* 2d or 3d tree Ed tree

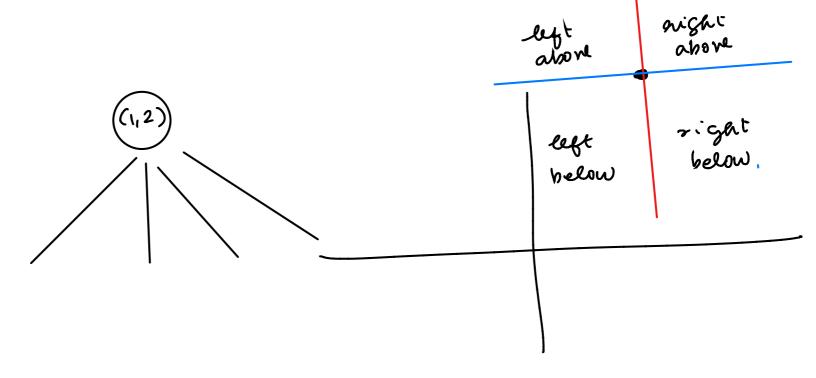
At each level of the thee, we compare with a ponticular dim

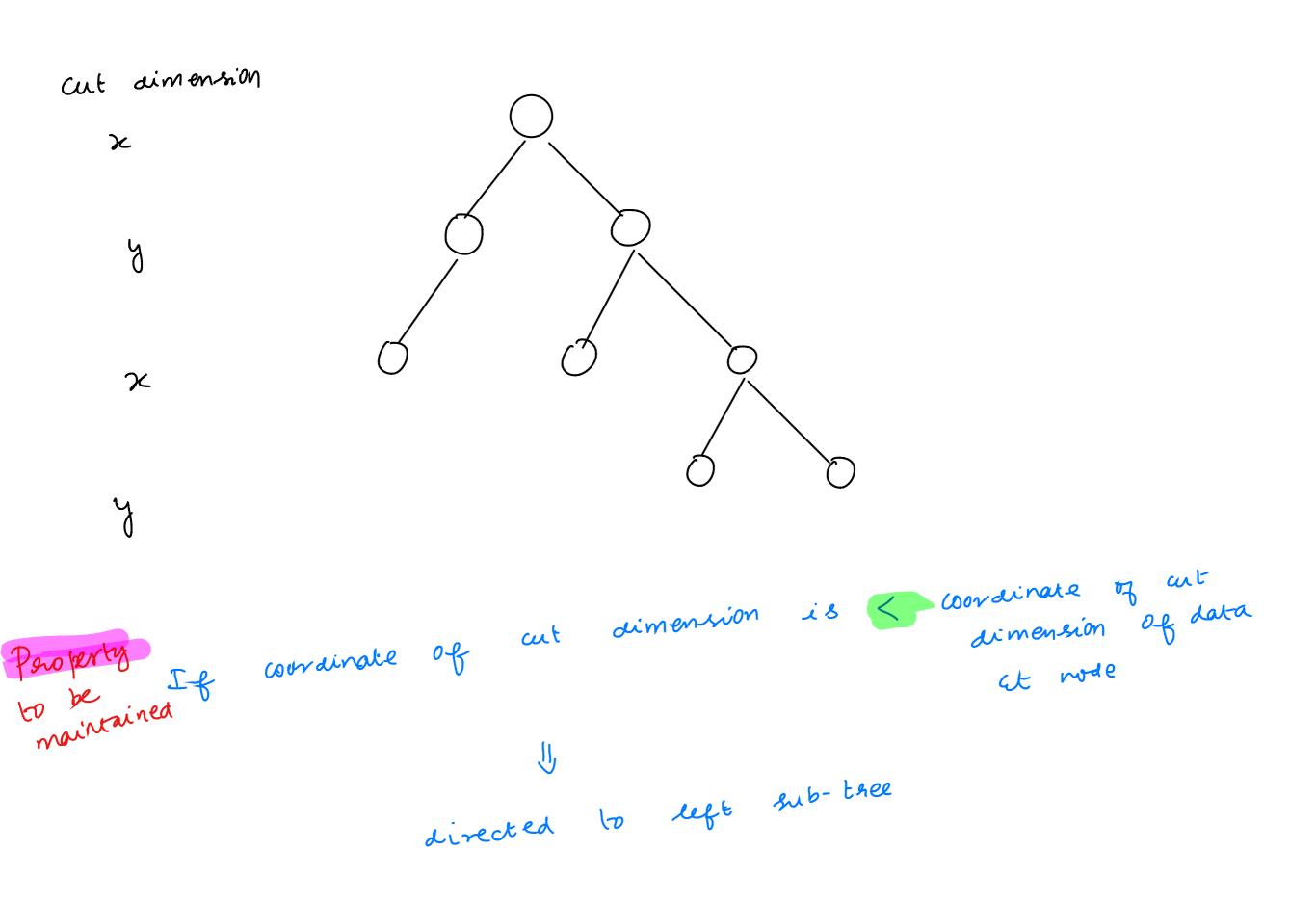
4 Typically one circulates dimensions (round robin)

