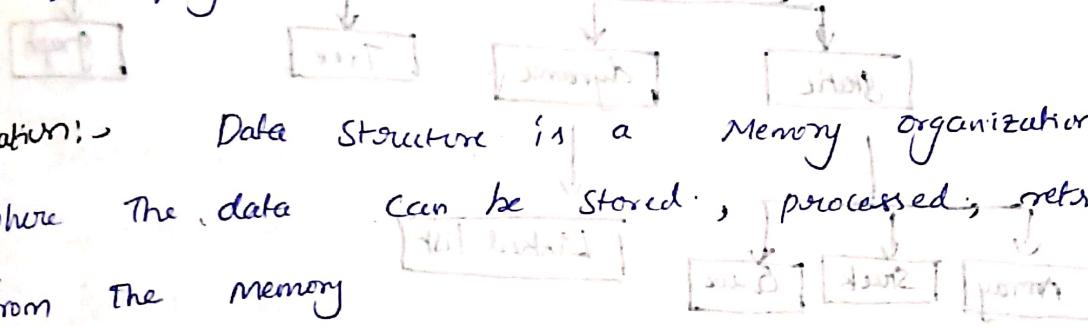


UNIT - 3 -

DATA STRUCTURES

Introduction to Data Structures:- When Large amount of data is organized in a memory, we need to maintain the data efficiently by using the concept of "Data Structures". In present days, Each and Every application process a large amount of data. Every application needs to apply the data structures technique.

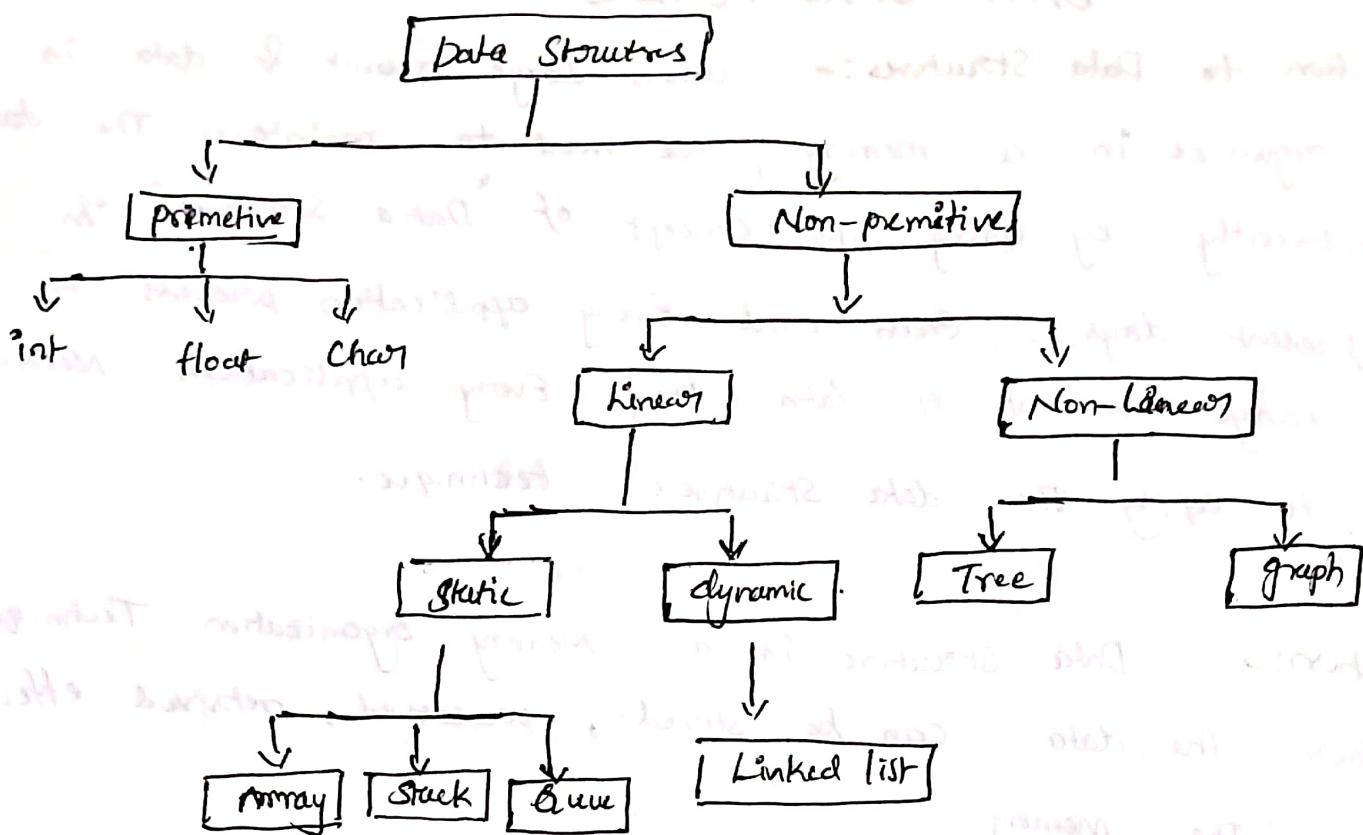


Definition:- Data Structure is a Memory organization Technique where the data can be stored, processed, retrieved efficiently from the memory.

Advantages :-

- 1] Data structures are used to efficiently organizing and accessing the data quickly.
- 2] Data structures are always used to reduce the time complexity of a Data processing.
- 3] In data structures data always organizing with dynamic approach rather than static approach.
- 4] Data structures organizing with the data with different operations like Searching, Sorting, Inserting, deleting, etc...
- 5) The following is the Hierarchy for data structures concept

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The term primitive data structure means where the data structure can hold any single value. This type of data structures are mainly used in programming language.

