity).

### IEEE 802.11 Architecture



→ IEEE 802.11 standard, popularly known as wi-fi the industry standard WLAN technology.

→ The components of an IEEE 802.11 architecture are as follows.

1. Station (sta) comprises all devices and equipment that are connected to the wireless land.

→ Station can be of two types (WAP or) AP (wireless access point (or) Access point)

→ Wireless Access points are generally wireless routers that form the base station (or) access.

→ Clients: Workstation, laptops, computers, printers etc.

2. Basic Service Set (BSS):

→ Group of stations communicating at physical layer level.

→ Infrastructure BSS - devices communicate with other devices through AP.

→ Independent BSS - devices communicate in a peer-to-peer basis in an adhoc manner.

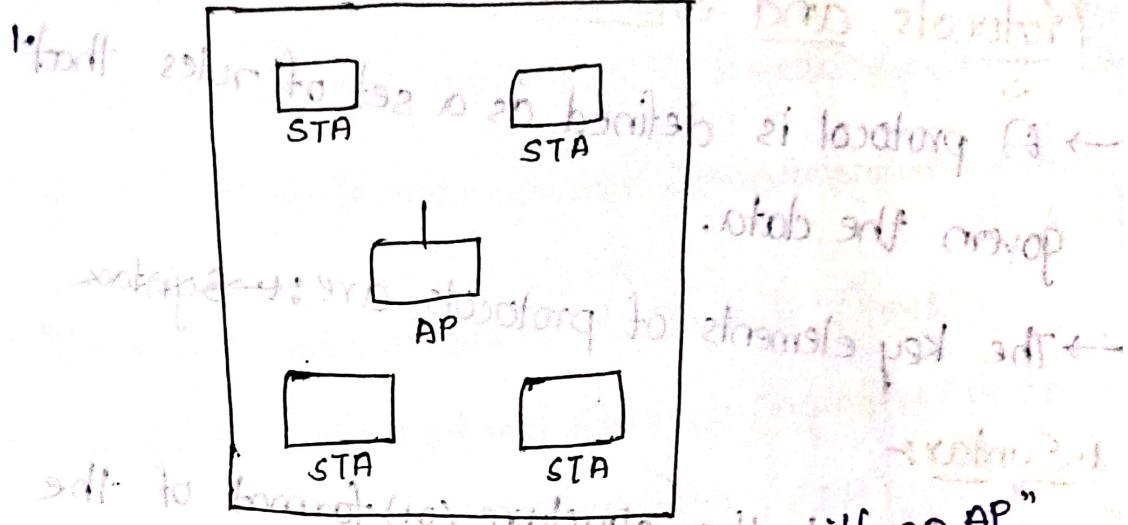
3. Extended Service Set (ESS)

→ Set of all connected BSS.

→ Distribution System (DS) connects access point in ESS.

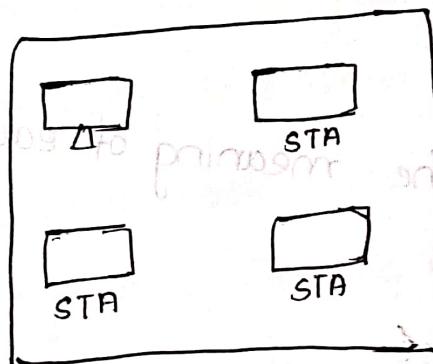
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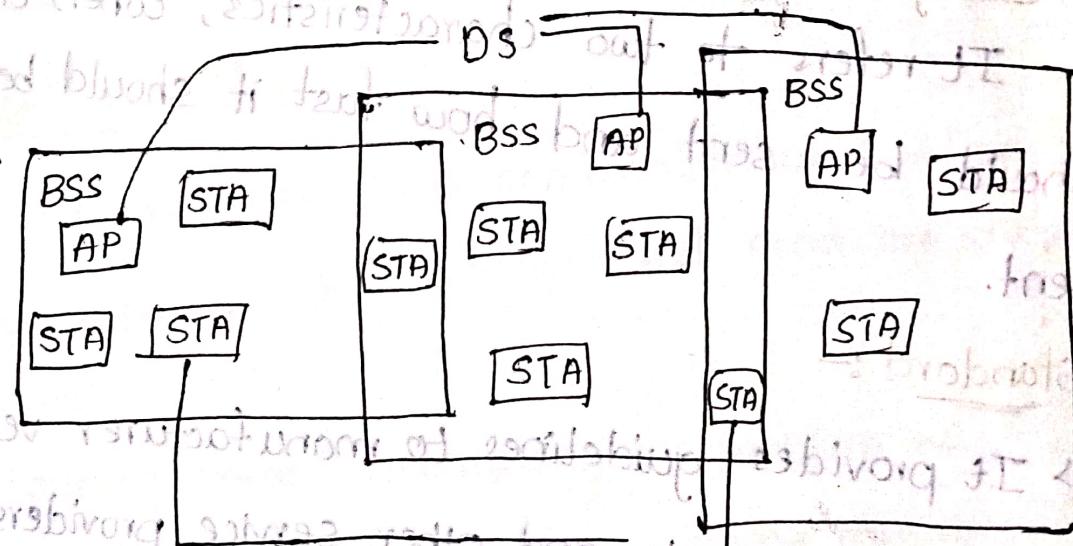
Infrastructure of Network "BSS's with an AP"

2.



stations across (sites) of band

3.



no base station connection or communication between BSS

multiple stations can connect to one AP

(co-channel PS (or) "foot PS") & (DIFS)

(co-channel PS (or) "wol PS") & (DIFS)

## Protocols and Standards

→ A protocol is defined as a set of rules that govern the data.

→ The key elements of protocols are :- Syntax

1. Syntax :-

It refers to the structure (or) format of the data.

2. Semantics :-

It refers to the meaning of each section of bits

3. Timing :-

It refers to two characteristics, when data should be sent and how fast it should be sent.

Standard :-

→ It provides guidelines to manufacturer vendor, government agencies and other service providers.

Data communication fall into two categories.

i. Defactor ("By fact" (or) By convention)

ii. De-jure ("By law" (or) By regulation)

the organisations which maintain standards:-

1. International organisational for standardisation (ISO)
2. International Telecommunication Union-telecommunication (ITU-T)
3. American National Standards Institute (ANSI)
4. Institute of electrical and Electronic Engineers (IEEE)
5. Electronic Industries Association (EIA)

### Internet Standards Organisation:-

1. Internet Engineering Task Force (IETF)
2. Internet Research Task Force (IRTF)
3. Internet Engineering Steering Group (IESG)
4. Request for Comments (RFC)

### Network Models:-

A network is a combination of hardware and software that receives and sends the data.

### Layer:-

Layering means decomposing the problem of building a network into more manageable components.

- Higher layer
- Middle layer
- Lower layer.

## Connection oriented and connectionless oriented services

Layers can offer two types of services

Connection-oriented service.

1. Connection-oriented service.

2. Connectionless oriented service.

1. Connection oriented service:-

- Establish the connection

- Use the connection

- Release the connection

2. connectionless oriented service

Connectionless service is much like same as

the postal service

3. Quality of service:-

- Reliable
- Unreliable

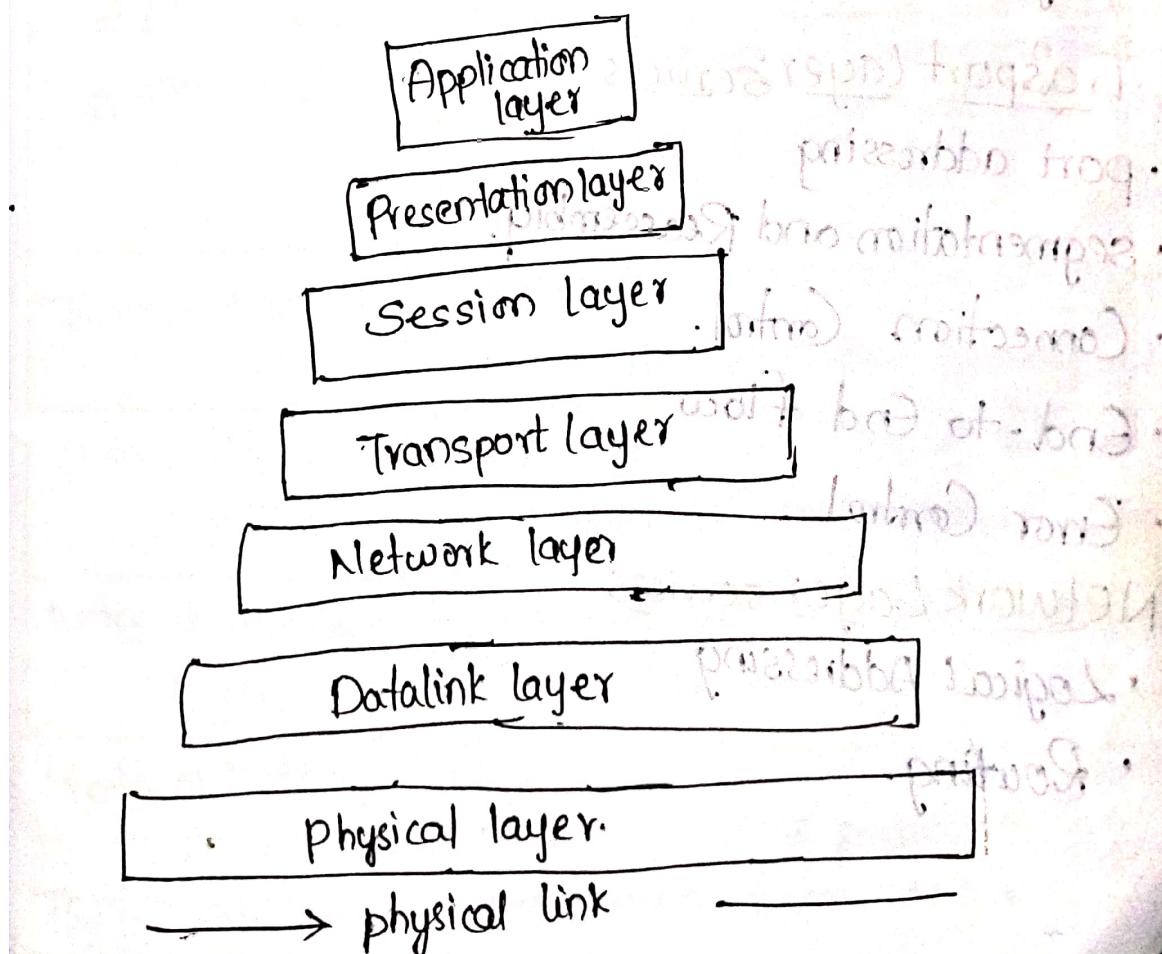
## Service Primitives:-

A primitive means operation. A service is described by a set of primitives.

- Listen
- Connect
- Send
- Receive
- Disconnect

## OSI Model

The purpose of the OSI Model is to facilitate communication between different systems without requiring changes to the logic of the underlying hardware and software.



## Services provided by the Application Layer

- ① FTAM (file Transfer and Access Management)
  - It allows to send the data to remote computer or receive the data from the remote computer.
- ② Mail services (allows to access email services)
- ③ Directory services (provides access to data globally).

## Presentation Layer Services :

Translation.

Encryption

Compression

Session Layer services

- Dialog control
- Synchronization.

## Transport Layer Services

- port addressing
- segmentation and Reassembly.
- Connection Control.
- End-to-End flow
- Error Control

## Network Layer Services

- Logical Addressing
- Routing.

## Datalink layer services

- Framing
- physical Addressing
- flow control
- Error control
- Access Control

## Physical layer services

physical characteristics of the media

- physical characteristics of bits
- Representation of bits
- Data rate (Transmission layer) Rate
- Synchronization of bits.
- Line configuration — point to point (or) point to multipoint
- physical topology (star, bus, ring, mesh, hybrid)
- Transmission mode (baseband, narrow band, broad band)

TCP/IP layers	Functionalities.
Application layer	provides interface between users and network to access network resources.
Transport layer.	Segmentation, process to process communication, error free delivery.
Internet layer	Moving packets from source to destination, provides routing.
Network interface layer.	Control the hardware and media that constitute the network.

## Protocol Data unit (PDU):—

PDU's are named according to the protocols of the TCP/IP suite, data, segment, packet, frame and bits.

Application layer - Data  
Transport layer - Segment  
Network layer - packet  
Datalink layer - frames  
Physical layer - Bits.

## DATA AND SIGNALS:-

### signals:-

- It is a function that represents the variation of a physical quantity with respect to time. Two types of signals:
  - Analog signal
  - Digital signal.

Analog signals have continuous electrical signals while digital signals have non-continuous electrical signals.

### 1. Analog signal:-

→ It is the signal that can take any value in the defined range.

→ All real life signals are analog in nature.

Ex: Human's voice.

### 2. Digital Signal:-

→ It is the signal that can take one of the finite values at any given time.

→ In case of digital signals, we discretize both time and magnitude.

Ex:- Files in a disc.

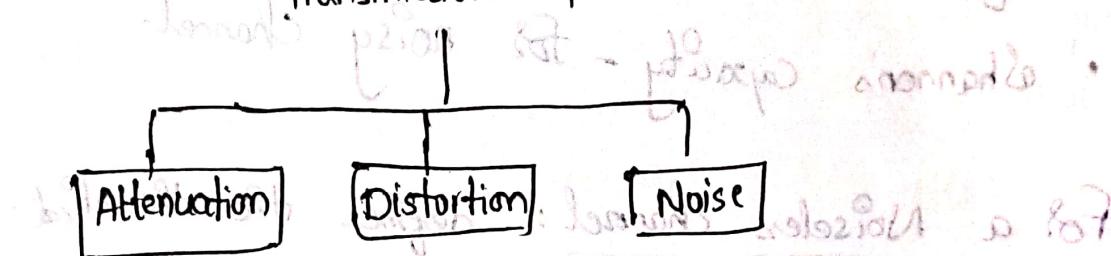


## Transmission impairment:-

When the received signal is not identical to the transmitted one due to the transmission impairment. The quality of the signals will get destroyed due to transmission impairment.