a cut aimension of that level

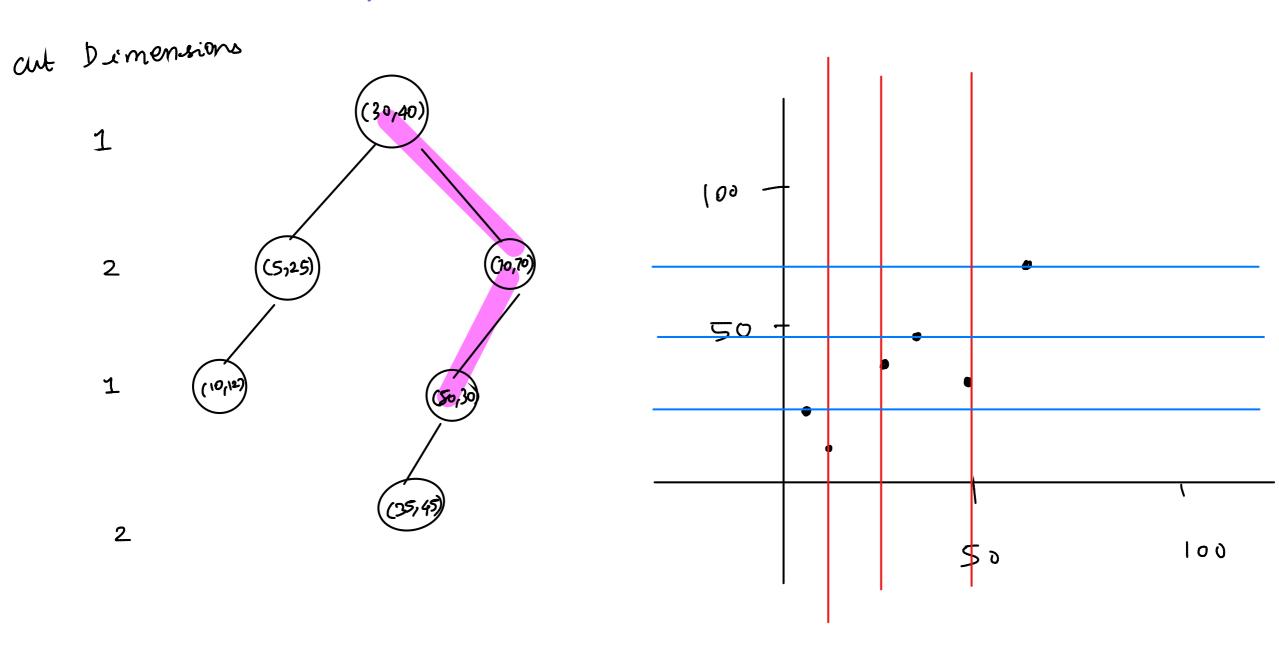
only 2 children for each

a If all armencions are cut we get 2d children





Data:
$$(30,40)$$
, $(5,25)$, $(10,12)$, $(70,70)$, $(50,30)$, $(35,45)$



insert. (x, urde t, cut-dim) if t = NVLL t = new mae (x) iq (x = = t. data) say it is duplicate if (x cont.dim) < t.data [cut-dim]) insert (2, t. left, (cut-dim+1)% total-dim) UX insert (2, t. right, (cut-dim ti) %. total-dim) evsl

Since there no complete ordering

Find Min (dim)

