Hopcroft-Karp Algorithm

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Introduction

Hopcroft-Karp Algorithm

The Hopcroft-Karp algorithm is a graph algorithm that finds the maximum cardinality matching in a bipartite graph.

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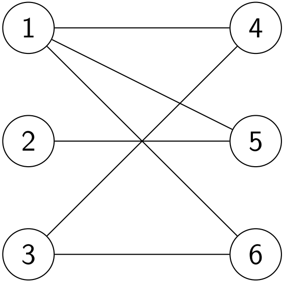

Figure: Bipartite graph

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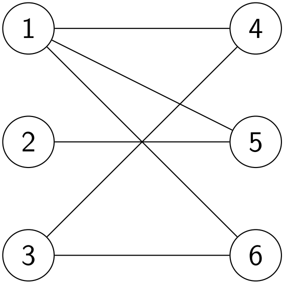

Figure: Bipartite graph

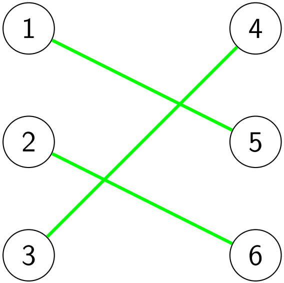

Figure: Maximum cardinality matching

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Definitions

Bipartite Graph

A graph is bipartite if its vertices can be divided into two disjoint sets such that every edge connects a vertex in one set to a vertex in the other set.

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Free vertex

A free vertex is a vertex with no matching edge connected to it.

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Algorithm

Hopcroft-Karp(G)

1: $M = \phi$

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  $P = \{p_1, p_2, \ldots, p_k\}$ //vertex-disjoint shortest augmenting paths

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  $M := M \oplus \{p_1 \cup p_2 \cup \cdots \cup p_k\}$

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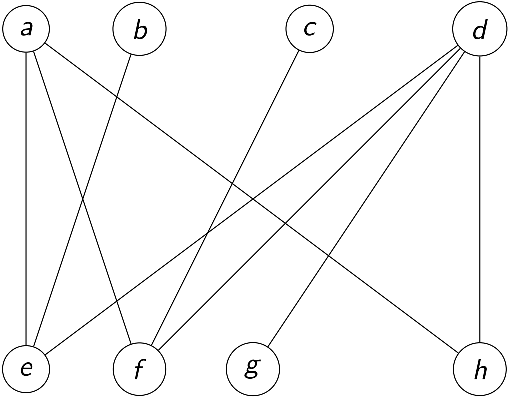
while $P \neq \phi$ **do**
    $P = \{p_1, p_2, \ldots, p_k\}$ //vertex-disjoint shortest augmenting paths
    $M := M \oplus \{p_1 \cup p_2 \cup \cdots \cup p_k\}$
end while
Output the matching M as the maximum cardinality matching.

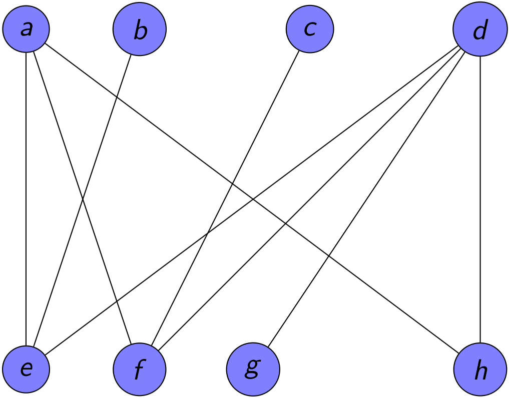
Example

A bipartite graph is given below where we will find its maximum cardinality mathching:



