

Experiment 7

Ques 1) Topological Sort of a directed acyclic graph (DAG)

Steps 1) Create the graph by calling `addEdge(a,b)`

2) Call the `topologicalSort()`

(a) Create a stack and a boolean array named as `visited[]`
 (b) Mark all the vertices as not visited i.e., initialize `visited[]` with 'false' value.

(c) Call the recursive helper function `topologicalSortUtil()` to store Topological Sort starting from all vertices one by one.

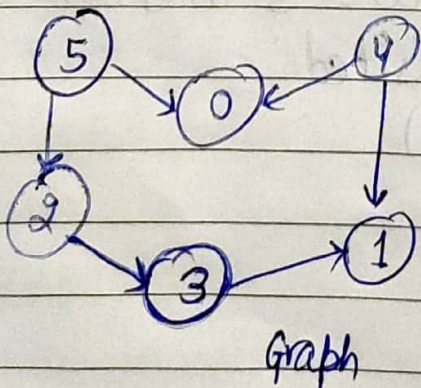
3) def `topologicalSortUtil(int v, bool visited[], stack<int> &Stack)`

a) Mark the current node as visited.

b) Recur for all the vertices adjacent to this vertex.

c) Push current vertex to stack which stores result.

4) At last after return from the utility function, print contents of stack.



1) `visited[] =`

F	F	F	F	F	F
0	1	2	3	4	5

2) $i=0$

T	F	F	F	F	F
0	1	2	3	4	5

Similarly for other values of i

3)

								5
						4		4
					2	2		2
	→		→	3	→	3	→	3
		1		1		1		1
0		0		0		0		0

4) Content of stack is pop out one by one and output is
5 4 2 3 1 0

Complexity :- Time Complexity

- a) Best $\Theta(|V| + |E|)$
- b) Average $\Theta(|V| + |E|)$
- c) Worst $\Theta(|V| + |E|)$

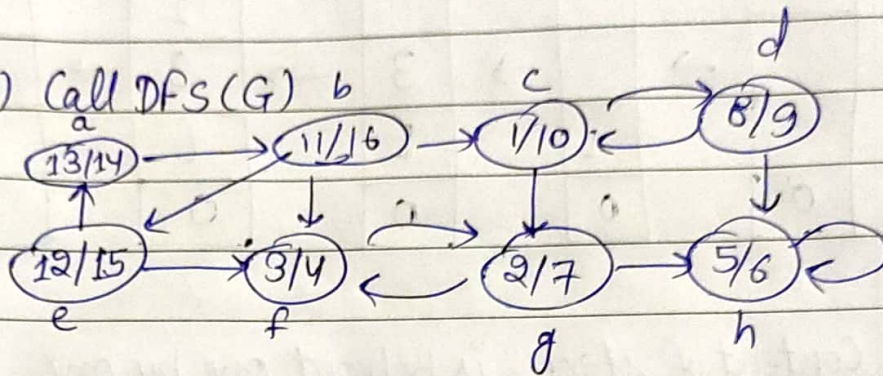
Space Complexity $\Theta(|V|)$

Ques 2) Compute the strongly connected components of a directed graph $G = (V, E)$ using two DFSs one on G and one on G^T .

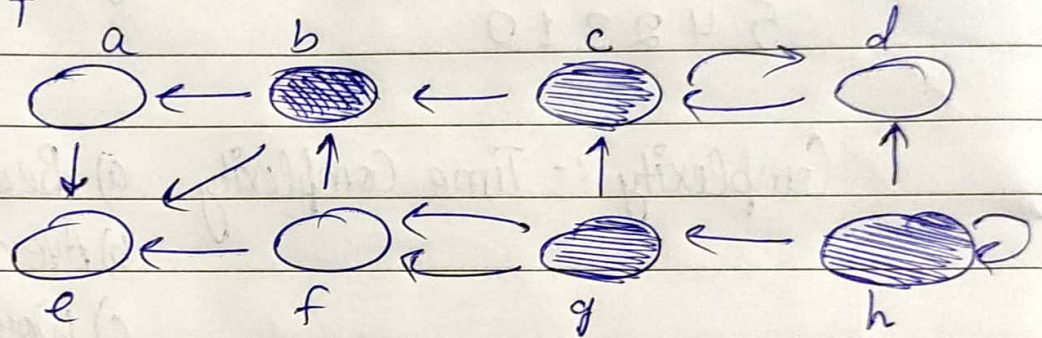
- Steps
- 1) Call DFS(G) to compute finishing times $f[u]$ for all u .
 - 2) Compute G^T
 - 3) Call DFS(G^T) but in the main loop, consider vertices in order of decreasing $f[u]$
 - 4) Output the vertices in each tree of the depth first forest formed in second DFS as a separate SCC.

Time taken linear time i.e., $\Theta(V+E)$ to compute SCC of a digraph G .

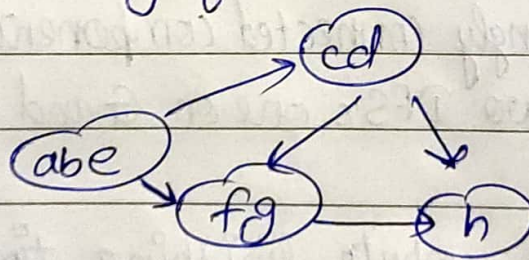
Example :- (1) Call DFS(G)



2) Compute G^T



3) Call DFS(G^T) but this time consider the vertices in order to decreasing finish time.



4) Output :- $\{a, b, e\}, \{c, d\}, \{f, g\}, \{h\}$