A Reenhirely browerse the thee

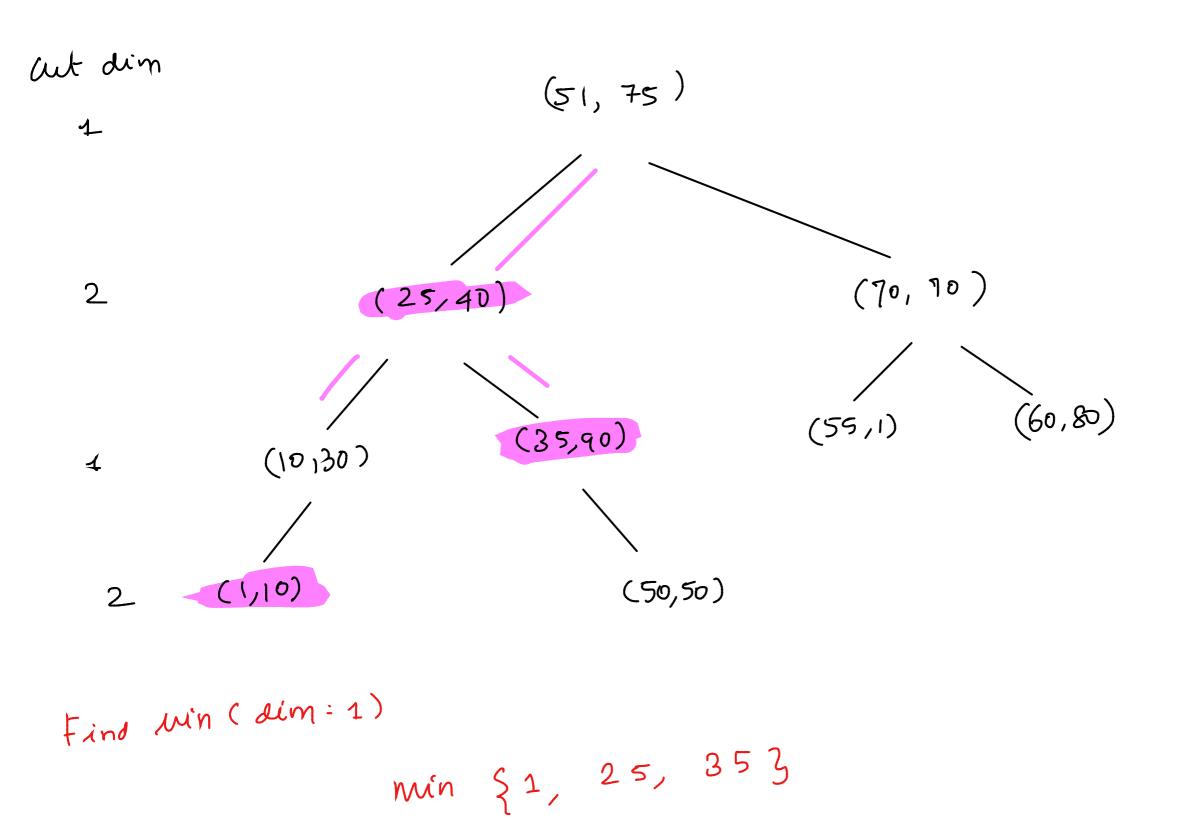
s at the current level.