## Causes of Transmission Impairment

### Transmission Impairment



### Attenuation:-

It means loss of energy.

### Distortion:-

Distortion occurs where there is a change in shape of signal.

Signal - when data is travelled over a transmission medium, some

Noise - when data is travelled over a transmission medium, some unwanted signal is added to it which creates the signal.

### Data rate:-

Refers to the speed of data transfer a channel. It is generally computed in bits per second (bps). Higher data rates are expressed as

Kbps ("kilo" bits per second i.e 1000 bps) Mbps

("Mega" bits per second, i.e, 1000 Kbps) Gbps ("Giga" bits per second)

1. Bandwidth of the channel

2. Number of levels of signals that are used.

3. Noise present in the channel.

- Data rate can be calculated using two theoretical formulas:
- Nyquist Bit Rate - for noiseless channel.
  - Shannon's Capacity - for noisy channel.

For a Noisless channel: defines the theoretical maximum bit rate.

- The theoretical formula for the maximum bit rate is: maximum bit rate =  $2 \times \text{Bandwidth} \times \log_2 V$
- Here, maximum bit rate is calculated in bps.
- Bandwidth (The difference between the maximum frequency and the minimum frequency) is the bandwidth of the channel.
  - $V$  is the number of discrete levels in the signal.
- called the shannon capacity to determine the theoretical highest data rate for a noisy channel
- $$\text{capacity} = \text{Bandwidth} \times \log_2 (1 + \text{SNR})$$

→ Here, capacity is the maximum data rate of the channel in bps.

Bandwidth is the bandwidth of the channel

SNR is the signal-to-noise ratio.

→ SNR is actually the ratio of what is wanted (signal-to) what is not wanted (noise) a high SNR means the signal is less corrupted by noise. A low SNR means the signal is more corrupted by the noise.

→  $\text{SNR} = \frac{\text{Average signal power}}{\text{Average noise power}}$

### Performance:-

The examination and review of collective network information and review of collective network information to describe the quality of services delivered by the underlying computer network is known as "network performance". Hence, to measure the performance of a network here are the major factors to be considered:

- Transit time
- Response time.
- Bandwidth.

### Band width in bps

Bandwidth = capacity.

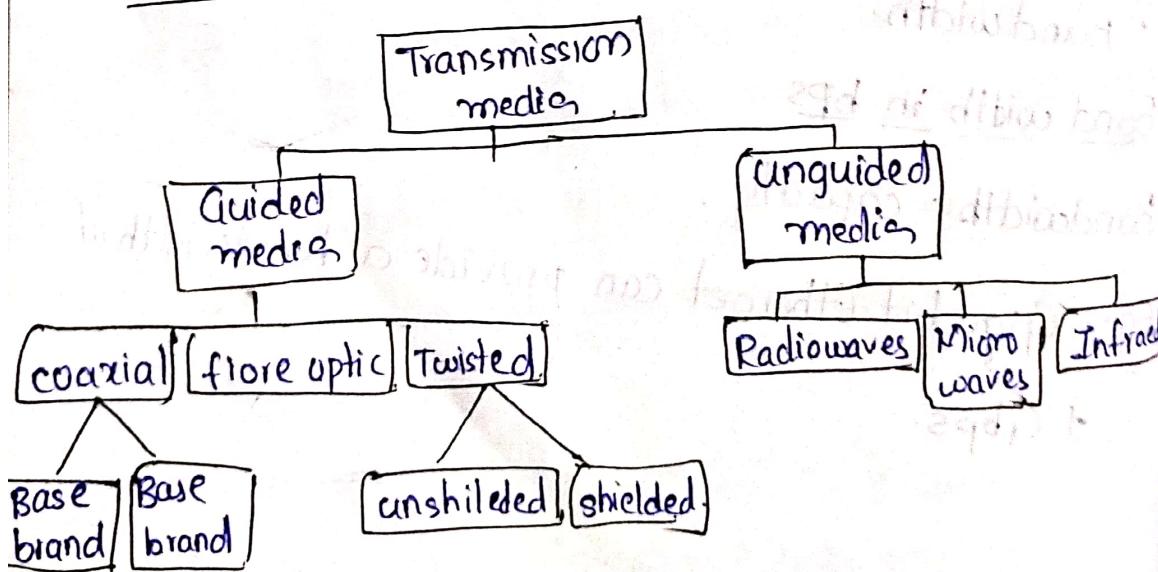
Eg:- Gigabit Ethernet can provide a bandwidth of 1 Gbps.

Bandwidth in Hertz is often referred to as frequency.  
→ A range of frequencies used to transmit signals which is measured in hertz.  
→ A link may have a bandwidth of 'B' bps but we can only send 'T' bps through this link with  $T \leq B$  always.

Delay / latency :- As we discussed, Throughput is the number of data packets successfully delivered in a given time. Delay is the measure of time taken to do the delivery.

The latency or delay defines how long it takes for an entire message to completely arrive at the destination from the time taken to the delivery.

Classification of Transmission media :-



## Guided Media:-

It is also referred to as wired or bounded transmission media signals being transmitted are directed and confined in a narrow pathway by using physical links.

### Features:-

→ High speed

→ Secure transmission

→ Used for comparatively shorter distance.

### Types of Guided Media:-

1. Twisted pair-

Copper wires are the most common wires used for transmitting signals because of good performance at low costs. However, if two or more wires are lying together they can interfere with each other's signals. To reduce

this electromagnetic interference, pair of copper wires are twisted together in a helical shape like a DNA molecule.

A twisted pair consists of two insulated copper

wires arranged in a regular spiral pattern.

### unshielded twisted pair (UTP)

UTP is an unshielded twisted pair cable used in computer and telecommunication medium.

### Advantages:-

→ It is cheap

→ Installation of the unshielded twisted pair is easy.

→ It can be used for high speed LAN.

### Disadvantages:-

- This cable can only be used for shorter distances because of attenuation.
- lower capacity and performance in comparison to STP

### Shielded twisted pair (STP)

A STP is a type of twisted pair cable that contains an extra wrapping foil or copper braid jacket to protect the cable from defects like cuts, noise and signal to the interference.

### Advantages:-

- It has lower noise and attenuation than UTP
- It is shielded with a plastic cover that protects the STP cable from a harsh environment and increases the data transmission rate.

### Disadvantages:-

- It is more expensive as compared to UTP and coaxial cable.
- It has a higher attenuation rate.

## Coaxial Cable:-

- Coaxial cable is very commonly used transmission media, for example, TV wire is a coaxial cable.
- The name of the cable is coaxial as it contains two conductors parallel to each other.
- It has a higher frequency limit than P.E.

Coaxial cable is two types:-

1. Base band transmission

a. Broad band transmission.

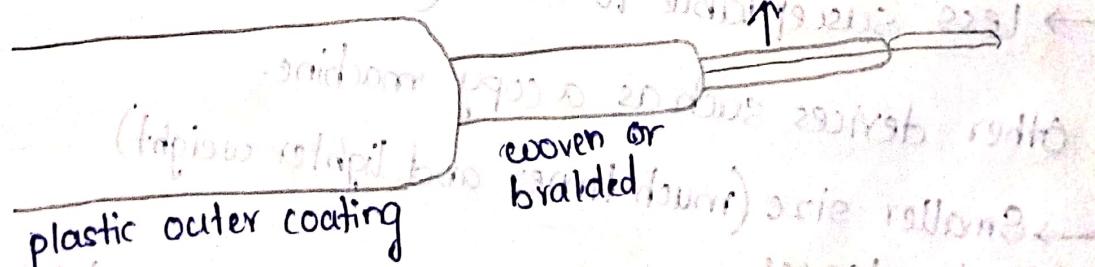
b. Baseband transmission:-

It is defined as the process of transmitting

a single signal at high speed.

2. Broad band transmission:-  
It is defined as the process of transmitting

multiple signals.



Advantages:-

- The data can be transmitted at high speed.
- It has better shielding as compared to twisted pair.

## Twisted pair cable:-

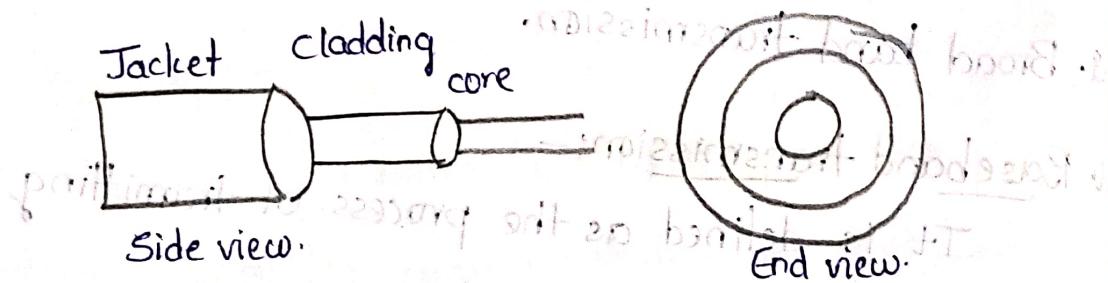
→ It provides higher bandwidth.

## Disadvantages:-

→ It is more expensive as compared to twisted pair cable.

→ If any fault occurs in the cable causes the failure in the entire network.

## Fibre Optic Cable:-



## Advantages:-

→ Capability of carrying significantly more signals than wire cable.

→ Faster data transmission.

→ Less susceptible to noise (interference) from other devices such as a copy machine.

→ Smaller size (much thinner and lighter weight)

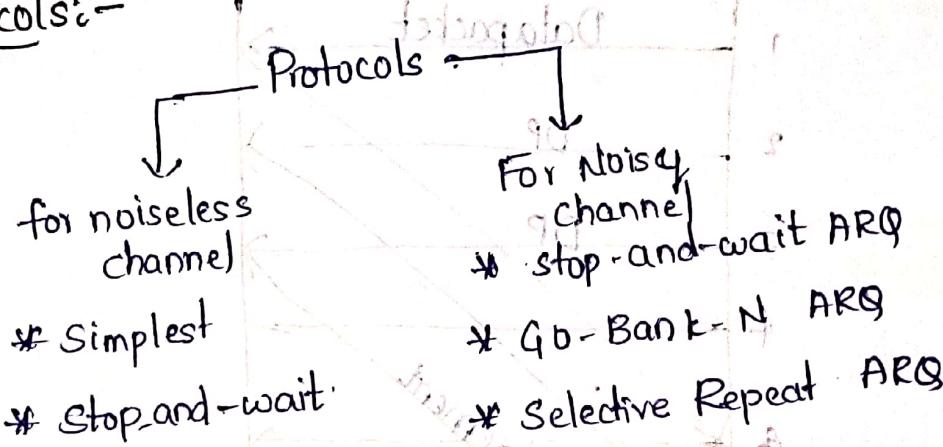
## Disadvantages:-

→ Fibre optic cables are more expensive than twisted pair or coaxial cable and can be difficult to install and modify.

## Error Control:-

→ Error control is a combination of both error detection and error correction it ensures that the data received at the receiver end in the same as the one send by the sender.

## Protocols:-



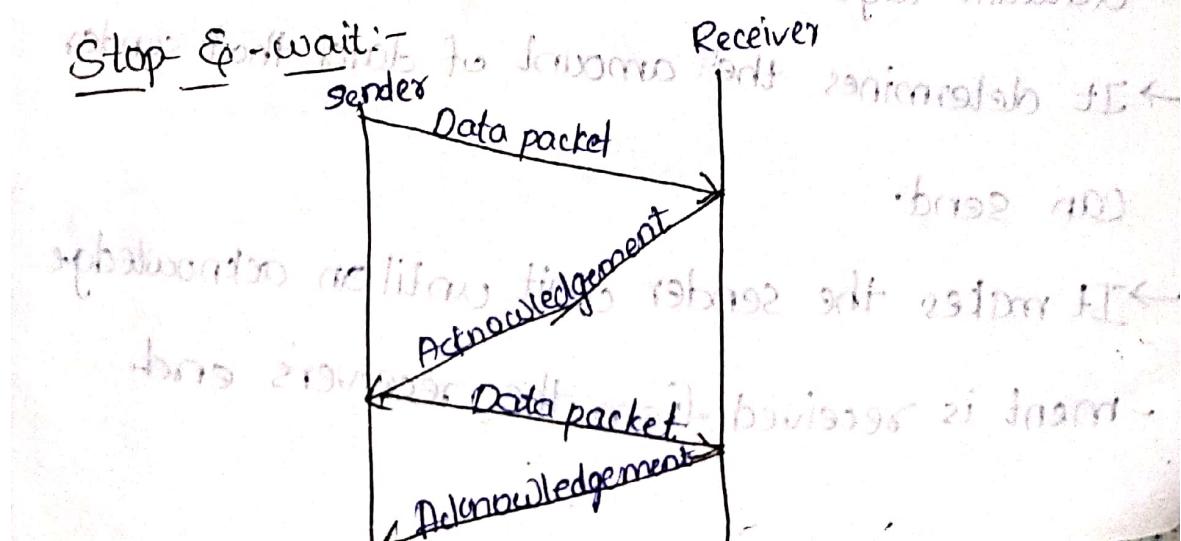
## ARQ (Automatic Repeat Request)

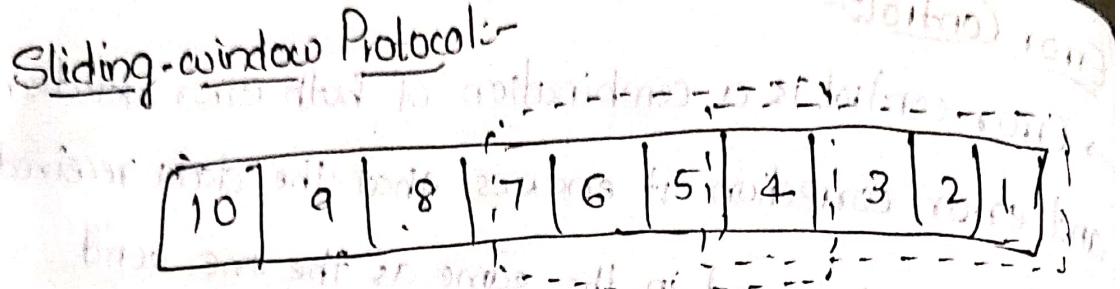
Fig:- Taxonomy of protocols.

