

Assignment

Ques 1. Insertion Sort

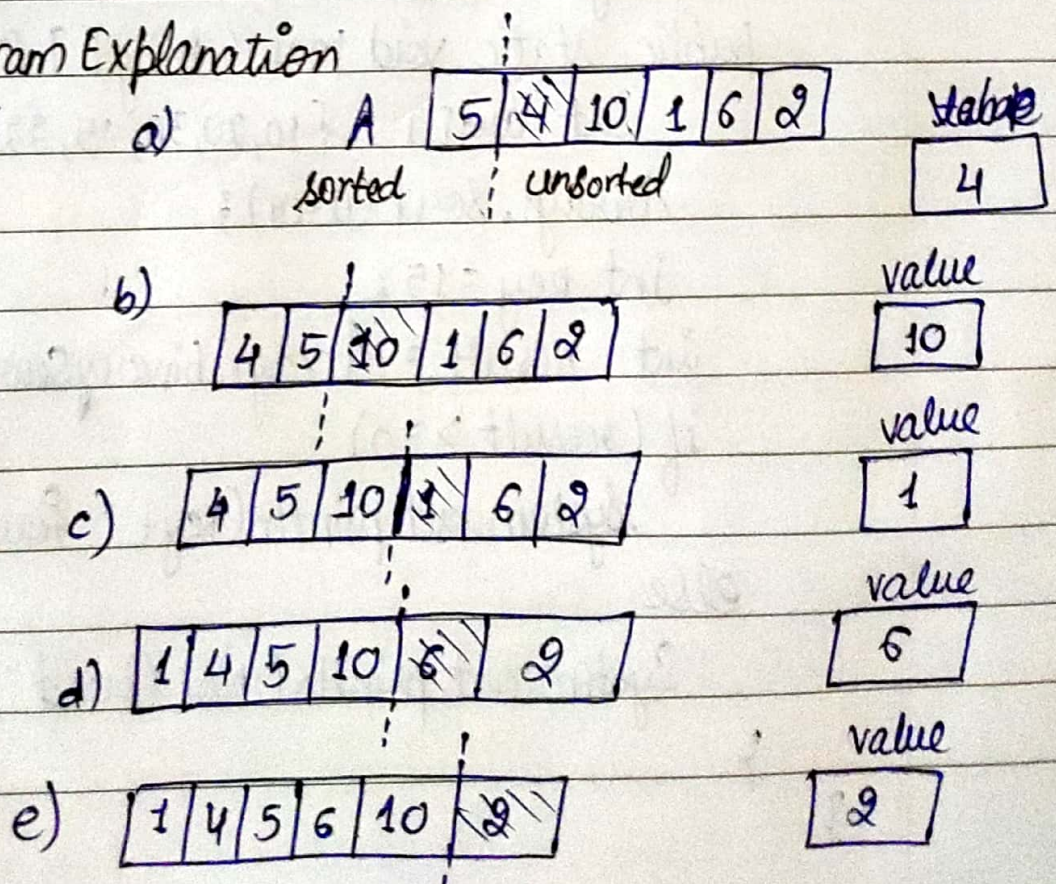
Insertion Sort (A, n)

```

{ for i ← 1 to n-1
  value ← A[i]
  hole ← i
  while (hole > 0 && A[hole-1] > value)
  { A[hole] ← A[hole-1]
    hole ← hole - 1
  }
  A[hole] ← value
}

```

Code Diagram Explanation



f)

1	2	4	5	6	10
---	---	---	---	---	----

sorted array.

(ii) Selection Sort

SelectionSort (A, n)

```

{
  for (i ← 0 to n-2)
  {
    iMin ← i
    for j ← i+1 to n-1
    {
      if (A[j] < A[iMin])
        iMin ← j
    }
    temp ← A[i]
    A[i] ← A[iMin]
    A[iMin] ← temp
  }
}

```

Code Diagram Explanation

	0	1	2	3	4	5
sorted	7	4	10	8	3	1
unsorted						

Pass 1:

1	4	10	8	3	7
---	---	----	---	---	---

Pass 2:

1	3	10	8	4	7
---	---	----	---	---	---

Pass 3:

1	3	4	8	10	7
---	---	---	---	----	---

Pass 4:-

1	3	4	7	10	8
---	---	---	---	----	---

Pass 5:-

1	3	4	7	8	10
---	---	---	---	---	----

sorted array.

Ques 2. Merge Sort

MergeSort(A, p, r)

if $p < r$

$q = \lfloor (p+r)/2 \rfloor$

MergeSort(A, p, q)

MergeSort(A, q+1, r)

Merge(A, p, q, r)

Merge(A, p, q, r)

$n_1 = q - p + 1$

$n_2 = r - q$

let L and R be new arrays

for $i = 1$ to n_1

$L[i] = A[p+i-1]$

for $j = 1$ to n_2

$R[j] = A[q+j]$

$L[n_1+1] = \infty$

$R[n_2+1] = \infty$

$i = 1$

$j = 1$
 for $k = p \text{ to } r$
 if $L[i] \leq R[j]$
 $A[k] = L[i]$
 $i = i + 1$
 else $A[k] = R[j]$
 $j = j + 1$

Diagram Explanation

