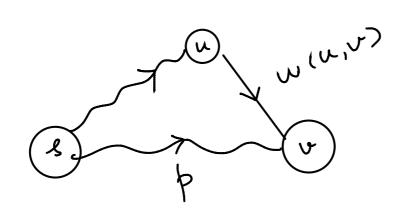
1) Taiongle Inequality

$$S(s,v) \in S(s,u) + \omega(u,v)$$



$$A \xrightarrow{2} B$$

$$A = 2$$

$$A = 2$$

$$A = 2$$

$$A = 2$$

$$A = 3$$

2) upper bound peroperty: During the course of any algorithm
for shortest pain via relaxing edges

$$d[v] \geq \delta(s,v)$$
 and once $d[v] = \delta(s,v)$
estimated

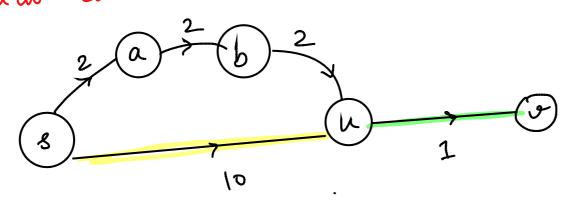
then it never changes.

4) convergence Property:

ら~りょ

if $d \in U = \delta(s, u)$ at any time point to according (u, v), then on relaxing (u, v), then on relaxing (u, v) you will have $d \in U = \delta(s, v)$ and it stays that way later as well.

