

## Experiment 4

1) Heap Sort : Not stable, but can be made stable  
 $T(n) = O(n \log n)$

Heapsort(A) {

    BuildHeap(A)

    for  $i \leftarrow \text{length}(A)$  downto 2 {

        exchange  $A[1] \leftrightarrow A[i]$

        heapsize  $\leftarrow$  heapsize - 1

        Heapify(A, 1)

    }

BuildHeap(A) {

    heapsize  $\leftarrow \text{length}(A)$

    for  $i \leftarrow \text{floor}(\text{length}/2)$  downto 1

        Heapify(A, i)

    }

Heapify(A, i) {

    lef  $\leftarrow \text{left}(i)$

    righ  $\leftarrow \text{right}(i)$

    if ( $\text{lef} \leq \text{heapsize}$ ) and ( $A[\text{lef}] > A[i]$ )

        largest  $\leftarrow \text{lef}$

    else

        largest  $\leftarrow i$

    if ( $\text{righ} \leq \text{heapsize}$ ) and ( $A[\text{righ}] > A[\text{largest}]$ )

        largest  $\leftarrow \text{righ}$



```
    if (largest != i) {  
        exchange A[i] ↔ A[largest]  
        Heapify (A, largest)  
    }
```

```
    }
```

```
2) Quicksort (A, p, r) {  
    if (p < r)  
        q ← Partition (A, p, r)  
        Quicksort (A, p, q)  
        Quicksort (A, q+1, r)  
    }
```

```
}
```



```

Partition(A, p, r) {
    pivot ← A[p]
    i ← p
    j ← r
    while (i < j) {
        while (A[i] ≤ pivot)
            i++;
        while (A[j] > pivot)
            j--;
        if (i < j)
            exchange A[i] ↔ A[j]
    }
    exchange A[p] ↔ A[j]
    return j;
} // end.

```

$T(n) \Rightarrow$  Best case  $O(n \log n)$   
 Worst case  $O(n^2)$

3) Implementation of quick sort in Java

```

int [] array = {1, 91, 18, 9, 27, 4, 29, 19}
quickSort(arr, 0, 7);

```



this func<sup>n</sup> is in java.util.Arrays package



4) Quick Sort in C++ (STL)

```
int arr[7] = {1, 8, 9, 10, 19, 4, 27};  
sort(arr, arr+7);
```

↓

```
#include <algorithm>
```