

# Introduction to Blockchain

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# What is Blockchain?

- A blockchain is considered as a chain of blocks that are created from several blocks and it potentially consists of information.
- By the words “block” and “chain”, we actually specify in the context of digital information (“block”) which is stored in a public domain say database (“chain”).
- Since the digital information is stored in the form of “block” and it is linked in a “chain” form, the linked blocks constitute a chain, and hence, the name “blockchain”.
- The blockchain’s first block is known as the **Genesis block**.

# What is Blockchain?

The reasons why the blockchain have gained so much admiration are that:

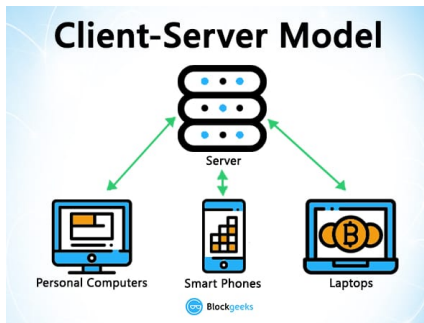
- It is not owned by a single entity, hence it is **decentralized**
- The data is **cryptographically** stored inside
- The blockchain is **immutable**, so no one can tamper with the data that is inside the blockchain
- The blockchain is **transparent** so one can track the data if they want to

# Three Pillars of Blockchain Technology

- Decentralization
- Transparency
- Immutability

# Pillar #1 : Decentralization

- Example of a centralized system is banks. They store all your money, and the only way that you can pay someone is by going through the bank.
- The traditional client-server model is a perfect example of this.
- When you google search for something, you send a query to the server who then gets back at you with the relevant information. That is simple client-server.



The centralized systems have treated us well for many years, however, they have several vulnerabilities.

- Firstly, because they are centralized, all the data is stored in one spot. This makes them easy target spots for potential hackers.
- If the centralized system was to go through a software upgrade, it would halt the entire system
- What if the centralized entity somehow shut down for whatever reason? That way nobody will be able to access the information that it possesses
- Worst case scenario, what if this entity gets corrupted and malicious? If that happens then all the data that is inside the blockchain will be compromised.

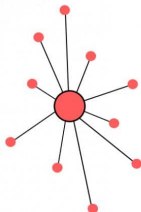
# Pillar #1 : Decentralization

- In a decentralized system, the information is not stored by one single entity. In fact, everyone in the network owns the information.
- In a decentralized network, if you wanted to interact with your friend then you can do so directly without going through a third party.
- That was the main ideology behind Bitcoins. You and only you alone are in charge of your money. You can send your money to anyone you want without having to go through a bank.



## The New Networks

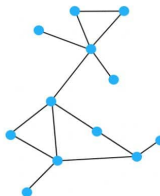
### Centralized



Centralized systems have a core authority that **dictates the truth** to the other participants in the network.

Only **privileged users** or institutions can access the history of transactions or confirm new transactions.

### Decentralized



Decentralized systems have **no core authority** to dictate the truth to other participants in the network.

**Every participant** in the network can access the history of transactions or confirm new transactions.



# Pillar #2 : Transparency

- While the person's real identity is secure, you will still see all the transactions that were done by their public address.
- This level of transparency has never existed before within a financial system.
- It adds that extra, and much needed, level of accountability which is required by some of these biggest institutions.

# Pillar #3 : Immutability

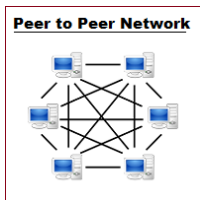
- Immutability, in the context of the blockchain, means that once something has been entered into the blockchain, it cannot be tampered with.
- The reason why the blockchain gets this property is that of cryptographic hash function.

| INPUT   | HASH  |
|---|---|
| Hi  | 3639EFCDD08ABB273B1619E82E78C29A7DF02C1051B1820E99FC395DCAA3326B8 |
| Welcome to blockgeeks. Glad to have you here. | 53A53FC9E2A03F9B6E66D84BA701574CD9CF5F01FB498C41731881BCDC68A7C8  |

| INPUT          | HASH   |
|----------------|--|
| This is a test | C7BE1ED902FB8DD4D48997C6452F5D7E509FBCDBE2808B16BCF4EDCE4C07D14E |
| this is a test | 2E99758548972A8E8822AD47FA1017FF72F06F3FF6A016851F45C398732BC50C |

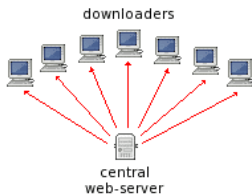
# Maintaining the Blockchain – Network and Nodes

- The blockchain is maintained by a peer-to-peer (P2P) network.
- The network is a collection of nodes which are interconnected to one another.
- Nodes are individual computers which take in input and performs a function on them and gives an output.
- The blockchain uses a special kind of network called *peer-to-peer network* which partitions its entire workload between participants, who are all equally privileged, called *peers*.
- There is no longer one central server, now there are several distributed and decentralized peers.