Constructs
Constructs

The term Non-primitive data structure means where the data structure can hold more than one value. This type of data structures are mainly used for both programming language constructs and application constructs.

Linear search with Recursion :- Linear Search is called as

Sequential Search. This is oldest and slow searching technique for finding the search element. The main objective of linear search is to search the element from whole list in sequential approach. That means One element after the ~~and~~^{the} element. In this searching process first we need to collect a input of elements then after we need to set a search element. For finding search element, we will compare each and every element with search element until the element is found. If the search element is not found. Then we can display. Search element is not existing in a list.

Ex:- Let us consider the following array of elements and find the search element.

$$a[5] = \boxed{10 \mid 20 \mid 30 \mid 40 \mid 50} \\ a[0] \quad a[1] \quad a[2] \quad a[3] \quad a[4]$$

Here Search element is : $se = 40$

if $(a[0] == se)$ $\therefore [a[0] = 10 == se]$
 $10 == 40$.

Here The search element is not found. Then control will go.
Next

if $(a[1] == se)$
 $20 == 40$.

\therefore Here The search element is not found.

if $(a[2] == se)$
 $30 == 40$.

\therefore Here The search element is not found.



if ($a[3] == \text{se}$)

$\text{pos} = 3$

break;

Then Search element is found.

programming with recursion:-

```
int lin_recursion(int [], int, int);
```

```
void main ()
```

```
{
```

```
int a[10], n, se, i, pos;
```

```
printf ("Enter the array size\n");
```

```
scanf ("%d", &n);
```

```
printf ("Enter elements in array\n");
```

```
scanf ("%d",
```

```
for (i=0; i<=n-1; i++)
```

```
{
```

```
scanf ("%d", &a[i]);
```

```
}
```

```
printf ("Enter search element\n");
```

```
scanf ("%d", &se);
```

```
pos = lin_recursion (a, n-1, se);
```

if ($pos \geq 0$)

printf ("Element is not found");

else

printf ("Search element %d is located at %d position\n", se, pos+1);



q1) lin- recursion (int a[10], int n, int se)

{
if ($n > 0$)

return -1;

if ($a[n] == se$)

return n;

else:

return lin- recursion (a, n-1, se);

}

III Binary Search Using Recursion:- Binary Search is efficient searching technique as compare to Linear Search.

Binary Search is working based on Divide and Conquer approach. In this approach The input of elements are partitioned into two different parts based on calculation of mid value. Once the mid value is return we can perform the following steps for finding searching element.

* Here $mid = \frac{low + high}{2}$

{ \therefore low represents starting position}

{ \therefore high represent ending position}

i) If $mid(a[mid] == se)$

return mid;

ii) else if ($a[mid] > se$)

return $high = \frac{high + mid - 1}{2}$;

iii) else

return $low = \frac{low + mid + 1}{2}$;



The above process will be continued until the search element is found.

Note: Binary searching is only working with ordered elements
i.e., That is in ascending order & descending order.

Ex:- Let us consider the following array and find a search element (Here search is 14).

2	4	6	8	10	12	14	16	18
0	1	2	3	4	5	6	7	8

To calculate $\text{Mid} = \frac{\text{low} + \text{high}}{2}$

$$\Rightarrow \frac{0+8}{2} = 4.$$

$$\text{low} = \text{mid} + 1 = 4 + 1 = 5$$

here $a[\text{mid}] < \text{se}$ satisfied That is $10 < 14$.

Then $\text{low} = \text{mid} + 1$.

12	14	16	18
5	6	7	8.

1
 low

high

$$\text{mid} = \frac{5+8}{2} \Rightarrow \frac{13}{2} = 6.5 \approx 6.$$

Here $a[\text{mid}] = \text{se}$ satisfied.



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ic: 14 == 14 Then. Search element is found.

```
#include < stdio.h>
#include < conio.h>

int binary_rec(int [], int, int, int);

Void main ()
{
    int a[20], n, i, se, pos;
    clrscr();

    printf("Enter The size of array in");
    scanf("%d", &n);

    printf("Enter The elements in array[n]");
    for (i=0; i<=n-1; i++)
    {
        scanf("%d", &a[i]);
    }

    printf("Enter The search element in");
    scanf("%d", &se);

    pos = binary_rec(a, 0, n-1, se);

    if (pos < 0)
        printf("Search element is not found");
}
```



```
else  
    printf (" search element %d occurs at %d location  
            (%d , %d);\n  
    getch();  
}
```

```
int binary_rec ( int a[], int l, int h, int se ).  
{  
    int mid;  
    if (l > h)  
    {  
        return -1;  
    }  
    mid = l + h / 2;  
    if (a[mid] == se)  
    {  
        return mid;  
    }  
    else if (a[mid] > se)  
    {  
        return binary_rec (a, l, mid - 1, se);  
    }  
    else  
    {  
        return binary_rec (a, mid + 1, h, se);  
    }  
}
```

Analysis of algorithms using Time complexity :-

In data structures every algorithm efficiency is calculated by using two factors. That is time complexity and space complexity.

Time complexity is defined as the amount of time required for executing the instructions of algorithm. The time complexity of algorithm can be calculated based on the following factors

- Type of processor used. [Single processor, (8) multiprocessor)
- Type of Architecture used [8 bit, (8), 32 bit, (8), 64-bit .. etc]
- The Unit time cost of time required for executing the operations of Algorithm.
i.e: Assignment operator, Arithmetic, logical, etc.

Ex:- Let us consider the single processor and 32 bit machine architecture. To calculate the time complexity for the following piece of code. (8) Algorithm.

