

Data Structures & Algorithms

Topic: Heap



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Introduction

- **Heap (MaxHeap):** A complete binary tree H with n elements is called a Heap if each node N of H has the following property:

“The value at N is greater than or equal to the value at each of the children of N .”

- If the value at N is less than or equal to the value at each of the children of N , then it is called MinHeap.
- Heaps are maintained in memory by using linear array TREE.



Insertion in a Heap

- INSERT_HEAP(TREE, N, ITEM)

1. Set $N = N+1$, and $PTR = N$.
2. Repeat Step 3 to 6 while $PTR > 1$
3. Set $PAR = \lfloor PTR / 2 \rfloor$
4. If $ITEM \leq TREE[PAR]$, then:
Set $TREE[PTR] = ITEM$ and Return.
5. Else: Set $TREE[PTR] = TREE[PAR]$.
6. Set $PTR = PAR$.
7. Set $TREE[1] = ITEM$.
8. Return.



Example

- Create a Heap from the following list of numbers:
40, 30, 50, 20, 60, 55, 70, 60, 65, 50



Deletion of Root of a Heap

- Assign Root R to some variable ITEM.
- Replace the deleted node R by the last node L of Heap H so that H is still complete tree, **but not a Heap**.
- Call MaxHeapify to maintain the heap property.



Build MaxHeap

BUILD_MAX-HEAP(A)

1. $\text{heapsize}[A] = \text{length}[A]$
2. Repeat for $i = \lfloor \text{length}[A]/2 \rfloor$ to 1
3. Call **MAX_HEAPIFY(A, i)**
4. Exit



Maintaining Heap Property

MAX_HEAPIFY(A, i)

1. Set: $l = \text{LEFT}(i)$
2. Set: $r = \text{RIGHT}(i)$
3. If $l \leq \text{heapsize}[A]$ and $A[l] > A[i]$, then:
4. $\text{largest} = l$.
5. Else: $\text{largest} = i$.
6. If $r \leq \text{heapsize}[A]$ and $A[r] > A[\text{largest}]$, then:
7. $\text{largest} = r$.
8. If $\text{largest} \neq i$, then:
9. Exchange $A[i] \longleftrightarrow A[\text{largest}]$.
10. MAX_HEAPIFY (A, largest).



Heap Sort

HEAP_SORT (A)

1. BUILD_MAXHEAP (A)
2. Repeat for $i = \text{length}[A]$ to 2
3. Exchange $A[1] \longleftrightarrow A[i]$.
4. Heapsize[A] = Heapsize [A] – 1.
5. Call MAX_HEAPIFY(A, 1).
6. Exit.



Complexity of Heap Sort

Average and Worst case Complexity of Heap sort = $O(n \log n)$.



Questions



Review Questions

- What do you mean by Heap?
- What is the complexity of Heap sort?
- Calculate the complexity of Heap sort.
- How will you insert an element in a heap?
- What is the difference between Maxheap and Minheap?