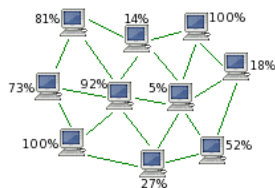
- One of the main uses of the peer-to-peer network is file sharing, also called torrenting.

**Traditional Centralized Downloading**



- Slow
- Single point of failure
- High bandwidth usage for server

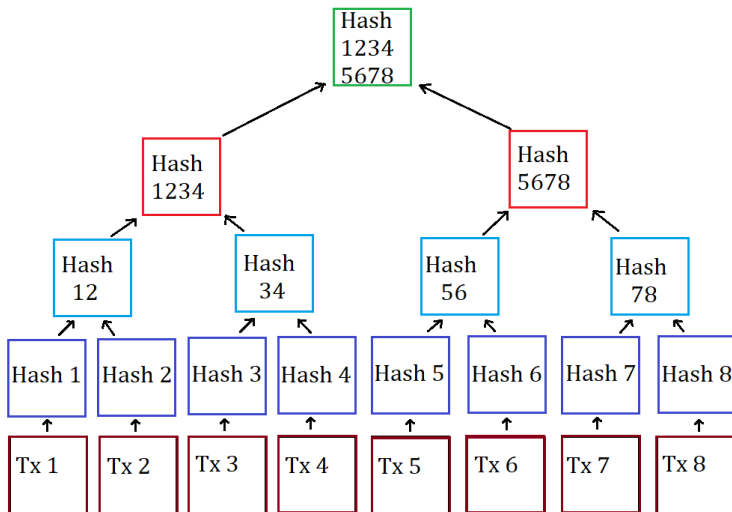
**Decentralized Peer-to-Peer Downloading**



- Fast
- No single point of failure
- All downloaders are also uploaders

- **Public (Permissionless) Blockchain:** Everyone has the right to join, access, send, verify and receive transactions of the blocks in the blockchain to create a consensus. One widely successful permissionless blockchain is *bitcoin*.
- **Private (Permissioned) Blockchain:** The owner of the network decides which node to assign the right to access, send, receive, join and verify the block for creating an agreement between the nodes. Example: Healthcare Applications
- **Consortium or Hybrid Blockchain:** Internet of Vehicles (IoV) application

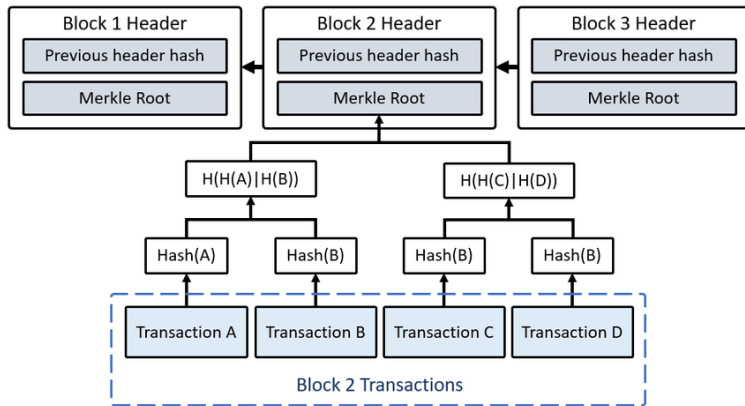
# How to Make a Merkle Tree?