## Flow Control:-

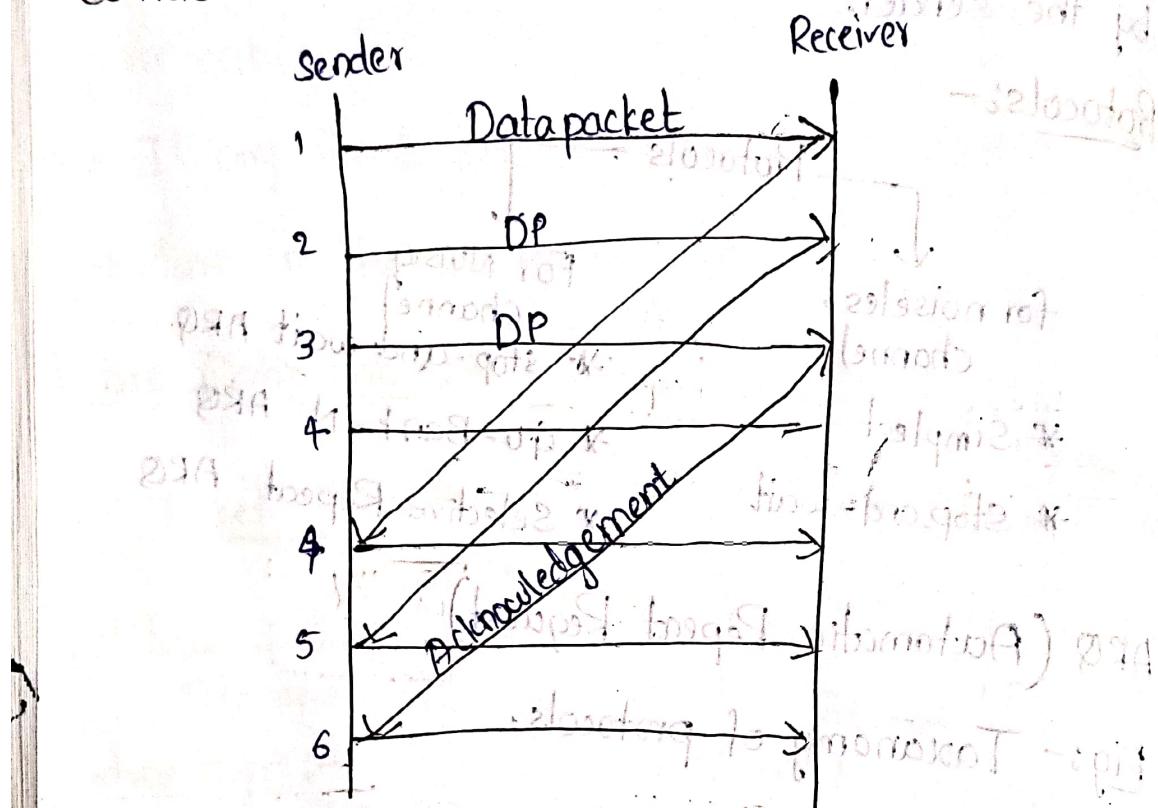
- stop and wait protocol (sends only one frame at a time)
- sliding window protocol (sends several frames at a time)

### Stop & wait:-





window size = 4



### Flow control:-

→ Flow control is essential function of the datalink layer.

→ It determines the amount of data that sender can send.

→ It makes the sender wait until an acknowledgement is received from the receiver's end.

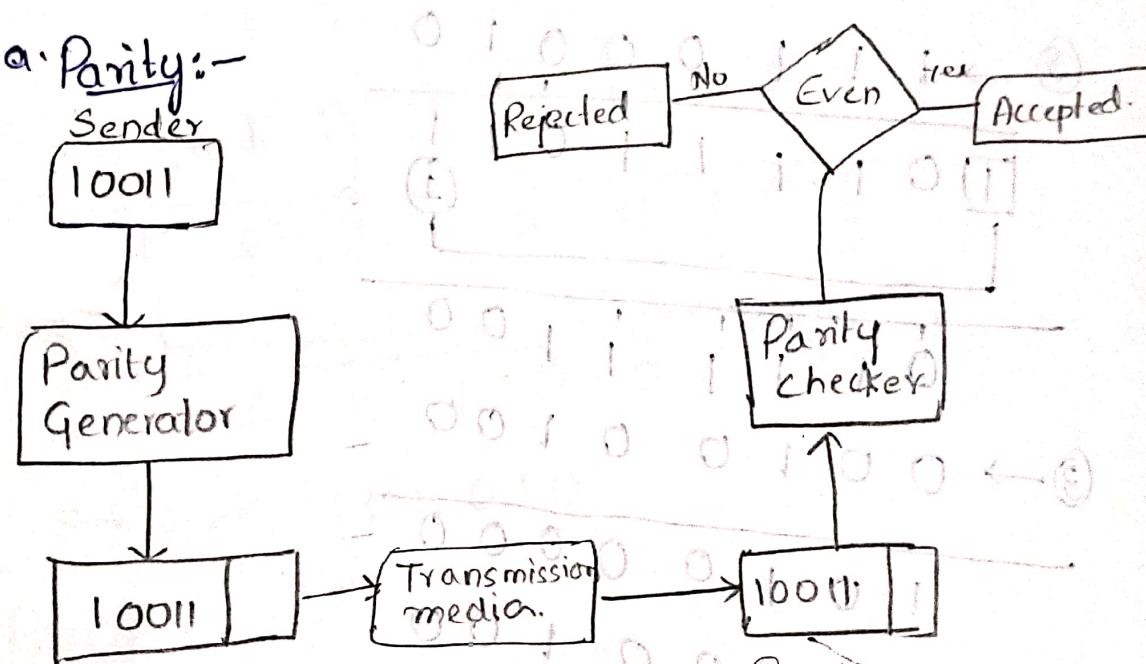
→ Methods of flow control are stop and wait, sliding window.

## Error Correction and Detection:-

## 1. Error Detection:-

- a. Parity
  - b. Checksum
  - c. CRC (cyclic Redundancy check)

a. Parity:-



Note:- Rule:  $P_1$  - check one bit, skip one bit

$p_i \rightarrow$  check 1 bit skip 1 bit

1, 3, 5, 7 - - -

$B_2 \rightarrow$  check 2 bit skip 2 bit

2, 3, 6, 7, 10, 11 . . .

$P_4 \rightarrow$  check 4 bit skip 4 bit

9, 5, 6, 7, 12, 13, 14, 15, . . .

Checksum: In what follows, we'll go through

10011001	11100010	00100100	10000100
----------	----------	----------	----------

$m=8$  (bits)

$k=4$  (parts)

Sender side :-

$$\textcircled{1} \rightarrow 1\ 0\ 0\ 1\ 1\ 0\ 0\ 0$$

$$\begin{array}{r} \textcircled{2} \rightarrow 1\ 1\ 1\ 0\ 0\ 0\ 1\ 0 \\ \hline 1\ 0\ 1\ 1\ 1\ 1\ 0\ 1 \\ \hline 0\ 1\ 1\ 1\ 1\ 1\ 0\ 0 \end{array}$$

$$\textcircled{3} \rightarrow 0\ 0\ 1\ 0\ 0\ 1\ 0\ 0$$

$$\begin{array}{r} \textcircled{4} \rightarrow 1\ 0\ 0\ 0\ 0\ 1\ 0\ 0 \\ \hline 1\ 0\ 1\ 0\ 0\ 0\ 0\ 0 \end{array}$$

$$\begin{array}{r} \textcircled{4} \rightarrow 1\ 0\ 0\ 0\ 0\ 1\ 0\ 0 \\ \hline 1\ 0\ 0\ 1\ 0\ 0\ 1\ 0 \end{array}$$

$$\text{sum} = 0\ 0\ 1\ 0\ 0\ 1\ 0\ 1$$

$$\text{checksum} = 1\ 1\ 0\ 1\ 0\ 0\ 1\ 0$$

$\textcircled{1} \rightarrow$	1	0	0	1	1	0	0	1
$\textcircled{2} \rightarrow$	1	1	1	0	0	0	1	0
$\boxed{\textcircled{1}}$	0	1	1	1	1	0	1	1
$\boxed{\textcircled{2}}$	1	0	1	1	1	0	1	1
$\textcircled{3} \rightarrow$	0	0	1	0	0	0	0	0
$\textcircled{4} \rightarrow$	1	0	0	0	0	1	0	0
$\boxed{\textcircled{1}}$	0	0	1	0	0	1	0	0
$\boxed{\textcircled{2}}$	1	0	0	1	0	0	1	0
$\boxed{\textcircled{3}}$	0	0	1	0	0	0	1	0
$\boxed{\textcircled{4}}$	1	0	0	0	0	1	0	0

## Receiver

checksum CRC		0	1	1	0	0	1	0
sender	101101000	0	0	0	0	1	1	0
	1001	↓	↓	↓	↓	↓	↓	↓
	01000							
	1001	↓	↓	↓	↓	↓	↓	↓
	0001100							
	1001	↓	↓	↓	↓	↓	↓	↓
	01010							
	1001	↓	↓	↓	↓	↓	↓	↓
	00111							
	101101000	0	0	1	0	0	1	0
	101101000	0	0	1	0	0	1	0
	101101011	0	0	1	0	0	1	0

## Error Corrections

### 1. Hamming Codes:-

D <sub>7</sub>	D <sub>6</sub>	D <sub>5</sub>	P <sub>4</sub>	D <sub>3</sub>	P <sub>2</sub>	P <sub>1</sub>
1	0	1	0	1	0	1

$$P_1 \rightarrow C_1, S_1$$

1, 3, 5, 7, ...

$$P_2 \rightarrow C_2, S_2$$

2, 3, 6, 7, 10, 11

$$P_4 \rightarrow C_4, S_4$$

4, 5, 6, 7, 12, 13, 14, 15, ...

$$P_1 \rightarrow D_3$$

D <sub>5</sub>	D <sub>7</sub>
1	1

$$P_1 \rightarrow 1$$

$$P_2 \rightarrow 0$$

$$P_4 \rightarrow 0$$

## HDLC

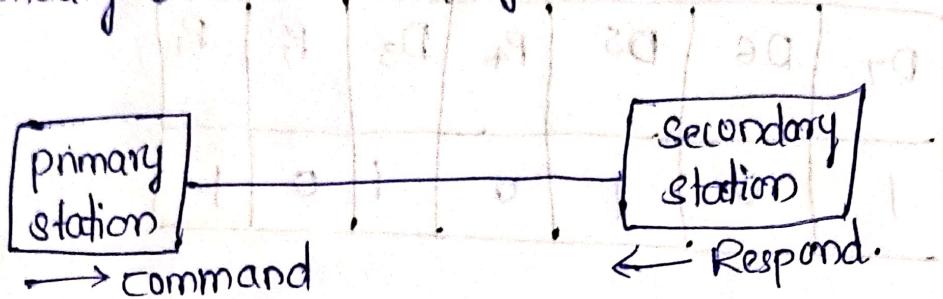
High level Data link control (HDLC)

→ HDLC is a bit oriented protocol for communication

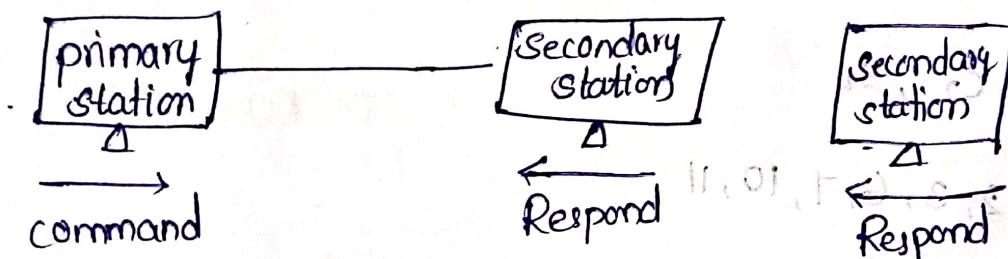
over point-to-point and multi-point links.

## point - to - point:-

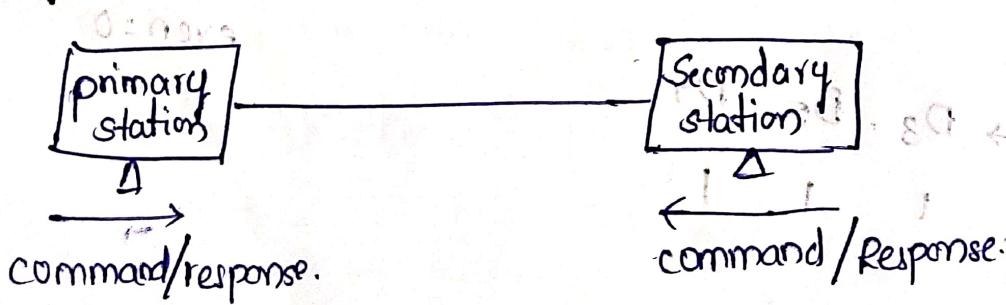
A primary station can send commands or  
secondary station can only respond.



