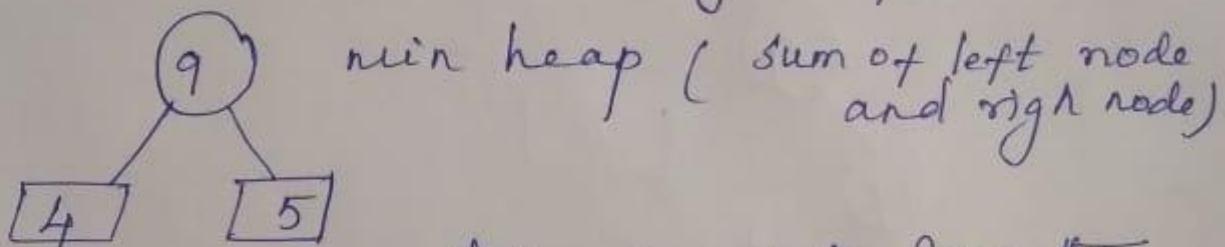


Consider the following messages m_1, \dots, m_7
with relative frequencies $(q_1, \dots, q_7) = (4, 5, 7, 8,$

Step 1: Arrange the data in ascending order
 $4, 5, 7, 8, 10, 12, 20$

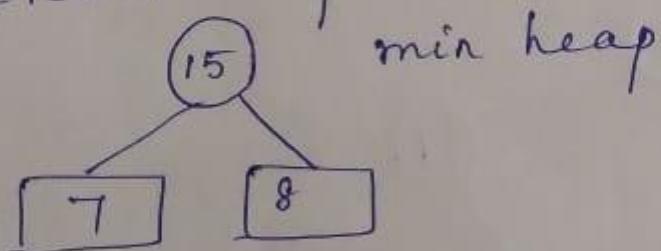
Step 2: Combine first two entries of a table
Create a minimum binary heap



Step 3: Remove the entries 4 and 5 from the table, insert 9 in appropriate position

$7, 8, 9, 10, 12, 20$

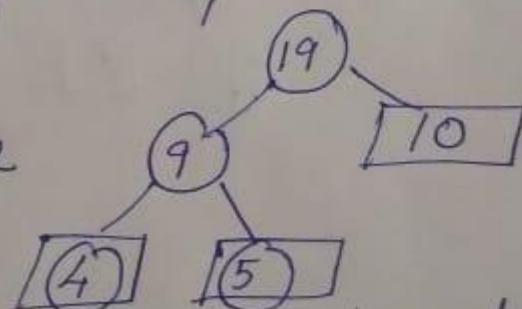
Combine minimum value of two blocks,
Create a parent node,



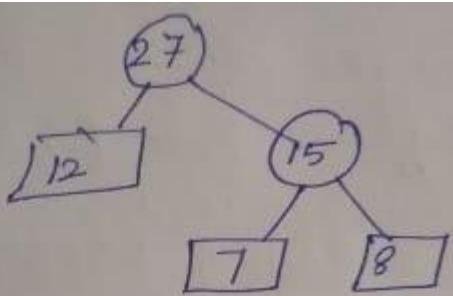
Step 4: Remove the entries 7 and 8 from table,
insert 15 at appropriate position