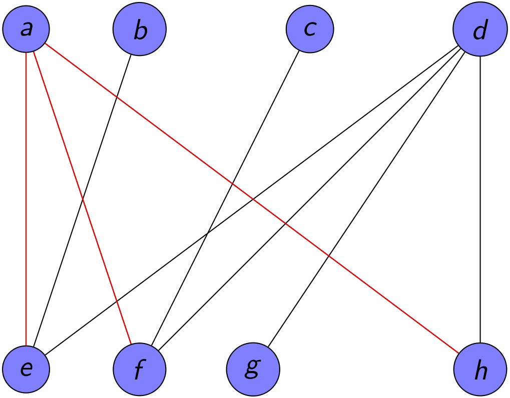
First iteration:BFS

BFS sources $= \{a, b, c, d\}$
$F = \{\}$



First iteration:BFS

BFS sources $= \{a, b, c, d\}$
$F = \{e, f, h\}$



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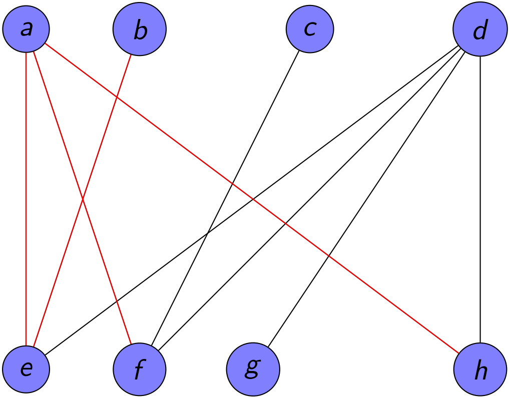

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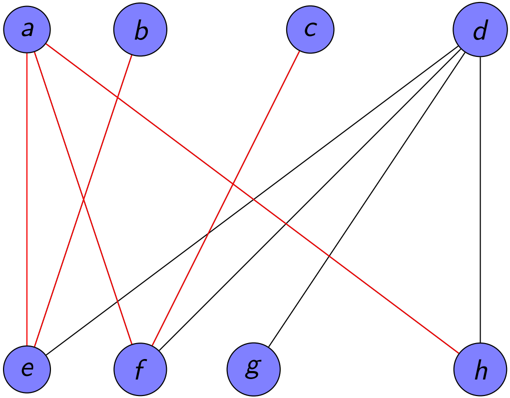

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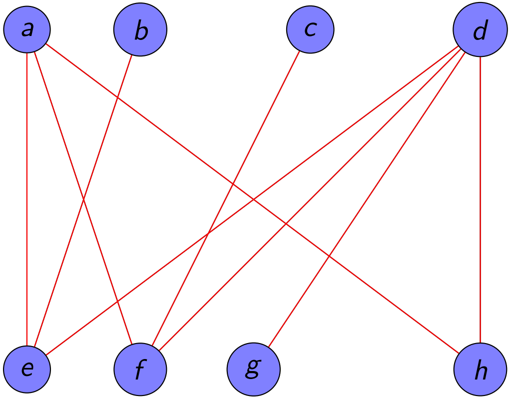
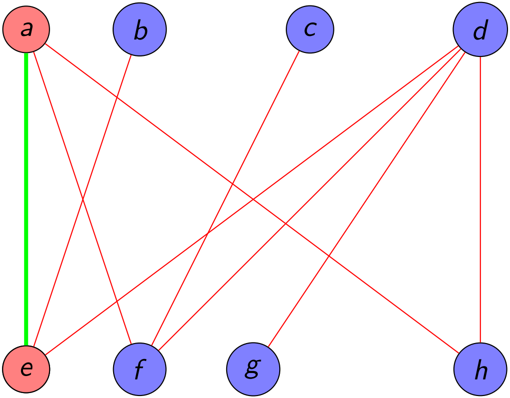


First iteration:DFS

$$F = \{e, f, h, g\}$$
$$P = \{(e - a)\}$$
$$M = \{(e - a)\}$$



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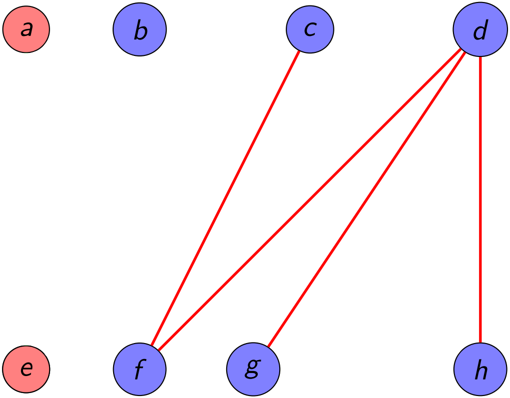