Sum-of-numbers (int a [], int n)

{

 int sum = 0;

 for (i=0; i < n; i++)

{

 sum = sum + a[i];

}

 return sum;

Cost	repetition	Total
1	*	1
1+1+1	* + (n+1)+n	2n+2
1	+ 2n	2n
1	* + 1	1

The above

The time complexity for the above code. $T(n) = 4n + 4$

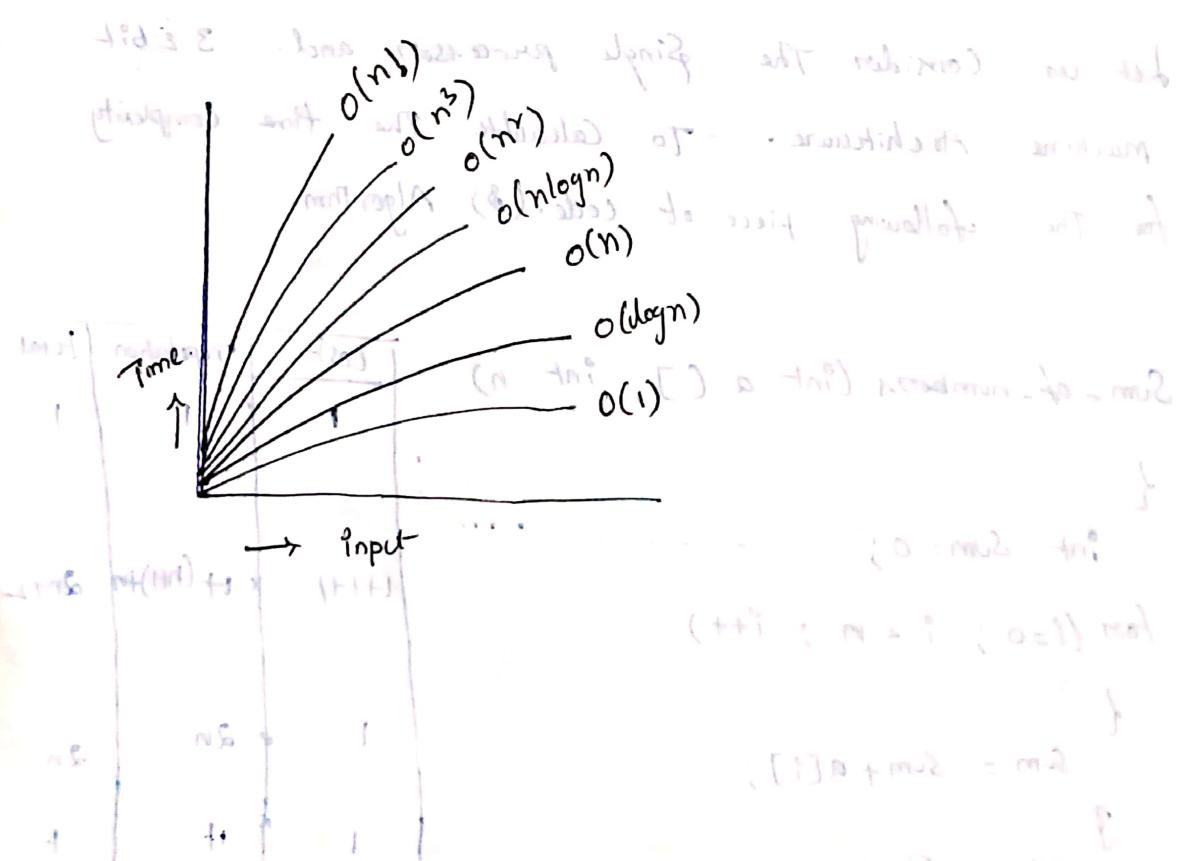
Generally the time complexity of Algorithm is measured by using Big-oh notation (O). According to notation, the time complexity is representing without polynomial terms.

$$\therefore T(n) = c \cdot n + c'$$

[where c, c' are constants for a polynomial.]

$$\therefore T(n) = O(n)$$

The following diagram shows the performance of analysis of algorithm with time complexity

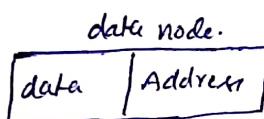


what is Linked List ? Types of Linked List:-

Linked List is also called as Linear List and which supports Linear data structure. In Linked List the elements are organized in memory dynamically with inter connection of addresses by using sequential approach. That's why Linked list is also called as Linear List.

In Linked List The elements are organizing in memory with a data-node representation, where The data node consists of two parts That is First part Storing a data and the second part Storing address of a node.

The following shows The structure of a data node.



There are Three types of Linked list

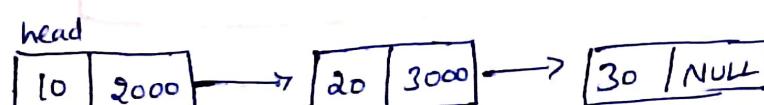
1) Single Linked list

2) Double Linked list

3) Circular linked list

1) Single Linked list :- This is a default linked list for implementing the list of elements. where The elements are organizing in a sequential fashion with connecting of address next node. In this type of linked list the elements are only traversing in unidirectional.

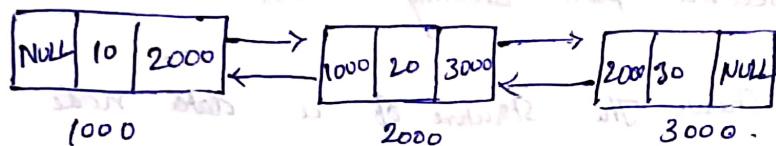
Ex:-



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2) Double Linked list :- In double linked list the elements are organized in memory with a representation of data node. but here the data node consists of three parts. The first part indicates storing the address of a previous node. The second part indicates data, the third part indicates storing address of next node. In this type of list the elements are inter connect to each other. in a by directional way.

