

UNINFORMED SEARCH

-BFS

-DFS

INT404

• **State-space search** is the process of searching through a state space for a solution by making explicit a sufficient portion of an implicit state-space graph to include a goal node.

- Hence, initially $V=\{S\}$, where S is the start node;

- when S is expanded, its successors are generated and those nodes are added to V and the associated arcs are added to E .

- This process continues until a goal node is generated (included in V) and identified (by goal test)

• During search, a node can be in one of the three categories:

- Not generated yet (has not been made explicit yet)

- OPEN**: generated but not expanded

- CLOSED**: expanded

- Search strategies differ mainly on how to select an OPEN node for expansion at each step of search

A General State-Space Search Algorithm

```
open := {S}; closed := { };  
  
repeat  
    n := select(open);           /* select one node from open for expansion */  
    if n is a goal  
        then exit with success; /* delayed goal testing */  
    expand(n)  
        /* generate all children of n  
        put these newly generated nodes in open (check duplicates)  
        put n in closed (check duplicates) */  
  
until open = { };  
exit with failure
```

Some Issues

- Search process constructs a search tree, where
 - **root** is the initial state S , and
 - **leaf nodes** are nodes
 - not yet been expanded (i.e., they are in OPEN list) or
 - having no successors (i.e., they're "deadends")
- Some important issue that arises
 - The direction in which conduct the search(forward vs. backward reasoning)
 - How to select applicable rules(matching).
 - How to represent each node of search process(the knowledge representation problem)
 - Search tree vs. search graph

Evaluating Search Strategies

- **Completeness**

- Guarantees finding a solution whenever one exists

- **Time Complexity**

- How long (worst or average case) does it take to find a solution? Usually measured in terms of the **number of nodes expanded**

- **Space Complexity**

- How much space is used by the algorithm? Usually measured in terms of the **maximum size that the “OPEN” list becomes** during the search

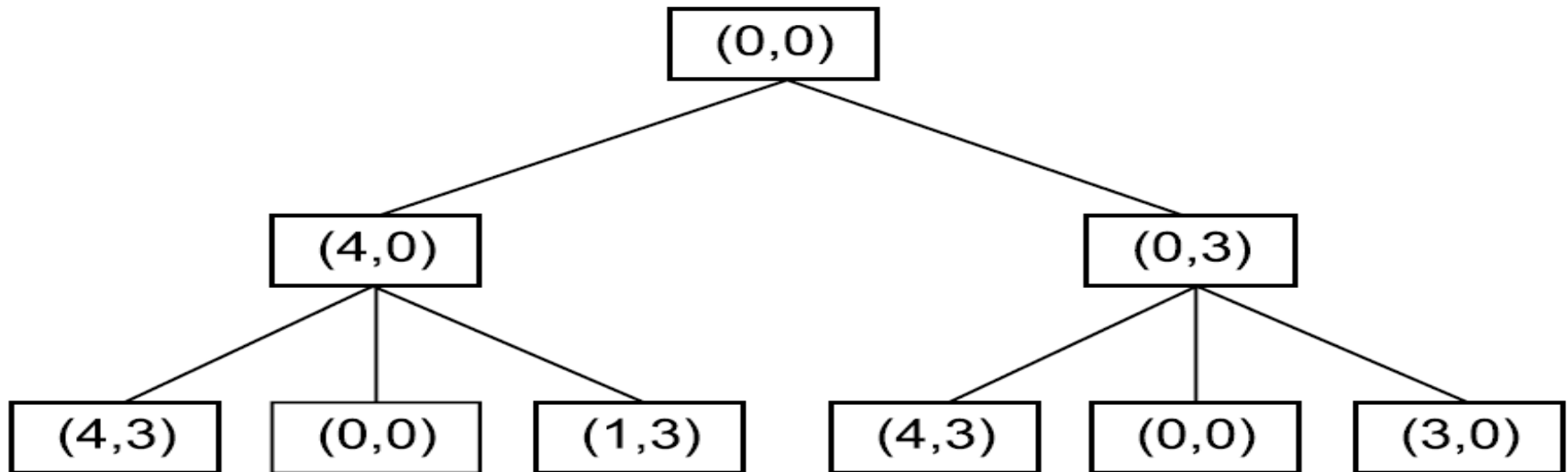
- **Optimality/Admissibility**

- If a solution is found, is it guaranteed to be an optimal one? For example, is it the one with minimum cost?

