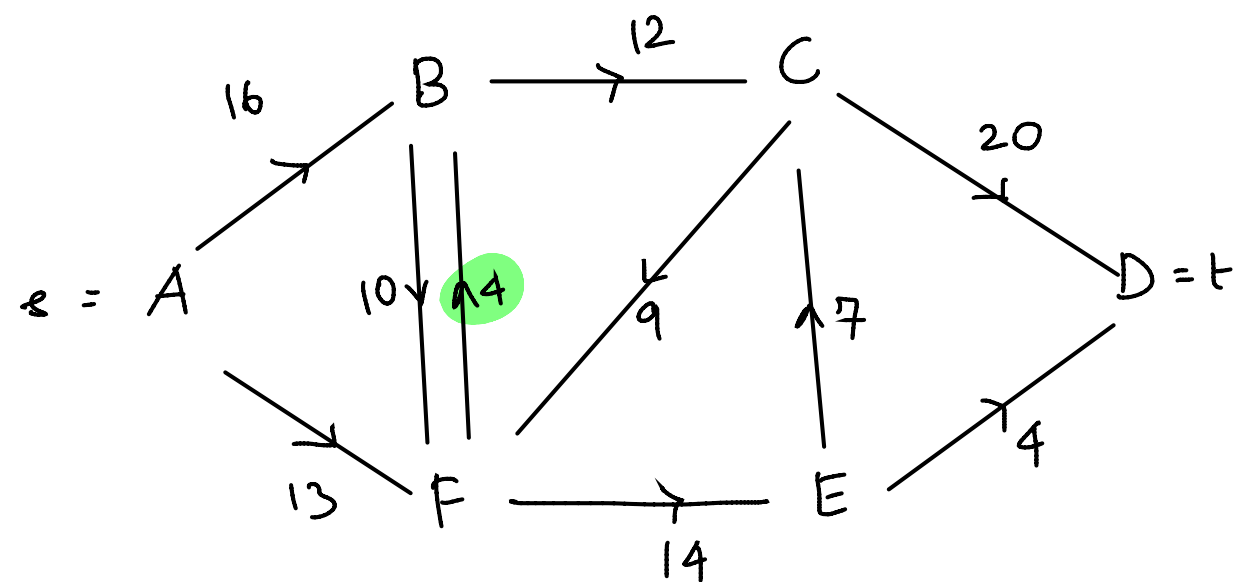


# Residual Network

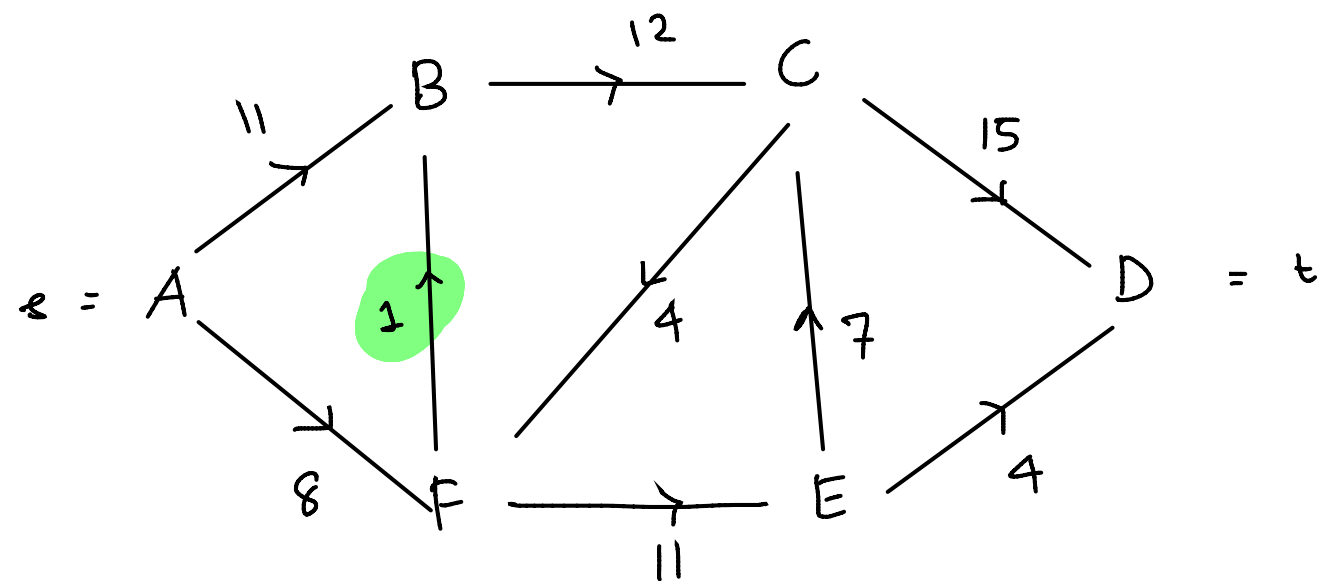
Residual capacity :