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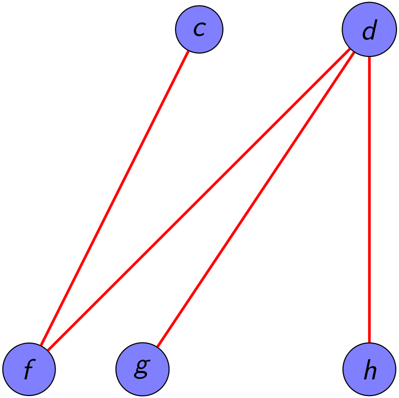

First iteration:DFS

$F = \{f, h, g\}$
$P = \{(e - a), (f - c)\}$
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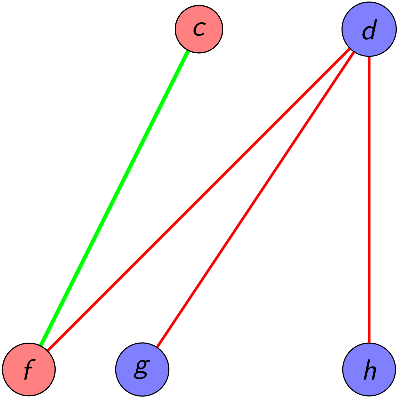

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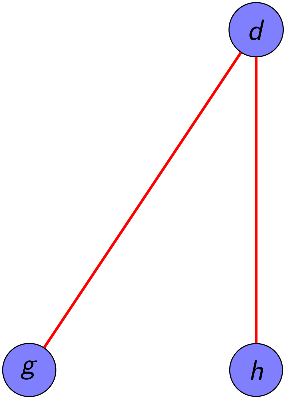

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$F = \{h, g\}$
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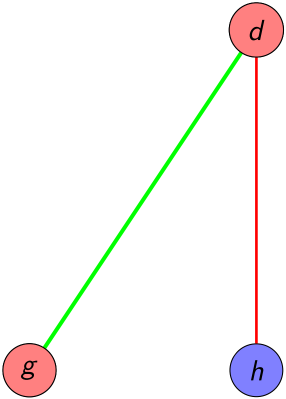
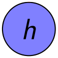


First iteration:DFS

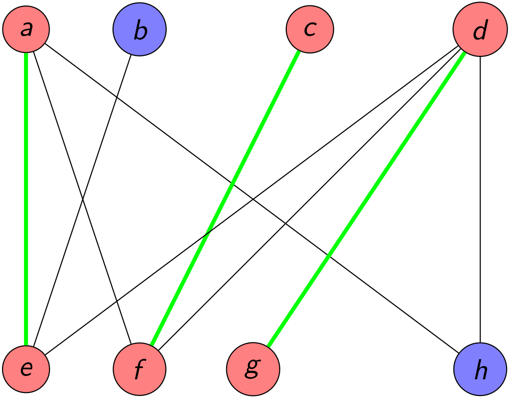
$$F = \{h\}$$
$$P = \{(e-a),(f-c),(g-d)\}$$
$$M = \{(e-a),(f-c),(g-d)\}$$



After first iteration

After running the DFS,

$$M = \{(e - a), (f - c), (g - d)\}$$



Second iteration:BFS

BFS sources $= \{b\}$
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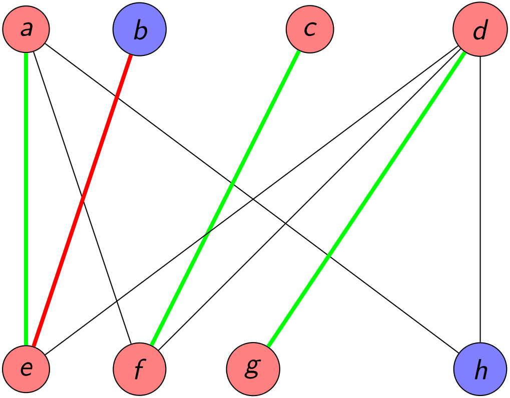
Augmenting path $= b- > e$



