

**Programme: B E – Computer Science and Engineering (AI&ML) &
Computer Science and Engineering (Cyber Security)
Internal Assessment – I**

TERM : 4 th OCT 2023 to 27 th JAN 2024	COURSE NAME: AUTOMATA THEORY AND COMPILER DESIGN
DATE : 01-12-2023 TIME: 11.00am-12.00pm	COURSE CODE : CY53/CI53
MAX MARKS: 30	PORTIONS : 2.5 UNITS



Mobile Phones are banned

Instructions to Candidates: **Answer any two full questions.**

Marks: 15x2=30

Q. NO	Questions	Blooms Levels (L1 to L6)*	CO	Marks
1.a	For the expression grammar $E \rightarrow E * F \mid F + E \mid F$ $F \rightarrow F - \mid id$ Obtain the following i. Leftmost derivation for the input string ' id-+id*id ' ii. Rightmost derivation for the input string ' id-+id*id ' iii. Is this Grammar ambiguous? If so, Eliminate the ambiguity.	L4	CO3	1+1 +2= 4
b	Consider the grammar given below. If not suitable for Predictive parser make necessary changes and construct predictive parsing table, M for the grammar given below. If LL(1), do parsing for the input " idid\$ " $S \rightarrow FR$ $R \rightarrow S \mid \epsilon$ $F \rightarrow id$	L3	CO2	5
c	Compute FIRST and FOLLOW of all Non-terminals for the grammar G. Show the traces of computation. i. $S \rightarrow ABC$ $A \rightarrow BS \mid b$ $B \rightarrow bS \mid CA \mid \epsilon$ $C \rightarrow \epsilon$ ii. $S \rightarrow ACB \mid cbB \mid Ba$ $A \rightarrow da \mid BC$ $B \rightarrow g \mid \epsilon$ $C \rightarrow h \mid \epsilon$	L3	CO4	6
2.a	i. Construct a transition diagram for her, she, he, him ii. Specify the input and output of each phase of a C compiler by translating the given assignment statement. (Assume all variables are of type int) val=sqrt(a*a+b*b);	L2, L3	CO1 , CO4	2+3 =5

B	Demonstrate how tokens are identified using “sentinels” technique. Identify the tokens and its appropriate lexemes for the code fragment. <pre> int fact(int n){ // computing factorial if(n==1) printf(“the factorial is 1”); if(n>1) return(n*fact(n-1)); } </pre>	L2	CO2	2+2 =4
C	i. What is an Augmented Grammar? Specify the reasons for the necessity of augmenting a grammar for LR Parsers. ii. Identify the initial item in the item set for LR parsers. iii. Construct LR(0) set of items for the following grammar G:- $S \rightarrow +SS \mid *SS \mid (S) \mid a$	L4	CO2	2+1 +3= 6
3.a	i. Describe the actions of a Shift Reduce Parser. ii. Show the moves made by shift reduce parser for the input string “bab\$”. Identify handles for the given input string. G: $S \rightarrow AaBb \mid BaAb$ $A \rightarrow b$ $B \rightarrow \epsilon$ iii. If $A \rightarrow \alpha$ is a production of a CFG and $ \alpha =n$. Then how many items can be generated using $A \rightarrow \alpha$.	L3 ,L4	CO2	1+3 +1= 5
B	Design a Predictive Parsing Table for the Grammar given below. Modify the Grammar if required. G: $S \rightarrow fES \mid wES \mid id$ $E \rightarrow E:T \mid E,T \mid \epsilon$ $T \rightarrow num$	L3	CO2	6
c	Write regular expressions for i. Accepting the Subject Code for semester 3,4,5,6 ii. Accepting Email id for gmail/yahoo	L4	CO1 , CO3	2+2 =4

* **L1** – Remember, **L2** – Understand, **L3**- Apply, **L4**- Analyze, **L5**-Evaluate, **L6**-Create

Course Outcomes (COs):

- Understand the core concepts in automata theory and Theory of Computation. (PO 1, 2, PSO 1,2)
- Design and develop lexical analyzers, parsers and code generators. (PO1,2, 3, 9, PSO 1,2)
- Design Grammars and Automata (recognizers) for different language classes and become knowledgeable about restricted models of Computation (Regular, Context Free) and their relative powers. (PO 1, 2,3,4, PSO 1,3)
- Gain the knowledge of the structure of a Compiler and Apply concepts automata theory and Theory of Computation to design Compilers. (PO 1,2,3,4,9, PSO 1, 2)

