

User Authentication in Wireless Sensor Networks

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Temporal credential-based three-factor user authentication for distributed wireless sensor networks

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Necessity for user authentication

- Most queries in wireless sensor network (WSN) applications are issued at the point of the base station or gateway node of the network.
- However, for critical applications of WSNs (e.g., battle field, healthcare application) there is a great need to access the real time data inside the WSN from the nodes, because the real-time data may no longer be accessed through the base station only.
- The real-time data can be given access directly to the external users (parties) those who are authorized to access data as and when they demand.
- The user authentication plays a vital role for this purpose.

Three factors used in the designed scheme

- Smart card
- Password
- Personal biometrics (for example, fingerprints, faces, irises, hand geometry and palm-prints, etc.)

- Uses the user's personal biometrics along with traditional password to design user authentication protocols in WSNs.
- The biometric verification allows one to confirm or establish an individual's identity.
- There are major advantages of using biometric keys (for example, fingerprints, faces, irises, hand geometry and palm-prints, etc.):
 - ▶ Biometric keys can not be lost or forgotten.
 - ▶ Biometric keys are very difficult to copy or share.
 - ▶ Biometric keys are extremely hard to forge or distribute.
 - ▶ Biometric keys can not be guessed easily.
 - ▶ Someone's biometrics is not easy to break than others.

- The output of a conventional hash function $h(\cdot)$ is sensitive and it may also return completely different outputs even if there is a little variation in inputs.
- The biometric information is prone to various noises during data acquisition, and the reproduction of actual biometric is hard in common practice.
- To avoid such problem, a fuzzy extractor method is preferred, which can extract a uniformly random string and a public information from the biometric template with a given error tolerance t .

Definition

The fuzzy extractor is a tuple (\mathcal{M}, l, t) , which is composed of the following two algorithms, called *Gen* and *Rep*:

- **Gen:** It is a probabilistic algorithm, which takes a biometric information $B_i \in \mathcal{M}$ as input, and then outputs a secret key data $\sigma_i \in \{0, 1\}^l$ and a public reproduction parameter τ_i , where $Gen(B_i) = \{\sigma_i, \tau_i\}$.
- **Rep:** This is a deterministic algorithm, which takes a noisy biometric information $B'_i \in \mathcal{M}$ and a public parameter τ_i and t related to B_i , and then it reproduces (recovers) the biometric key data σ_i . In other words, we have $Rep(B'_i, \tau_i) = \sigma_i$ provided that the condition $d(B_i, B'_i) \leq t$ is met.

- The probability to guess the biometric key data $\sigma \in \{0, 1\}^l$ by an attacker is approximately $\frac{1}{2^l}$, where $l = m - 2 \log(\frac{1}{\epsilon}) + O(1)$, where ϵ is the statistical distance between two given probability distributions, and m is the min-entropy given as follows. the min-entropy $H_\infty(A)$ of a random variable A is $-\log(\max_a \Pr[A = a])$.

Vanga Odelu, Ashok Kumar Das, and Adrijit Goswami. “A Secure Biometrics-Based Multi-Server Authentication Protocol using Smart Cards,” in *IEEE Transactions on Information Forensics and Security*, Vol. 10, No. 9, pp. 1953 - 1966, 2015, DOI: 10.1109/ TIFS.2015.2439964. (2021 SCI Impact Factor: 7.231) [This article is one of the top 50 most frequently downloaded documents for Popular Articles (June - November 2015)]

Threat Model In the following we consider the three types of models:

- **Honest-but-Curious adversary model:** This model [HCAM] is a passive adversarial model where the adversary \mathcal{A} will behave like a legitimate entity and follow the specified protocol. However, \mathcal{A} can read all the transmitting information between the corrupted entities in the network.
- **Dolev-Yao (DY) threat model:** This model is known as the DY model [DYM]. In the DY model, an adversary \mathcal{A} has the potential ability to eavesdrop, intercept, modify and delete messages that are being communicated among various agents through a wireless network.
- **Canetti and Krawczyk's model:** This model is also known as the “CK-adversary model” [CKM]. Keeping all the fundamental assumptions used in the DY model, the CK-adversary model empowers \mathcal{A} to compromise secret keys, secret credentials, and session states through the session hijacking attacks. Thus, leakage of the short term secrets from the UE node's memory can lead to disclosure of session key and other secrets.