## Multi-points:-

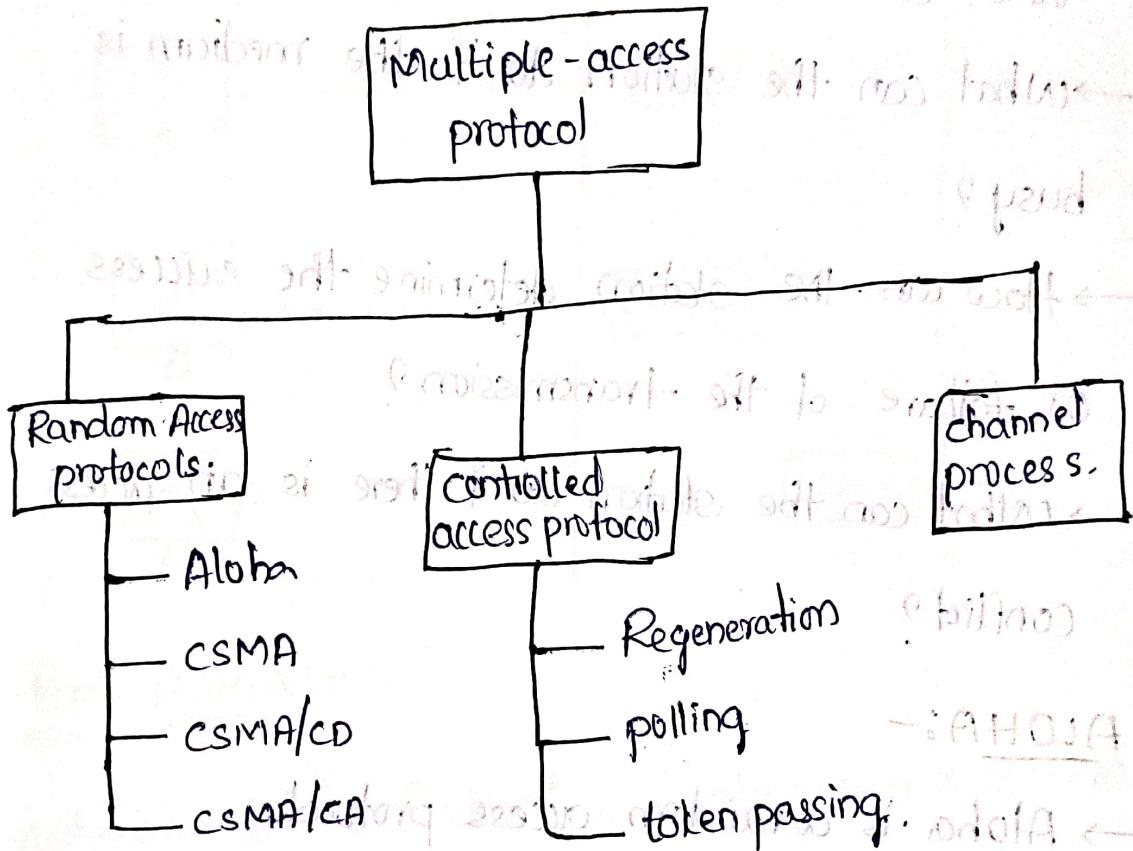


## Asynchronous balanced mode:-



## Multiple Access Schemes, wired LAN and

### wireless LAN

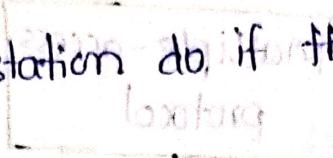
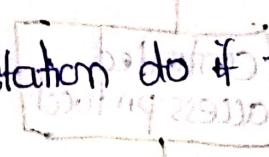


### Random Access:

→ In this, all stations have same superiority that is no station has more priority than another. Any station can send data depending on medium state (idle or busy).

→ Each station has the right to the medium without being controlled by any other station.

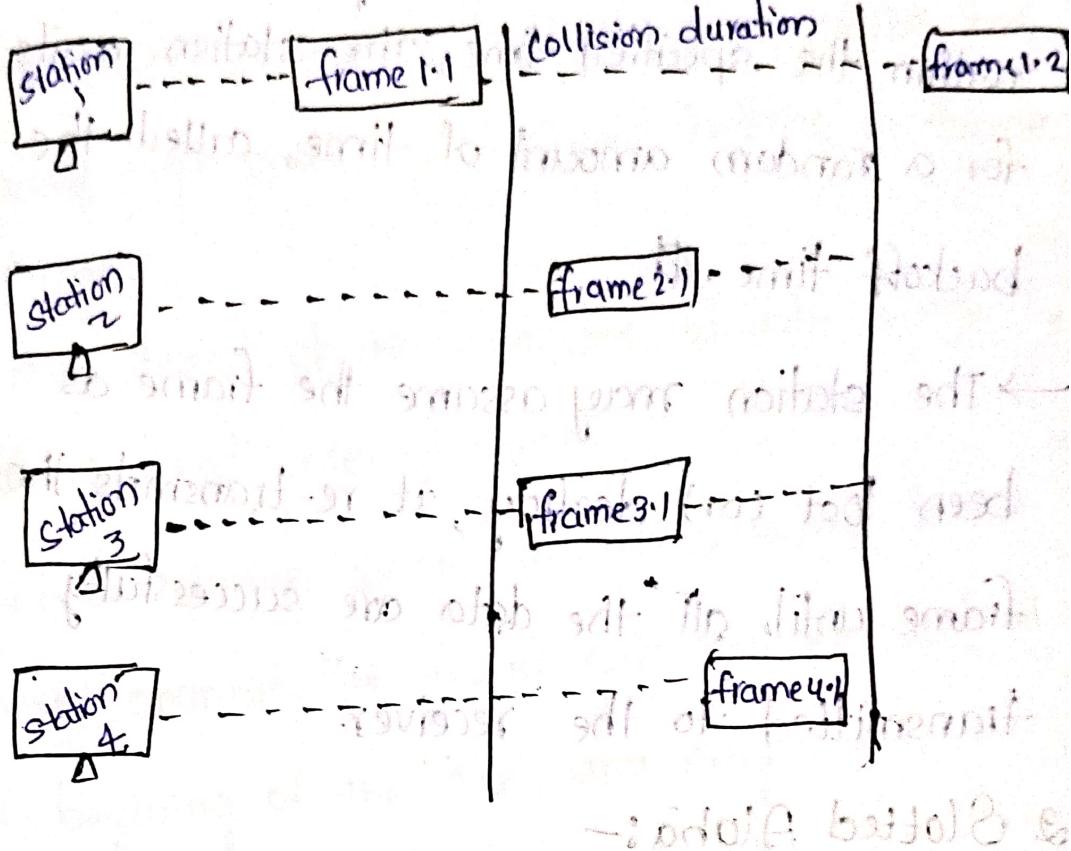
To avoid access conflict, each station follows a processor:-

- When can the station access the medium?
- What can the station do if the medium is busy?  

- How can the station determine the success or failure of the transmission?
- What can the station do if there is an access conflict?  


### ALOHA:-

- Aloha is a random access protocol.
- It was actually designed for UNIBAN; But it is also applicable for share medium.
- In this multiple station can transmit data at the same time and can hence lead to collision and data been garbled.
- It is of two types: 1. pure Aloha 2. Slotted Aloha.

## 1. Pure Aloha:-



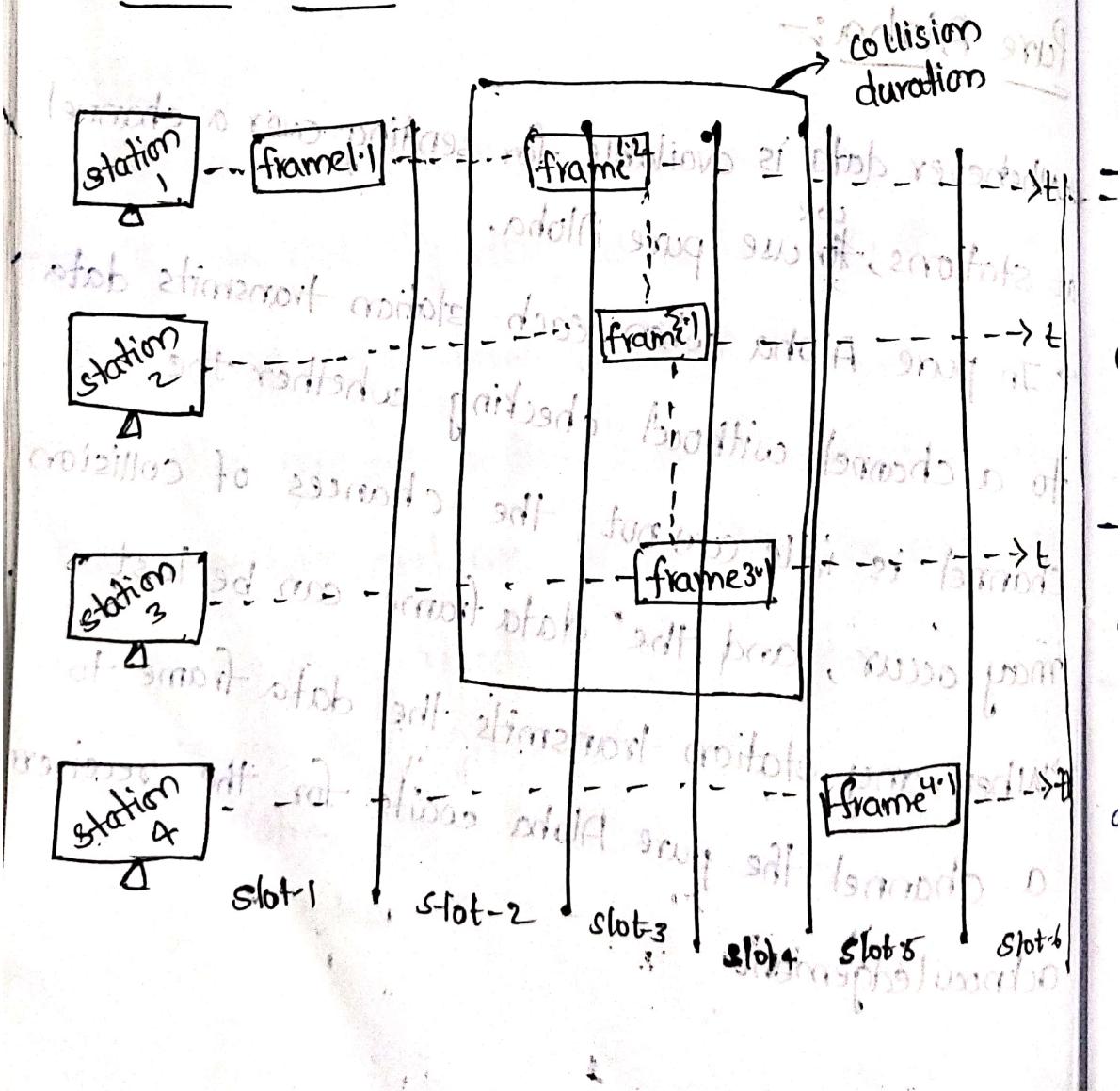
## Pure Aloha:-

- Whenever data is available for sending over a channel at stations, we use pure Aloha.
- In pure Aloha when each station transmits data to a channel without checking whether the channel is idle or not, the chances of collision may occur, and the data frame can be lost.
- When any station transmits the data frame to a channel the pure Aloha waits acknowledgement.

→ If it does not acknowledge the receiver end within the specified time, the station waits for a random amount of time, called the backoff time.

→ The station may assume the frame as been lost (or) destroyed, it re-transmits the frame until all the data are successfully transmitted to the receiver.

## 2 Slotted Aloha:



→ slotted aloha is designed to overcome the pure aloha efficiency because pure aloha has a very high possibility of frame hitting or frame collapse.

→ The shared channel is divided into slots that is fixed interval time.

→ If a station wants to send a frame to a shared channel, the frame can only be sent at the beginning of the slot, and only one frame is allowed to be sent to each slot.

→ If a station misses out the allowed time it must wait for the next slot.

CSMA :-

carrier sense multiple access protocol

There are two types of CSMA.

1. CSMA/CD

Collision Detection

2. CSMA/CA

Collision Avoidance

- principle of CSMA:
- \* "sense before transmit" prevents collisions by waiting to transmit if conflict occurs.
  - "listen before talk"
  - \* Carrier busy = Transmission is taking place.
  - \* carrier idle = No transmission is currently taking place.