Sl.No.	SCHEME AND SOLITIONS	MARKS																																																																
1 a.	<p>$E \rightarrow E * F \mid F + E \mid F$ $F \rightarrow F - \mid id$</p> <p>Obtain the following</p> <p>i. Leftmost derivation for the input string ‘id-+id*id’ $E \rightarrow F + E \rightarrow F - + E \rightarrow id - + E \rightarrow id - + E * F \rightarrow id - + F * F \rightarrow id - + id * F \rightarrow id - + id * id$</p> <p>ii. Rightmost derivation for the input string ‘id-+id*id’ $E \rightarrow F + E \rightarrow F + E * F \rightarrow F + E * id \rightarrow F + F * id \rightarrow F + id * id \rightarrow F - + id * id \rightarrow id - + id * id$</p> <p>iii. Is this Grammar ambiguous? If so, Eliminate the ambiguity. Yes. $E \rightarrow FT \mid F + ET$ $T \rightarrow *FT \mid \epsilon$ $F \rightarrow idR$ $R \rightarrow -R \mid \epsilon$</p>	1M 1M 2M																																																																
b.	<p>$S \rightarrow FR$ $R \rightarrow S \mid \epsilon$ $F \rightarrow id$</p> <p>OR ELR</p> <p>$S \rightarrow FR$ $R \rightarrow FR \mid \epsilon$ $F \rightarrow id$</p> <p>The grammar is suitable for PP. -----</p> <table><tr><th>NT</th><th>FIRST</th><th>FOLLOW</th></tr><tr><td>S</td><td>{id}</td><td>{ \$ }</td></tr><tr><td>F</td><td>{id}</td><td>{id, \$ }</td></tr><tr><td>R</td><td>{id, ε}</td><td>{ \$ }</td></tr></table> <p>PARSE TABLE M</p> <table><tr><th>NT/INPUT SYMBOL</th><th>id</th><th>\$</th></tr><tr><td>S</td><td>$S \rightarrow FR$</td><td></td></tr><tr><td>F</td><td>$F \rightarrow id$</td><td></td></tr><tr><td>R</td><td>$R \rightarrow S$</td><td>$R \rightarrow \epsilon$</td></tr></table> <p>Parsing on input: idid</p> <table><tr><th>Matched string</th><th>Stack</th><th>Input</th><th>Action</th></tr><tr><td></td><td>S\$</td><td>idid\$</td><td></td></tr><tr><td></td><td>FR\$</td><td>idid\$</td><td>Reduce $S \rightarrow FR$</td></tr><tr><td></td><td>idR\$</td><td>idid\$</td><td>Reduce $F \rightarrow id$</td></tr><tr><td>id</td><td>R\$</td><td>id\$</td><td>Matched id</td></tr><tr><td></td><td>S\$</td><td>id\$</td><td>Reduce $R \rightarrow S$</td></tr><tr><td></td><td>FR\$</td><td>id\$</td><td>Reduce $S \rightarrow FR$</td></tr><tr><td></td><td>idR\$</td><td>id\$</td><td>Reduce $F \rightarrow id$</td></tr><tr><td>idid</td><td>R\$</td><td>\$</td><td>Reduce $R \rightarrow \epsilon$</td></tr><tr><td>idid</td><td>\$</td><td>\$</td><td>Accept</td></tr></table>	NT	FIRST	FOLLOW	S	{id}	{ \$ }	F	{id}	{id, \$ }	R	{id, ε}	{ \$ }	NT/INPUT SYMBOL	id	\$	S	$S \rightarrow FR$		F	$F \rightarrow id$		R	$R \rightarrow S$	$R \rightarrow \epsilon$	Matched string	Stack	Input	Action		S\$	idid\$			FR\$	idid\$	Reduce $S \rightarrow FR$		idR\$	idid\$	Reduce $F \rightarrow id$	id	R\$	id\$	Matched id		S\$	id\$	Reduce $R \rightarrow S$		FR\$	id\$	Reduce $S \rightarrow FR$		idR\$	id\$	Reduce $F \rightarrow id$	idid	R\$	\$	Reduce $R \rightarrow \epsilon$	idid	\$	\$	Accept	1M 2M 2M
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NT	FIRST	FOLLOW
S	{b}	{\$,b}
A	{b}	{b,\$}
B	{b, ε }	{b,\$}
C	{ε }	{b,\$}