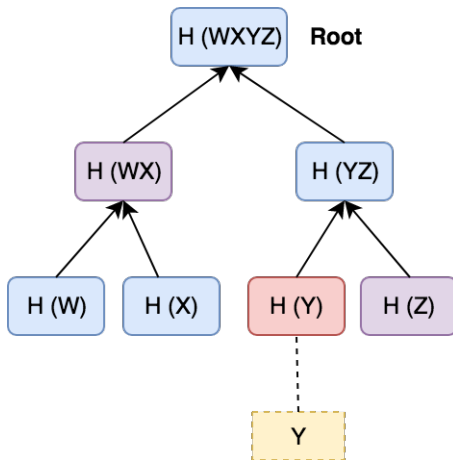
$Tx_i$ :  $i^{th}$  transaction;  $H_i$ : Hash of  $i^{th}$  transaction ( $H_{Tx_i} = H(Tx_i)$ );

$H_{Tx_1 Tx_2} = H_{Tx_1} \oplus H_{Tx_2}$ ;  $H_{12345678}$ : Merkle root

# Use of Merkle Tree in Blockchain



# Verifying Transactions Using the Merkle Root



**To confirm transaction  $Y$ , one only needs to know:**  
 **$H(WX)$ ,  $H(Y)$ ,  $H(Z)$  and  $H(WXYZ)$ ;  $H(\cdot)$ : hash function**



# Types of blockchain

| Block Header                    |               |
|---------------------------------|---------------|
| Block Version                   | BV            |
| Previous Block Hash             | PBHash        |
| Merkle Tree Root                | MTR           |
| Block Type                      | Public        |
| Timestamp                       | TS            |
| Owner of Block                  | $ES_i$        |
| Public key of signer ( $ES_i$ ) | $Pub_{ES_i}$  |
| Block Payload (Transactions)    |               |
| Transaction #1                  | $T_{x_1}$     |
| Transaction #2                  | $T_{x_2}$     |
| $\vdots$                        | $\vdots$      |
| Transaction # $n_t$             | $T_{x_{n_t}}$ |
| Current Block Hash              | CBHash        |
| Signature on block using ECDSA  | BSign         |

| Block Header                           |                               |
|--|-------------------------------|
| Block Version                          | BV                            |
| Previous Block Hash                    | PBHash                        |
| Merkle Tree Root                       | MTR                           |
| Block Type                             | Private                       |
| Timestamp                              | TS                            |
| Owner of Block                         | $ES_i$                        |
| Public key of signer ( $ES_i$ )        | $Pub_{ES_i}$                  |
| Block Payload (Encrypted Transactions) |                               |
| Encrypted Transaction #1               | $E_{Pub_{ES_i}}(T_{x_1})$     |
| Encrypted Transaction #2               | $E_{Pub_{ES_i}}(T_{x_2})$     |
| $\vdots$                               | $\vdots$                      |
| Encrypted Transaction # $n_t$          | $E_{Pub_{ES_i}}(T_{x_{n_t}})$ |
| Current Block Hash                     | CBHash                        |
| Signature on block using ECDSA         | BSign                         |

| Block Header                    |                               |
|---------------------------------|-------------------------------|
| Block Version                   | BV                            |
| Previous Block Hash             | PBHash                        |
| Merkle Tree Root                | MTR                           |
| Block Type                      | Hybrid                        |
| Timestamp                       | TS                            |
| Owner of Block                  | $ES_i$                        |
| Public key of signer ( $ES_i$ ) | $Pub_{ES_i}$                  |
| Block Payload                   |                               |
| Encrypted Transaction #1        | $E_{Pub_{ES_i}}(T_{x_1})$     |
| Transaction #2                  | $T_{x_2}$                     |
| $\vdots$                        | $\vdots$                      |
| Encrypted Transaction # $n_t$   | $E_{Pub_{ES_i}}(T_{x_{n_t}})$ |
| Current Block Hash              | CBHash                        |
| Signature on block using ECDSA  | BSign                         |

a) Formation of a block on public blockchain

b) Formation of a block on private blockchain

c) Formation of a block on consortium blockchain

Consensus mechanisms are used to verify transactional data between the nodes in a P2P network.

- **Byzantine Fault Tolerance (BFT):** An agreement protocol which helps to tolerate the Byzantine failures in a network. BFT maintains the reliable record of transactions in a transparent and tamper-proof way, as long as the number of traitors does not exceed one-third of the general network nodes.
- **Practical Byzantine Fault Tolerance (PBFT):** This consensus mechanism is used when BFT fails to tolerate the faults in a network system. The algorithm for PBFT works in asynchronous systems and is optimized to achieve high performance along with an impressive overhead runtime.

