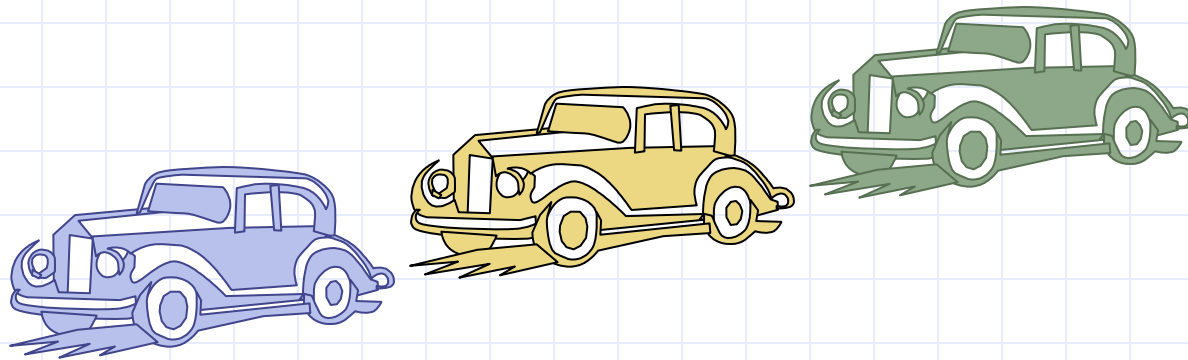


Queues



The Queue ADT

- The **Queue** ADT stores arbitrary objects
- Insertions and deletions follow the first-in first-out scheme
- Insertions are at the rear of the queue and removals are at the front of the queue
- Main queue operations:
 - **enqueue**(object): inserts an element at the end of the queue
 - object **dequeue**(): removes and returns the element at the front of the queue
- Auxiliary queue operations:
 - object **first**(): returns the element at the front without removing it
 - integer **len**(): returns the number of elements stored
 - boolean **is_empty**(): indicates whether no elements are stored
- Exceptions
 - Attempting the execution of dequeue or front on an empty queue throws an **EmptyQueueException**

Example

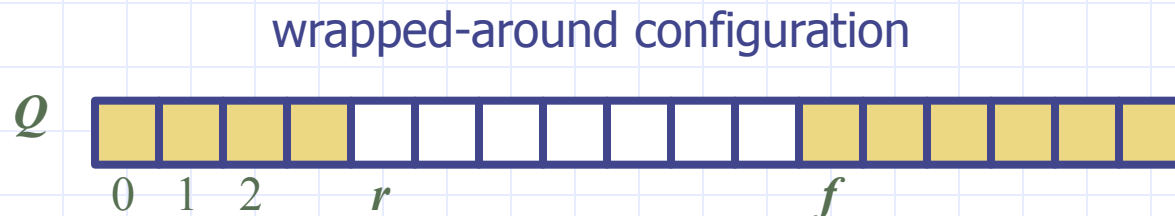
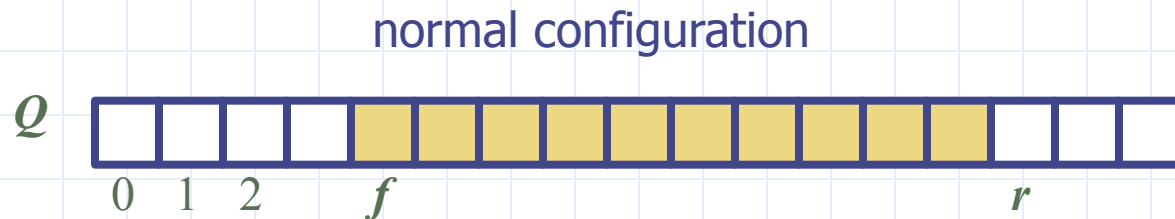
Operation	Return Value	first \leftarrow Q \leftarrow last
Q.enqueue(5)	—	[5]
Q.enqueue(3)	—	[5, 3]
len(Q)	2	[5, 3]
Q.dequeue()	5	[3]
Q.is_empty()	False	[3]
Q.dequeue()	3	[]
Q.is_empty()	True	[]
Q.dequeue()	“error”	[]
Q.enqueue(7)	—	[7]
Q.enqueue(9)	—	[7, 9]
Q.first()	7	[7, 9]
Q.enqueue(4)	—	[7, 9, 4]
len(Q)	3	[7, 9, 4]
Q.dequeue()	7	[9, 4]