$$c_f(u, v) = c(u, v) - f(u, v)$$

Edge capacity

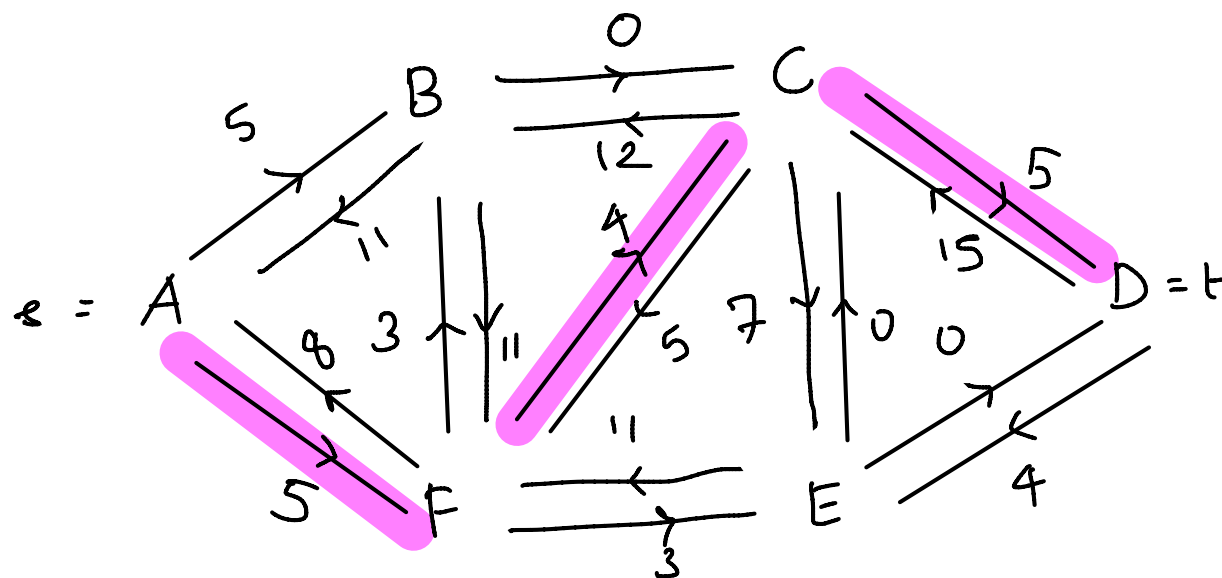


Flow f



## Residual Graph

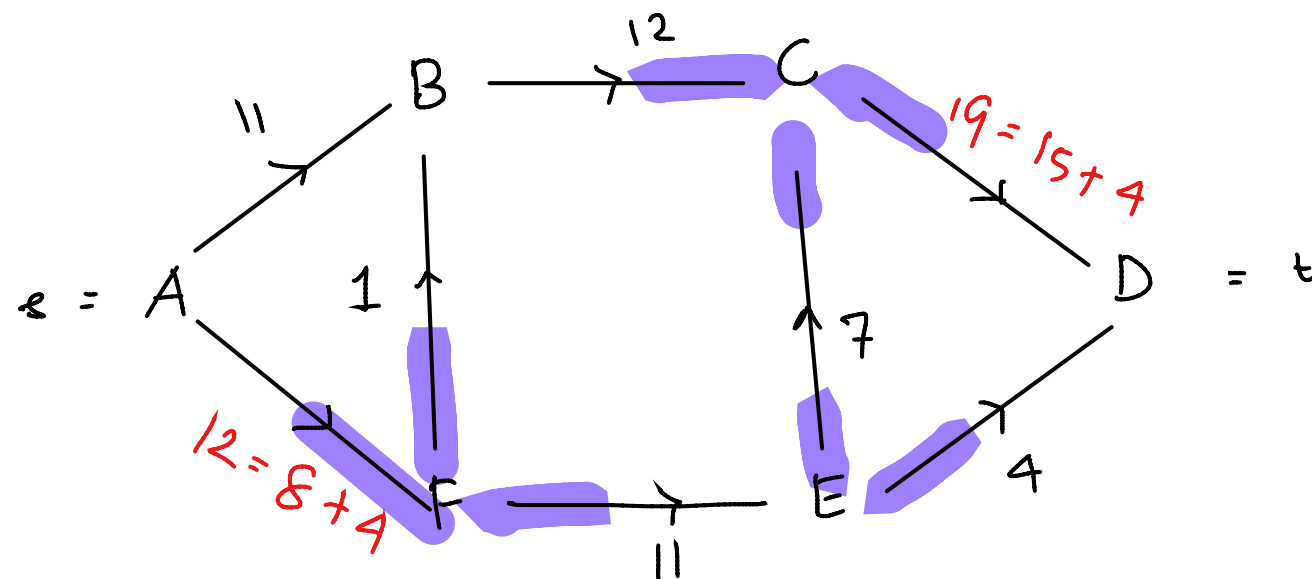
$$G_f = (V, c_f)$$



$$c_f(A, F) = 13 - 8 = 5, \quad c_f(F, A) = 