Binary Search Tree: Need not be balanced (BST)

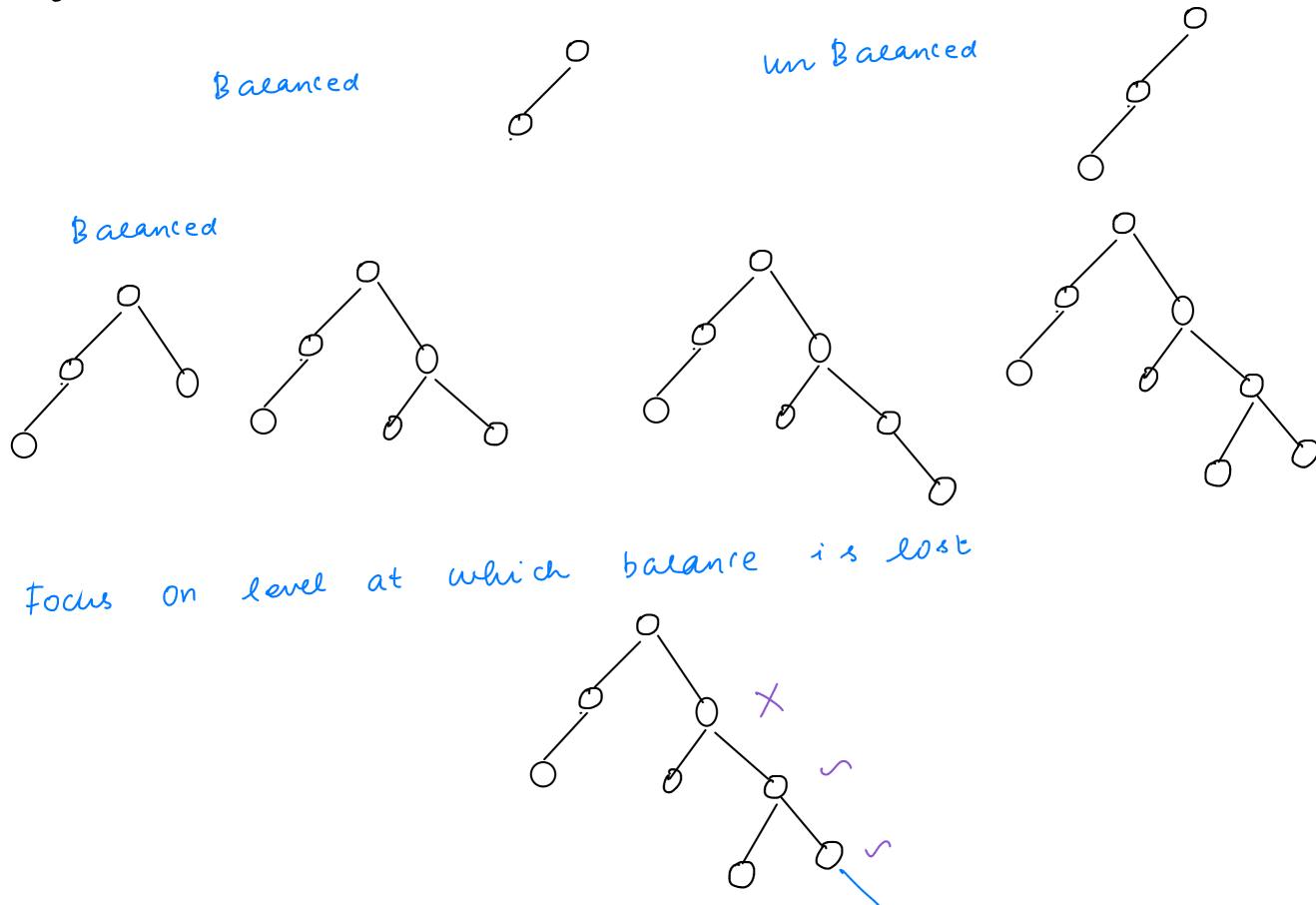
. Balance Factor:

BF(n) = Height of enight subtree - Height of Left subtree

BF(n) 70: sight heavy

BF(n) < 0: left heavy

- · AVL Tree which are BST, |BF(n)| < 1
 Adellon, Yelsky, Landis
- . Peroperty $h = O(log(n)) \leq C log(n)$ (paove it later)