Applications of Queues

- ❑ Direct applications
 - Waiting lists, bureaucracy
 - Access to shared resources (e.g., printer)
 - Multiprogramming
- ❑ Indirect applications
 - Auxiliary data structure for algorithms
 - Component of other data structures

Array-based Queue

- Use an array of size N in a circular fashion
- Two variables keep track of the front and rear
 - f index of the front element
 - r index immediately past the rear element
- Array location r is kept empty

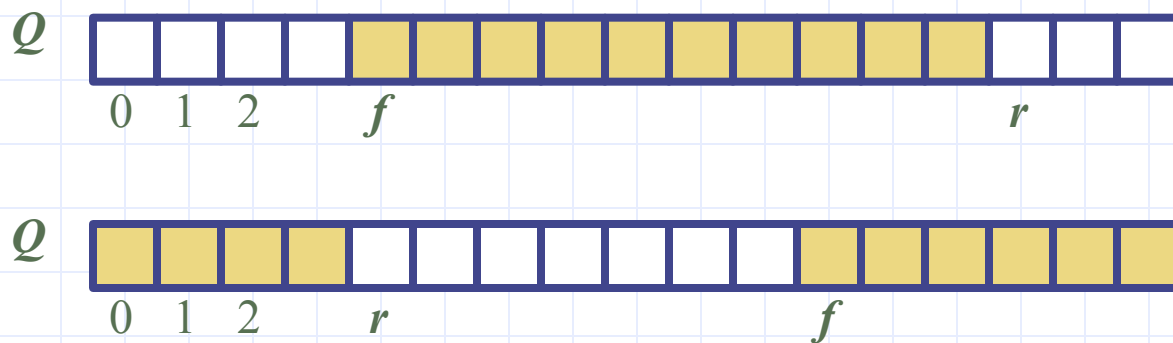


Queue Operations

- We use the modulo operator (remainder of division)

Algorithm *size()*
return $(N - f + r) \bmod N$

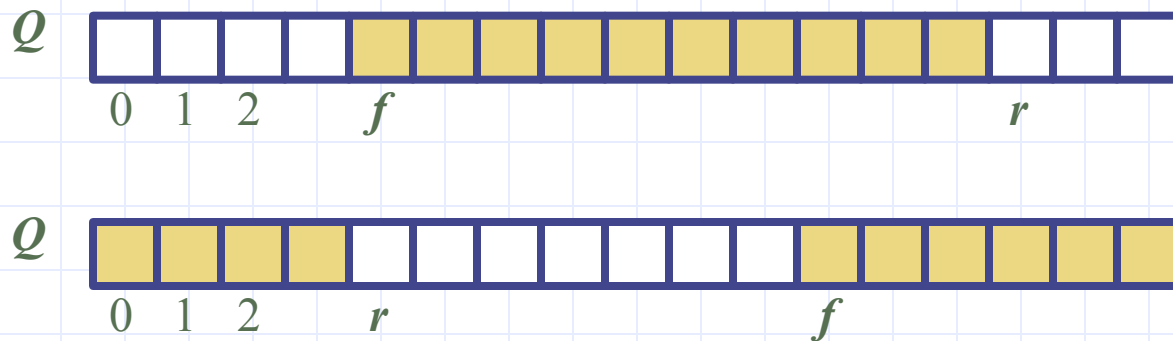
Algorithm *isEmpty()*
return $(f = r)$



Queue Operations (cont.)

