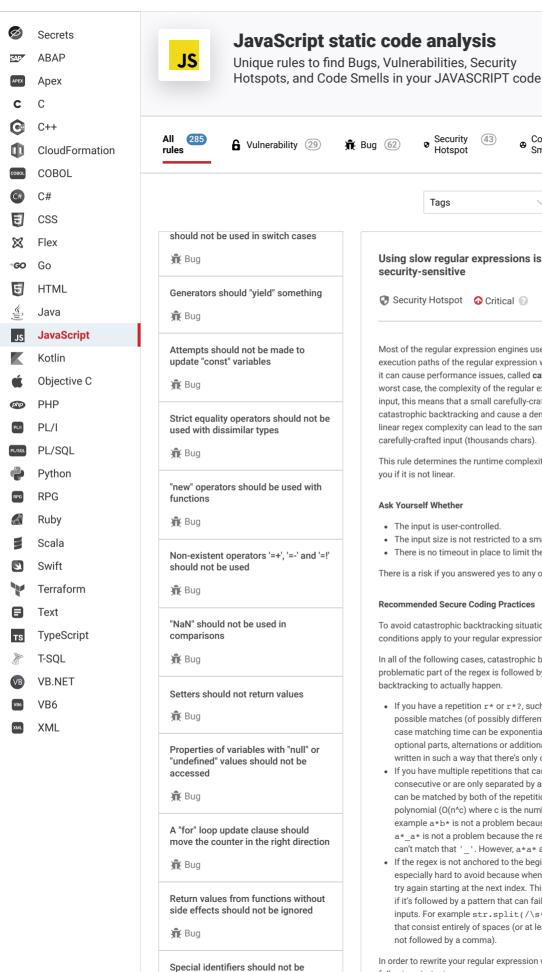
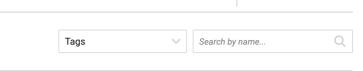


Products ✓





⊗ Code

Smell

(151)

Using slow regular expressions is security-sensitive

Security

Hotspot

(43)

Analyze your code

Quick 41 Fix

Security Hotspot Oritical cwe owasp regex

Most of the regular expression engines use backtracking to try all possible execution paths of the regular expression when evaluating an input, in some cases it can cause performance issues, called catastrophic backtracking situations. In the worst case, the complexity of the regular expression is exponential in the size of the input, this means that a small carefully-crafted input (like 20 chars) can trigger catastrophic backtracking and cause a denial of service of the application. Superlinear regex complexity can lead to the same impact too with, in this case, a large carefully-crafted input (thousands chars).

This rule determines the runtime complexity of a regular expression and informs you if it is not linear.

Ask Yourself Whether

- . The input is user-controlled
- The input size is not restricted to a small number of characters.
- There is no timeout in place to limit the regex evaluation time.

There is a risk if you answered yes to any of those questions.

Recommended Secure Coding Practices

To avoid catastrophic backtracking situations, make sure that none of the following conditions apply to your regular expression.

In all of the following cases, catastrophic backtracking can only happen if the problematic part of the regex is followed by a pattern that can fail, causing the backtracking to actually happen.

- If you have a repetition r* or r*?, such that the regex r could produce different possible matches (of possibly different lengths) on the same input, the worst case matching time can be exponential. This can be the case if r contains optional parts, alternations or additional repetitions (but not if the repetition is written in such a way that there's only one way to match it).
- If you have multiple repetitions that can match the same contents and are consecutive or are only separated by an optional separator or a separator that can be matched by both of the repetitions, the worst case matching time can be polynomial ($O(n^{-}c)$) where c is the number of problematic repetitions). For example a*b* is not a problem because a* and b* match different things and a*_a* is not a problem because the repetitions are separated by a '_' and can't match that ' $_$ ' . However, a*a* and $.*_.*$ have quadratic runtime.
- If the regex is not anchored to the beginning of the string, quadratic runtime is especially hard to avoid because whenever a match fails, the regex engine will try again starting at the next index. This means that any unbounded repetition, if it's followed by a pattern that can fail, can cause quadratic runtime on some inputs. For example ${\tt str.split(/\s^*,/)}$ will run in quadratic time on strings that consist entirely of spaces (or at least contain large sequences of spaces, not followed by a comma).

In order to rewrite your regular expression without these patterns, consider the following strategies:

 If applicable, define a maximum number of expected repetitions using the bounded quantifiers, like {1,5} instead of + for instance.

bound or assigned

📆 Bug

Values should not be uselessly incremented

R Bug

Related "if/else if" statements should not have the same condition

Rug

Objects should not be created to be dropped immediately without being

Bug

Identical expressions should not be used on both sides of a binary operator

- Refactor nested quantifiers to limit the number of way the inner group can be
 matched by the outer quantifier, for instance this nested quantifier situation
 (ba+)+ doesn't cause performance issues, indeed, the inner group can be
 matched only if there exists exactly one b char per repetition of the group.
- Optimize regular expressions by emulating possessive quantifiers and atomic arouping.
- Use negated character classes instead of . to exclude separators where applicable. For example the quadratic regex .*_.* can be made linear by changing it to [^]* .*

Sometimes it's not possible to rewrite the regex to be linear while still matching what you want it to match. Especially when the regex is not anchored to the beginning of the string, for which it is quite hard to avoid quadratic runtimes. In those cases consider the following approaches:

- Solve the problem without regular expressions
- Use an alternative non-backtracking regex implementations such as Google's RE2 or node-re2.
- Use multiple passes. This could mean pre- and/or post-processing the string
 manually before/after applying the regular expression to it or using multiple
 regular expressions. One example of this would be to replace
 str.split(/\s*,\s*/) with str.split(",") and then trimming the
 spaces from the strings as a second step.
- It is often possible to make the regex infallible by making all the parts that could fail optional, which will prevent backtracking. Of course this means that you'll accept more strings than intended, but this can be handled by using capturing groups to check whether the optional parts were matched or not and then ignoring the match if they weren't. For example the regex x*y could be replaced with x*(y)? and then the call to str.match(regex) could be replaced with matched = str.match(regex) and matched[1] !== undefined.

Sensitive Code Example

The regex evaluation will never end:

Compliant Solution

Possessive quantifiers do not keep backtracking positions, thus can be used, if possible, to avoid performance issues. Unfortunately, they are not supported in JavaScript, but one can still mimick them using lookahead assertions and backreferences:

See

- OWASP Top 10 2017 Category A1 Injection
- MITRE, CWE-400 Uncontrolled Resource Consumption
- MITRE, CWE-1333 Inefficient Regular Expression Complexity
- owasp.org OWASP Regular expression Denial of Service ReDoS
- stackstatus.net Outage Postmortem July 20, 2016
- regular-expressions.info Runaway Regular Expressions: Catastrophic Backtracking
- <u>docs.microsoft.com</u> Backtracking with Nested Optional Quantifiers

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