Threat Model

- [HCAM]: B. Narwal and A. K. Mohapatra, “A survey on security and authentication in wireless body area networks,” *Journal of Systems Architecture*, vol. 113, p. 101883, 2021.
<https://www.sciencedirect.com/science/article/pii/S1383762120301600>
- [DYM]: D. Dolev and A. Yao, “On the security of public key protocols,” *IEEE Transactions on Information Theory*, vol. 29, no. 2, pp. 198208, 1983.
<https://ieeexplore.ieee.org/document/1056650>
- [CKM]: R. Canetti and H. Krawczyk, “Universally Composable Notions of Key Exchange and Secure Channels,” in *International Conference on the Theory and Applications of Cryptographic Techniques (EUROCRYPT'02)*, Amsterdam, The Netherlands, 2002, pp. 337–351. https://link.springer.com/chapter/10.1007/3-540-46035-7_22

Threat Model

- Due to the hostile environments in the deployment field, nodes can be physically captured by an attacker.
- Sensor nodes as well as cluster heads can be compromised or captured by an attacker. Usually, nodes are not equipped with tamper-resistant hardware due to cost constraints and hence we assume that once a node is captured by an attacker, all the stored sensitive data as well as cryptographic information are revealed to the attacker.
- In any case, the *GWN* will not be compromised by an attacker.
- Finally, we make use of the famous Dolev-Yao threat model in which two communicating parties (nodes) communicate over an insecure channel. We adopt the similar threat model for WSNs where the channel is insecure and the end-points (users, sensor nodes) cannot in general be trustworthy.

Table: Notations used

Symbol	Description
GWN	WSN gateway node (base station)
U_i	i^{th} user
SC_i	Smart card of U_i
ID_i	Identity of user U_i
PW_i	Password of user U_i
B_i	Biometric information of U_i
K	1024-bit secret number known to U_i only
$h(\cdot)$	Secure collision-free one-way hash function
X_s	1024-bit secret master key of GWN
SN_j	j^{th} sensor node in WSN
ID_{SN_j}	Identity of SN_j
TE_i	Expiration time of U_i 's temporal credential
TS_X	Current timestamp of an entity X
$Gen(\cdot)$	Fuzzy generator function
$Rep(\cdot)$	Fuzzy reproduction function
t	Error tolerance threshold used in fuzzy extractor
ΔT	Maximum transmission delay
$A \oplus B$	Bitwise XORed of data A with data B
$A B$	Data A concatenates with data B

Pre-Deployment Phase

Before deployment of nodes in the network, the *GWN* does the following steps.

- Step PD1. For each deployed sensor node SN_j , the GWN selects a unique identifier ID_{SN_j} .
- Step PD2. The GWN generates randomly a large 1024-bit number K_{GWN-S} , which is considered as the GWN's private key only known to the GWN. After that for each deployed sensor node SN_j , the GWN computes $TC_j = h(K_{GWN-S} || ID_{SN_j})$, which is the temporal credential for SN_j .
- Step PD3. Finally, each deployed sensor node SN_j is pre-loaded with the information TC_j as its temporal credential prior to its deployment in the target field.

Pre-Deployment Phase

ID_{SN_j}	$TC_j = h(K_{GWN-S} ID_{SN_j})$
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Figure: Pre-loaded information into SN_j 's memory.

Registration Phase

- Before accessing data from a particular sensor node in the sensor network, the user U_i needs to register with the GWN of the network.
- U_i first selects a unique identity ID_i and chooses a password PW_i .
- U_i generates randomly a large 1024-bit secret number K . U_i computes the masked password $RPW_i = h(ID_i || K || PW_i)$ and sends the registration request message $\langle ID_i, RPW_i \rangle$ to the GWN via a secure channel.
- The remaining steps are summarized in the following table.

User authentication in DWSNs

User (U_i)/Smart Card (SC_i)	GWN
<p>Inputs ID_i, PW_i, B_i. Generates a random secret number K. Computes $RPW_i = h(ID_i K PW_i)$. $\langle ID_i, RPW_i \rangle$ <div style="text-align: center;">\downarrow</div> (via a secure channel)</p> <p>Computes $Gen(B_i) = (\sigma_i, \tau_i)$, $e_i = h(ID_i \sigma_i) \oplus K$, $f_i = h(ID_i RPW_i \sigma_i)$, $r_i^* = r_i \oplus h(ID_i K)$. Replaces r_i with r_i^* in smart card. Stores $e_i, f_i, Gen(\cdot), Rep(\cdot), t$ and τ_i in smart card.</p>	<p>Generates private key K_{GWN-U}. Computes $TC_i = h(K_{GWN-U} ID_i TE_i)$, $PTC_i = TC_i \oplus RPW_i$. Generates secret information X_s and computes $r_i = h(ID_i X_s)$. Selects temporary identity TID_i of U_i and initializes it. Stores the tuple (TID_i, ID_i, TE_i) in its verification table. $\langle SmartCard(h(\cdot), TID_i, TE_i, PTC_i, r_i) \rangle$ <div style="text-align: center;">\downarrow</div> (via a secure channel)</p>

Registration Phase

$$h(\cdot), TID_i, TE_i, PTC_i, r_i^*, f_i, e_i, Gen(\cdot), \\ Rep(\cdot), t, \tau_i.$$

Figure: Information stored into SC_i 's memory.