

KATHMANDU UNIVERSITY

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Mini Project
[Code No: COMP 202]

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In the program ascending priority queue has been implemented by using min heap. Priority queue uses the insert and delete function of heap as min heap works in the same way as ascending priority queue. Deletion from the root/front which is always the minimum element.

5, 4, 10, 16, 2, 11, 1 are inserted in order.
They are stored in the vector as following

5						
---	--	--	--	--	--	--

4	5					
---	---	--	--	--	--	--

4	5	10				
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4	5	10	16			
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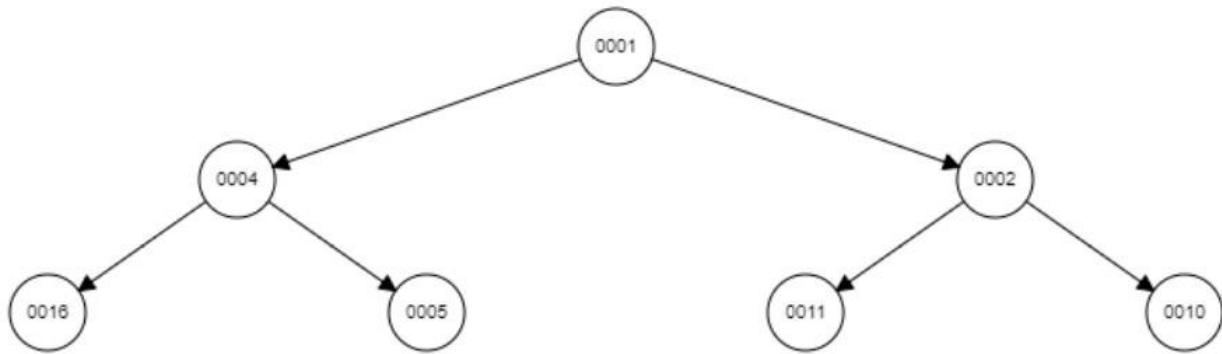
2	4	10	16	5		
---	---	----	----	---	--	--

2	4	10	16	5	11	
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1	4	2	16	5	11	10
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While inserting elements into the heap, if the inserted element is lower than its parent's element then we swap the two. This process is done to maintain the property of the heap

After insertion of the elements the heap is as shown below:



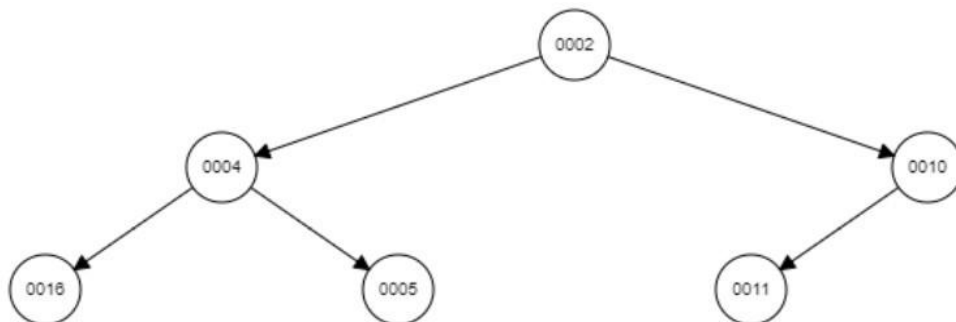
Since the data are stored starting from 0th index, the left child of an index i becomes $(2*i + 1)$ and similarly right child index becomes $(2*i + 2)$. The parent of an index i becomes $(i-1)/2$.

Upon deleting the min-element (root node), the last element takes the place on root then, it is compared with its lowest children (among the right and left children) then swapped if it is lower.

Upon deleting minimum element i.e. 1

2	4	10	16	5	11	
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The heap becomes



In this way we can implement ascending priority queue using min heap. The element with lowest priority (i.e., minimum element) is removed. Insert an element with arbitrary priority. And return an element with minimum priority.

Time Complexity of the program

For the time complexity of heap implementation, we have the following major operations

Insertion – The time complexity of inserting an element to the heap is $O(1)$ for the best case when the inserted element doesn't violate the heap property. If it does violate the heap property, `reheap_up` must be called which makes the time complexity $O(\log n)$.

Deletion – Deletion removes the minimum element from the node. Doing so violates the heap property, then calling `reheap_down` makes the time complexity $O(\log n)$.

Top – Top returns the minimum element/root in the heap, so it completes in $O(1)$.

Display() – For display, it traverse through the whole vector so, it's time complexity is $O(n)$ where n is the number of elements in the heap.

So, the overall time complexity of the program becomes $O(n)$.

Similarly for priority queue

For enqueue, `InsertKey` is used with time complexity $O(\log n)$

For dequeue, `DeleteKey` is used with time complexity $O(\log n)$

And for top/peek, `top` is used with time complexity $O(1)$

So, the overall time complexity of the program becomes $O(\log n)$.

Output of Program

```
Inserting 5,4,10,16,2,11,1 to the priority Queue  
The size is 7  
The elements in priority Queue are  
1 4 2 16 5 11 10  
  
The top element is 1  
Deleting ROOT  
  
The top element is 2  
Deleting ROOT  
  
The top element is 4  
Deleting ROOT  
  
The top element is 5
```

[GitHub Repository](#)