0 - f(F, A) = 0 - (-f(A, F)) = f(A, F)$$

An Augmenting path is  $A \rightarrow F \rightarrow C \rightarrow D$ , path capacity = 4

New Flow:

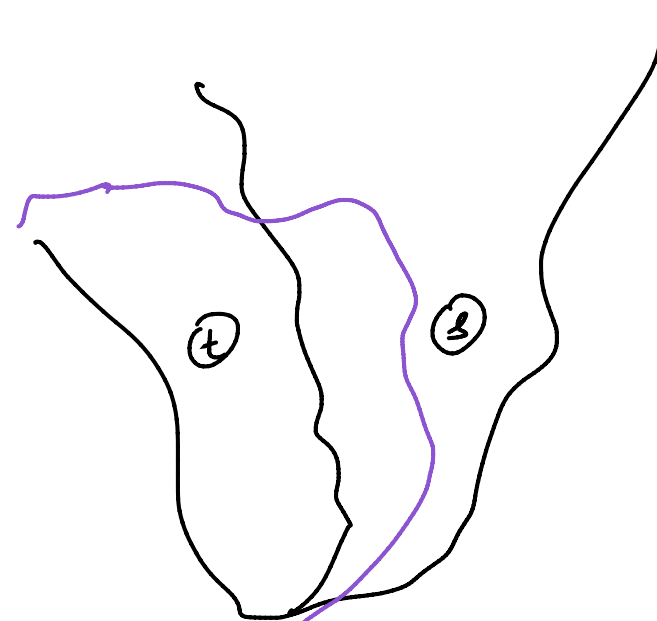
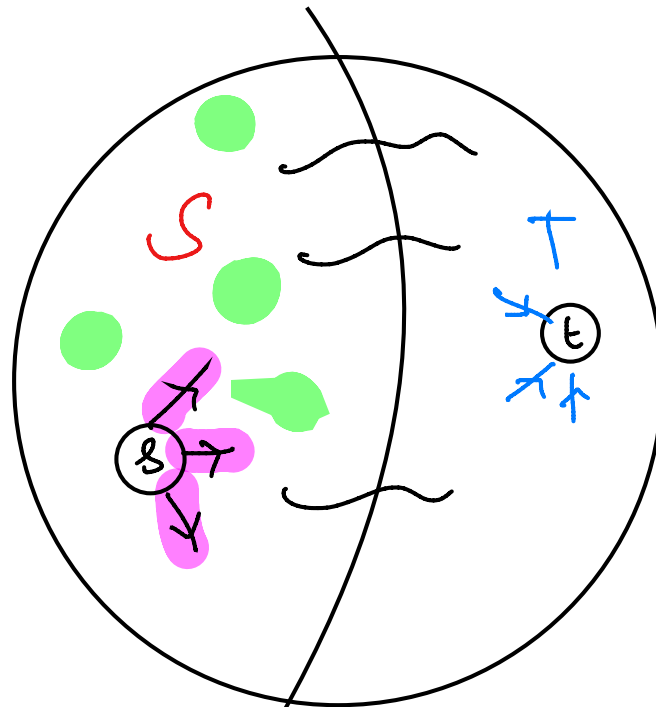