- ❑ Operation enqueue throws an exception if the array is full
- ❑ This exception is implementation-dependent

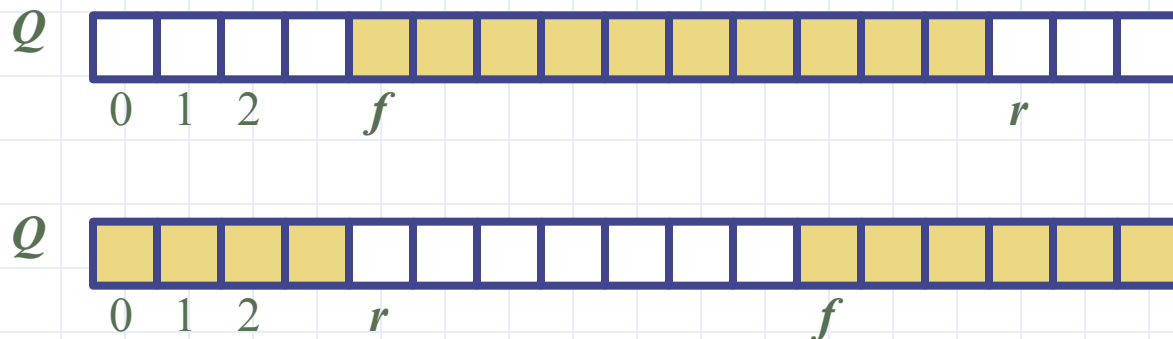
```
Algorithm enqueue(o)  
  if  $size() = N - 1$  then  
    throw FullQueueException  
  else  
     $Q[r] \leftarrow o$   
     $r \leftarrow (r + 1) \bmod N$ 
```



Queue Operations (cont.)

- ❑ Operation `dequeue` throws an exception if the queue is empty
- ❑ This exception is specified in the queue ADT

```
Algorithm dequeue()  
  if isEmpty() then  
    throw EmptyQueueException  
  else  
     $o \leftarrow Q[f]$   
     $f \leftarrow (f + 1) \bmod N$   
    return  $o$ 
```



Queue in Python

- Use the following three instance variables:
 - `_data`: is a reference to a list instance with a fixed capacity.
 - `_size`: is an integer representing the current number of elements stored in the queue (as opposed to the length of the data list).
 - `_front`: is an integer that represents the index within data of the first element of the queue (assuming the queue is not empty).

Queue in Python, Beginning

```
1 class ArrayQueue:
2     """FIFO queue implementation using a Python list as underlying storage."""
3     DEFAULT_CAPACITY = 10          # moderate capacity for all new queues
4
5     def __init__(self):
6         """Create an empty queue."""
7         self._data = [None] * ArrayQueue.DEFAULT_CAPACITY
8         self._size = 0
9         self._front = 0
10
11    def __len__(self):
12        """Return the number of elements in the queue."""
13        return self._size
14
15    def is_empty(self):
16        """Return True if the queue is empty."""
17        return self._size == 0
18
```

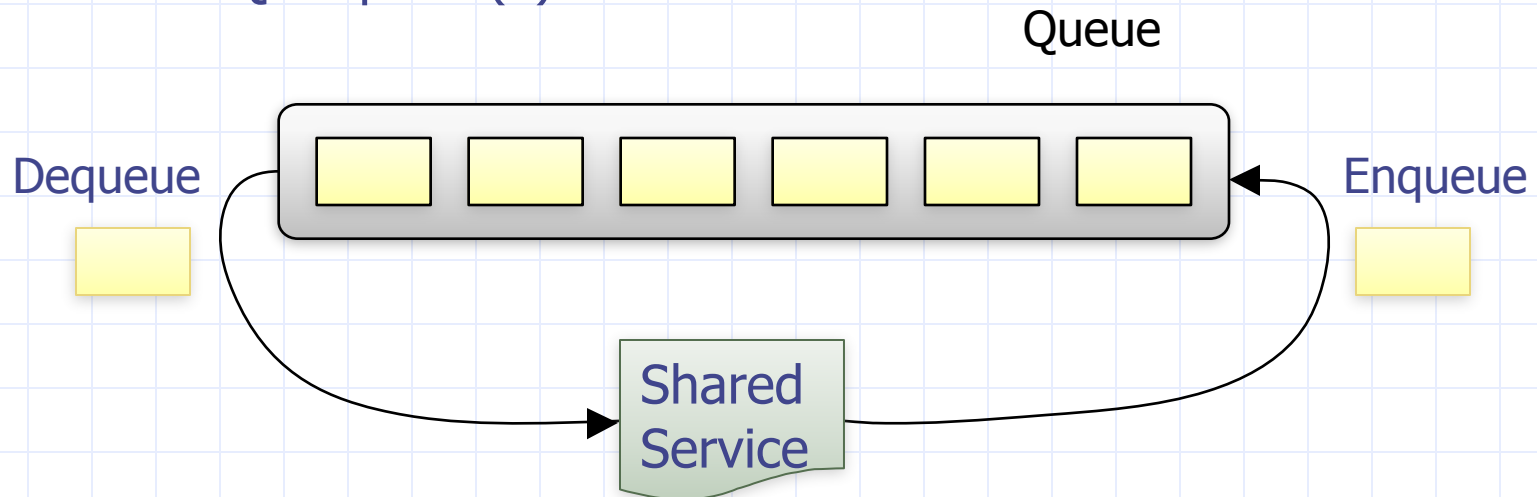
```
19    def first(self):
20        """Return (but do not remove) the element at the front of the queue.
21
22        Raise Empty exception if the queue is empty.
23        """
24        if self.is_empty():
25            raise Empty('Queue is empty')
26        return self._data[self._front]
27
28    def dequeue(self):
29        """Remove and return the first element of the queue (i.e., FIFO).
30
31        Raise Empty exception if the queue is empty.
32        """
33        if self.is_empty():
34            raise Empty('Queue is empty')
35        answer = self._data[self._front]
36        self._data[self._front] = None          # help garbage collection
37        self._front = (self._front + 1) % len(self._data)
38        self._size -= 1
39        return answer
```

