here 
$$x \gg 0$$
 ...  $E(X) = \sum_{x \gg 0}$ 

$$P_{P}(x>0) = P_{P}(x=1) + P_{P}(x=2) + P_{P}(x=3) + P_{P}(x=3) + P_{P}(x=3)$$

in of 1 Pr (x>1) = Por in ed DO

$$Pr(x=1)+Pr(x=2)+Pr(x=3)+---$$

puting x=2

putting x=K

in ed 0

$$PP(X>0) + PP(X>1) + PP(X>2) = 1. PP(X=1) + 2. PP(X=2) + 3. PP(X=3)$$

$$+ + - \cdot \cdot + k. PP(X=K=1) + (K+1) PP(X=K+1)$$

+ - - . 0 ⇒ Em(x>x)

XNO

$$= \sum_{x \ge 0.1} x \cdot PP(x = x) + 0 \cdot PP(x = 0)$$

$$= \sum_{x \ge 0.1} x \cdot PP(x = x) = E(x) [from @]$$

 $: | E(X) = \sum Pr(X > X) | (boroved)$ 

Criven a permutation IT of {1,2,3,...,n}, i ∈ {1,2,3,...,n} is said to be a fixed point of TT if TT(i)=i. Let o be a grandom permutation of [1,2,3,...,n] X is a pandom variable corresponding to the number of fixed points in 5. Let's consider Xi a pardom variable such that  $Xi = \begin{cases} 1 & \text{if } TT(i) = i \text{ (th point is fixed)} \\ 0 & \text{if } TT(i) \neq i \text{ (ith point is not fixed)} \end{cases}$ obability of  $X_i = P(X_i) - \Gamma(n-1)I$ 

Then, Probability of  $X_i = P(X_i) = \left\{ \frac{(n-1)!}{n!} = \frac{1}{n} ; \text{ if } X_i = 1 \right\}$ The number of [ 1 - 1 ; if Xi=0]

ways are n

points can be points can be arranged of : (100)

:. Expectation of 
$$Xi = E(Xi) = \sum XiP(Xi)$$

$$= 1 \cdot \frac{1}{n} + o(1 - \frac{1}{n}) = \frac{1}{n}.$$

 $X = X_1 + X_2 + X_3 + X_4 + \cdots + X_n$ . Now we can definitely say that

AS A POSTU

Hence, from linearity of expectation we can say that,

 $E(X) = E(X_1) + E(X_2) + E(X_3) + --- + E(X_n)$  $= n \cdot \frac{1}{n} = 1$ total  $= n \cdot \frac{1}{n} = 1$  = 1 = 1 = 1 = 1

 $\therefore E(x) = 1 \quad \text{(for (Ans))}$ 

(b) We need to find PMF of x. Let's first book at the formula for derangement. A permutation of the elements of a set where no element permains at the same place is called denangement. It is denoted by Dr. The number of derangements of a set with n objects is given by the formula  $D_n = n! \sum_{i=0}^n \frac{(-1)^{i}}{i!}$ Now, Probability that no. of fixed points in o 18 K Those K fixed only

A points have to -> Rest (n-15) points MCK · (1) · (Desargement of (n-K) points) should be arranged in Such a way that, no one choose K points should be among n points all possible arrangement at their which will be correct pos. considered as fixed of n points. 1.e.π(i) ≠ i to maintain the fact that loco not multer be sound at (L) exactly K points are fixed putting n = n-k in ed (1)  $\therefore PMF(X) = P(X = K) = \frac{1}{K!} \sum_{i=0}^{n-K} \frac{(-1)^i}{i!} \left( \underbrace{An8} \right)$ 

We have to find expected number of samps in a uniformly nandom permutation  $\sigma$ .

Let's define a nandom variable  $Y_{ij}$  in such a way that  $Y_{ij} = \begin{cases} 1 & \text{if } TT(i) = j \text{ and } TT(j) = i \text{ i.e. i.i.s. injth position and} \\ j & \text{is in th position.} \end{cases}$ Otherwise two points in carbon or are in place j; there is no points of j are in place j; if j are in place j; if j and j are in place j; if j are in place j are in place j; if j are in place j and j are in place j; if j are in place j and j a

$$\frac{1}{n!} \cdot \mathbb{E}(\lambda^{n}) = 1 \cdot \frac{1}{n(n-1)} + 0 \cdot \left(1 - \frac{n(n-1)}{n(n-1)}\right)$$

.. It will be equal to number of ways we can choose 2 points from n points. 
$$= {}^{n}C_{2} = \frac{n(n-1)}{2}$$

Now, let y be a grandom variable which denotes expected no. of swaps.