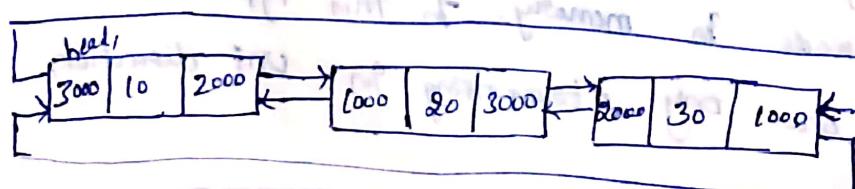
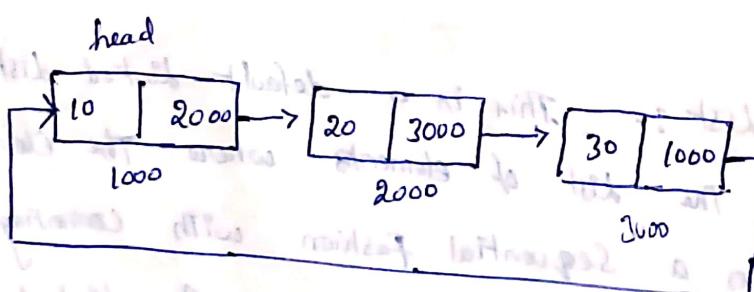
Ex:-

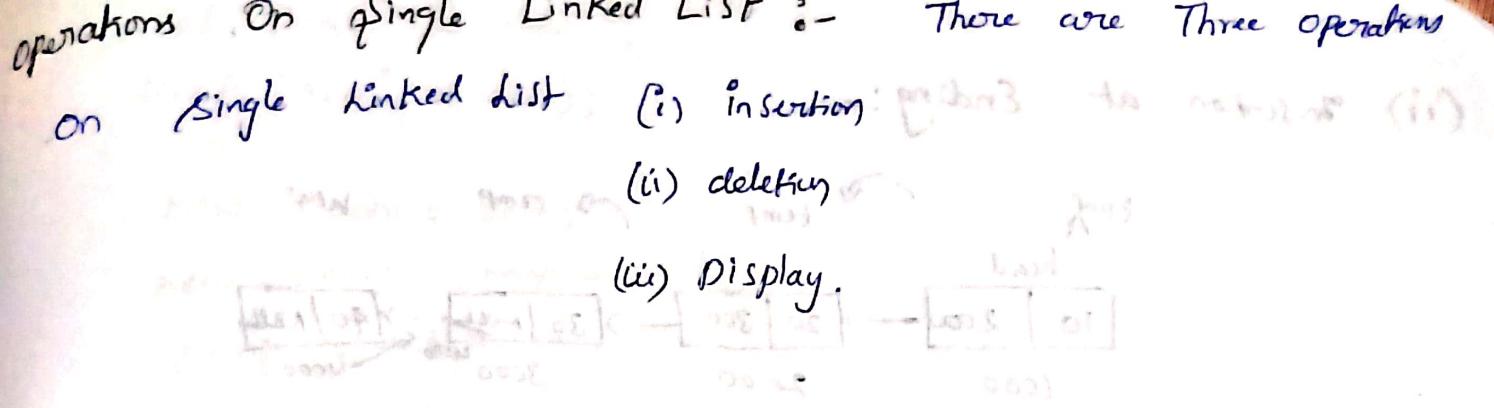


3) Circular linked list :-

The circular linked list also represented with circular linked list and circular double linked list in circular single linked list. The starting address of node is connected to last node. In a circular fashion. In circular double linked list the last and first address node are mutually inter connected. in a circular fashion.

Ex:-



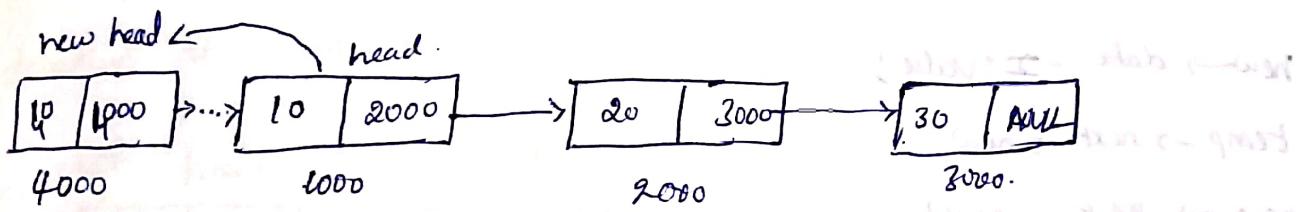


1) insertion :- There are Three ways of Insertion.

(i) insertion at beginning.

(ii) insertion at Ending

(iii) insertion at any specific position.



new = (struct node*) malloc (size of (struct node));

printf("Enter The Value.");

scanf("%d", &Value);

new->data = Value;

new->next = head;

new = head;



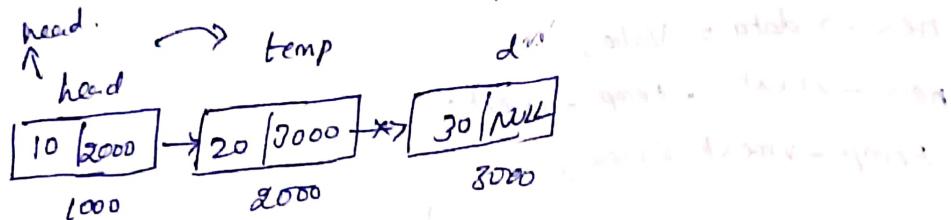
Struct node.