Algorithm : Breadth-First Search

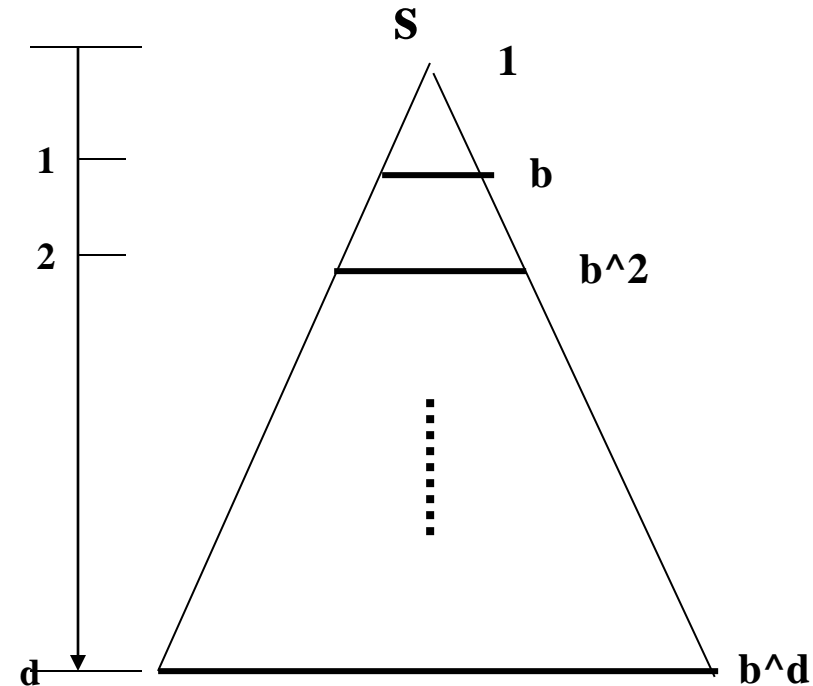
1. Create a variable called *NODE-LIST* and set it to the initial state.
2. Until a goal state is found or *NODE-LIST* is empty:
 - (a) Remove the first element from *NODE-LIST* and call it *E*. If *NODE-LIST* was empty, quit.
 - (b) For each way that each rule can match the state described in *E* do:
 - (i) Apply the rule to generate a new state,**
 - (ii) If the new state is a goal state, quit and return this state.**
 - (iii) Otherwise, add the new state to the end of NODE-LIST.**

Two Levels of a Breadth-First Search Tree



Breadth-First

- A complete search tree of depth d where each non-leaf node has b children, has a total of $1 + b + b^2 + \dots + b^d = (b^{d+1} - 1)/(b - 1)$ nodes
- Time complexity (# of nodes generated): $O(b^d)$
- Space complexity (maximum length of OPEN): $O(b^d)$



- For a complete search tree of depth 12, where every node at depths 0, ..., 11 has 10 children and every node at depth 12 has 0 children, there are $1 + 10 + 100 + 1000 + \dots + 10^{12} = (10^{13} - 1)/9 = O(10^{12})$ nodes in the complete search tree.
 - BFS is suitable for problems with shallow solutions