Insertion

Case 1
Invert
to either
left or
right of
left

Care 2:

auready

one child

is there

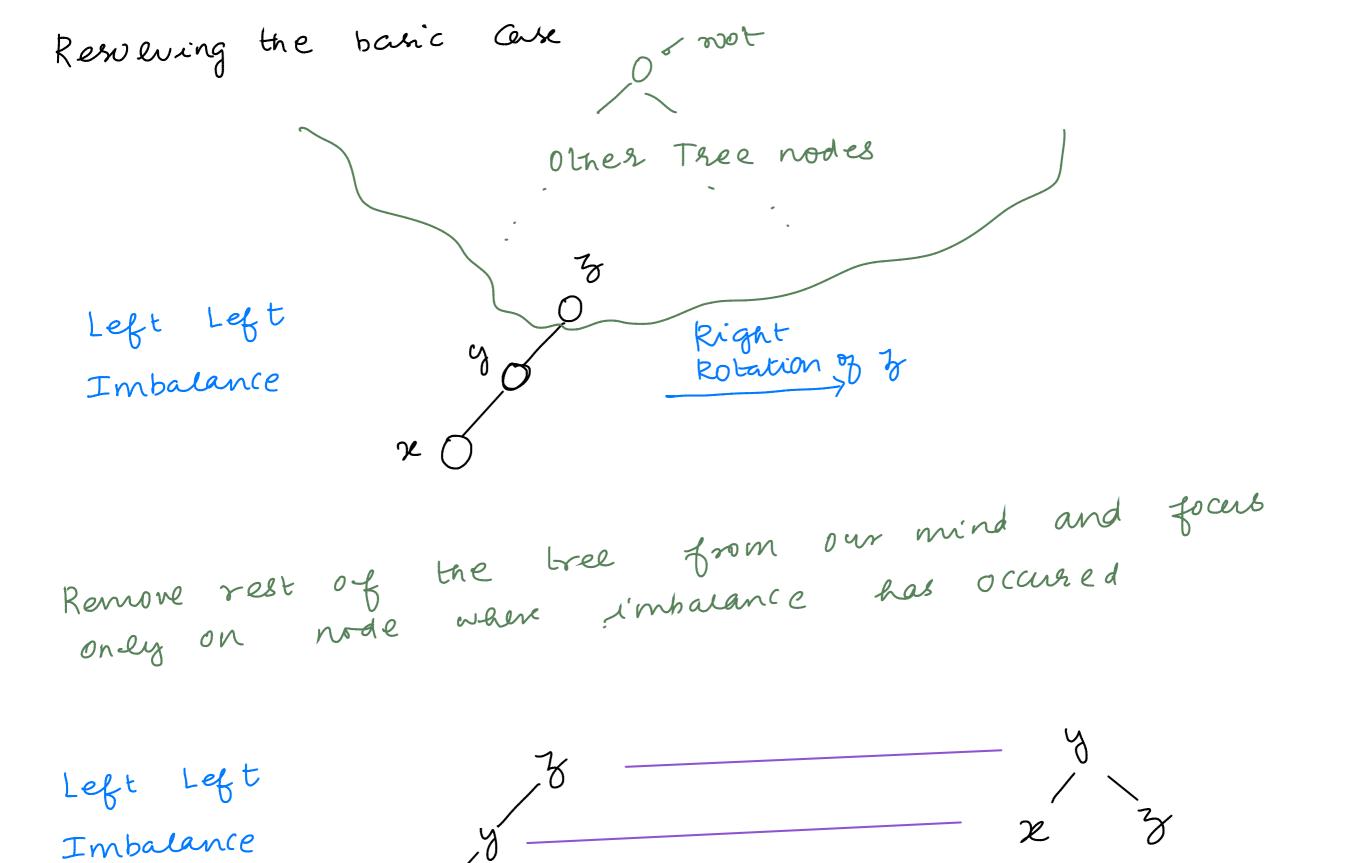
and we

are adding

a second child

Immediate parent to which the new node is interted will not have impalance

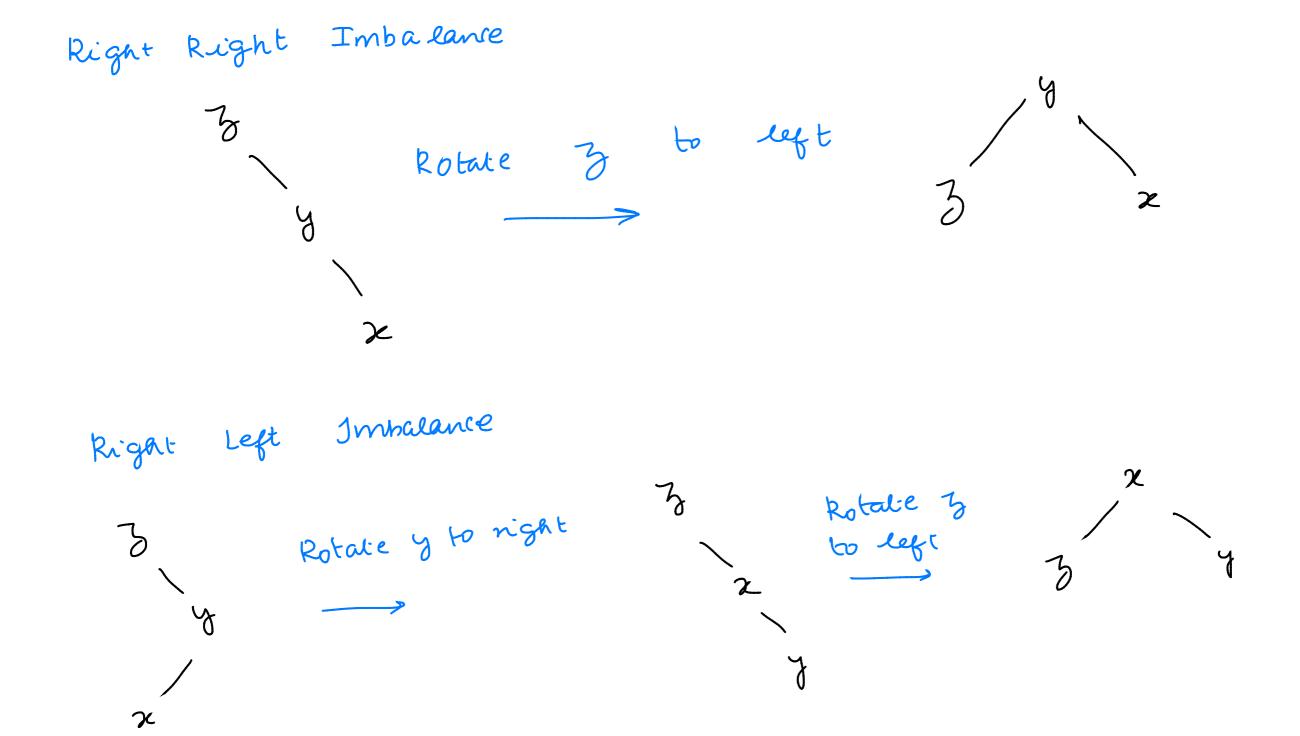
Simplest Configuration in which Barric Coure: Adding New node causes imbalance Before adding node no imbalance imbalance Immediate Grand parent has Banic Case:



Right

Robation 3 3

Left Right Imbalance left left impalance by doing First reduce it to left notation of y X Right
Robation of (3)



great grand parent of

grand parent of

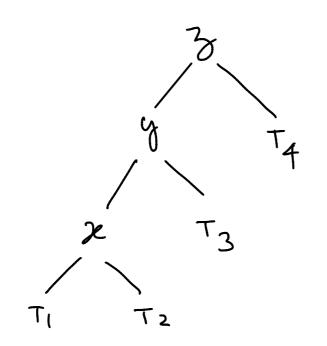
parent of

parent of

new note 0

xyz snifted up one level

General Casi. Xyz saifts up by arbitrary levels



Left Left Imbalance
(Property I) Before addition are nodes are balanced

(The property I) Before addition are nodes are balanced

(The property I) Before addition are nodes are balanced

(property)

After addition nodes x,y are balanced and y is imbalanced

After addition

y Ty

To new

To Te

To is To after adding New Node

Property I ⇒ \times · $\mathcal{L}(T_2) = \mathcal{L}(T_1) - 1$ \times • $h(T_2) = h(T_1)+1$

(not possible; cannot change height at lexel 2 by adding new node to T2) (not possible; because after adding new node to T2, level 2 it self becomes unb danced => contradicts property II) $\sqrt{\cdot}$ $\lambda (T_2) = \lambda (T_1)$ $\Rightarrow h(+_2^{\text{naw}}) = h(T_1) + 1$

Property IT

After adding, imbalance is at level 3; T4 is lighter

$$\mathcal{L}(T_2^{\text{new}}) + 2 \qquad \mathcal{L}(T_4^{\text{new}}) + 2 - \mathcal{L}(T_4^{\text{new}})$$