Queue in Python, Continued

```
40 def enqueue(self, e):
41     """ Add an element to the back of queue. """
42     if self._size == len(self._data):
43         self._resize(2 * len(self._data))    # double the array size
44     avail = (self._front + self._size) % len(self._data)
45     self._data[avail] = e
46     self._size += 1
47
48 def _resize(self, cap):                      # we assume cap >= len(self)
49     """ Resize to a new list of capacity >= len(self). """
50     old = self._data                         # keep track of existing list
51     self._data = [None] * cap                # allocate list with new capacity
52     walk = self._front
53     for k in range(self._size):              # only consider existing elements
54         self._data[k] = old[walk]            # intentionally shift indices
55         walk = (1 + walk) % len(old)         # use old size as modulus
56     self._front = 0
```

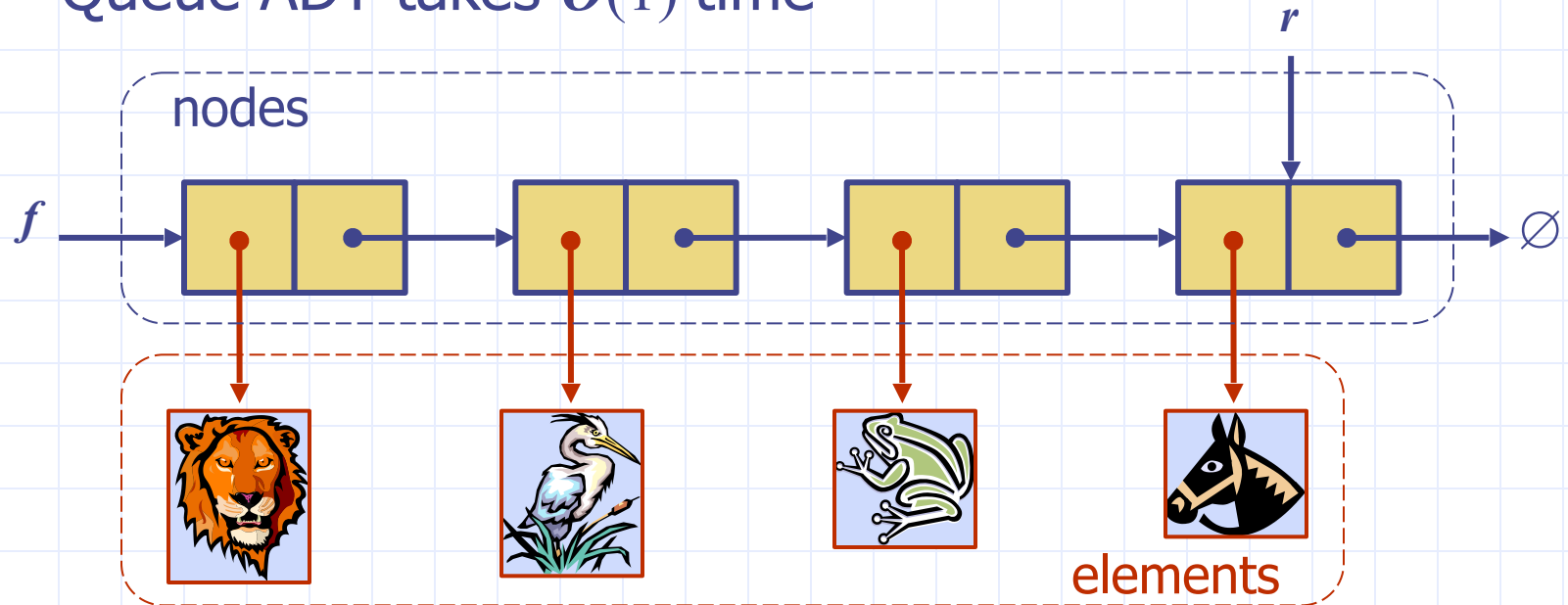
Application: Round Robin Schedulers

- We can implement a round robin scheduler using a queue Q by repeatedly performing the following steps:
 1. $e = Q.dequeue()$
 2. Service element e
 3. $Q.enqueue(e)$



Queue as a Linked List

- ◆ We can implement a queue with a singly linked list
 - The front element is stored at the first node
 - The rear element is stored at the last node
- ◆ The space used is $O(n)$ and each operation of the Queue ADT takes $O(1)$ time

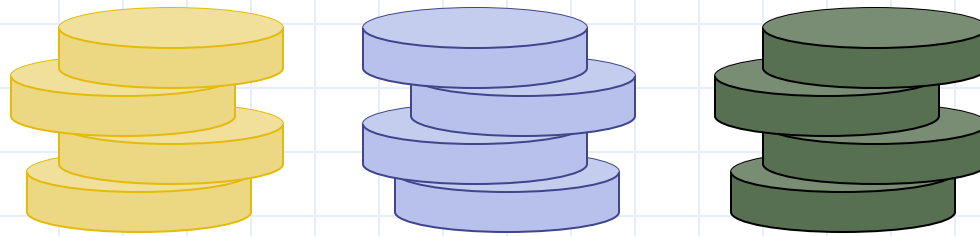


Linked-List Queue in Python

```
1 class LinkedQueue:
2     """FIFO queue implementation using a singly linked list for storage."""
3
4     class _Node:
5         """Lightweight, nonpublic class for storing a singly linked node."""