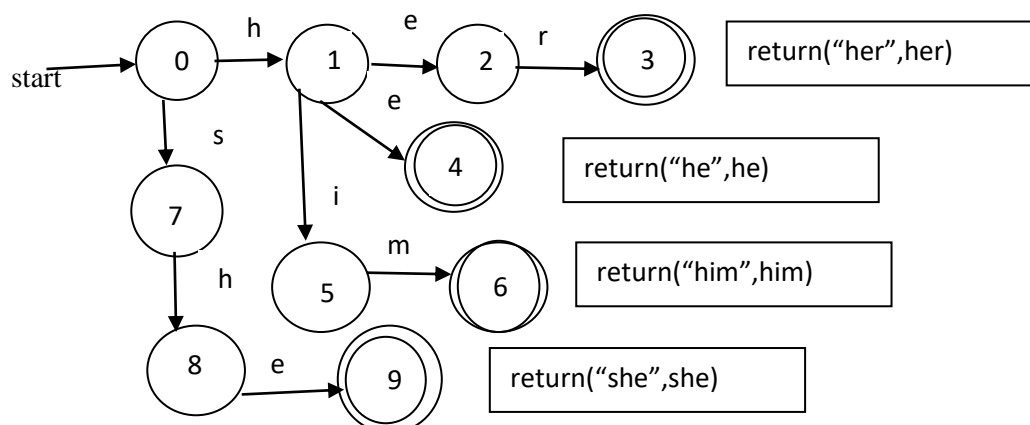
- ii. $S \rightarrow ACB \mid cbB \mid Ba$
 $A \rightarrow da \mid BC$
 $B \rightarrow g \mid \epsilon$
 $C \rightarrow h \mid \epsilon$

NT	FIRST	FOLLOW
S	{a,c,g,d,h, ε }	{\$}
A	{d,g,h, ε }	{h,g,\$}
B	{g, ε }	{a,h,g,\$}
C	{h, ε }	{h,g,\$}

1.5+1.5=3M

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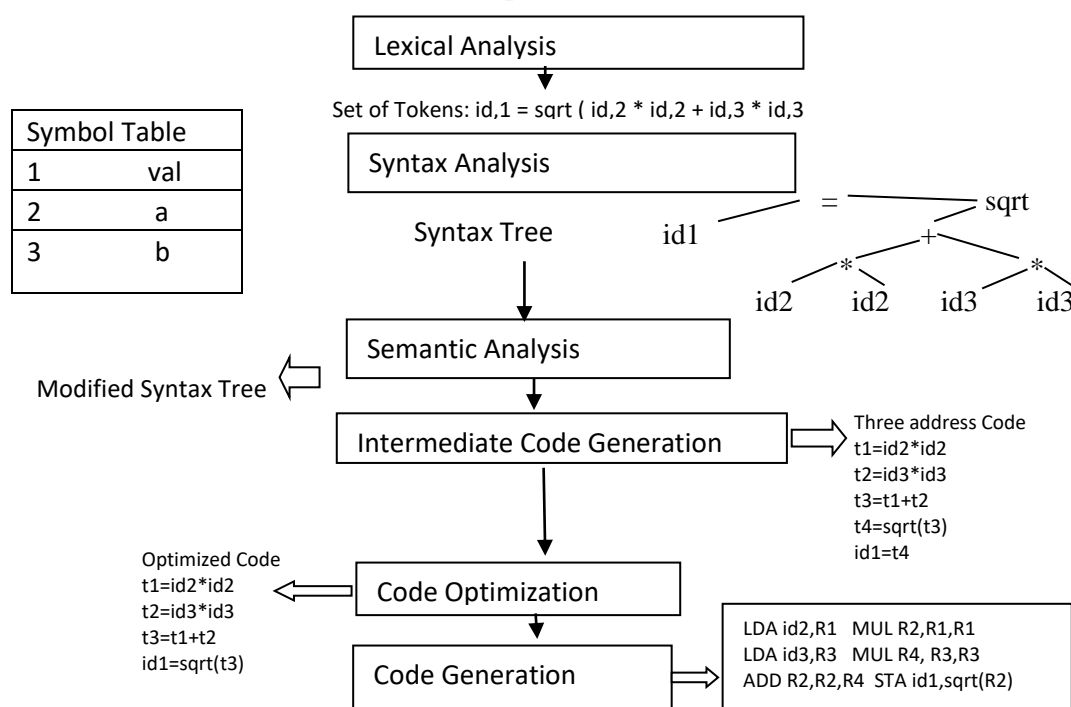
- 2.a. i. Construct a transition diagram for **her, she, he, him**



2M

- ii.

val=sqrt(a*a+b*b);