{ int data;

struct node *next;

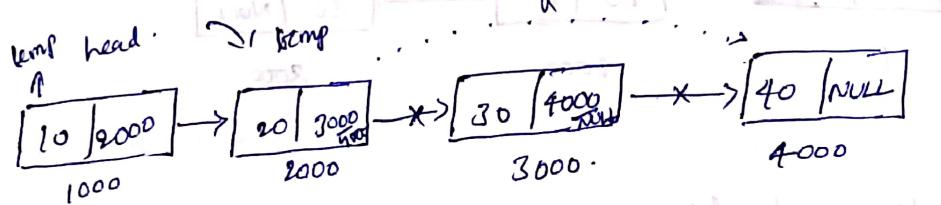
{ *new, *head;

Deletion at ending:-



```
temp = head;  
while (temp->next != NULL)  
{  
    temp = temp->next;  
}  
d = temp->next;  
temp->next = NULL  
free(d);
```

Deletion at any specific position:-

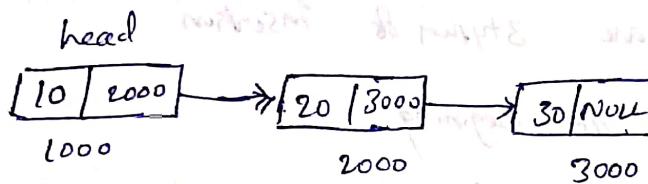


```
int pos, i;  
temp = head;  
printf("Enter a position you want to delete");  
scanf("%d", &pos);  
for (i=0; i< pos - 1 && i++)
```

```
{  
    temp = temp->next;  
}
```

$d_2 \rightarrow \text{temp} \rightarrow \text{next};$
 $\text{temp} \rightarrow \text{next} = d \rightarrow \text{next};$
 $d \rightarrow \text{next} = \text{NULL};$
 $\text{free}(d);$

Display :-



if (head == NULL)

{

printf ("List is Empty");
exit(0);

}

else.

{ temp

temp = head;

while ((temp → next → next != NULL))

{
printf ("%d\n", temp → data);

temp = temp → next;

}

}

Start node.

```

{
    int data;
    struct node *next;
}
struct node {
    int data;
    struct node *next;
}
temp, head;
    
```

(2000) →

(3000) →

(3000) →

(NULL)

(2000) →

(3000) →

(3000) →

(NULL)

(2000) →

(3000) →

(3000) →

(NULL)



VII operation on double Linked List :- There are 3 operations
in double linked list

(i) insertion

(ii) deletion

(iii) Display.

(i) Insertion:- There are 3 types of insertion.

(i) insertion at Beginning

(ii) insertion at ending

(iii) insertion at any specific position.

(i) insertion at Beginning :-

new

head



int Value;

new = (Struct node*) malloc (Sizeof (Struct node));

printf ('Enter Value\n');

.scanf ('%d', &value);

new → data = value;