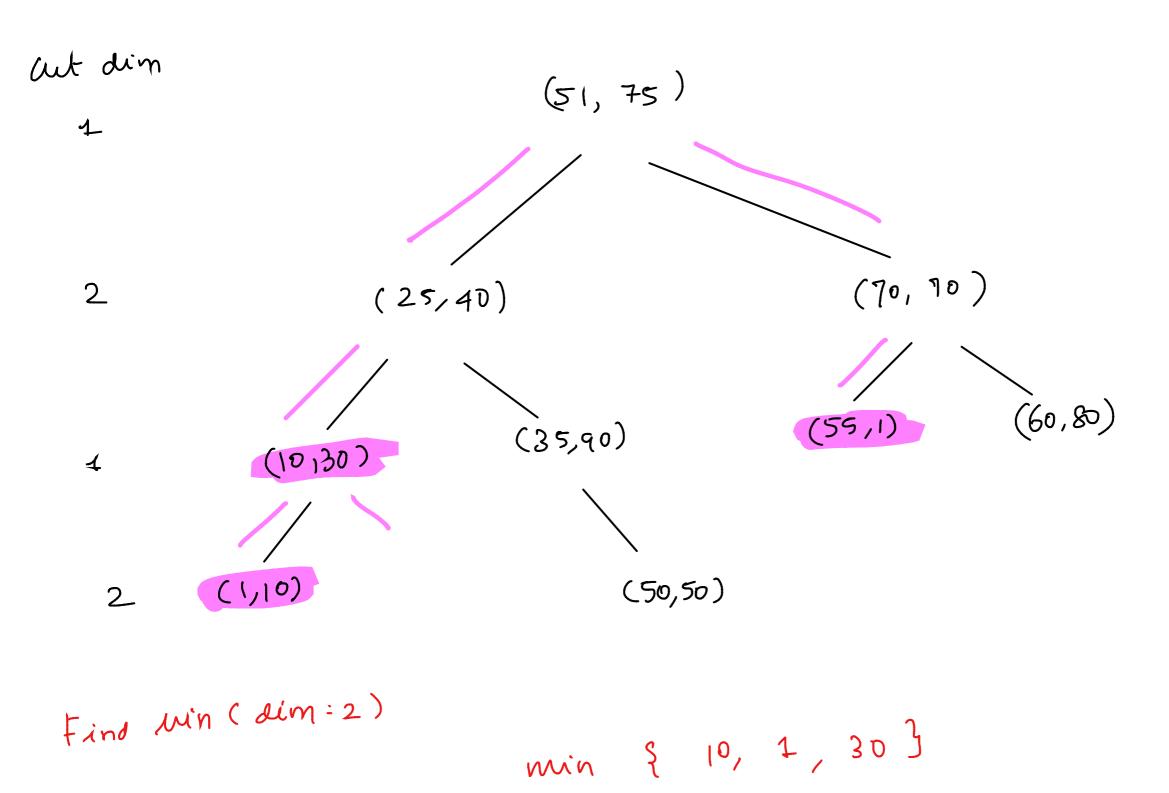
if at-din = dim

we can choose left

else

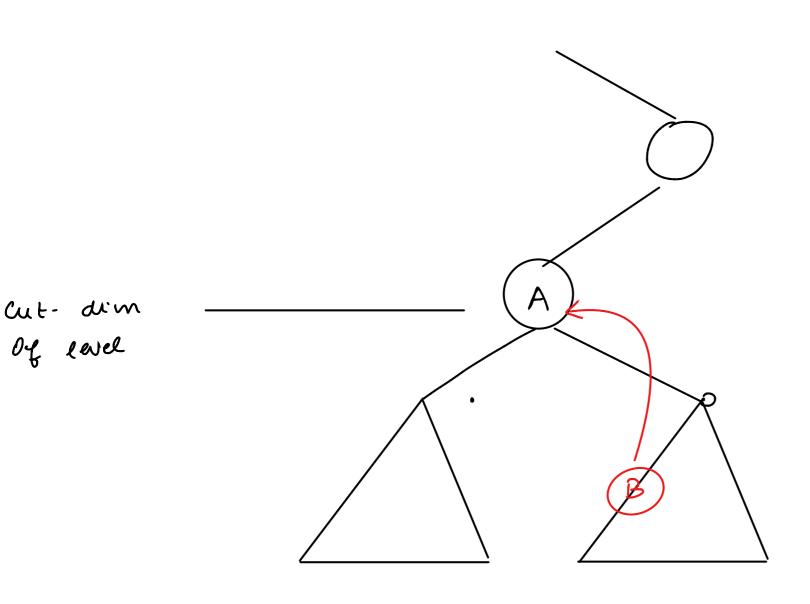
ue have to search both sides





Find Min (mode t, dim, out-dim) cd = - dim if t. left = = NULL: neturn t. data use seturn Find Min (t. left, dim, (ut.dim+1)% total-dim) else return min (t. data, Find win (t. left, dim, (cut-dim H)% total-dim), Find Min (t. augat, dim, (cut-dim H)% total-dim) Say we want to delete A, if night suppress is not empty

Find min (t. night, cut-dim at level A)



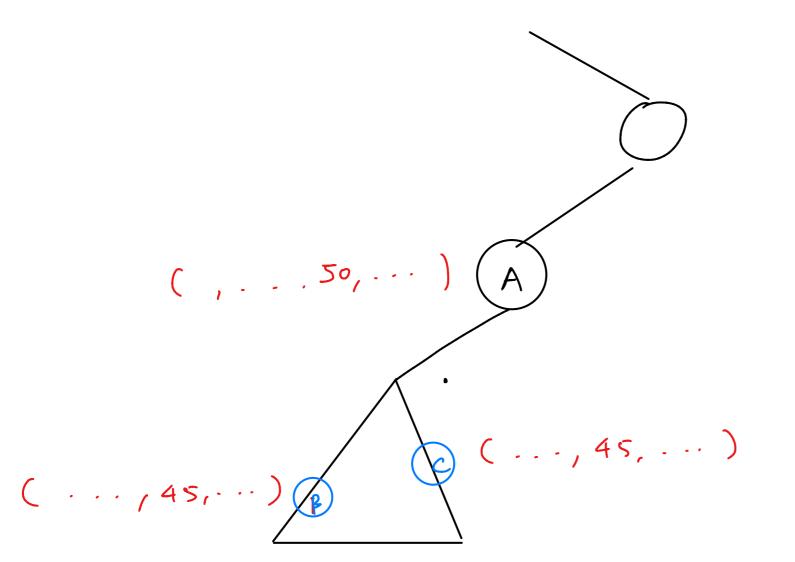
Of level

B replaces A Now Bis empty

of deletes until a leaf gets deleted

this initiates a societ

Say the night subtree is empty, search for the maximum of the left subtree



Say benak the tie somehow, say 13 was chosen to replace, we vio lake the < property

· make the left sub-tree the right sub-tree

trind min of the enight subtree

(, . . . 50, . . .) A

(--.., 1, ...