Can calculate a new residual capacity

Fact:  $|f|$  was for  $G$ ,  $|f'|$  was for  $G_f$

$$|f + f'| = |f| + |f'|$$

cut :  $(S, T)$   
 $s \in S$   $t \in T$



$$f(S, V) = f(S, V - S + S) = f(S, T) + \underset{0}{f(S, S)}$$

$$= f(S, V - S)$$

$$f(S, T) = f(S, V - S)$$

$$= f(S, V) - \overset{0}{f(S, S)}$$

$$= \overset{=0}{f(S, V)} + f(\overset{=0}{S - s}, V)$$

$$= |f|$$

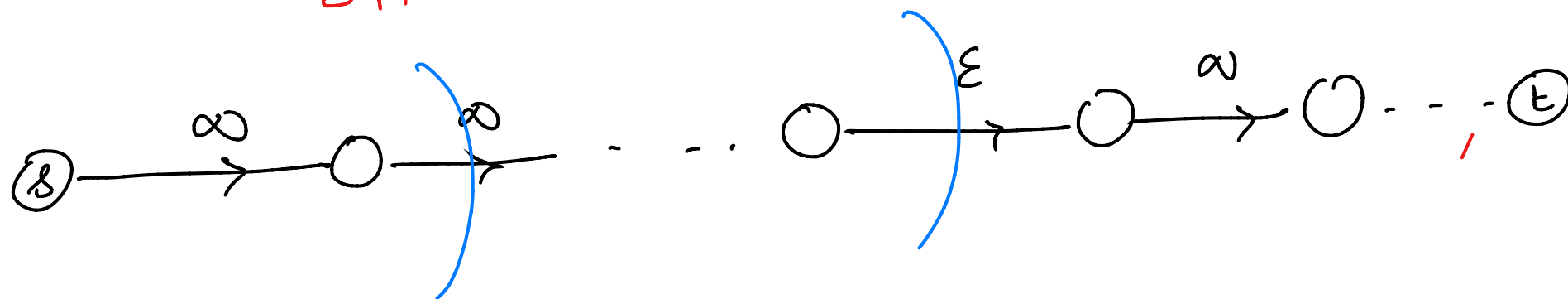
Fact :  $f(s, t) = \sum_{u \in S} \sum_{v \in T} f(u, v)$

