

CSE322 ALGEBRIC METHODS USING ARDEN'S THEORM

Lecture #8

ALGEBRIC METHOD ISUNG ARDEN'S THEORM



The following method is an extension of the Arden's theorem (Theorem 5.1). This is used to find the r.e. recognized by a transition system.

The following assumptions are made regarding the transition system:

- (i) The transition graph does not have Λ -moves.
- (ii) It has only one initial state, say v_1 .
- (iii) Its vertices are $v_1 \dots v_n$.
- (iv) V_i the r.e. represents the set of strings accepted by the system even though v_i is a final state.
- (v) α_{ij} denotes the r.e. representing the set of labels of edges from v_i to v_j . When there is no such edge, $\alpha_{ij} = \emptyset$. Consequently, we can get the following set of equations in $V_1 \dots V_n$:

$$\mathbf{V}_{1} = \mathbf{V}_{1}\boldsymbol{\alpha}_{11} + \mathbf{V}_{2}\boldsymbol{\alpha}_{21} + \cdots + \mathbf{V}_{n}\boldsymbol{\alpha}_{n1} + \Lambda$$

$$\mathbf{V}_{2} = \mathbf{V}_{1}\boldsymbol{\alpha}_{12} + \mathbf{V}_{2}\boldsymbol{\alpha}_{22} + \cdots + \mathbf{V}_{n}\boldsymbol{\alpha}_{n2}$$

$$\vdots$$

$$\mathbf{V}_{n} = \mathbf{V}_{1}\boldsymbol{\alpha}_{1n} + \mathbf{V}_{2}\boldsymbol{\alpha}_{2n} + \cdots + \mathbf{V}_{n}\boldsymbol{\alpha}_{nn}$$

By repeatedly applying substitutions and Theorem 5.1 (Arden's theorem), we can express V_i in terms of α_{ii} 's.

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Consider the transition system given in Fig. 5.13. Prove that the strings recognized are $(\mathbf{a} + \mathbf{a}(\mathbf{b} + \mathbf{a}\mathbf{a})^*\mathbf{b})^* \mathbf{a}(\mathbf{b} + \mathbf{a}\mathbf{a})^*\mathbf{a}$.

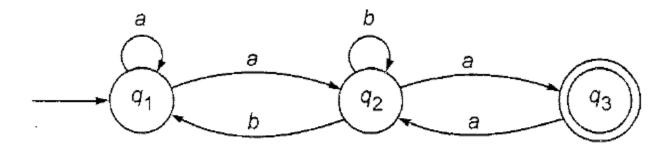


Fig. 5.13

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The three equations for q_1 , q_2 and q_3 can be written as

$$q_1 = q_1a + q_2b + \Lambda,$$
 $q_2 = q_1a + q_2b + q_3a,$ $q_3 = q_2a$

It is necessary to reduce the number of unknowns by repeated substitution. By substituting q_3 in the q_2 -equation, we get by applying Theorem 5.1

$$q_2 = q_1 a + q_2 b + q_2 a a$$

= $q_1 a + q_2 (b + a a)$
= $q_1 a (b + a a)^*$

Substituting \mathbf{q}_2 in \mathbf{q}_1 , we get

$$\mathbf{q}_{l} = \mathbf{q}_{l}\mathbf{a} + \mathbf{q}_{l}\mathbf{a}(\mathbf{b} + \mathbf{a}\mathbf{a})*\mathbf{b} + \Lambda$$
$$= \mathbf{q}_{l}(\mathbf{a} + \mathbf{a}(\mathbf{b} + \mathbf{a}\mathbf{a})*\mathbf{b}) + \Lambda$$

Hence,

$$q_1 = \Lambda(a + a(b + aa)*b)*$$
 $q_2 = (a + a(b + aa)*b)* a(b + aa)*$
 $q_3 = (a + a(b + aa)*b)* a(b + aa)*a$

Since q_3 is a final state, the set of strings recognized by the graph is given by

$$(a + a(b + aa)*b)*a(b + aa)*a$$