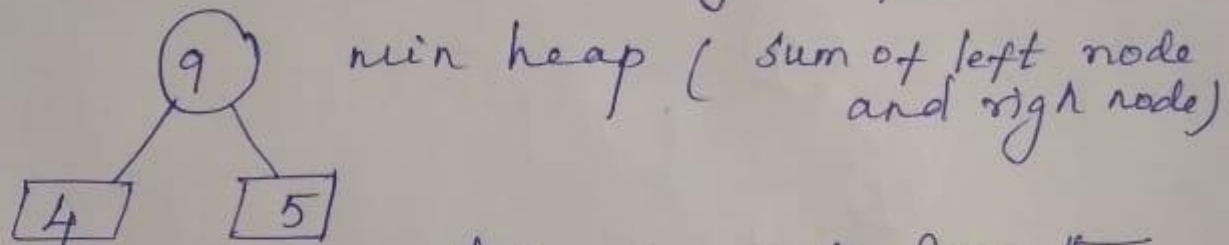


Consider the following messages m_1, \dots, m_7 with relative frequencies $(f_1, \dots, f_7) = (4, 5, 7, 8, 10, 12, 20)$

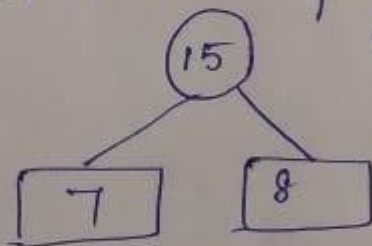
Step 1: Arrange the data in ascending order
4, 5, 7, 8, 10, 12, 20

Step 2: Combine first two entries of a table
Create a minimum binary heap



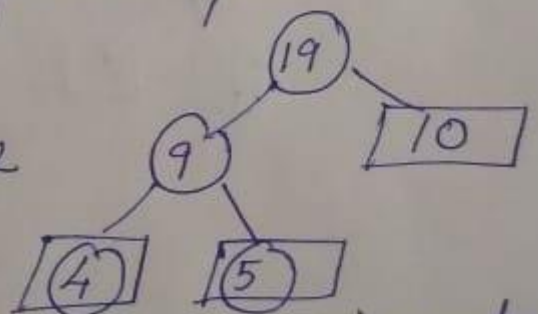
Step 3: Remove the entries 4 and 5 from the table, insert 9 in appropriate position
7, 8, 9, 10, 12, 20

Combine minimum value of two blocks,
Create a parent node,
min heap

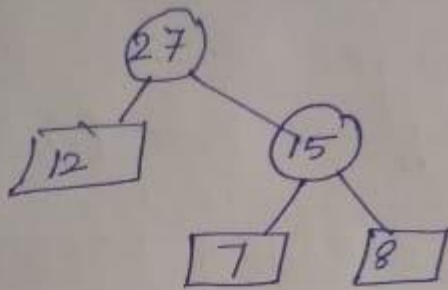


Step 4: Remove the entries 7 and 8 from table,
insert 15 at appropriate position
9, 10, 12, 15, 20