Host



Searcher



Host

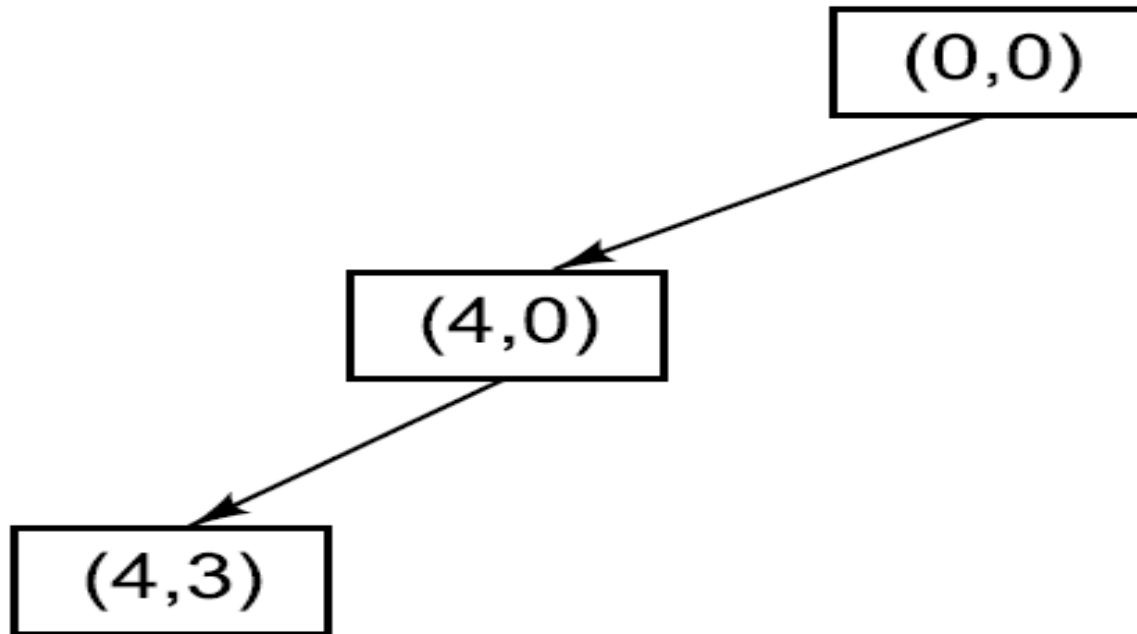


Host

Algorithm : Depth-First Search

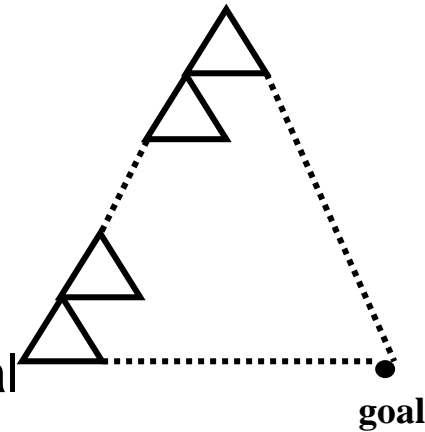
1. If the initial state is a goal state, quit and return success.
2. Otherwise, do the following until success or failure is signaled:
 - (a) **Generate a successor, E , of the initial state. If there are no more successors, signal failure.**
 - (b) **Call Depth-First Search with E as the initial state.**
 - (c) **If success is returned, signal success. Otherwise continue in this loop.**

A Depth-First Search Tree



Depth-First (DFS)

- Algorithm outline:
 - Always select from the OPEN the node with the greatest depth for expansion, and put all newly generated nodes into OPEN
 - OPEN is organized as **LIFO** (last-in, first-out) list.
 - Terminate if a node selected for expansion is a goal
- **May not terminate** without a "depth bound," i.e., cutting off search below a fixed depth D (How to determine the depth bound?)
- **Not complete** (with or without cycle detection, and with or without a cutoff depth)
 - **Exponential time**, $O(b^d)$, but only **linear space**, $O(bd)$, required
 - Can find **deep solutions quickly** if lucky
- When search hits a deadend, can only back up one level at a time even if the "problem" occurs because of a bad operator choice near the top of the tree. Hence, only does "chronological backtracking"



Thank You !!!