Types of CSMA:

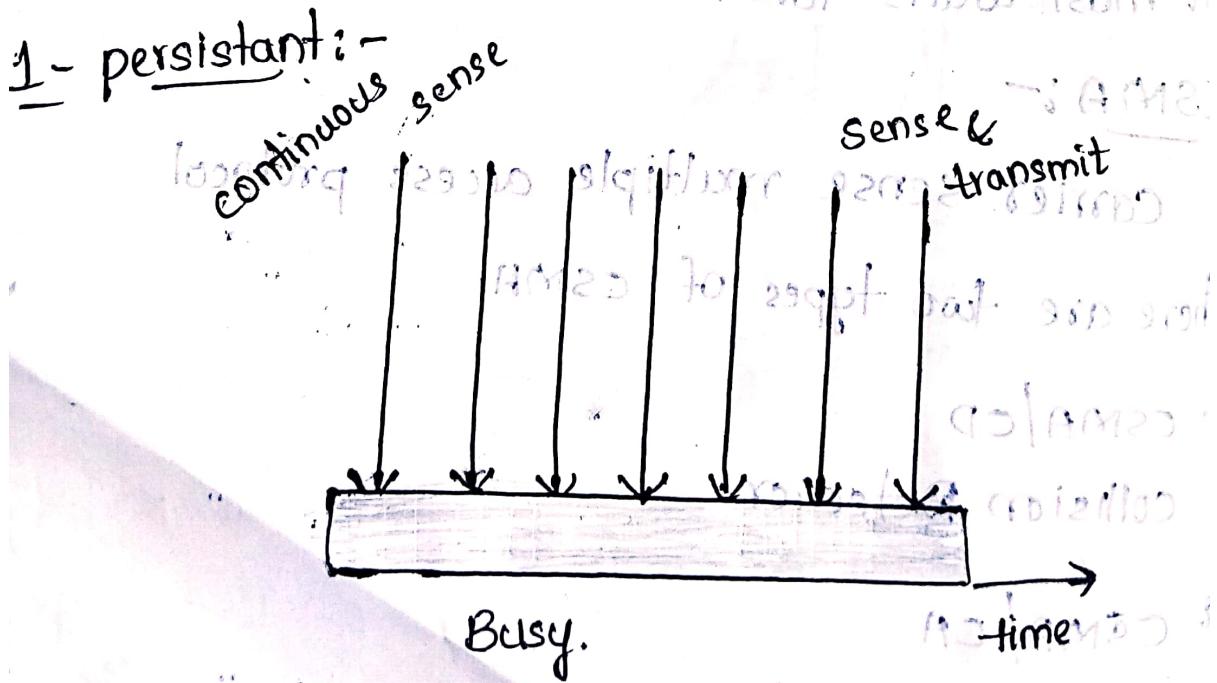
1 - persistent bus idle soft to principle of

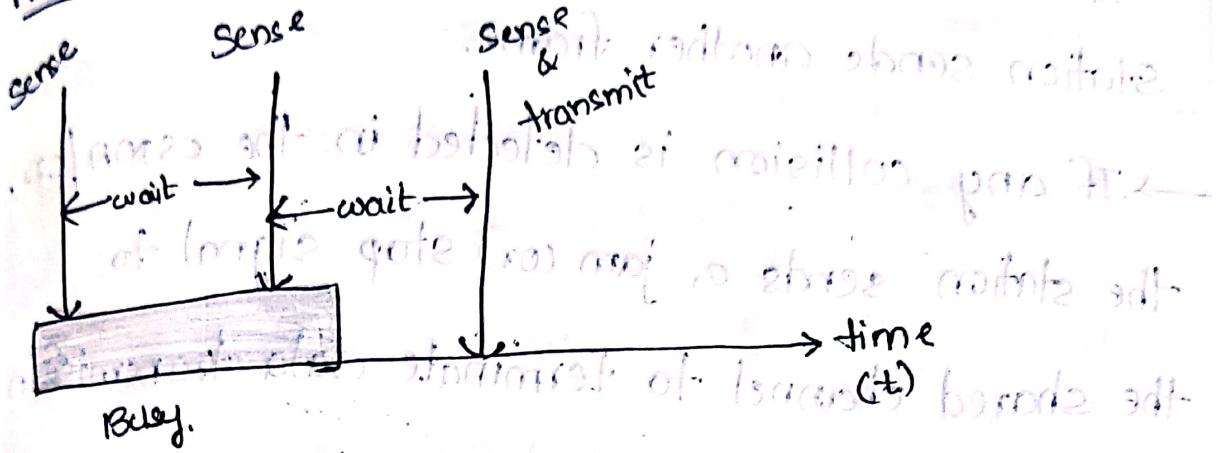
Non-persistent loss of bus idle a

2 - non-persistent bus idle two attempts in IEEE

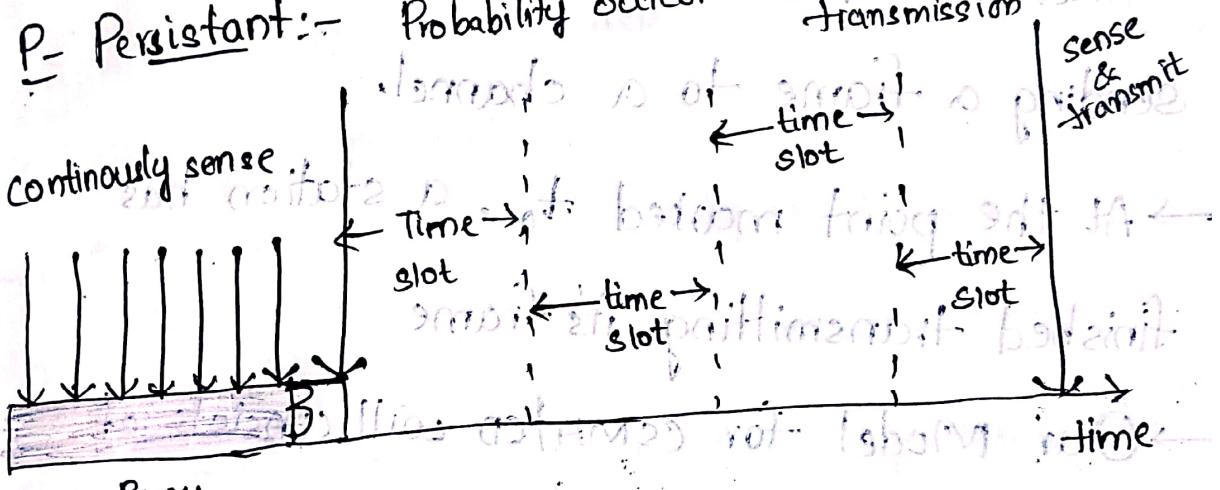
3 - p-persistent bus idle 50% chance to win

1 - persistent:





P- Persistent:- Probability outcomes doesn't follow transmission sense & fit



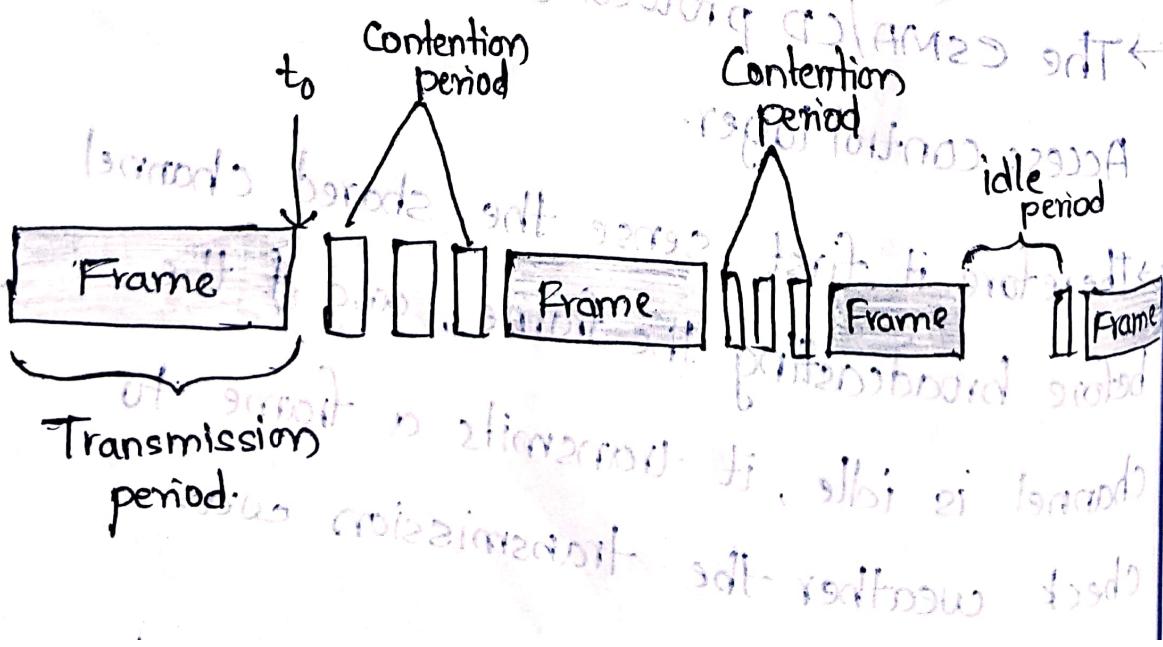
CSMA/CD uses CSMA/CD (Carrier Sense Multiple Access/Collision Detection).

→ The CSMA/CD protocol works with a shared medium  
Access control layer.

Access control layer  
→ Therefore it first sense the shared channel before broadcasting the frame, and if the

before broadcasting . . .  
channel is idle , it transmits a frame to  
check whether the transmission was successful.

- If the frame is successfully received, the station sends another frame.
  - If any collision is detected in the CSMA/CD, the station sends a jam (or) stop signal to the shared channel to terminate data transmission.
  - After that it waits for a random time before sending a frame to a channel.
  - At the point marked  $t_0$ , a station has finished transmitting its frame.
  - Our Model for CSMA/CD will consists of alternating contention and transmission period, with idle period occurring when all stations are idle.



## CSMA/CA

- It is a Carrier sense multiple Access/collisions Avoidance network protocol for carrier transmission of data frames. It is a protocol that works with a medium access control layer.
- When a data frame is sent to a channel, it receives an acknowledgement to check whether the channel is clear.
- If the station receives only a single acknowledgement (own), that means the data frame has been successfully transmitted to the receiver.
- Following are the methods used in the CSMA/CA to avoid the collision.

1. Interframe space:- In this method the station waits for the channel to become idle, and if it gets the channel is idle, it does not immediately send the data. It waits for some time. And, this time period is called interframe gap.

2. Contention Window:- In this contention window, the total time is divided into different slots.

3. Acknowledgement:- In the acknowledgement method the sender station sends the data frame, to the shared channel if the acknowledgement is not received a header of time.

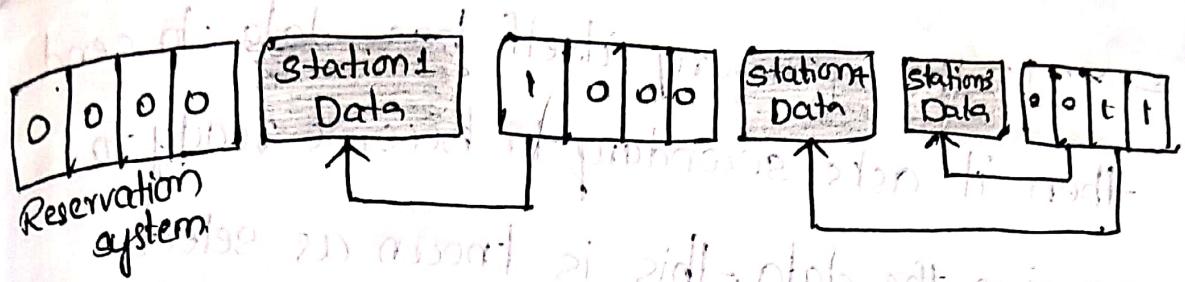
The diagram illustrates the CSMA/CA mechanism. It shows a sequence of frames over time. A station sends a frame, receives an ACK, and then sends another frame. If no ACK is received, it waits for a random duration before retransmitting. The random duration is determined by a binary exponential backoff algorithm. The diagram also shows the CSMA/CA process, where a station waits for the channel to be idle before sending a frame.

## Controlled Access

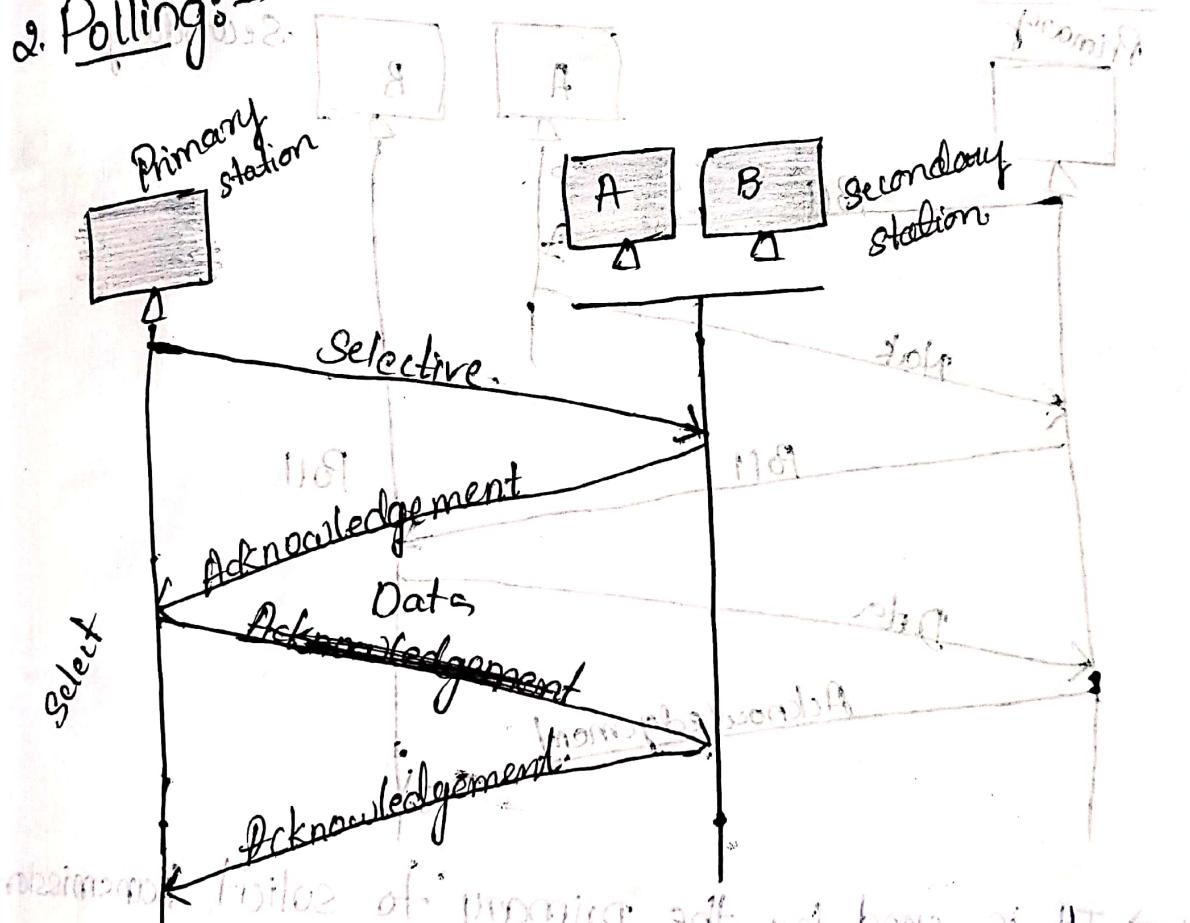
- In controlled Access protocol a station cannot send data unless it has right to send it.
- There are three popular controlled access methods
  - 1. Reservation system.
  - 2. Polling.
  - 3. Token passing

### i. Reservation System:-

- A station needs to send a frame had to reserve the channel before sending the frame. If there are  $N$  stations then the reservation frame consists of exactly  $N$  number of mini slots in the frame.