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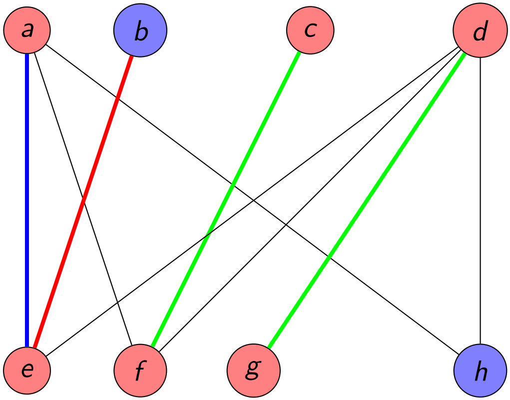
Augmenting path $= b- > e- > a$



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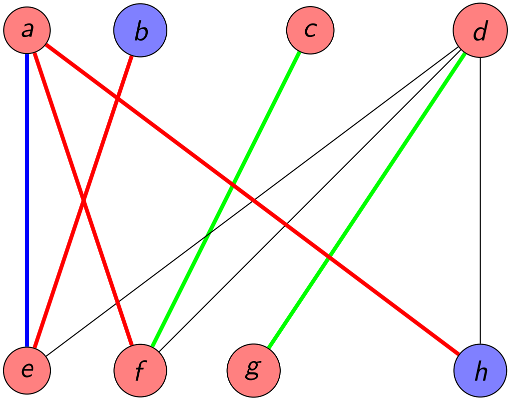

Second iteration:DFS

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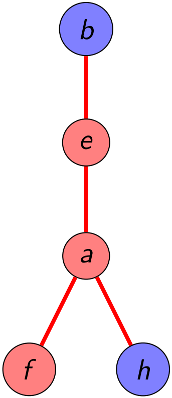

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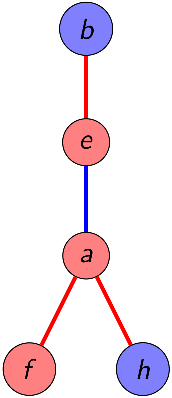

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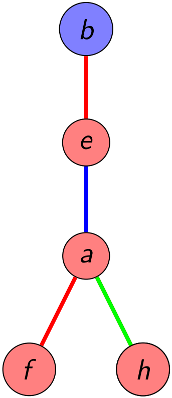

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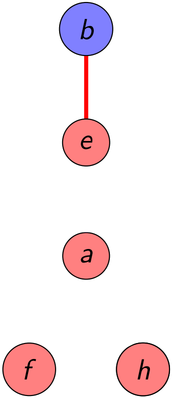

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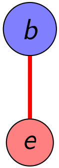

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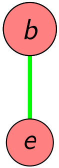

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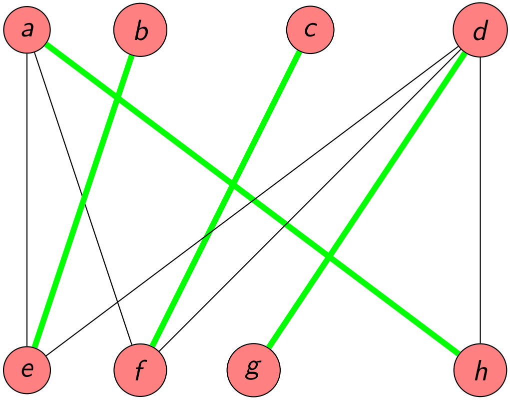
Algorithm termination

As no more free vertex is available,the algorithm terminates and finally we get the following graph with maximum cardinaliy 4 where $M = \{(h - a), (e - b), (f - c), (g - d)\}$



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Time Complexity

- The time complexity of the algorithm is O(E * sqrt(V)), where E is the number of edges and V is the number of vertices in the graph. This means that the time required to run the algorithm increases linearly with the number of edges, but is also influenced by the square root of the number of vertices.

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- The time complexity of the algorithm is considered efficient for most bipartite graphs, but may not be the best choice for extremely large graphs.

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- Online advertising - matching ads with potential viewers based on demographic and behavioral data.