6         (omitted here; identical to that of LinkedStack._Node)
7
8     def __init__(self):
9         """Create an empty queue."""
10        self._head = None
11        self._tail = None
12        self._size = 0                # number of queue elements
13
14    def __len__(self):
15        """Return the number of elements in the queue."""
16        return self._size
17
18    def is_empty(self):
19        """Return True if the queue is empty."""
20        return self._size == 0
21
22    def first(self):
23        """Return (but do not remove) the element at the front of the queue."""
24        if self.is_empty():
25            raise Empty('Queue is empty')
26        return self._head._element    # front aligned with head of list
```

```
27    def dequeue(self):
28        """Remove and return the first element of the queue (i.e., FIFO).
29
30        Raise Empty exception if the queue is empty.
31        """
32        if self.is_empty():
33            raise Empty('Queue is empty')
34        answer = self._head._element
35        self._head = self._head._next
36        self._size -= 1
37        if self.is_empty():           # special case as queue is empty
38            self._tail = None         # removed head had been the tail
39        return answer
40
41    def enqueue(self, e):
42        """Add an element to the back of queue."""
43        newest = self._Node(e, None)   # node will be new tail node
44        if self.is_empty():
45            self._head = newest        # special case: previously empty
46        else:
47            self._tail._next = newest
48            self._tail = newest        # update reference to tail node
49            self._size += 1
```

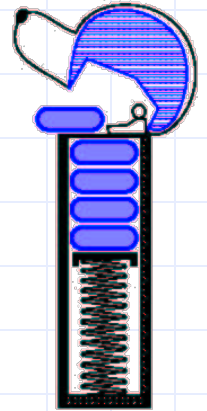
Stacks



Abstract Data Types (ADTs)