## Q. Polling:-

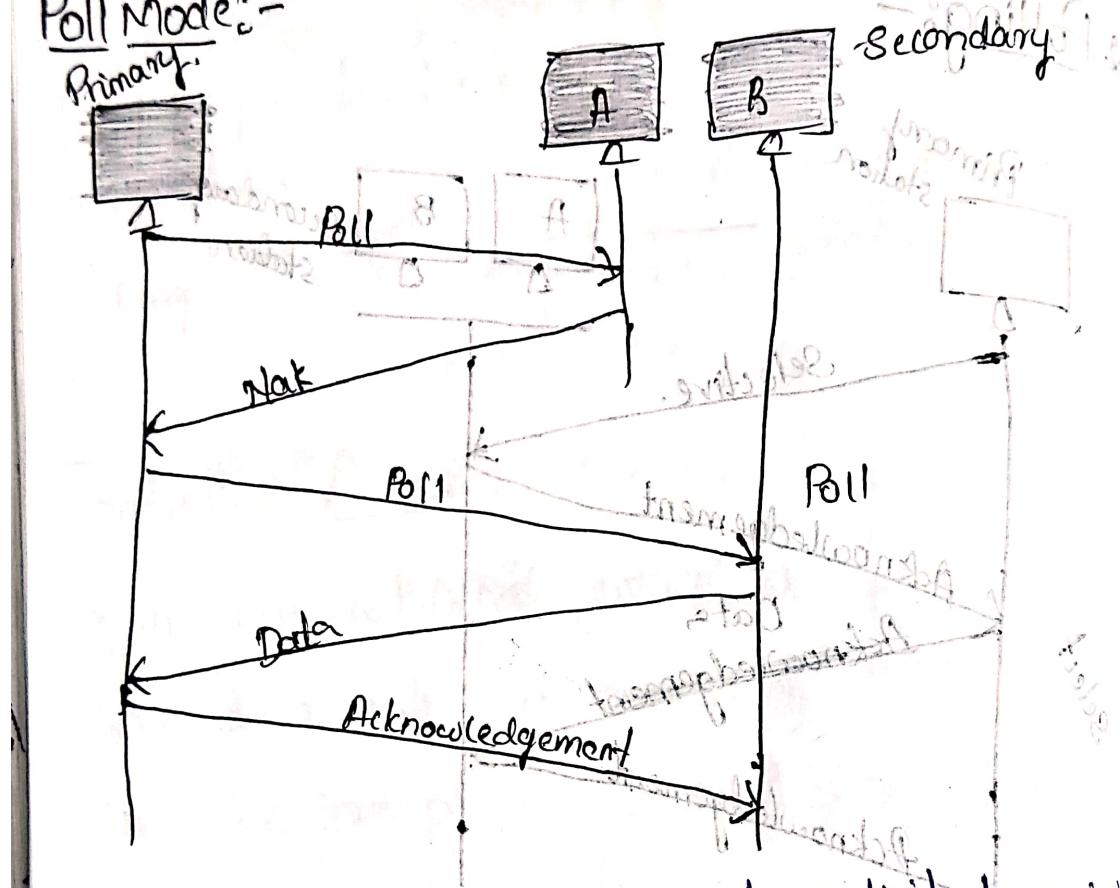


→ polling works with the system in which there is a primary device and others are secondary devices.

→ Data transfer occurs to the primary device. The primary device asks each of the secondary devices whether it has any data to send. This is known as polling.

→ If the primary itself has data to send then it asks secondary to become ready to receive the data. This is known as select.

### Poll Mode:-



→ It is used by the primary to solicit transmission from the secondary.

→ When the primary is ready to receive data it polls a secondary whether it has any data to send.

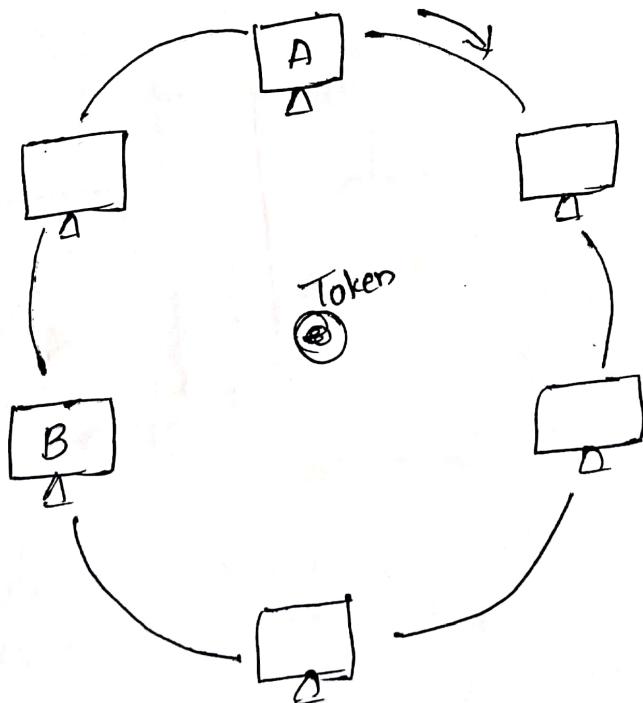
→ If the secondary has no data to send then it sends a negative acknowledgement (NAK) to the primary.

- The primary then polls the next secondary. If the secondary has data to send then it sends the data to the primary.
- primary sends positive acknowledgement in response.

### Select Mode:-

- When the primary device has to send, it sends an SEL(select) command to the secondary to make it know about the incoming data.

### TOKEN PASSING:- (page: 138)



## Channelization:-

→ It is a multiple access method in which the available band width of a link shared in time, frequency, (or) through code, between different stations.

## Multiplexing:-

→ In computer networking, multiple signals are combined together thus travel simultaneously in a shared medium.

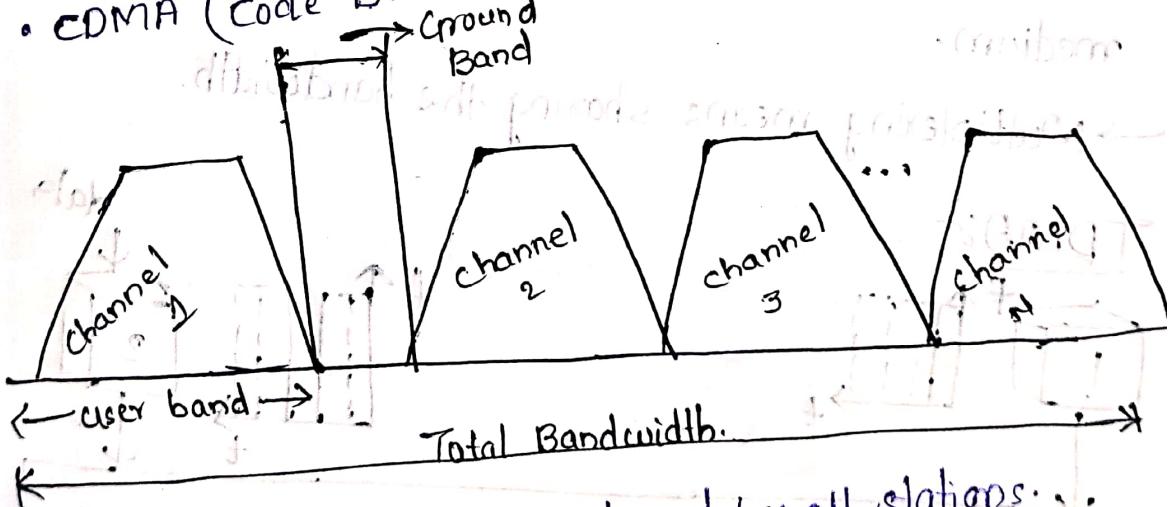
→ Multiplexing means sharing the band width.

→ It consists of three protocols.

• FDMA (frequency Division Multiple Access)

• TDMA (Time Division Multiple Access)

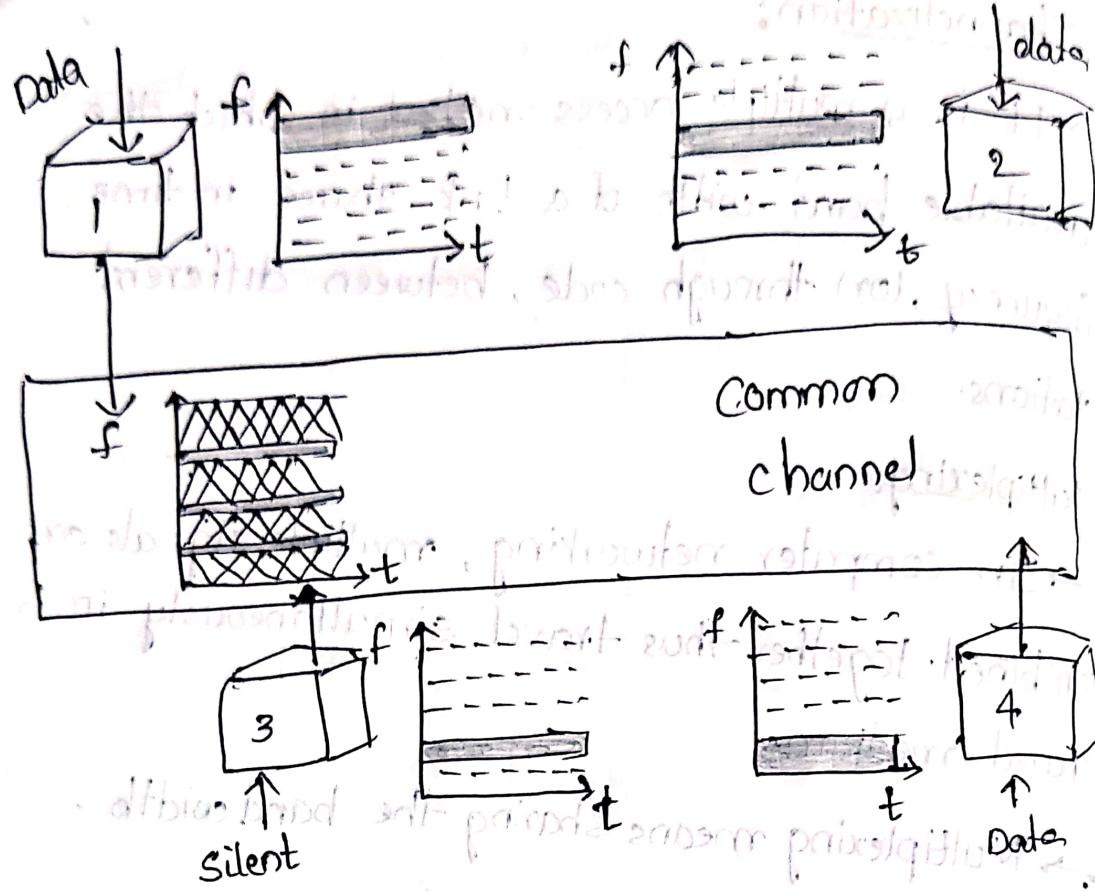
• CDMA (Code Division Multiple Access)



→ The available bandwidth is shared by all stations.

→ The FDMA is a datalimit protocol that uses, FDM (frequency division multiplexing) at the physical layer.



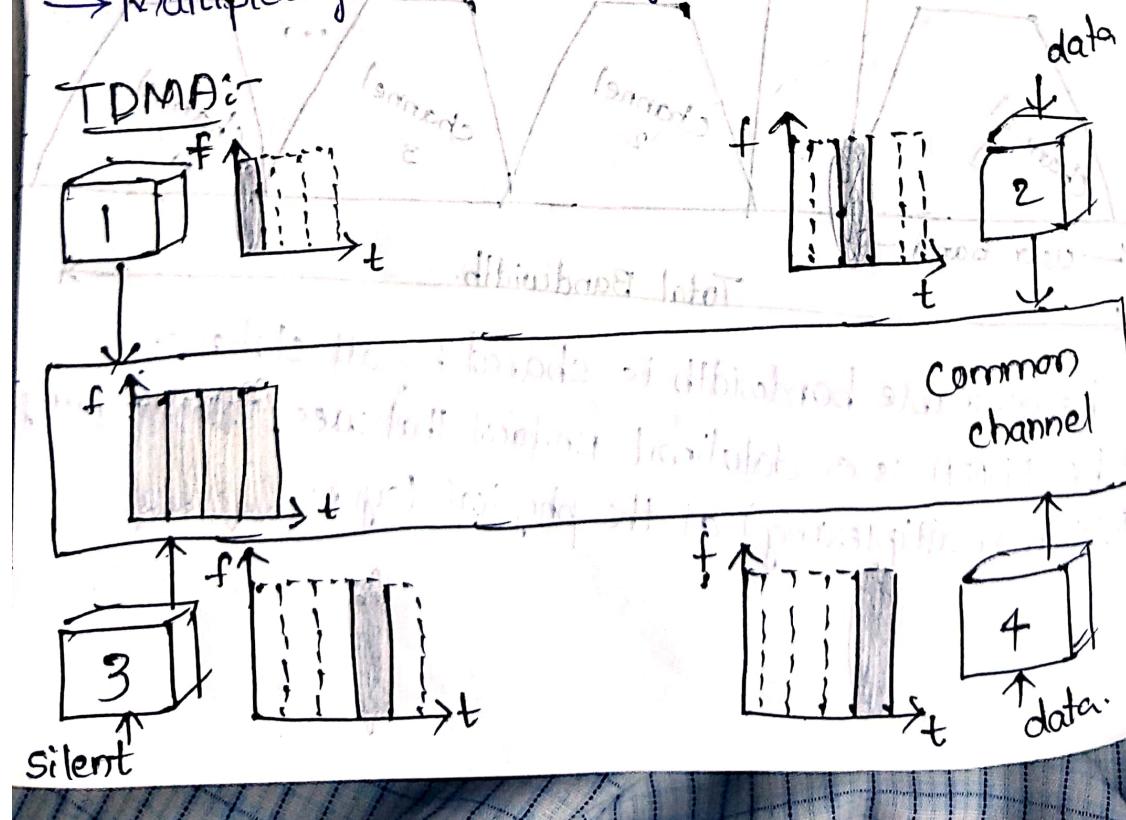


Multiplexing: → Allocating bandwidth to stations.