Shortest Pata to v is via u



Step 1: relax edge $(8,u) \Rightarrow d[u]=10$

2 ' " d[u]=11

3: " (b, n) =) d[n] = 10

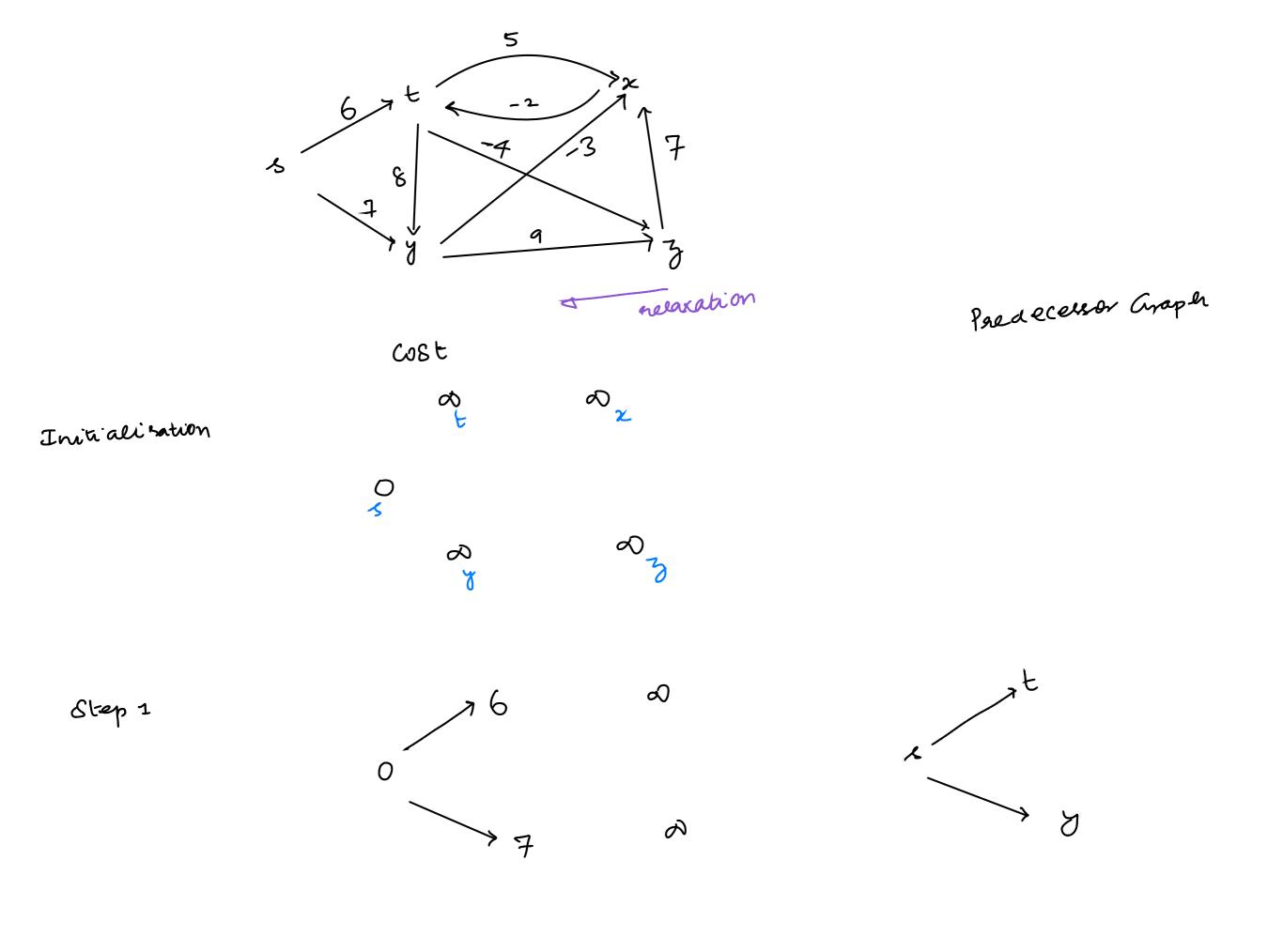
4: (a.b) =) d[b] =0

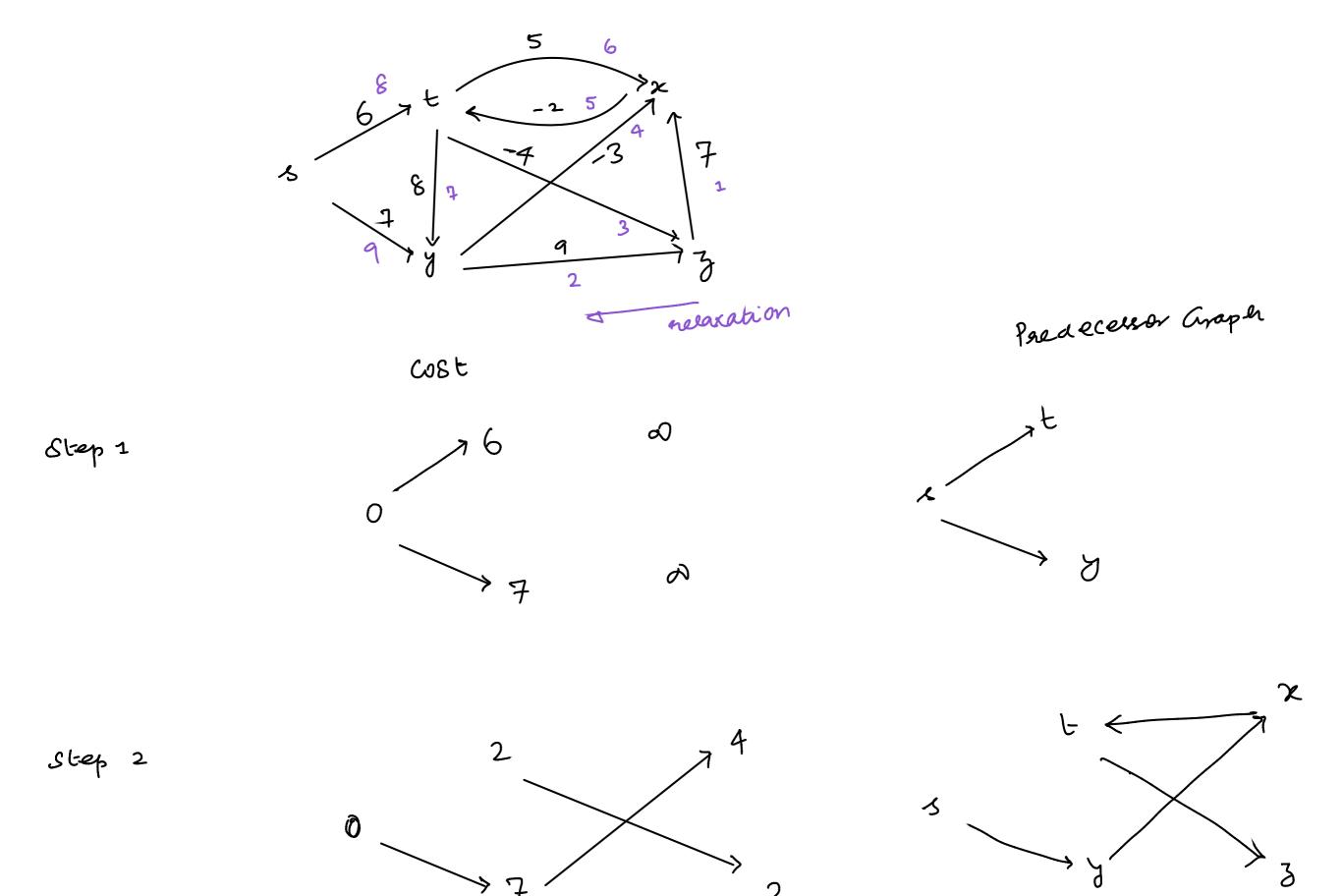
• $p = \langle v_0, v_1, \dots, v_k \rangle$ is shortest path from b to v, and the edges are relaxed in the following order (σ_{0}, σ_{1}) , (σ_{1}, σ_{2}) , ... $(\sigma_{k-1}, \sigma_{k})$ (inrespective of other relaxations that happen)

 $dc\sigma_{R}J = \delta(s,\sigma_{R})$

d[v] = S(s,v) for all $v \in V$, predecessor graph is the Shortest Path thee gooted at s.

```
BELLMAN-FORD ALGORITHM ( Negative Edge weights and even negative by a es)
                                                     Can actect
      INIT- SINGLE - SOURCE (G, S)
       for 1 = 1 to |V| - 1
              do for each edge (u,v) E E
                   do Relax (u, v, w)
                 if d[v] > d [u] + W(u,v) \ \
        for each edge (u,v) E E
                     then seturn FALSE
                                  graph has negative cycles
                 graph hus no negative aples => d[v] = S(s,v)
         selven TRUE
We could possibly be relaxing in the worst possible order
        e3, e2, e1
                      e,
```





Step3

O

7

Time: VE