# **EXPERIMENT - 12**

# **Advanced Transaction Simulation: Deadlocks, MVCC, and Concurrency**

This document demonstrates advanced transaction scenarios in MySQL using InnoDB: Deadlock simulation, MVCC-based concurrency, and comparison of locking vs MVCC. Use two concurrent sessions (Session A and Session B) in Nimbus or MySQL CLI to reproduce behavior.

## Setup (Run Once)

DROP TABLE IF EXISTS StudentEnrollments;

```
CREATE TABLE StudentEnrollments (
student_id INT NOT NULL,
student_name VARCHAR(100) NOT NULL,
course_id VARCHAR(10) NOT NULL,
enrollment_date DATE NOT NULL,
PRIMARY KEY (student_id)
) ENGINE=InnoDB;
```

INSERT INTO StudentEnrollments (student\_id, student\_name, course\_id, enrollment\_date) VALUES

```
(1, 'Ashish', 'CSE101', '2024-06-01'), (2, 'Smaran', 'CSE102', '2024-06-01'), (3, 'Vaibhav', 'CSE103', '2024-06-01');
```

## **Part A: Simulating a Deadlock Between Two Transactions**

Two concurrent transactions attempt to update overlapping rows in reverse order. MySQL detects a deadlock and rolls back one transaction (Error 1213).

## **Session A**

```
START TRANSACTION;

UPDATE StudentEnrollments

SET enrollment_date = '2024-07-01'

WHERE student_id = 1;

SELECT SLEEP(5);

UPDATE StudentEnrollments
```

```
SET enrollment_date = '2024-07-02'
WHERE student_id = 2;
COMMIT;
```

#### **Session B**

```
START TRANSACTION;

UPDATE StudentEnrollments

SET enrollment_date = '2024-08-01'

WHERE student_id = 2;

SELECT SLEEP(5);

UPDATE StudentEnrollments

SET enrollment_date = '2024-08-02'

WHERE student_id = 1;

COMMIT;
```

Expected Output: One transaction fails with ERROR 1213 (Deadlock found...). The database automatically rolls back the victim transaction.

② To avoid deadlocks: always lock rows in the same order, keep transactions short, and retry on deadlock errors.

## Part B: Applying MVCC to Prevent Conflicts During Concurrent Reads/Writes

MVCC allows non-blocking reads under Repeatable Read isolation. Session A reads a snapshot while Session B updates concurrently.

#### Session A (Reader)

```
SET SESSION TRANSACTION ISOLATION LEVEL REPEATABLE READ;
START TRANSACTION;
SELECT enrollment_date FROM StudentEnrollments WHERE student_id = 1;
-- Expect: 2024-06-01
SELECT SLEEP(5);
SELECT enrollment_date FROM StudentEnrollments WHERE student_id = 1;
COMMIT;
```

## Session B (Writer)

```
START TRANSACTION;

UPDATE StudentEnrollments

SET enrollment_date = '2024-07-10'

WHERE student_id = 1;
```

#### COMMIT;

Expected Output: Session A sees old snapshot (2024-06-01) until commit, while Session B successfully updates to 2024-07-10.

## Part C: Comparing Behavior With and Without MVCC

This demonstrates blocking behavior with explicit locking (SELECT FOR UPDATE) vs non-blocking reads with MVCC.

### Scenario 1: With Locking

#### Session A:

START TRANSACTION:

SELECT \* FROM StudentEnrollments WHERE student\_id = 1 FOR UPDATE;

SELECT SLEEP(6);

UPDATE StudentEnrollments SET enrollment\_date = '2024-09-10' WHERE student\_id = 1; COMMIT;

#### Session B:

START TRANSACTION;

UPDATE StudentEnrollments SET enrollment\_date = '2024-09-20' WHERE student\_id = 1; COMMIT:

-- Session B blocks until A commits

## Scenario 2: With MVCC (Non-blocking Reads)

#### Session A:

SET SESSION TRANSACTION ISOLATION LEVEL REPEATABLE READ;

START TRANSACTION;

SELECT enrollment\_date FROM StudentEnrollments WHERE student\_id = 1;

SELECT SLEEP(6);

SELECT enrollment\_date FROM StudentEnrollments WHERE student\_id = 1; COMMIT;

#### Session B:

START TRANSACTION;

UPDATE StudentEnrollments SET enrollment\_date = '2024-10-01' WHERE student\_id = 1; COMMIT;

## **Expected Output:**

- With Locking: readers/writers block each other.
- With MVCC: readers continue using snapshot; writers commit new versions.

# **Sample Output Image**

# Part A: Insert Multiple Fee Payments in a Transaction

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## START TRANSACTION Intennne

payment_id	student_name	arriers 11	amoı	payment.Jate
1	Ashish	5000.00	0 50	2024-06-01
2	Smaran	4500.00	030	2024-06-02
3	Valbhav	5500.00	083	2024-06-03

COMMIT

# Part B: Demonstrate ROLLBACK for Failed Payment Insertion

#### START TRANSACTION

payment_id	student_name	amount	amn	payment_date
4	Kiran	4000.00	400	2024-06-04
5	Smdran	-100.00	-100	2024-06-05

ROLLBACK

## Part C: Simulate Partial Failure and Ensure Consistent State

## START TRANSACTION

payment_id	student_name	amount	amn	payment_date
5	Nidhi	3000.00	000	2024-06-08
6	Smaran	2500.00	000	2024-06-07
3	Valbhav	5500.00	000	2024-06-07

ROLLBACK

# Part D: Verify ACID Compliance with Transaction Flow

#### START TRANSACTION

REGET TRANSACTION SELE		Inalict tra	Inalict transactions for 22 024 -06-05			
payment_id	Vaibhav	5500	5500.00	22	354-0085	2024-06-03
payment_id	Smaran	4600	4500.00	MΩ	254-0035	2024-06-02
payment_id	Ashish	5000	5000.00	FL	233-0765	2024-06-01