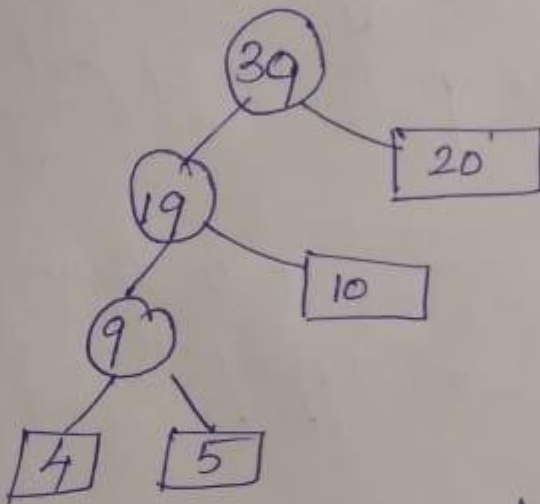
Combine minimum value



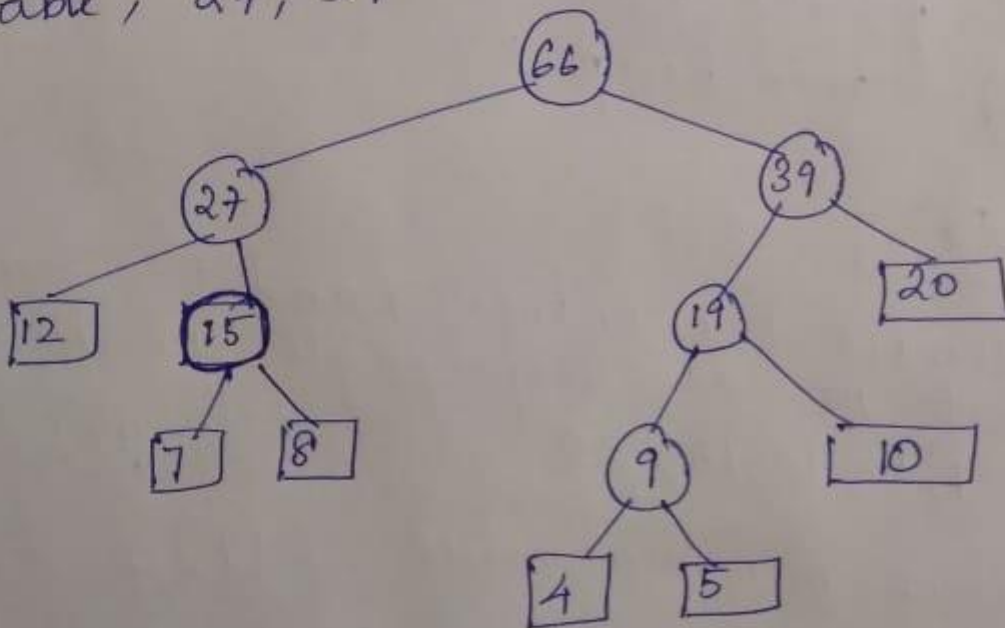
Step 5: Remove entries 9 and 10, insert 19 at its proper position 12, 15, 19, 20



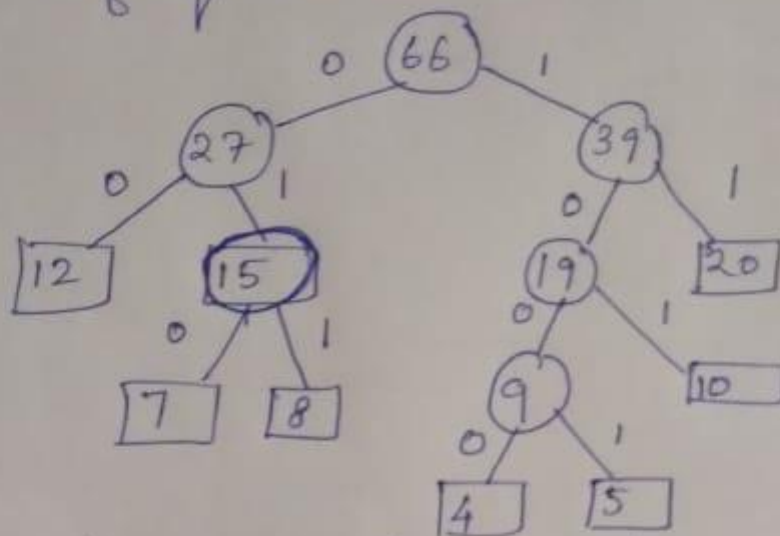
Step 6: Remove entries 15 and 12 from table, insert 27 at appropriate position
19, 20, 27



Step 7: Remove entries 19 and 20, insert 39 in table, 27, 39



Now assign left child 0, right child 1, to encode frequencies



Codes for given frequencies,

Frequencies

Codes

| | |
|----|------|
| 4 | 1000 |
| 5 | 1001 |
| 7 | 010 |
| 8 | 011 |
| 10 | 101 |
| 12 | 00 |
| 20 | 11 |

Variable
length
Codeword

Huffman Encoding is not applicable for fixed length code, The character that has highest frequency will has Codeword of minimum length.

Time Complexity: Construction of minheap
 \therefore Complexity of minheap $O(n \log n)$
 where n , is number of characters