$9, 10, 12, 15, 20$

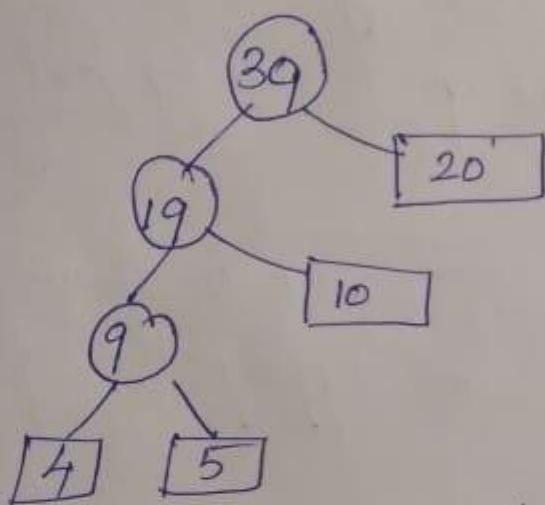
Combine minimum value



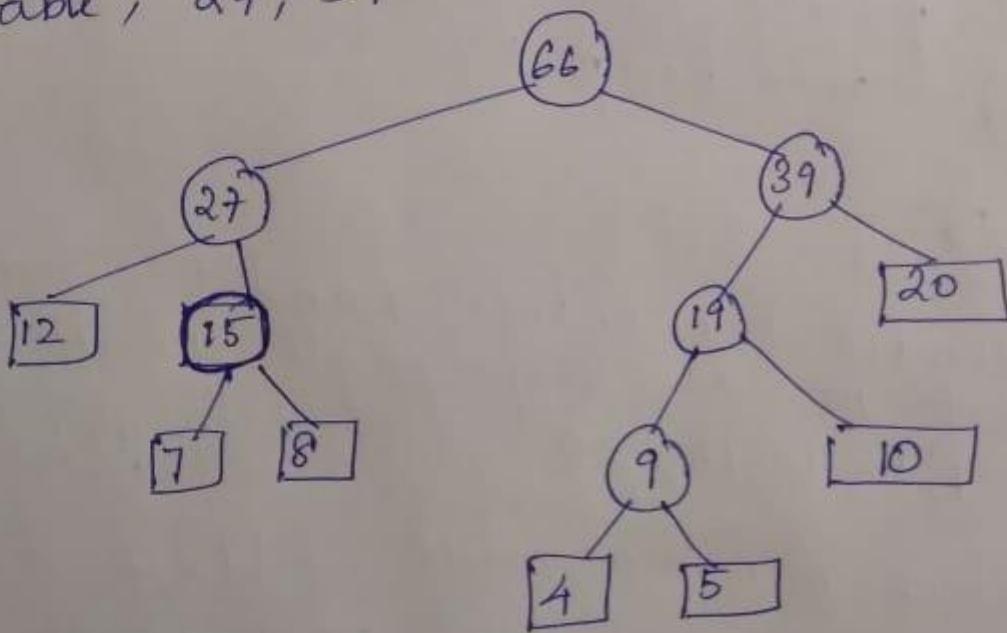
Step 5: Remove entries 9 and 10, insert 19 at its proper position
 $12, 15, 19, 20$



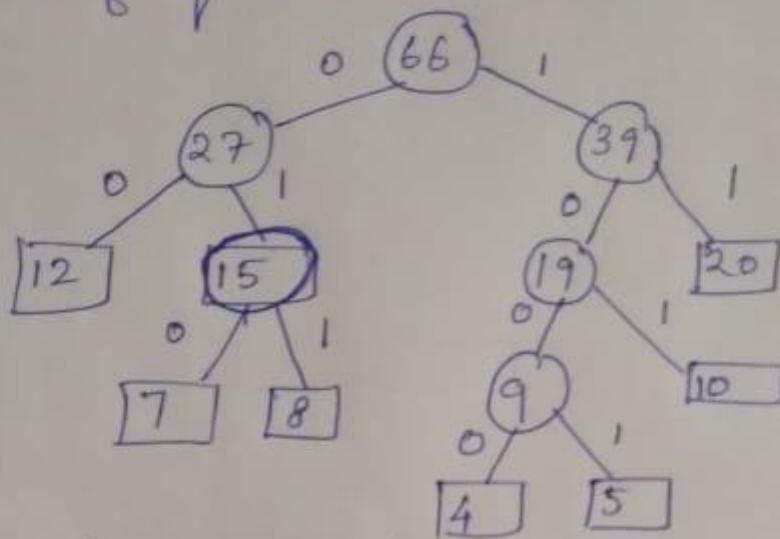
Step 6: Remove entries 15 and 12 from table, insert 27 at appropriate position
1.9, 20, 27



Step 7: Remove entries 19 and 20, insert 39 in table, 27, 39



Now assign left child 0, right child 1 to encode frequencies



Codes for given frequencies,

Frequencies	Codes
4	1000
5	1001
7	010
8	011
10	101
12	00
20	11

} Variable length Codeword

Huffman Encoding is not applicable for fixed length code, the character that has highest frequency will has codeword of minimum length.

Time Complexity : Construction of minheap
 \therefore Complexity of minheap $O(n \log n)$
 where n , is number of characters