→ Multiplexing means combining multiple signals on a shared medium.

→ Multiplexing means sharing the bandwidth.

TDMA:



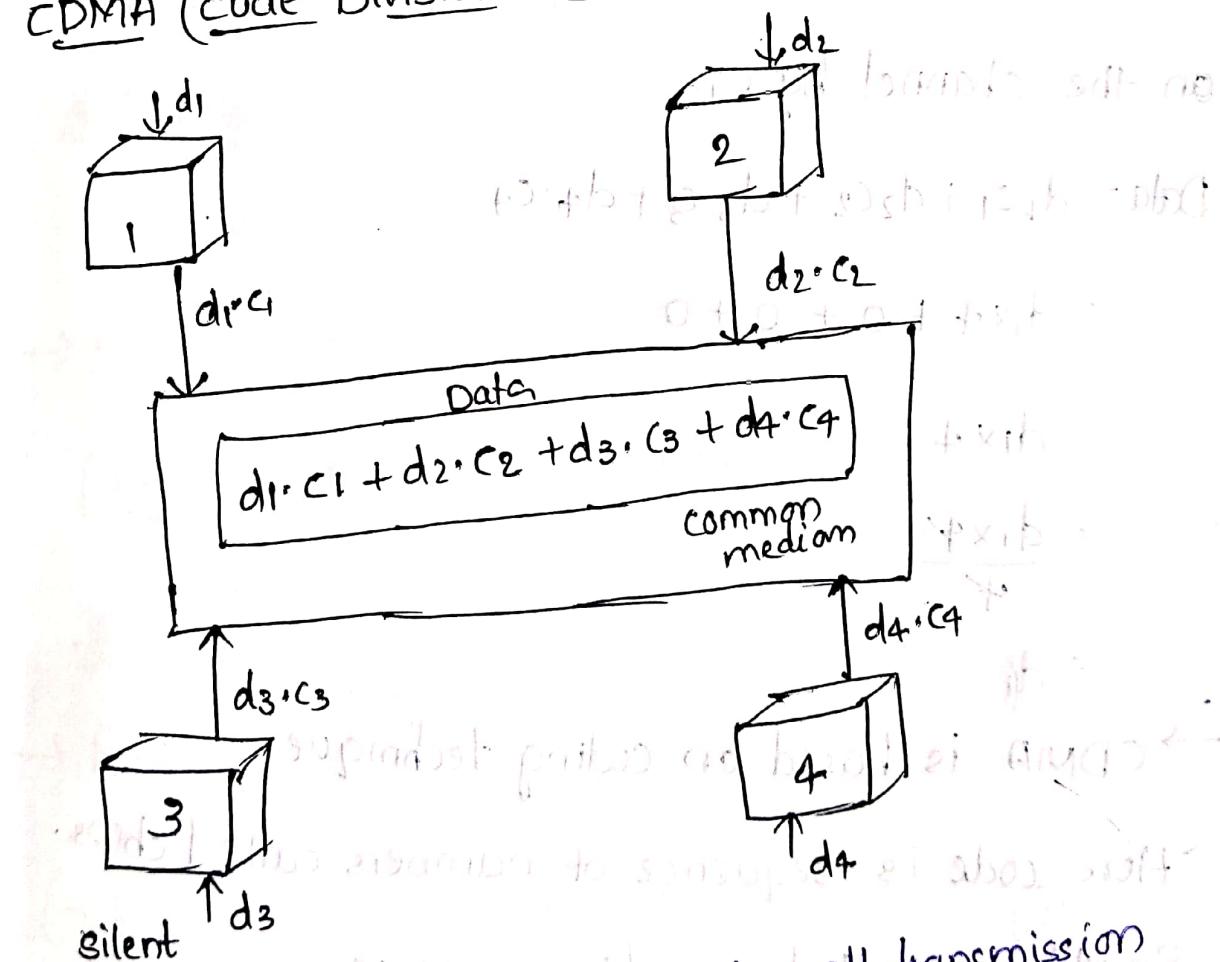
→ the time division multiple access (TDMA) method of multiple access is suitable for the digitally modulated signals.

→ In TDMA, the bandwidth of the channel is shared by the stations in time.

→ Each station can send data during its allocated time slot.

→ The beginning of the slot and the location of the slot must be known by each station.

CDMA (Code Division Multiple Access)



→ In CDMA, one channel carries all transmission simultaneously.

→ CDMA differs from FDMA because only one channel occupies the entire bandwidth of the link.

→ The assigned codes have two properties.

1. If we multiply each code by another, we get zero(0)

2. If we multiply each code by itself, we get four(4)

Example:

Let us assume that if station 2 wants to send data on the channel by  $c_1$ ,

$$\text{Data} = d_1 \cdot c_1 + d_2 \cdot c_2 + d_3 \cdot c_3 + d_4 \cdot c_4$$

$$= d_1 \cdot 4 + 0 + 0 + 0$$

$$= d_1 \cdot 4$$

$$= \frac{d_1 \cdot 4}{4}$$

$$= d_1$$

→ CDMA is based on coding technique.

→ Here code is sequence of numbers called chips.

So this is called as chip sequence.

→ The chip sequences are also called as orthogonal sequences.

# Chip Sequences & Properties

$$[+1 +1 +1 +1]$$

$c_1$

$$[+1 -1 +1 -1]$$

$c_2$  Bidirectional

$$[+1 +1 -1 -1]$$

$c_3$

$$[+1 -1 -1 +1]$$

$c_4$  Unidirectional

Four different chip sequences. Each pass to site 1.

→ Each sequence is made of  $N$  elements (N means number of stations)

→ Multiply sequence by a number, add bias and

$$2 [+1 +1 -1 -1] \xrightarrow{\text{bias} = 0} (+2 +2 -2 -2) = 0$$

$c_3$

→ Multiply two different sequences → 0.

$$(+1 +1 -1 +1) \otimes (+1 +1 +1 +1) = 0$$

$c_3$

→ Adding 2 sequences.

$$(+1 +1 -1 -1) + (+1 +1 +1 +1) =$$

$c_3$

$$(+2 +2 0 0) \rightarrow \text{phase out of site 1}$$



## Data representation in CDMA

Data bit 0 → +1

Data bit 1 → -1

Silence → 0

## Ethernet:-

→ It is a way connecting computers together in a LAN. The basic idea of its design is that multiple computers can allow to access it & can send the data at anytime.

→ There are 2 kinds of ethernet:

i. classic ethernet

& Switched ethernet.

### i. Classic Ethernet:-

→ It is a single long cable to which all the computers are attached.

→ The first variety is thin ethernet, which is sent more easily and made connections using industry standard BNC connectors.

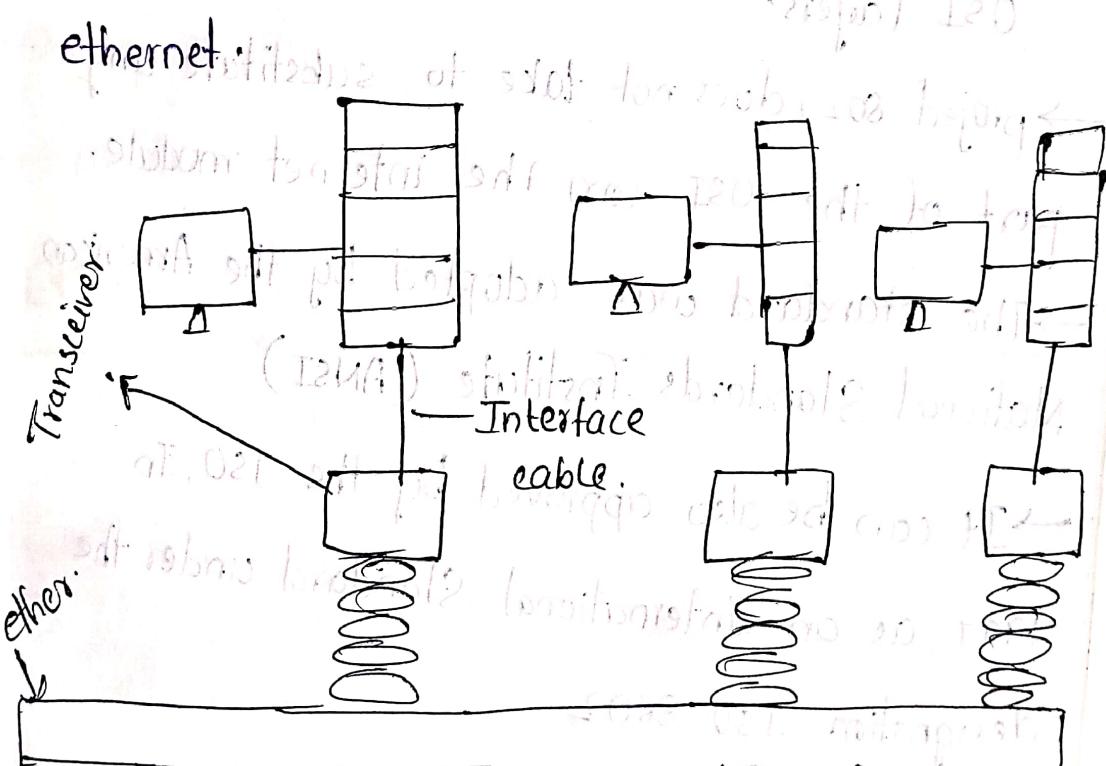
→ Thin ethernet is cheaper & easy to install

→ It is run only upto 185 meter per segment.

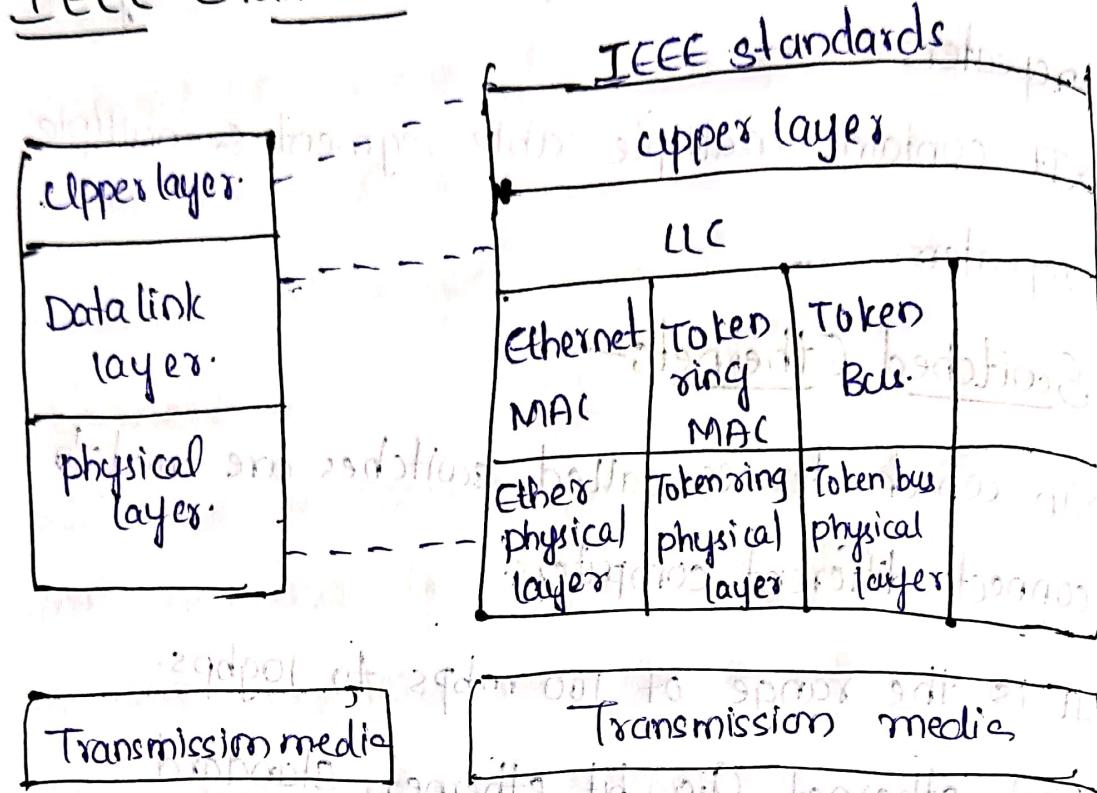
- A series of cables segments connected by repeaters.
- It contains multiple cable segments & multiple repeaters.