- An abstract data type (ADT) is an abstraction of a data structure
- An ADT specifies:
 - Data stored
 - Operations on the data
 - Error conditions associated with operations
- Example: ADT modeling a simple stock trading system
 - The data stored are buy/sell orders
 - The operations supported are
 - ◆ order **buy**(stock, shares, price)
 - ◆ order **sell**(stock, shares, price)
 - ◆ void **cancel**(order)
 - Error conditions:
 - ◆ Buy/sell a nonexistent stock
 - ◆ Cancel a nonexistent order

The Stack ADT



- ❑ The **Stack** ADT stores arbitrary objects
- ❑ Insertions and deletions follow the last-in first-out scheme
- ❑ Think of a spring-loaded plate dispenser
- ❑ Main stack operations:
 - **push**(object): inserts an element
 - object **pop**(): removes and returns the last inserted element
- ❑ Auxiliary stack operations:
 - object **top**(): returns the last inserted element without removing it
 - integer **len**(): returns the number of elements stored
 - boolean **is_empty**(): indicates whether no elements are stored

Example

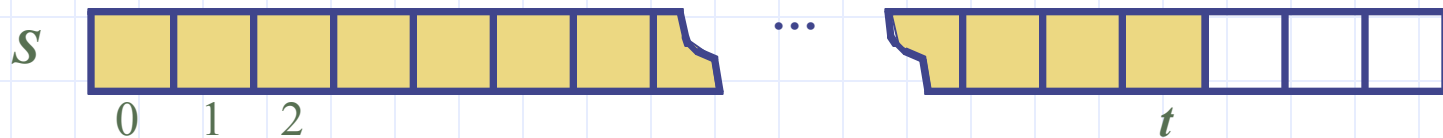
Operation	Return Value	Stack Contents
S.push(5)	—	[5]
S.push(3)	—	[5, 3]
len(S)	2	[5, 3]
S.pop()	3	[5]
S.is_empty()	False	[5]
S.pop()	5	[]
S.is_empty()	True	[]
S.pop()	“error”	[]
S.push(7)	—	[7]
S.push(9)	—	[7, 9]
S.top()	9	[7, 9]
S.push(4)	—	[7, 9, 4]
len(S)	3	[7, 9, 4]
S.pop()	4	[7, 9]
S.push(6)	—	[7, 9, 6]
S.push(8)	—	[7, 9, 6, 8]
S.pop()	8	[7, 9, 6]

Applications of Stacks

- ❑ Direct applications
 - Page-visited history in a Web browser
 - Undo sequence in a text editor
 - Chain of method calls in a language that supports recursion
- ❑ Indirect applications
 - Auxiliary data structure for algorithms
 - Component of other data structures

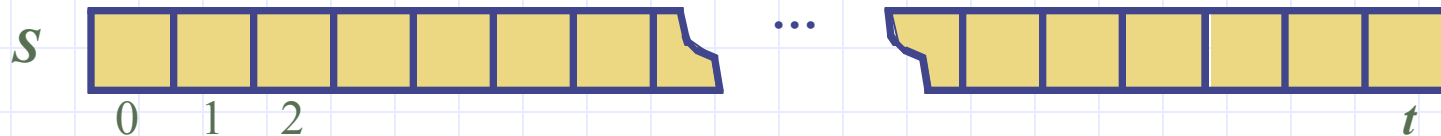
Array-based Stack

- ❑ A simple way of implementing the Stack ADT uses an array
- ❑ We add elements from left to right
- ❑ A variable keeps track of the index of the top element



Array-based Stack (cont.)

- ❑ The array storing the stack elements may become full
- ❑ A push operation will then need to grow the array and copy all the elements over.