$$\leq \sum_{u \in S} \sum_{v \in T} c(u, v)$$

$$= c(S, T)$$

$\Rightarrow |f| \leq c(S, T)$   
 $\uparrow$   
 varies

$\Rightarrow |f| \leq \min_{S, T} c(S, T)$



## Max-Flow - Min-Cut

The following are equivalent

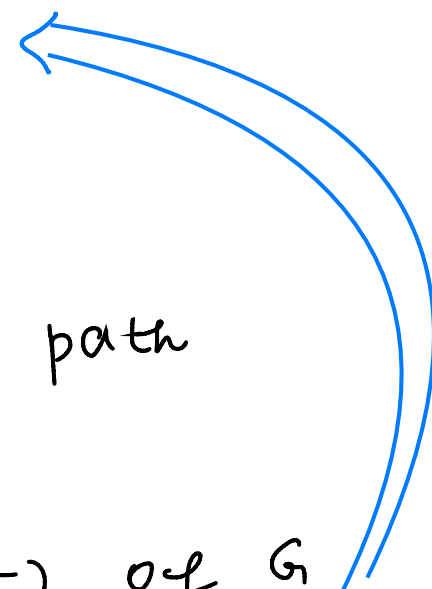
- $f$  is a maximum flow in  $G$



- residual  $G_f$  has no augmenting path



- $|f| = C(S, T)$  for some cut  $(S, T)$  of  $G$



(1)  $\Rightarrow$  (2)

(2)  $\Rightarrow$  (1) say there exists an augmenting path in  $G_f$

then  $|f + f'| = |f| + |f'| > |f| \Rightarrow$

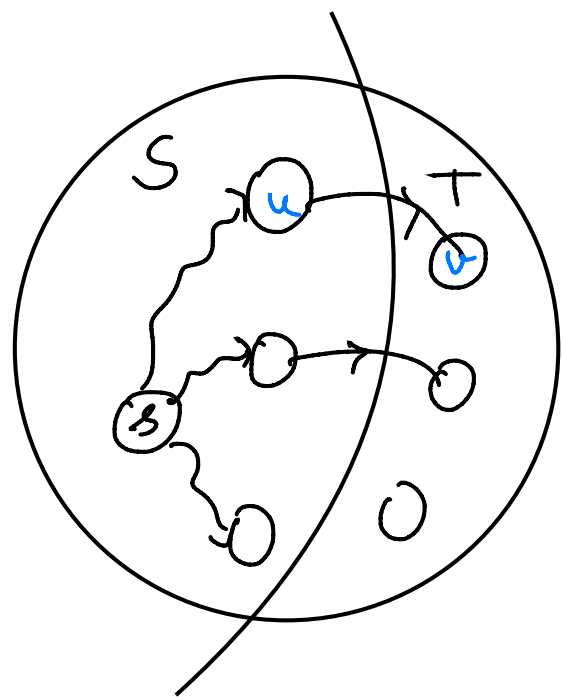
$f$  is not maximum

(2)  $\Rightarrow$  (3)

Push as much tree boundary

$S = \{ v \in V : \text{there is a path from } s \text{ to } v \text{ in } G_f \}$

$S \neq V$  ( because no augmenting path,  $t$  cannot be reached by any path. So definitely  $t$  is not in  $S$  )



$u \in S, v \in T$

$$c_f(u, v) = 0$$

$$c_f(u, v) = c(u, v) - f(u, v)$$

$$\Rightarrow f(u, v) = c(u, v)$$

$$|f| = f(S, T) = c(S, T)$$

(3)  $\Rightarrow$  (1)

equality is achieved  $\Rightarrow$  flow is maximum

## FORD - FULKERSON $(G, s, t)$

for each edge  $(u, v) \in E$   
do  $f(u, v) \leftarrow 0$   
 $f(v, u) \leftarrow 0$

while there exists a path from  $s$  to  $t$  in  $G_f$   
do  $C_f(p) \leftarrow \min \{C_f(u, v) : (u, v) \text{ is in } p\}$

for each edge  $(u, v)$  in  $p$ :

do  $f(u, v) \leftarrow f(u, v) + C_f(p)$   
 $f(v, u) \leftarrow -f(u, v)$