## 2. Switched Ethernet:-

- In which devices called switches are used to connect different computers.
- If is the range of 100 mbps to 10gbps.
- Fast ethernet, Giga bit ethernet, standard ethernet.



## IEEE Standards:-



OSI layers.

- project 802, does not take to substitute any part of the OSI (or) the internet module.
- The standard was adopted by the American National Standards Institute (ANSI)
- It can be also approved by the ISO, In 1987 as an International Standard under the designation ISO 8802

### Data link layer:-

- the datalink layer in the IEEE standard can be divided into two suppliers.
  - i, logic link control
  - ii, Medium access control.

## Logic Link Control

The data link layer provides the functions are framing, flow control and error control.

→ In IEEE project 802, flow control, error control and part of the framing duties are together into one sublayer called the logic link control.

## Ethernet Evolution

- Standard Ethernet - 10 Mbps
- Fast ethernet - 100mbps
- Gigabit ethernet - 1 Gbps
- 10 Gigabit ethernet - 10 Gbps

## Ethernet Frame Format

start frame delimiter							
Preamble	SFD	Destination Address	Source Address	length or type	Data padding	CRC	frame
7bytes	1byte	6 bytes	6 bytes	min=46 bytes max=1500bytes			

Protocol id cannot be 0 because no padding

## Preamble

- It is of 7bytes alternating zeros and ones (0's & 1's)

used to used to alert the receiver to say that we are sending the data so get ready to receive it.

## SFD (start Frame Delimiter)

- It is the pattern of 1010101
- The last two (one) bits tell the receiver that the rest of the frame is about to start.
- If the receiver is not ready to accept the data, again we will find the SFD to say that we are going to send the next field will be with destination address, so get ready to receive it.

## Destination Address

It represents the link layer address of receiver.

## Source Address(SA)

It represents the link layer address of sender.

## Length (or) type

Depending on whether the frame is standard ethernet or IEEE 802.3

Standard frame size is 1500 bytes, IEEE frame size is 64 bytes.

IEEE frame size is 64 bytes, IEEE frame size is 1500 bytes.

IEEE frame size is 64 bytes, IEEE frame size is 1500 bytes.

## Data :-

- It contains the actual data.
- It is the range of 46 bytes to 1500 bytes.
- If more than 1500 bytes data it will be included in the next frame.
- If it is less than 46 bytes, we have to padding with zeroes (0's)

## CRC (cyclic Redundancy Check) :-

- It is for error reduction.

## Minimum (or) Maximum frame length :-

Destination Address	Source Address	Length Type	Data padding	C.R.C
6 bytes.	6 bytes.	(2bytes)	Min-46 bytes Max-1500bytes	4 bytes

→ Min frame length

$$= 6 + 6 + 2 + 46 + 4$$

$$= 64 \text{ bytes.}$$

→ Max frame length

$$= 6 + 6 + 2 + 1500 + 4$$

$$= 1518 \text{ bytes.}$$

## Addressing:-

→ Among the 16 bits, the first four bits provides identification to the fragment and the rest of 12 bits contain the sequence number that increments with each transmission.

It is of 3 types.

- Unicast
- Multi cast
- Broad cast

## Physical Layer

→ The standard ethernet defines several physical layer implementation. Here 4 of the most common are described.

## Ethernet Cabling:- (page: 152)

Name	Cable	Max segment	Nodes	Advantages
10 Base 5	thick coaxial	500 mts	100	Original cable
10 Base 2	Thin coaxial	185 mts	30	No hub needed
10 Base T	Twisted-pair	100 mts	1024	Cheapest system
10 Base F	fibre-optics	2000 mts	1024	best between buildings

Eg:- 10 Base 5  
Collision domain  $\rightarrow$  500 mts  
Mbps Bandwidth

### Changes in the Standard:-

The 10 Mbps standards ethernet has gone through several changes before moving to the higher data rates.

- Bridge ethernet
- Switched ethernet
- Full Duplex ethernet

### Bridged Ethernet :-

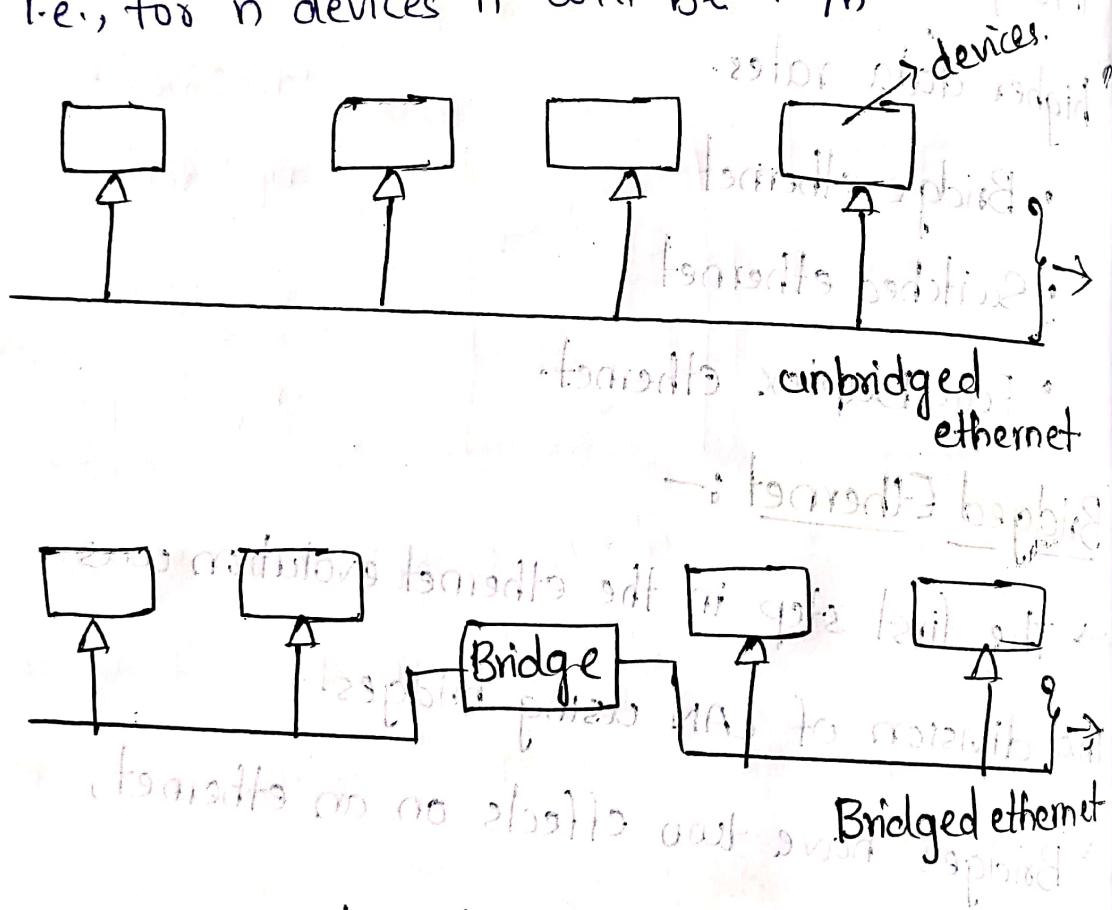
→ The first step in the ethernet evolution was the division of LAN using bridges.

→ Bridges have two effects on an ethernet,

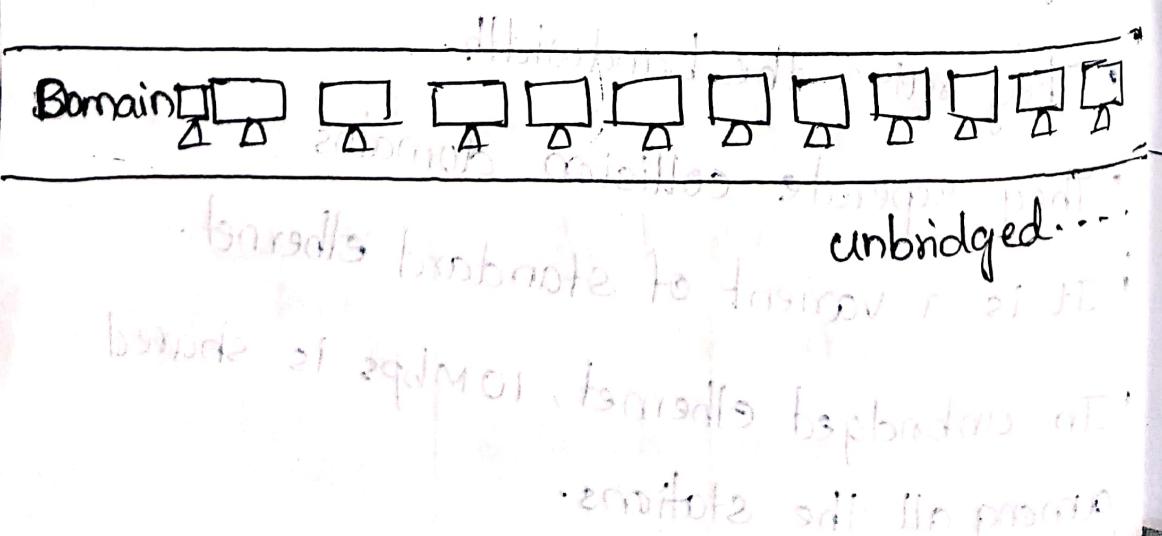
they are as follows:

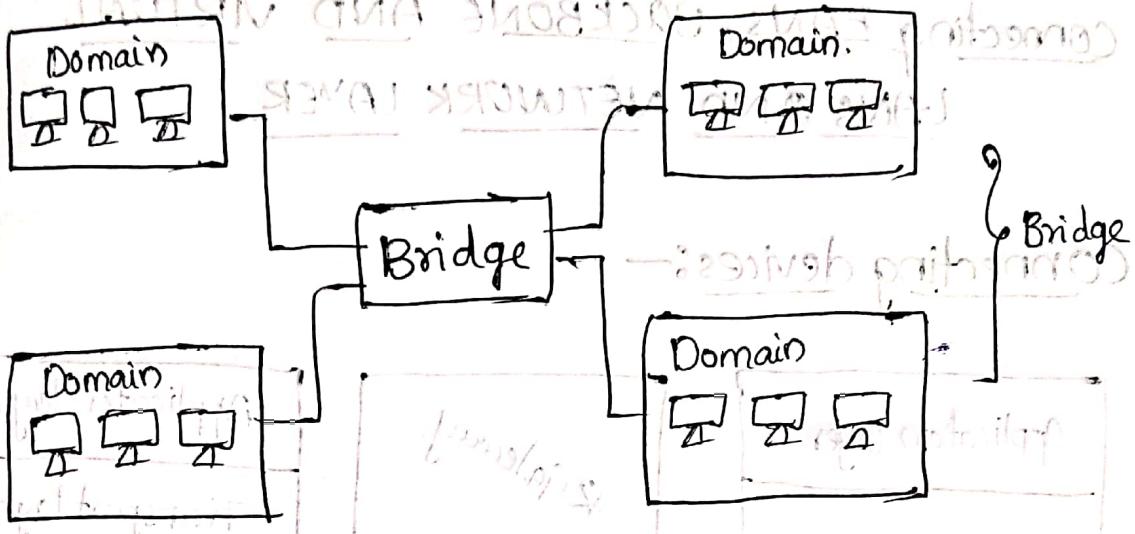
- They raise the bandwidth.
- They separate collision domains.
- It is a variant of standard ethernet.
- In unbridged ethernet, 10 Mbps is shared among all the stations.

- A bridge divides the network into two or more networks (raising the bandwidth)
- Without bridge, if a device is there, it can have 10mbps, for two devices, it will be 5Mbps, for 4 devices, it will be 2.5Mbps i.e., for  $n$  devices it will be  $10/n$

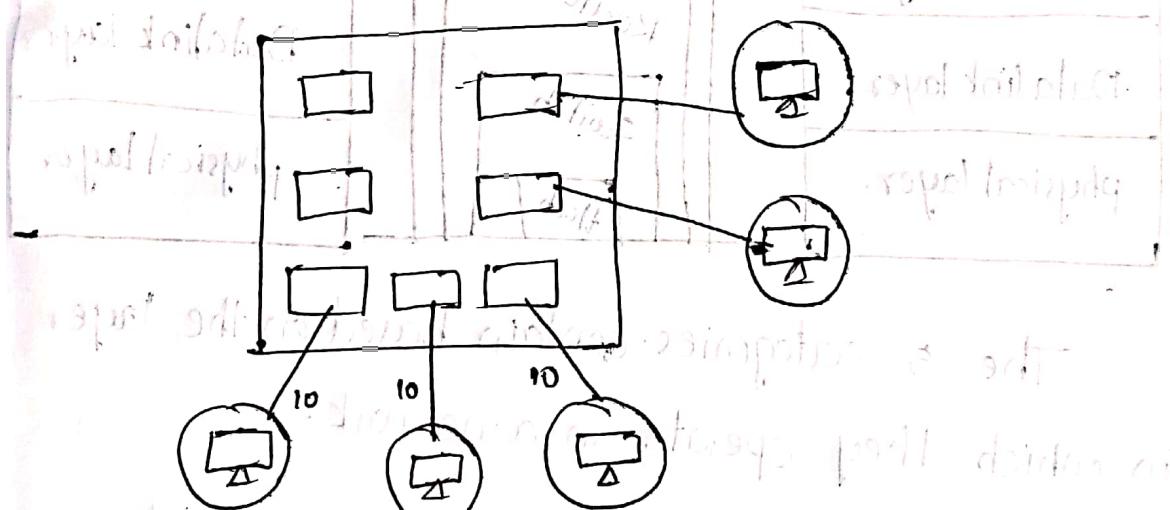


### Collision Separating domain





## 2. Switched Ethernet :- page 157



## 3. Full-Duplex Ethernet :-

