

## Recursion Tree Method

1. A tree where each node represents the cost of a certain recursive sub-problem
2. Sum up the values in each node to get the cost of entire algorithm.

I: Draw a recursion tree based on given recurrence relation.

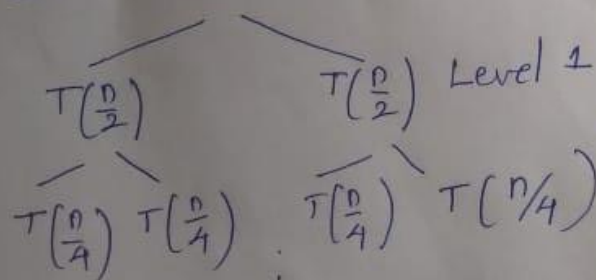
II: Cost of each level  
Total no. of levels in recursion tree  
No. of nodes in last level  
Cost of last level.

III: Add cost of all levels of recursion tree and simplify the expression so obtained in terms of asymptotic notation.

Problem-1: Solve the following recurrence relation using recursion tree method:

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

1:  $T(n)$  Level 0



\* A problem of size  $n$  will get divided in 2 subproblem of size  $\frac{n}{2}$ .

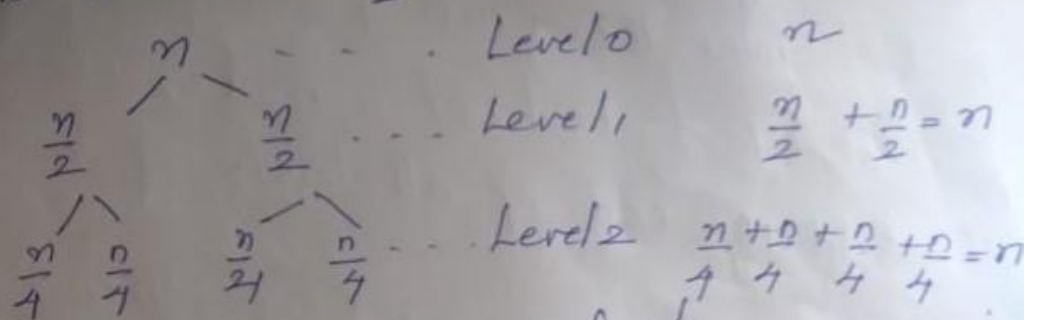
\* Then, each sub-problem of size  $\frac{n}{2}$  will get divided into 2 sub-problems of size  $\frac{n}{4}$ .

\* At the bottom layer the size of subproblem will reduce to 1.

The given recurrence relation shows:

\* The cost of dividing a problem of size  $n$  into 2 sub-problems and combining its solution is  $n$ .

\* Cost of dividing a problem of size  $n/2$  into 2 sub-problems and then combining its solution is  $n/2$  . . . . . (2) Cost



(3) Determine the total no. of levels

$$\text{Level} - 0 = \frac{n}{2^0}$$

$$\text{Level} - 1 = \frac{n}{2^1}$$

$$\text{Level} - 2 = \frac{n}{2^2}$$

Continuing we have  
size of subproblem at level  $i = \frac{n}{2^i}$

the last level  $x$ , size of sub-problem becomes 1.

$$\frac{n}{2^x} = 1$$

$$= 2^x = n$$

Logarithm on both sides  $x = \log_2 n$   
Total no. of levels in recursion tree

Determine the no. of nodes in last level

Level 0 has  $2^0$  i.e., 1 node

Level 1 has  $2^1$  i.e., 2 nodes

Level 2 has  $2^2$  i.e., 4 nodes

⋮

Level  $\log_2 n$  has  $2^{\log_2 n}$  nodes i.e.,  $n$  nodes

To Determine the cost of last level,

$$\text{Cost of last level} = nT(1) = \Theta(n)$$

6) Add the cost of all levels of recursion

Tree

$$T(n) = \underbrace{\{n + n + n + \dots\}}_{\log_2 n \text{ Levels}} + \Theta(n)$$

$$= n \log_2 n + \Theta(n)$$

$$= n \log_2 n + \Theta(n)$$

$$= \Theta(n \log_2 n)$$