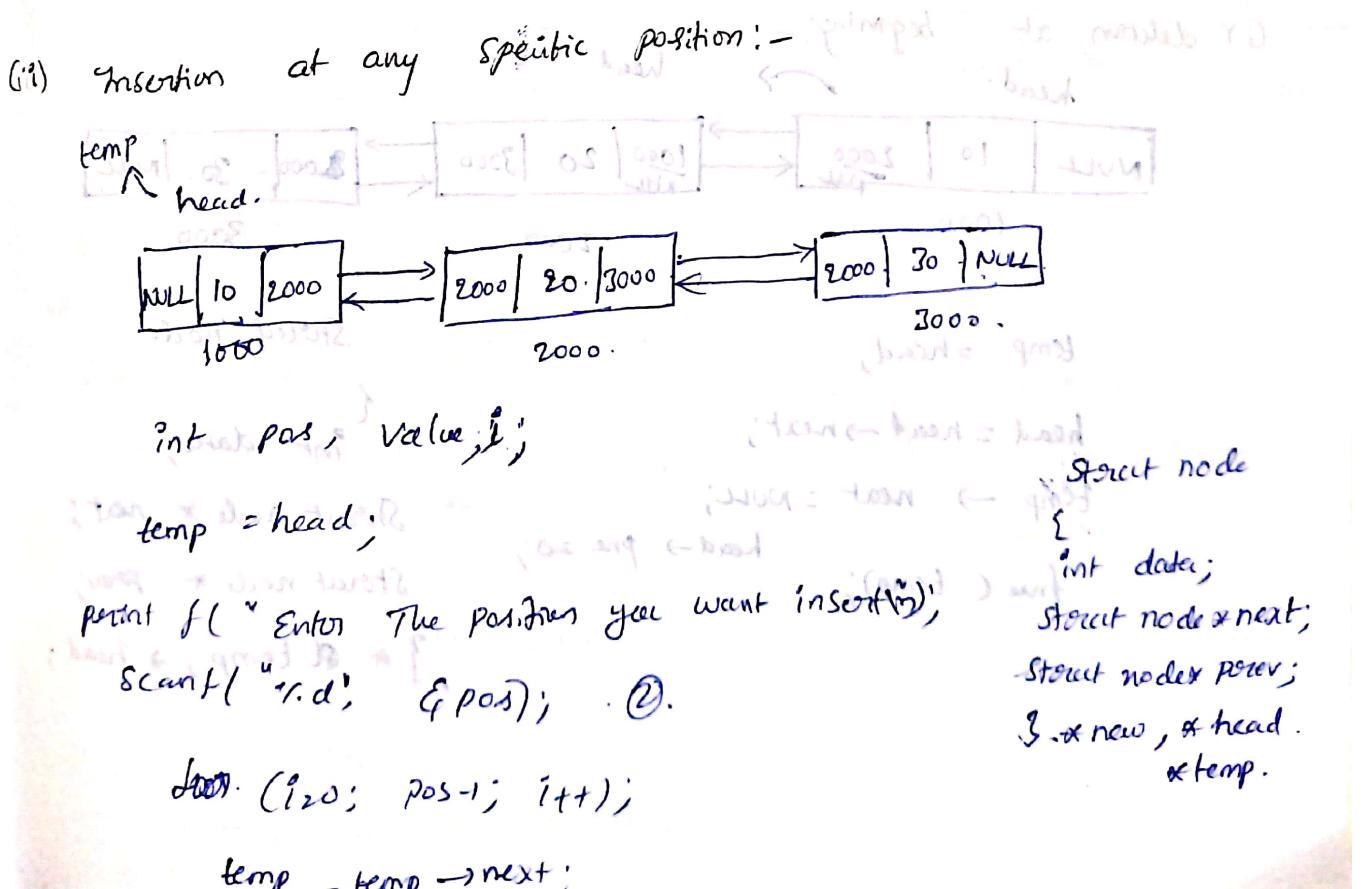
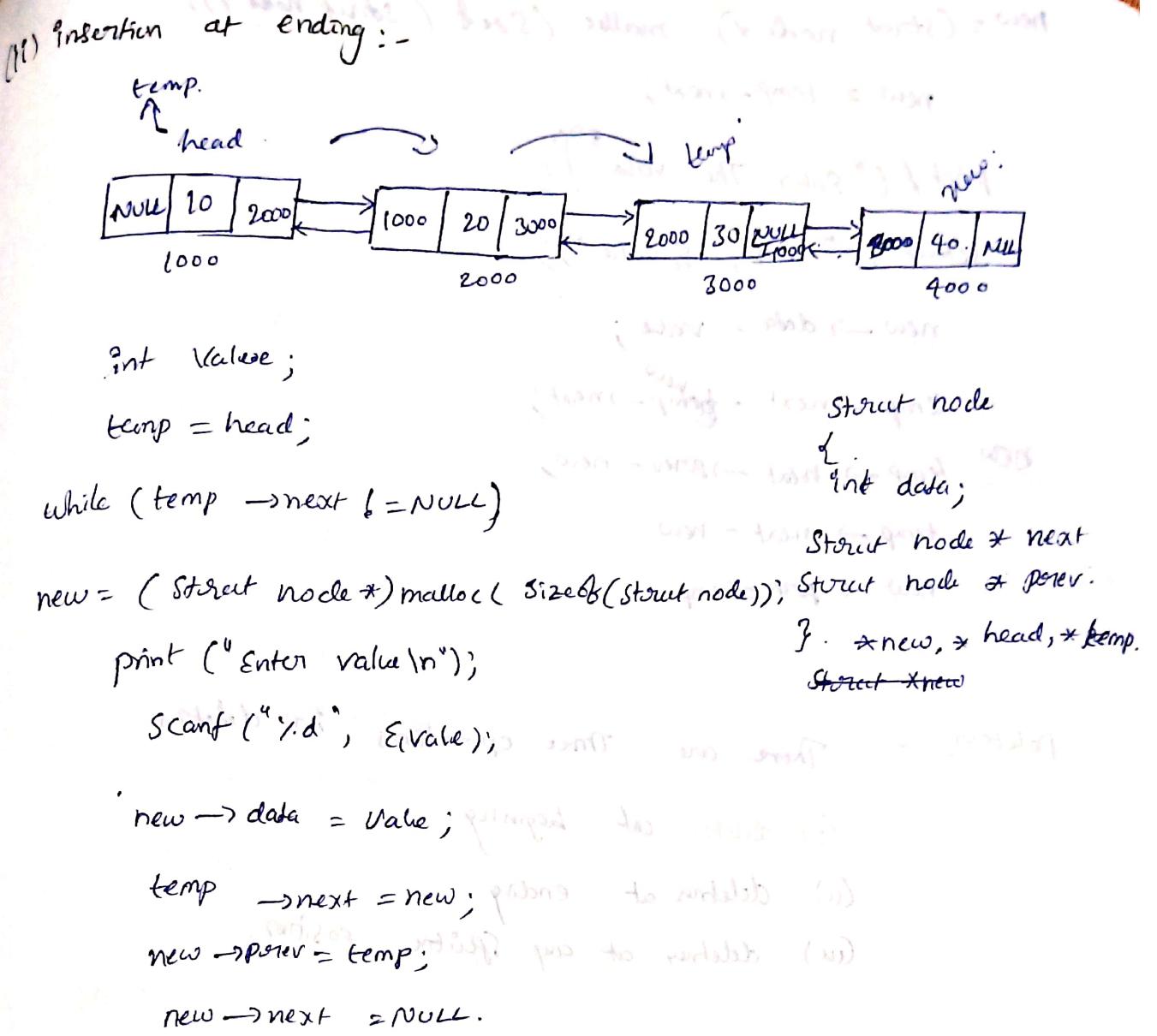
new → next = head;

head → prev = new;

new → prev = NULL;

head = new;





new = (struct node*) malloc (sizeof (struct node));

new->temp->next;

printf ("Enter The value: ");

scanf ("%d", &value);

new->data = value;

temp->next = new;

new->temp->next->prev = new;

temp->next = new;

new->prev = temp->next; // deletion at end of list

free (temp);
temp = NULL;

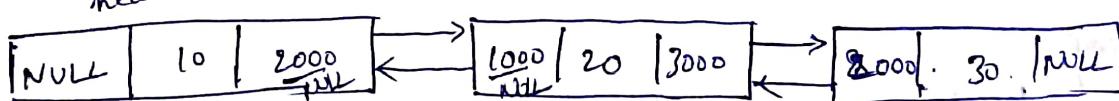
Deletion :- There are three operations in deletion

(i) deletion at beginning; // head = start - next

(ii) deletion at end; // start - last - next

(iii) deletion at any specific position.