Hence we can say that 
$$y = \sum_{i < j} y_{ij}$$
  
 $i < j$   
 $1 \le i, j \le n$ .

PMF(X)=

$$E(y) = E\left(\sum_{\substack{i < i \\ i \neq i, j \neq n}} y_{ij}\right)$$

$$= \frac{N(N-1)}{2} \cdot E(Y_{ij})$$
total no. 
$$= \frac{N(N-1)}{2} \cdot \frac{1}{N(N-1)}$$
of such totals.

d) We have to show that 
$$P(X>10) \leq \frac{1}{10}$$

We have already proved that, E(X)=1 [ where X is a r.v. corresponding to the no From Markov's inequality, we know that

$$P(X \geqslant K) \leq \frac{E(X)}{K}$$

$$\Rightarrow P(X \geqslant 10) \leq \frac{1}{10} \left[ \text{butting } K = 10 \right]$$

We know that if A 18 a Subset of B i.e. ACB

Then 
$$P(A) \leq P(B)$$
.



For any edge e E E. Let Xibe pandom variable Which is I when eis monochomatic and o otherwise.

The warm work of extension ment ..

We have to show that set of random variable [Xe] eff are painwise independent. Show that they are not independent.

$$\begin{array}{c|c}
(2) & xe=0 \\
Xe=1 & R & xe=1 & R \\
(2) & Xe=0 & B \\
Xe=1 & B
\end{array}$$

$$\begin{array}{ll} C_{2} & xe=0 \\ Xe=1 & R & xe=1 R \\ C_{2} & xe=0 \end{array} \qquad \begin{array}{ll} Xe=\left\{ \begin{array}{ll} 1 & \text{if eis monochromatic} \\ 0 & \text{if eis non-chromatic} \end{array} \right\}. \end{array}$$

Xe=1) B Let's define a sample space which denotes the two endpoints colour of any edge

$$P(Xe) = \begin{cases} 3/9 = 1/3 & \text{if } Xe = 1 \\ 16/9 = 2/3 & \text{if } Xe = 0 \end{cases}$$

[ case 1] ( Two edges are not connected)

(hovered)

[ case I] (Two edges are not connected)

i. J. We need to show, 
$$P(Xe_i = 1, Xe_j = 1) = P(Xe_i = 1) P(Xe_j = 1)$$

No of ways s.t. (colous(i) = colous(i2))

Xei  $Xe_j$  LHS  $P(Xe_i = 1, Xe_j = 1) = \frac{x \text{ no of ways s.t.}(colous(j,) = colous(j,2))}{x \text{ no of all possible colourings}}$ 

of  $\int i_1, i_2, j_3, j_2$ 

$$=\frac{1}{3^2}=\frac{1}{9}$$

P(
$$Xe_i = 1$$
) =  $\frac{1}{3}$  (from  $\emptyset$ )

P( $Xe_j = 1$ ) =  $\frac{1}{3}$  (from  $\emptyset$ )

P( $Xe_i = 1$ ). P( $Xe_j = 1$ ) =  $\frac{1}{3} \cdot \frac{1}{3} = \frac{1}{4}$ 

P( $Xe_i = 1$ ,  $Xe_j = 1$ ) = P( $Xe_i = 1$ ). P( $Xe_j = 1$ )

We need to Show P( $Xe_i = 1$ ,  $Xe_j = 1$ ) = P( $Xe_i = 1$ ) P( $Xe_j = 1$ )

Let's consider  $Xe_i$  more charactic and Let accoun ( $i_1$ ) = eclass ( $i_2$ ) = R

Hence, if  $Xe_j$  has to be more charactic than is colar ( $i_1$ ) has to R

P( $Xe_i = 1$ ,  $Xe_j = 1$ )

To of ways (colars( $i_1$ ) = colars( $i_2$ ) = accoun( $i_3$ ) =  $i_1$  and  $i_2$  proof ways of eccansing of ( $i_1$ ,  $i_2$ ,  $i_3$ )

P( $Xe_i = 1$ ) =  $i_1$  . P( $Xe_i = 1$ ). P( $Xe_j = 1$ ) =  $i_2$  . If  $i_3$  =  $i_4$  P( $Xe_i = 1$ ) =  $i_3$  . P( $Xe_i = 1$ ) P( $Xe_i = 1$ ) P( $Xe_i = 1$ ) P( $Xe_i = 1$ )

P( $Xe_i = 1$ ,  $Xe_j = 1$ ) = P( $Xe_i = 1$ ) P( $Xe_i = 1$ )