# Practical Byzantine Fault Tolerance (PBFT)

For adding a block in the blockchain, the following procedures in PBFT are required:

- A leader acting as a miner will select by the leader selection algorithm for adding a block.
- Leader receives a block with block adding request from any nodes (or client) into the blockchain.
- Leader sends this block to every node in the network for verifying the transactions.
- After successful verification of the transaction in the block, each received node sends a valid reply for adding that block.
- Leader counts the received reply and checks the number of counts (say,  $RCount$ ) if it is greater than the twice number of the faulty nodes, i.e.,  $RCount > 2n_f + 1$  or the two-third nodes give the same reply. The  $2n_f + 1$  non-faulty or valid replies provide the liveness of the system, that is, the message delay need to be bounded in due course. If this condition is satisfied, the leader will add the block into the blockchain and broadcasts a commit for adding the block into their respective blockchain for backup purposes.

- **Proof-of-Work (PoW):** This is the original consensus mechanism used to verify the transactions and produce new blocks in the blockchain. Mining is a complex process, and miners need to demonstrate that they can validate the transaction block. Here, the miners are financially rewarded if they perform verification. Thus, as the complexity increases in the mining process over time, the power consumption also increases. In other words, PoW is a costly process as the miners compete with each other to solve a mathematical problem.
- **Proof-of-Stake (PoS):** In 2017, Ethereum began the process of switching from a PoW mechanism to a PoS system. The latter was designed to mitigate the limitations of PoW, in terms of energy, cost, and processing time. Specifically, it adopts a forging process rather than the mining process to validate transaction blocks.

- **Delegated Proof-of-Stake (DPoS):** This is a fast, efficient, flexible and most decentralized consensus mechanism. DPoS holds the power of stakeholder for the approval of voting and resolving the consensus issues in an honest and representative way. The deterministic selection of witnesses allows the transactions to be confirmed on an average of just 1 second. This consensus mechanism is designed to protect all the participants in a free, fair, and transparent environment.
- **Proof-of-Burn (PoB):** An alternative consensus protocol for PoS and PoW. In PoB mechanism, the miners prove that they burn one cryptocurrency to create another currency, i.e., they are sent to a bitcoin address which is unspendable. Its significance depends on the burning tokens in an unrecoverable manner. As comparative to PoW/PoS, it is easily verifiable and hard to undo.

- It is a voting based consensus algorithm proposed in the literature in order to achieve high accuracy of a correct agreement for adding a block into a blockchain over unreliable distributed network.
- RPCA achieves an agreement in the voting process.
- Every participant node in the voting process maintains a unique node list (UNL), where each node in the list is considered as trusted one.
- The protocol executes with the help of the following steps [Wang et al. [2019]]:
  - ▶ In voting process, every participant node constantly receives the transaction. If the transaction is valid, the node integrates the transactions into a set or list, called a “candidate set”.
  - ▶ Every participant node dispatches its own candidate set to other participant node as a proposal.

# Ripple Protocol Consensus Algorithm (RPCA)

## (Continued...)

- The remaining steps are as follows [Wang et al. [2019]]:
  - ▶ The participant node receives the proposal from other nodes. The node will then check whether the sender node belongs to its UNL list or not. If it is there, the node will verify the transactions with its own local candidate set. If all are valid, the transactions will gain a vote. Only when the transaction gets more than 50% of vote, the transaction will enter into the next round.
  - ▶ Next, the participant node sends the transaction that it gains more than 50% of the votes than the others and if it increases to 60% of vote, it needs to wait until it reaches to the threshold of 80% of the votes.
  - ▶ Finally, the participant node records the transaction confirmed by the 80% UNL nodes to be added into its ledger data.

The transaction is accepted only if 80% of the votes in the UNL of a participant node agrees with it. Thus, 80% of the UNL is honest, that is, the percentage of the faulty nodes in the UNL is less than that for the 20% of the UNL. When the UNL contains  $d$  number of nodes in the network, and the RPCA will maintain its correctness as long as  $n_f \leq \frac{d-1}{5}$ , i.e.,  $d \geq 5n_f + 1$  where  $n_f$  is the Byzantine failure persisted nodes in the network.

