



CSE408

Knapsack problem

Lecture # 35

Dynamic Programming



Dynamic Programming is a general algorithm design technique for solving problems defined by or formulated as recurrences with overlapping sub instances

- Invented by American mathematician Richard Bellman in the 1950s to solve optimization problems and later assimilated by CS
- “Programming” here means “planning”
- Main idea:
 - set up a recurrence relating a solution to a larger instance to solutions of some smaller instances
 - solve smaller instances once
 - record solutions in a table
 - extract solution to the initial instance from that table

Knapsack Problem by DP



Given n items of

integer weights: $w_1 \ w_2 \ \dots \ w_n$

values: $v_1 \ v_2 \ \dots \ v_n$

a knapsack of integer capacity W

find most valuable subset of the items that fit into the knapsack

Consider instance defined by first i items and capacity j ($j \leq W$).

Let $V[i,j]$ be optimal value of such an instance. Then

$$V[i,j] = \begin{cases} \max \{V[i-1,j], v_i + V[i-1,j-w_i]\} & \text{if } j-w_i \geq 0 \\ V[i-1,j] & \text{if } j-w_i < 0 \end{cases}$$

Initial conditions: $V[0,j] = 0$ and $V[i,0] = 0$

Knapsack Problem by DP (example)



Example: Knapsack of capacity $W = 5$

item	weight	value
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1	2	\$12
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2	1	\$10
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3	3	\$20
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4	2	\$15
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capacity j

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	0	12	12	12
2	0	10	12	22	22	22
3	0	10	12	22	30	32
4	0	10	15	25	30	37

$w_1 = 2, v_1 = 12$ 1

$w_2 = 1, v_2 = 10$ 2

$w_3 = 3, v_3 = 20$ 3

$w_4 = 2, v_4 = 15$ 4

Backtracing
finds the actual
optimal subset,
i.e. solution.

Knapsack Problem by DP (pseudocode)



Algorithm DPKnapsack($w[1..n]$, $v[1..n]$, W)

var $V[0..n, 0..W]$, $P[1..n, 1..W]$: int

for $j := 0$ to W do

$V[0, j] := 0$

for $i := 0$ to n do

$V[i, 0] := 0$

for $i := 1$ to n do

for $j := 1$ to W do

if $w[i] \leq j$ and $v[i] + V[i-1, j-w[i]] > V[i-1, j]$ then

$V[i, j] := v[i] + V[i-1, j-w[i]]$; $P[i, j] := j-w[i]$

else

$V[i, j] := V[i-1, j]$; $P[i, j] := j$

return $V[n, W]$ and the optimal subset by backtracing

Running time and space:
 $O(nW)$.



Thank You !!!