Performance and Limitations

□ Performance

- Let n be the number of elements in the stack
- The space used is $O(n)$
- Each operation runs in time $O(1)$ (amortized in the case of a push)

Array-based Stack in Python



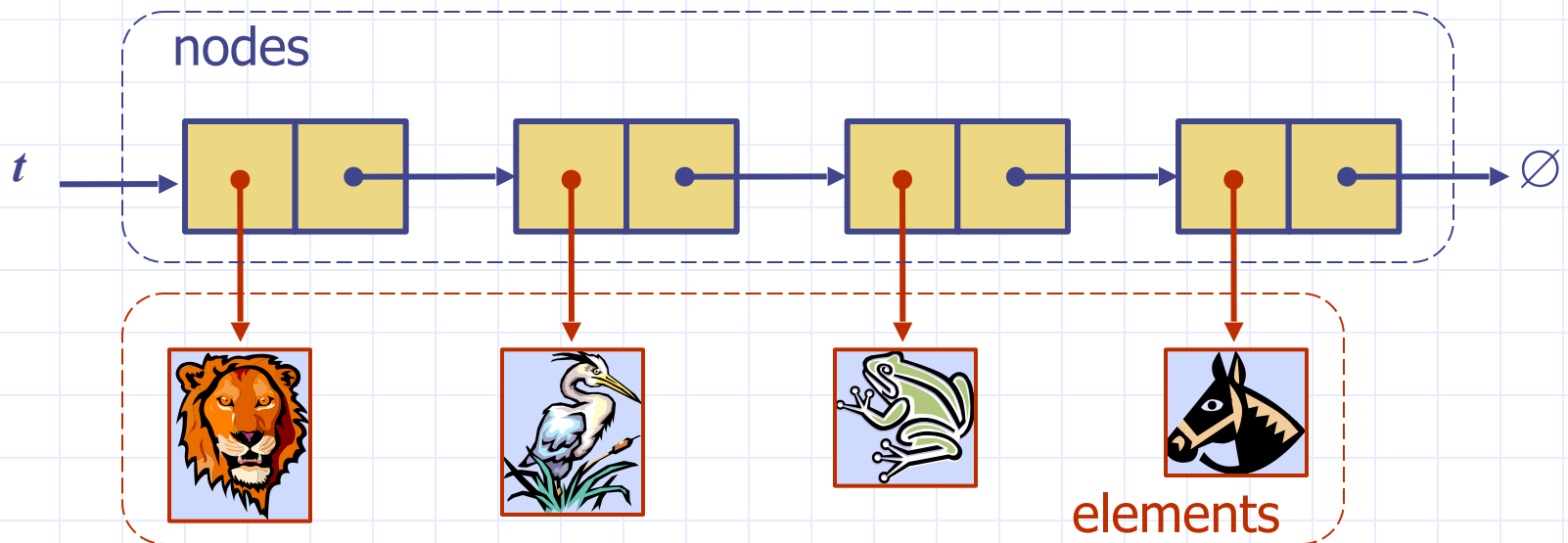
```
1 class ArrayStack:
2     """LIFO Stack implementation using a Python list as underlying storage."""
3
4     def __init__(self):
5         """Create an empty stack."""
6         self._data = [ ]           # nonpublic list instance
7
8     def __len__(self):
9         """Return the number of elements in the stack."""
10        return len(self._data)
11
12    def is_empty(self):
13        """Return True if the stack is empty."""
14        return len(self._data) == 0
15
16    def push(self, e):
17        """Add element e to the top of the stack."""
18        self._data.append(e)        # new item stored at end of list
19
```

```
20    def top(self):
21        """Return (but do not remove) the element at the top of the stack.
22
23        Raise Empty exception if the stack is empty.
24        """
25        if self.is_empty():
26            raise Empty('Stack is empty')
27        return self._data[-1]       # the last item in the list
28
29    def pop(self):
30        """Remove and return the element from the top of the stack (i.e., LIFO).
31
32        Raise Empty exception if the stack is empty.
33        """
34        if self.is_empty():
35            raise Empty('Stack is empty')
36        return self._data.pop( )    # remove last item from list

```

Stack as a Linked List

- ◆ We can implement a stack with a singly linked list
- ◆ The top element is stored at the first node of the list
- ◆ The space used is $O(n)$ and each operation of the Stack ADT takes $O(1)$ time