(i) deletion at beginning:-



temp = head;

head = head->next;

temp->next = NULL;

free (temp);

head->prev = NULL;

Struct node * next;

Struct node * prev;

} * & temp, * head;

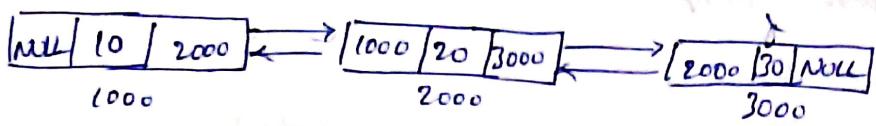
(free (temp); free (head);)



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(ii) Deletion at ending :-

head.



temp = head;

while (temp → next → next != NULL)

{

temp = temp → next;

}

d = temp → next;

temp → next = NULL;

d → prev = NULL.

free (d);

Struct node.

{ int data;

Struct node * next;

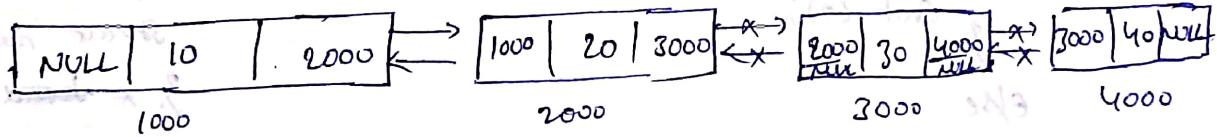
Struct node * prev;

} * temp, * head, * d;



(ii) Deletion at Specific position! -

head:



int, pos, ?

temp = head;

printf ("Enter position you want ");

scanf ("%d", & pos);

Struct node.

for (i = 0; i < pos-1; i++) {

 * /

 i < 1

}

temp = temp → next;

}

d = temp → next;

temp → next = d → next;



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$d \rightarrow next \rightarrow pow = temp;$

$d \rightarrow next \rightarrow pow = temp.$

$d \rightarrow next = null;$

$d \rightarrow pow = null;$

$free(d);$

$return;$

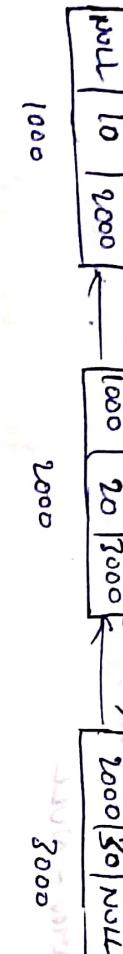
($head = 1000$ → 1000 → 2000) address

last node address

(1000 → 2000 → 3000) address

$Display:$ → 1000 → 2000 → 3000

$head:$



$if (head == 0 (NULL))$

{

printf ("List is Empty");

short node *temp;

int data;



$else$ → 1000

{

temp = head;

while (temp)

while (temp → next != NULL)

→ 2000 → 3000

{

printf ("Node %d", temp → data);

temp = temp → next;

→ 3000

{

return 0;



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`temp = temp -> next;`

`printf("%d", temp -> data);`

3.

circular linked list :- (Circular Singularity) :-
Data + Next

There are three operations in circular linked list

(i) insertion

(ii) deletion

(iii) Display.



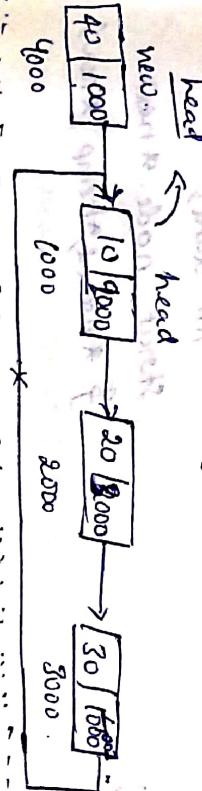
(a) insertion at beginning

(b) insertion at ending

(c) insertion at specific position.

insertion steps

insertion of at beginning:



int value;

new = (struct node *) malloc (sizeof (struct node)); struct node

temp = head;

while (temp->next != head)

i.

temp = temp -> next;

3

temp -> next = new;



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printf("Enter The value");

scanf("%d", &value);

new->data = value;

new->next = head;

head = new;

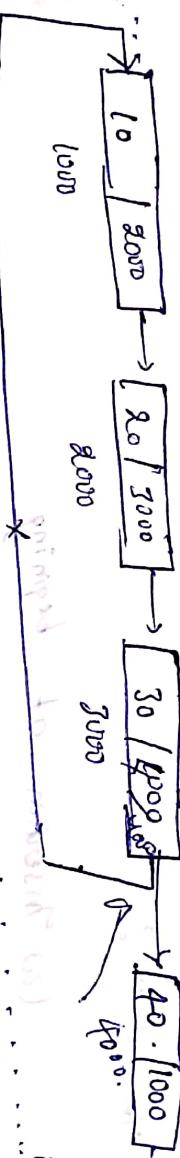
(ii) Insertion at end :-

temp.

head.

tail->next = new;

new



int Value;

temp = head;

while (temp->next != head)

Struct node

{printf("Enter the value");

scanf("%d", &data);

Struct node *next;

Struct node *next;

{temp->next->next = head;

g.

new2 = (Struct node*) malloc (sizeof(Struct node));

new2->data = value;

new2->next = head;

printf("Enter The Value");

scanf("%d", &value);

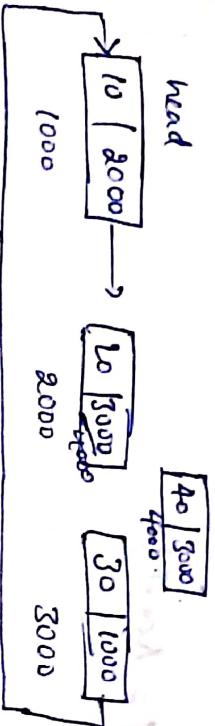
new->data = value;

temp->next = new;



