Chosen-ciphertext Attack (CCA)

Assumptions: I've used **output feedback mode (OFB)** to design CPA and **Cipher block chaining (CBC)** to design MAC.

Input format:

- 1. It will ask for a safe prime number p in integer format.
- 2. It will ask for a generator (primitive root) of that prime.
- 3. It will then ask for input the block length.
- 4. It will then ask for two keys k1 and k2 in binary form of block length.
- 5. It will ask to enter the data in binary form.
- 6. It will also ask for seed for generating initial vector via PRG.

Output Format:

- 1. It will first generate the ciphertext by OFB mode and CPA.
- 2. Then it will generate the MAC TAG by CBC mode on the ciphertext.
- 3. On the verification end, the MAC TAG will be verified first.
- 4. If the MAC matches, then the ciphertext will be decrypted to see if we get back the actual data. If yes, then it is successful.

```
PS C:\Users\Sudipta Halder\Desktop\IIITH ASSIGNMENTS\POIS> python .\5_cca_encryption.py
Enter the prime number(The prime should be such that p-1/2 should also be prime. Sophie Germain Prime)(1907): 31 Enter the generator(Primitive root for the prime)(987): 31 Enter (block_size = key_size) for the data: 4 Enter the key k1 in binary of length 4: 1001 Enter the key k2 in binary of length pref diff from k1 4: 1010 Enter the data to in binary(preferably of length multiple of 4): 1100011001001011001
Round #1
The Initial Vector is: 1111
The PRF is : 1111
 The data is : 1100
 Round #2
The Initial Vector is: 1111
The PRF is : 1111
The data is : 0110
Round #3
 The Initial Vector is: 1111
The PRF is: 1111
 The data is : 0100 Round #4
 The Initial Vector is: 1111
 The PRF is : 1111
 The data is : 1010
 Round #5
 The Initial Vector is: 1111
The PRF is: 1111
The data is: 0101
 Round #6
The Initial Vector is: 1111
The PRF is: 1111
The data is: 1001
The encrypted data is: 001110011011010110100110
```

```
Received CBC MAC TAG: 1111
Received cipher text: 001110011011010110100110
Round #1
ti: 1111, mi: 0011
t_xor_mi: 1100
prf_res: 1111
Round #2
ti: 1111, mi: 1001
t_xor_mi: 0110
prf_res: 1111
Round #3
ti: 1111, mi: 1011
t_xor_mi: 0100
prf_res: 1111
Round #4
ti: 1111, mi: 0101
t_xor_mi: 1010
prf_res: 1111
Round #5
ti: 1111, mi: 1010
t_xor_mi: 0101
prf_res: 1111
Round #6
ti: 1111, mi: 0110
t_xor_mi: 1001
prf_res: 1111
CBC-MAC: 1111
Mac Tag is verified. Sucess
```

```
*******************Decryption Stage**************
['0011', '1001', '1011', '0101', '1010', '0110']
Round #1
The Initial Vector is: 1111
The PRF is : 1111
The data is : 0011
Round #2
The Initial Vector is: 1111
The PRF is : 1111
The data is: 1001
Round #3
The Initial Vector is: 1111
The PRF is : 1111
The data is: 1011
Round #4
The Initial Vector is: 1111
The PRF is : 1111
The data is: 0101
Round #5
The Initial Vector is: 1111
The PRF is: 1111
The data is: 1010
Round #6
The Initial Vector is: 1111
The PRF is : 1111
The data is: 0110
Decrypted data we got after verifying: 110001100100101001011001
MSG and MAC TAG both verified. Successful
```

Working Flow:

- 1. Upon getting the input, the function `generate_cpa_ofb(prime, generator, data, block size, key)` will be called.
- 2. It will generate the cipher text c using key k1.
- 3. Then the ciphertext will be passed to `generate_cbc_mac(prime, generator, key, data, block_size) `and it will give the MAC TAG t using key k2.
- 4. After that, `verify_mac_tag_encrypted_data(prime, generator, k1, k2, block_size, encrypted_data, initial_vector, cbc_mac_tag) ` will be called to verify the correctness of the tag.
- 5. Upon successful verification of the MAC tag, `cpa_ofb_decryption(encrypted_data, initial_vector, block_size, prime, generator, key)` will be called to verify the ciphertext. It will be decrypted and matched with the original data to see the correctness.
- 6. Upon successful verification, we would respond `MSG and MAC TAG both verified. Successful`.

CONSTRUCTION 4.17 CCA-secure encryption.

Define a CCA-secure encryption scheme as follows:

- $\operatorname{\mathsf{Gen}}'(1^n)$: upon input 1^n , choose $k_1, k_2 \leftarrow \{0, 1\}^n$
- $\operatorname{Enc}'_k(m)$: upon input key (k_1, k_2) and plaintext message m, compute $c = \operatorname{Enc}_{k_1}(m)$ and $t = \operatorname{Mac}_{k_2}(c)$ and output the pair (c, t)
- $\mathsf{Dec}'_k(c,t)$: upon input key (k_1,k_2) and ciphertext (c,t), first verify that $\mathsf{Vrfy}_{k_2}(c,t)=1$. If yes, then output $\mathsf{Dec}_{k_1}(c)$; if no, then output \perp .