Linked-List Stack in Python

```
1 class LinkedStack:
2     """LIFO Stack implementation using a singly linked list for storage."""
3
4     #----- nested _Node class -----
5     class _Node:
6         """Lightweight, nonpublic class for storing a singly linked node."""
7         __slots__ = '_element', '_next' # streamline memory usage
8
9         def __init__(self, element, next): # initialize node's fields
10             self._element = element # reference to user's element
11             self._next = next # reference to next node
12
13     #----- stack methods -----
14     def __init__(self):
15         """Create an empty stack."""
16         self._head = None # reference to the head node
17         self._size = 0 # number of stack elements
18
19     def __len__(self):
20         """Return the number of elements in the stack."""
21         return self._size
22
```

```
23     def is_empty(self):
24         """Return True if the stack is empty."""
25         return self._size == 0
26
27     def push(self, e):
28         """Add element e to the top of the stack."""
29         self._head = self._Node(e, self._head) # create and link a new node
30         self._size += 1
31
32     def top(self):
33         """Return (but do not remove) the element at the top of the stack.
34
35         Raise Empty exception if the stack is empty.
36         """
37         if self.is_empty():
38             raise Empty('Stack is empty')
39         return self._head._element # top of stack is at head of list

```

```
40     def pop(self):
41         """Remove and return the element from the top of the stack (i.e., LIFO).
42
43         Raise Empty exception if the stack is empty.
44         """
45         if self.is_empty():
46             raise Empty('Stack is empty')
47         answer = self._head._element
48         self._head = self._head._next # bypass the former top node
49         self._size -= 1
50         return answer

```

Parentheses Matching

- Each “(”, “{”, or “[” must be paired with a matching “)”, “}”, or “]”
 - correct: ()(()){([())}
 - correct: ((())(()){([())}
 - incorrect:)(()){([())}
 - incorrect: ({ []})
 - incorrect: (

Parentheses Matching Algorithm

Algorithm ParenMatch(X, n):

Input: An array X of n tokens, each of which is either a grouping symbol, a variable, an arithmetic operator, or a number

Output: **true** if and only if all the grouping symbols in X match

Let S be an empty stack

for $i=0$ to $n-1$ **do**

if $X[i]$ is an opening grouping symbol **then**

$S.push(X[i])$

else if $X[i]$ is a closing grouping symbol **then**

if $S.is_empty()$ **then**

return false {nothing to match with}

if $S.pop()$ does not match the type of $X[i]$ **then**

return false {wrong type}

if $S.isEmpty()$ **then**

return true {every symbol matched}

else return false {some symbols were never matched}

Parentheses Matching in Python

```
1 def is_matched(expr):
2     """Return True if all delimiters are properly match; False otherwise."""
3     lefty = '([{' # opening delimiters
4     righty = ')]}' # respective closing delims
5     S = ArrayStack()
6     for c in expr:
7         if c in lefty:
8             S.push(c) # push left delimiter on stack
9         elif c in righty:
10            if S.is_empty():
11                return False # nothing to match with
12            if righty.index(c) != lefty.index(S.pop()):
13                return False # mismatched
14    return S.is_empty() # were all symbols matched?
```

Evaluating Arithmetic Expressions

$$14 - 3 * 2 + 7 = (14 - (3 * 2)) + 7$$

Operator precedence

* has precedence over +/−

Associativity

operators of the same precedence group
evaluated from left to right

Example: $(x - y) + z$ rather than $x - (y + z)$

Idea: push each operator on the stack, but first pop and perform higher and equal precedence operations.

Algorithm for Evaluating Expressions

Two stacks:

- opStk holds operators
- valStk holds values
- Use \$ as special “end of input” token with lowest precedence

Algorithm **doOp()**

```
x ← valStk.pop();  
y ← valStk.pop();  
op ← opStk.pop();  
valStk.push( y op x )
```

Algorithm **repeatOps(refOp)**:

```
while ( valStk.size() > 1 ∧  
        prec(refOp) ≤  
        prec(opStk.top())  
    doOp()
```

Algorithm **EvalExp()**

Input: a stream of tokens representing
an arithmetic expression (with
numbers)

Output: the value of the expression

while there's another token z

if isNumber(z) **then**

valStk.push(z)

else

repeatOps(z);

opStk.push(z)

repeatOps(\$);

return valStk.top()

Algorithm on an Example Expression

14 ≤ 4 - 3 * 2 + 7

Operator ≤ has lower
precedence than +/−