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(ii) insertion at any specific position:-



```
int value, pos, i;
```

Struct node.

```
printf("Enter the position you want");
```

{
int data.

```
scanf("%d", &pos);
```

Struct node *next;

```
temp = head;
```

```
for (i=0; i<pos-1, i++)
```

temp = temp->next;

```
for (i=0; i<pos, i++)
```

(Or) next

```
temp = temp->next;
```

```
new = (struct node*) malloc(sizeof(struct node));
```

new: New

```
printf("Enter the value");
```

Value

```
scanf("%d", &value);
```

Value

```
new->data = value;
```

```
new->next = temp->next;
```

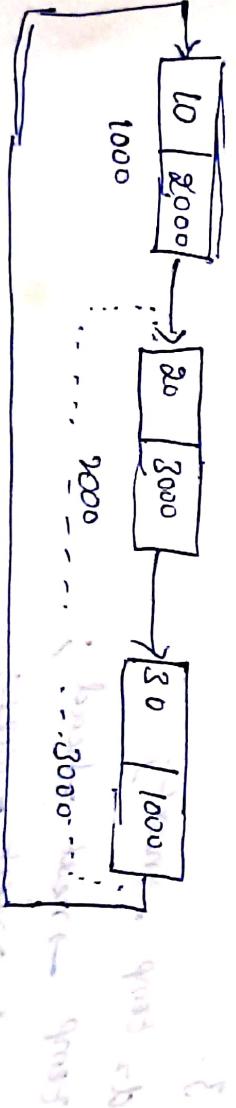
Next

```
temp->next = new;
```

Next

Problem :- (a) deletion at beginning.

head.



(b) with Struct node.

```
int data;
```

```
Struct node *next;
```

```
Struct Temp, *head, *xd;
```



$\text{temp} = \text{head};$

while ($\text{temp} \rightarrow \text{next} \neq \text{head}$)

{

$\text{temp}_0 = \text{temp} \rightarrow \text{next};$

3.

$d = \text{head};$

$\text{head} = \text{head} \rightarrow \text{next};$

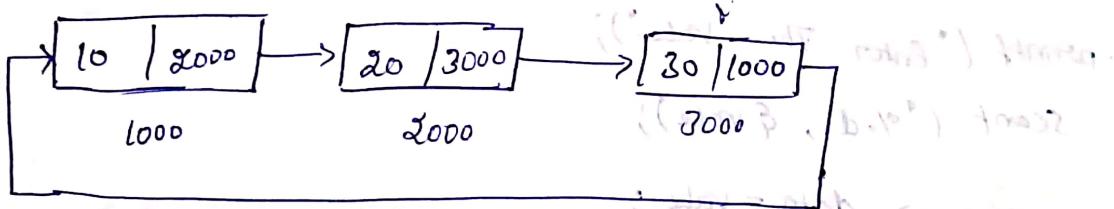
$\text{temp} \rightarrow \text{next} = \text{head};$

$d \rightarrow \text{next} = \text{NULL};$

$\text{free}(d);$

(ii) ~~to~~ detection of deletion at ending point :-

$\text{temp} = \text{head}.$



$\text{temp} = \text{head};$

while ($\text{temp} \rightarrow \text{next} \rightarrow \text{next} \neq \text{head}$)

{

$\text{temp} = \text{temp} \rightarrow \text{next};$

3.

$d = \text{temp} \rightarrow \text{next};$

$\text{temp} \rightarrow \text{next} = \text{head};$

$d \rightarrow \text{next} = \text{NULL};$

$\text{free}(d);$

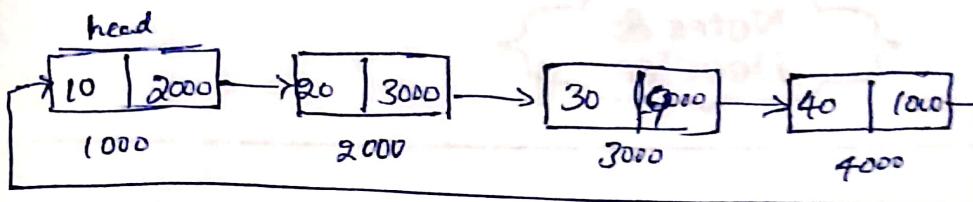
then free head;

else handle error;



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(iii) Deletion at any specific position:-



```
int pos, i;
temp = head;
printf ("Enter position you want\n");
scanf ("%d", &pos);
for (i=0; i<pos-1; i++)
{
    temp = temp ->next;
}
d = temp ->next;
temp ->next = d ->next;
d ->next = NULL;
free (d);
```

- returns pointer

Struct node

{

int data;

Struct node * next;

} *temp * head * d *t

3.

d = temp -> next;

temp ->next = d ->next;

d ->next = NULL;

free (d);

(head = (Struct *) malloc (sizeof (Struct)))

((Struct *) malloc (sizeof (Struct)))

temp = head;

((Struct *) malloc (sizeof (Struct)))

Notes & Doodles

display operation:-

10	12000
----	-------

20	3000
----	------

30	1000
----	------

data structure

if (`head == NULL`)

```
printf("List is Empty.");
exit(0);
```

else

{

`temp = head;`

{

`printf("%d", temp->data);`

`temp = temp->next; (temp = temp->next)`

{

`printf("%d", temp->data);`

`temp = temp->next;`

{

`printf("%d", temp->data);`

`temp = temp->next;`

}{

Refer Note 2 for 4 unit