P( $Xe_i = 1$ ,  $Xe_j = 1$ ) = P( $Xe_i = 1$ ) P( $Xe_i = 1$ ) P( $Xe_i = 1$ )

## Proving non-independence (Omo A) L= (1- 10x)9 To a prove non-independence Consider eyele C3 We need to prove at least one of these Xei Xej 3 1) $P(Xe_i=1, Xe_j=1) \neq P(Xe_i=1) P(Xe_j=1)$ ii) P(Xei=1, Xex=1) + P(Xei=1) P(Xex=1) ii) P(Xej=1, Xex=1) = P(Xej=1) P(Xex=1) P(XCC=1) P(XG=1) in) P(Xei=1, Xej=1, Xex=1) + P(Xei=1) P(Xej=1) P(Xex =1) Lets try to prove iv P(Xei=1, Xej=1, Xex=1) Let's consider Xei mono-chaomatic Hence colour (1)=colour(2) - 0/9/11/20 50, to if Xej and Xex wants to be mono-chromatic then $colour(1) = colour(2) = colour(3) = {B, (2, R}.$ LHS .. P(Xei=1, Xej=1, Xex=1) = no of ways (cdan(1)=colous(2)=colous(3)) no of ways of colouring [1,23] $z = \frac{3}{3^3} = \frac{1}{3^2} = \frac{1}{9}$ P(Xei=1)===== ... P(Xei=1). P(Xej=1) P(Xex=1) =1 P(Xej=1) = = = = = = $P(Xex=1)=\frac{1}{3}$ [.. P(Xei=1, Xej=1, Xek=1) + P(Xei=1) P(Xej=1) P(Xek=1)

... Non-independence provid

Y is the random variable corresponding to the number of non-monochromatic edges.

Let's define pandom variable Yi Such that

$$P(Yi) = \begin{cases} 6/9 = 2/3 & \text{if } Yi = 1 \\ 3/9 = 1/3 & \text{if } Yi = 0. \end{cases}$$

$$RR, BB, G2(2, RB, BR, RG, GR, RB)$$

$$RG, GB$$

$$E(Yi) = \sum_{i} Yi P(Yi)$$

$$= 1 \times \frac{2}{3} + 0.\frac{1}{3} = \frac{2}{3}$$

Now, we can say that Y = & Yi

From linearity of expectation we can write

$$E(Y) = E(\Sigma Y i)$$

$$= \sum E(Y i) = \sum E(Y i) \times E(Y i)$$

$$= |E| \times \frac{2}{3}$$
 where 
$$= \frac{2|E|}{3}$$
 |  $|E| = \text{total mof edges}$  in Ozraph

We have to show that there can't be any graph for which all 3-colour assignments make $\langle \frac{2 E }{3}   edges \rangle$
Let's try to prove it by contradiction  Elet's suppose that there can be some graph for which all  3-colour assignments make < 2/El edges non-monochromati
If that's true then $E(y) < \frac{2 E }{3}$ (Since all passion and the state of those numbers are less than x, then the day average of those numbers
But, in 4.2 we have proved that $E(y) = \frac{2 E }{3}$ Hence, our assumption was wrong.
There can't be any graph for which all 3-colour assignments make < 2/El edges non-monochromatic (broved)

(4) We have to show that  $P(y \gg \frac{|E|}{2}) \gg \frac{1}{3}$ 

We know that y is the random variable corresponding to the

number of ma non-monochromatic edges.

We have already proved in 4.2 that E(y) = 2|E| [Where |E| = total no of eages]

NOW. Let X be the grandom variable corresponding to the

number of monochromatic edges. if edge is chromotic Let also Xe be the r.v. Such that Xe= 1 if ,, is nonmono chromatic.

$$P(Xe) = \begin{cases} 3/4 = 1/3 & \text{if } Xe = 1 \\ 6/9 = 21/3 & \text{if } Xe = 0 \end{cases}$$
  

$$\therefore E(Xe) = 1 \times 1/3 + 0 \times 21/3 = 1/3$$

$$X = \sum X e$$

$$E(X) = E(X) = |E| \cdot E(X) = |E| \cdot |S| = \frac{|E|}{3}$$

$$[|E| = \text{total no of edges}]$$

NOW, P(Y> [E])  $P(Y) = \frac{1}{2}$   $= 1 - P(X) = \frac{1}{2}$ [ as an edge is either mono-chromatic or non mono-chromatic

$$= 1 - \left( \frac{E(x)}{\frac{|E|}{2}} \right) \left[ \begin{array}{c} From Markov's inequality \\ E(x > K) \leq \frac{E(x)}{K} \end{array} \right]$$

$$= 1 - \left( \frac{|E|}{\frac{|E|}{2}} \right) \left[ \begin{array}{c} \vdots E(x) = \frac{|E|}{3} \\ \hline \end{array} \right]$$

$$= 1 - \left( \frac{|E|}{3} \right) > \frac{1}{3} \left( \begin{array}{c} proved \\ proved \end{array} \right)$$