3M
Each other than code optimization 1M

Symbol Table 1M

b.	<p>Sentinels: //eof char should be the sentinel char. Lexeme Begin and Forward pointer-----</p> <table><tr><td></td><td></td><td>a</td><td>+</td><td>eof</td></tr></table> <table><tr><td>b</td><td>+</td><td>c</td><td>eof</td><td></td><td>eof</td></tr></table> <table><tr><th>Lexemes</th><th>Tokens</th></tr><tr><td>Chars i,n,t</td><td>int</td></tr><tr><td>Chars f,a,c,t</td><td>id,1</td></tr><tr><td>Char (</td><td>(</td></tr><tr><td>Char n</td><td>id,2</td></tr><tr><td>Char)</td><td>)</td></tr><tr><td>Char {</td><td>{</td></tr><tr><td>Char i,f</td><td>if</td></tr><tr><td>Char ==</td><td>==</td></tr><tr><td>1</td><td>1</td></tr><tr><td>Char p,r,i,n,t,f</td><td>printf</td></tr><tr><td>“the factorial is 1”</td><td>literal</td></tr><tr><td>Char ></td><td>></td></tr><tr><td>Char r,e,t,u,r,n</td><td>return</td></tr><tr><td>Char *</td><td>*</td></tr><tr><td>Char -</td><td>-</td></tr></table>			a	+	eof	b	+	c	eof		eof	Lexemes	Tokens	Chars i,n,t	int	Chars f,a,c,t	id,1	Char ((Char n	id,2	Char))	Char {	{	Char i,f	if	Char ==	==	1	1	Char p,r,i,n,t,f	printf	“the factorial is 1”	literal	Char >	>	Char r,e,t,u,r,n	return	Char *	*	Char -	-	1M 1M 2M
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c.	<p>i. What is an Augmented Grammar? Specify the reasons for the necessity of augmenting a grammar for LR Parsers.</p> <p>For Bottom Up Parsing, LR Parsers will stop parsing only after getting start symbol as handle (on top of stack). For specifying start symbol as handle</p> <p>ii. Identify the initial item in the item set for LR parsers.</p> <p>Initial Item: $S' \rightarrow .S$</p> <p>iii. Construct LR(0) set of items for the following grammar</p> <p>G:- $S \rightarrow +SS \mid *SS \mid (S) \mid a$</p> <div><div><p>I0:</p><p>$S' \rightarrow .S$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I1:</p><p>$S' \rightarrow S.$</p></div><div><p>I2:</p><p>$S \rightarrow +.SS$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I3:</p><p>$S \rightarrow *.SS$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I4:</p><p>$S \rightarrow (.S)$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I5:</p><p>$S \rightarrow a.$</p></div><div><p>I6:</p><p>$S \rightarrow +S.S$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I7:</p><p>$S \rightarrow *S.S$</p><p>$S \rightarrow .+SS$</p><p>$S \rightarrow .*SS$</p><p>$S \rightarrow .(S)$</p><p>$S \rightarrow .a$</p></div><div><p>I8:</p><p>$S \rightarrow (S.)$</p></div><div><p>I9:</p><p>$S \rightarrow +SS.$</p></div><div><p>I10:</p><p>$S \rightarrow *SS.$</p></div><div><p>I11:</p><p>$S \rightarrow (S).$</p></div></div> <div><div><p>I2,+=I2</p><p>I2,*=I3</p><p>I2,(=I4</p><p>I2,a=I5</p><p>I3,+=I2</p><p>I3,*=I3</p></div><div><p>I4,(=I4</p><p>I6,+=I2</p><p>I6,*=I3</p><p>I6,(=I4</p><p>I6,a=I5</p><p>I3,(=I4</p></div><div><p>I7,+=I2</p><p>I7,a=I5</p></div></div>	2M 1M 3M																																											

3 a.	<div><div><div>i. Describe the actions of a Shift Reduce Parser. Actions of a Shift reduce parser: Shift, Reduce, Accept, Error</div><div>ii. Show the moves made by shift reduce parser for the input string “bab\$”. Identify handles for the given input string. G: $S \rightarrow AaBb \mid BaAb$ $A \rightarrow b$ $B \rightarrow \epsilon$</div></div><table><tr><th>Stack</th><th>Input</th><th>Handle</th><th>Action</th></tr><tr><td>\$</td><td>bab\$</td><td></td><td></td></tr><tr><td>\$b</td><td>ab\$</td><td>b</td><td>shift</td></tr><tr><td>\$A</td><td>ab\$</td><td></td><td>Reduce $A \rightarrow b$</td></tr><tr><td>\$Aa</td><td>b\$</td><td></td><td>shift</td></tr><tr><td>\$AaB</td><td>b\$</td><td>ϵ</td><td>Reduce $B \rightarrow \epsilon$</td></tr><tr><td>\$AaBb</td><td>\$</td><td></td><td>shift</td></tr><tr><td>\$S</td><td>\$</td><td>AaBb</td><td>Reduce $S \rightarrow AaBb$</td></tr><tr><td></td><td></td><td></td><td>Accept</td></tr></table><div>iii. If $A \rightarrow \alpha$ is a production of a CFG and $\alpha =n$. Then how many items can be generated using $A \rightarrow \alpha$. Answer: $n+1$</div></div> <div>1M <</div>	Stack	Input	Handle	Action	\$	bab\$			\$b	ab\$	b	shift	\$A	ab\$		Reduce $A \rightarrow b$	\$Aa	b\$		shift	\$AaB	b\$	ϵ	Reduce $B \rightarrow \epsilon$	\$AaBb	\$		shift	\$S	\$	AaBb	Reduce $S \rightarrow AaBb$				Accept
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