

Experiment 9.

Ans1 Dijkstra Algorithm :- is a solution to the single source shortest path problem in graph theory.

Works on both directed and undirected graphs. However all edges must have non-negative weights.

Approach :- Greedy

Input :- Weighted graph $G = \{E, V\}$ and source vertex $v \in V$ such that all edge weights are non-negative

Output :- lengths of shortest paths from a given source vertex $v \in V$ to all other vertices.

Pseudocode :-

Dijkstra(G, w, s)

InitializeSingleSource(G, s)

$S = \phi$

$Q = G.V$

while $Q \neq \phi$

$u = \text{ExtractMin}(Q)$

$S = S \cup \{u\}$

for each vertex $v \in G.Adj[u]$

Relax(u, v, w)

Initialize Single Source (G, s)

for each vertex $v \in G.V$

$v.d = \infty$

$v.JT = \text{NIL}$

$s.d = 0$

Relax(u, v, w)

if $v.d > u.d + w(u, v)$

$v.d = u.d + w(u, v)$

$v.JT = u$

Bellmanford Algorithm helps us find the shortest path from a vertex to all other vertices of a weighted graph.

It is similar to Dijkstra's algorithm but it can work with graphs in which edges can have negative weights.

Pseudocode

function bellmanford(G, s)

for each vertex V in G

distance[V] $\leftarrow \infty$

previous[V] $\leftarrow \text{NULL}$

distance[s] $\leftarrow 0$

for each vertex V in G

for each edge (u, v) in G

tempDistance \leftarrow distance[u] + edge-weight(u, v)

if tempDistance < distance[v]

$\text{distance}[v] \leftarrow \text{tempDistance}$

$\text{previous}[v] \leftarrow u$

for each edge (u, v) in G

if $\text{distance}[u] + \text{edgeWeight}(u, v) < \text{distance}[v]$

Error: Negative Cycle Exists

return $\text{distance}[]$, $\text{previous}[]$

Time Complexity:- Best Case Complexity $O(E)$
 Av. Case Complexity $O(VE)$
 Worst Case complexity $O(VE)$

Space Complexity $O(V)$