(5) We have to devise a method that can find an assignment for which the number of non-monochromatic edges is at least |E1/2 with probability at least 99/100. Lots try to devise the agorithm first. Algorithm (input: any Geraph X) Step 1: Let X = {X1, X2, X3, -... Xx} be a let of K randomly beleated assignments. Step 2: Lot B Xi = ith assignment of graph 5tep 3: If the assignment Xi contains no number of non-momehomatic edges > [E] return the assignment Xi ELSE check for next assignment Xi+1 Now let's try to find how many assignments we need to check to achieve probability at least 99/100. Let Czi be a nandom variable Such that  $G_{i} = \begin{cases} 1 ; \text{ if ith assignment } X_{i} \text{ has } \Rightarrow \frac{1EI}{2} \text{ mono-chomatric} \end{cases}$ edges = 1-6(x>.5 Now from result of 44 P (ai = 1) = 2/3 : E (ai) = 1×2/3 = 2/3

Now let's define a random variable or such that

$$G_2 = \frac{\sum_{i=1}^{K} G_{2i}}{K} = \frac{G_1 + G_2 + G_3 + \dots + G_2 K}{K}$$

From linearity of expectation,

wity of expectation;
$$E(\alpha) \doteq \frac{E(\alpha_1 + \alpha_2 + \dots + \alpha_K)}{K}$$

$$\Rightarrow E(\alpha) = \frac{E(\alpha_1) + E(\alpha_2) + \dots + E(\alpha_K)}{K}$$

$$\Rightarrow E(\alpha) = \frac{K \cdot E(\alpha_1)}{K}$$

$$\Rightarrow E(\alpha) = \frac{K \cdot E(\alpha_1)}{K}$$

$$\Rightarrow E(\alpha) \leq \frac{2}{3} \quad \left[ : E(\alpha_1) \leq \frac{2}{3} \right]$$

(2=1=> (21, (22, -1) ax all are I irdividually

i.e all assignments { (21, (22, -., (2x) have mono-chromatic edges > IEI

have non-monochromatic Which implies all " edges L | El

which implies our above stated algorithm failed to find any assignment (2: with non-monochromatic edges > [E]

NOW, from S.  $P(\Omega=1) \leq (1-\frac{99}{100})$ 

Now. 62 can hold max. value =1

$$P(\alpha > 1) = P(\alpha = 1)$$

$$P(\alpha > 1) \leq 0.01$$

$$\Rightarrow P(\alpha > = 1) = P(\alpha - E(\alpha)) = 1 - E(\alpha)$$

$$\Rightarrow P(\alpha > = 1) = P(\alpha - E(\alpha)) > = P(\alpha - E(\alpha))$$

$$\Rightarrow P(\alpha > = 1) = P(\alpha - E(\alpha)) \Rightarrow P(\alpha - E(\alpha)) = 1 - E(\alpha)$$

$$\Rightarrow P(\alpha - E(\alpha)) > = 1 - E(\alpha)) \Rightarrow P(\alpha - E(\alpha)) > = 1 - E(\alpha)$$

$$\Rightarrow P(\alpha - E(\alpha)) > = 1 - E(\alpha)) \leq P(\alpha - E(\alpha)) > = 1 - E(\alpha)$$

Now from charyster we know  $P(|X-u|\gg K) \neq \frac{Var(X)}{K^2}$  $= P(|(\alpha - E(\alpha))| > = 1 - E(\alpha)) \le \frac{Van(\alpha)}{(1 - E(\alpha))^2}$ £ 4K(1-E(6))2  $\frac{1}{4\kappa(1-\frac{2}{3})^2}$ (1) = (30) 3 2 ] = (4 × (1/3)2 - , as alk are I individually & 4 14Kendari danlar Now, P((2>=1) \le 0.01  $\Rightarrow P(|\alpha - E(\alpha)|) \leq 0.01$ Non from S. P (00) = (100) = (100) => 4K > 000 1 0.01 SX X 7 8000 => 4K > 100 25 => K > 25×9 which weams we need to aleast check 225 assignments of graph to achieve probability > K7 225 E(62) & P(1 7/ 100