**Table:** Blockchain consensus mechanisms and their applications

| Consensus mechanism | Concept                   | Resource     | Applications  |
|---------------------|---------------------------|--------------|---|
| Ripple              | Voting in multiple rounds | No resource  | XRP ledger xrp [2014]   |
| PBFT                | Voting                    | No resource  | Tendermint Kwon [2014]  |
| PoW                 | Hashing                   | Computations | Bitcoin Nakamoto [2009]<br>Ethereum Wood et al. [2014]  |
| PoS                 | Digital signatures        | Currency     | PeerCoin King and Nadal [2012]<br>SnowWhite Bentov et al. [2016]<br>Ouroboros Kiayias et al. [2017] |
| DPoS                | Voting                    | Currency     | BitShared<br>Ark<br>EOS   |
| PoB                 | Address suspension        | Currency     | SlimCoin sli [2014]   |

**Ref.** Anusha Vangala, Ashok Kumar Das, Neeraj Kumar, and Mamoun Alazab. "Smart Secure Sensing for IoT-Based Agriculture: Blockchain Perspective," in *IEEE Sensors Journal*, 2020, DOI: 10.1109/JSEN.2020.3009382. (2019 SCI Impact Factor: 3.073)

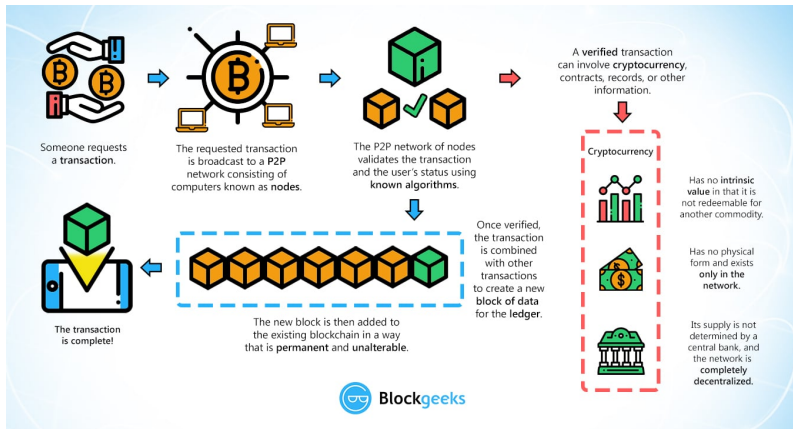


**Table:** Attacks on consensus mechanisms

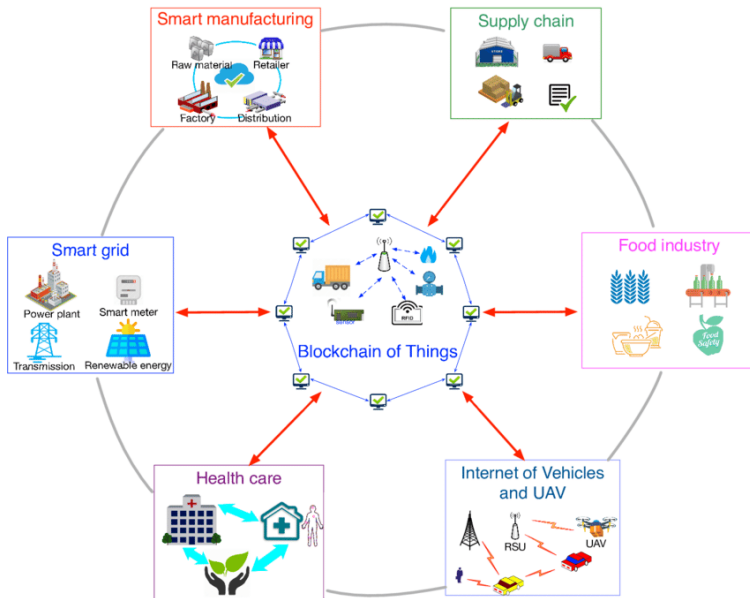
| <b>Attack</b>         | <b>Affected consensus protocols</b> | <b>Description</b>  |
|-----------------------|-------------------------------------|---|
| Double spending       | Most protocols                      | Repeated usage of token   |
| Selfish mining        | PoW                                 | Gain profits by generating blocks privately in a mining pool  |
| Nothing at stake      | PoS                                 | Blocks added to all branches in a fork  |
| Bribe attack          | PoS                                 | Honest nodes are given incentive to add blocks on private fork  |
| Stake bleeding attack | PoS                                 | Broadcast transactions copied from main chain onto private fork to earn extra fees and increase stake |
| Fake stake attack     | PoS                                 | Increase the smaller valued stakes to higher valued stakes  |

**Ref.** S. Zhang and J. Lee, “A Group Signature and Authentication Scheme for Blockchain-Based Mobile-Edge Computing,” *IEEE Internet of Things Journal*, vol. 7, no. 5, pp. 4557-4565, 2020. (2019 SCI Impact Factor: 9.936).

# Block creation and addition using consensus mechanism in a blockchain



# Applications of Blockchain Technology



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