Hopcroft-Karp Algorithm

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Introduction

Hopcroft-Karp Algorithm

The Hopcroft-Karp algorithm is a graph algorithm that finds the maximum cardinality matching in a bipartite graph.

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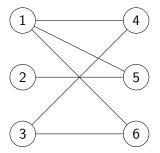


Figure: Bipartite graph

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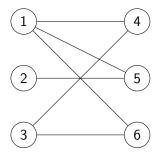


Figure: Bipartite graph

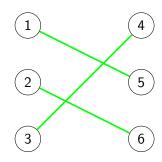


Figure: Maximum cardinality matching

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Definitions

Bipartite Graph

A graph is bipartite if its vertices can be divided into two disjoint sets such that every edge connects a vertex in one set to a vertex in the other set.

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Free vertex

A free vertex is a vertex with no matching edge connected to it.

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Hopcroft-Karp(G)

1: **M** = φ

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- 3: $P = \{p_1, p_2, p_k\}$ //vertex-disjoint shortest augmenting paths

Hopcroft-Karp(G)

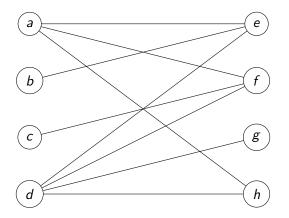
- 1: $M = \phi$
- 2: while $P \neq \phi$ do
- 3: $P = \{p_1, p_2, p_k\}$ //vertex-disjoint shortest augmenting paths
- 4: $M := M \oplus \{p_1 \cup p_2 \cup \cdots \cup p_k\}$

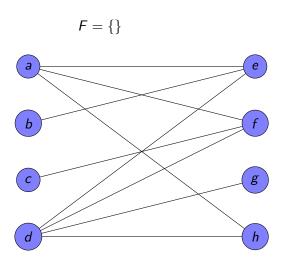
Hopcroft-Karp(G)

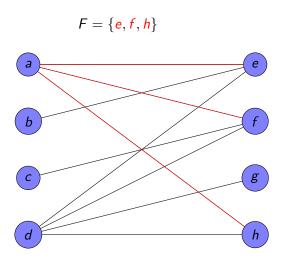
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- 5: end while
- 6: Output the matching M as the maximum cardinality matching.

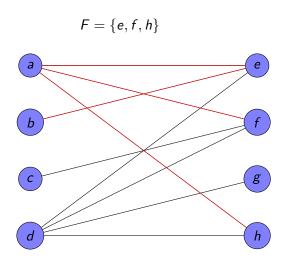
Example

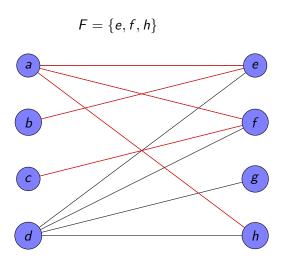
A bipartite graph is given below where we will find its maximum cardinality mathching:

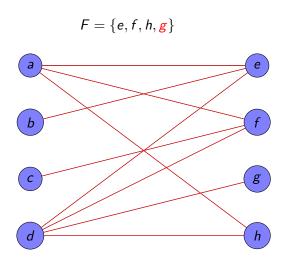




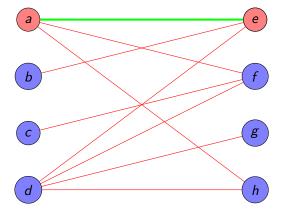








$$F = \{e, f, h, g\}$$
$$P = \{(e - a)\}$$



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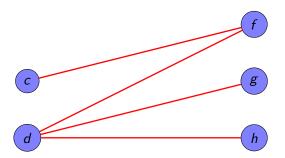
$$b$$

$$f$$

$$c$$

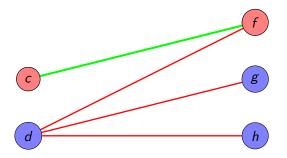
$$g$$

$$F = \{f, h, g\}$$
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$$F = \{f, h, g\}$$

 $P = \{(e - a), (f - c)\}$

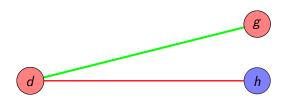


$$F = \{h, g\}$$

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$$F = \{h, g\} P = \{(e - a), (f - c), (g - d)\}$$



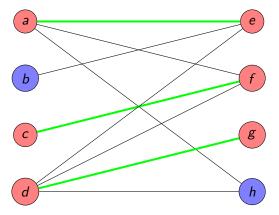
$$F = \{h\} P = \{(e-a), (f-c), (g-d)\}$$



After first iteration

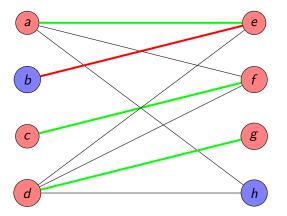
After running the DFS on the remaining part,

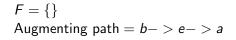
$$M = \{(e-a), (f-c), (g-d)\}$$

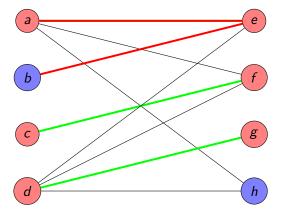


$$F = \{\}$$

 $Augmentingpath = b -> e$

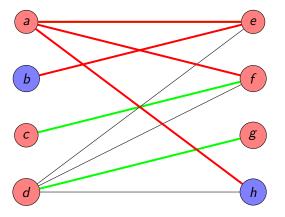


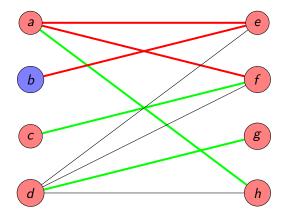


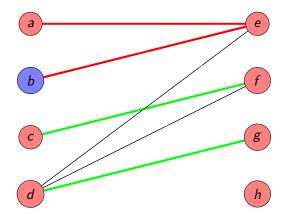


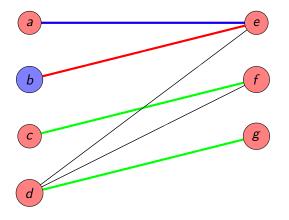
$$F = \{ \frac{H}{H} \}$$

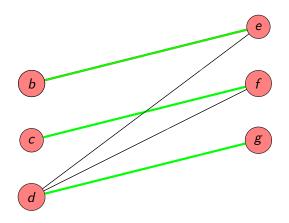
Augmenting path $= b - > e - > a - > h$

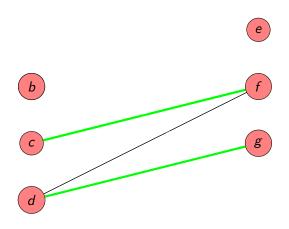


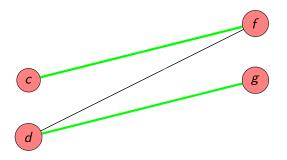












Algorithm termination..

As no more free vertex is available, the algorithm terminates and finally we get the following graph with maximum cardinaliy 4 where $M = \{A - H, B - E, C - F, D - G\}$

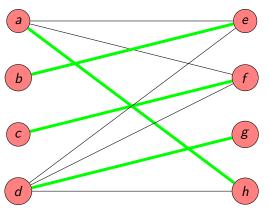


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Time Complexity

The time complexity of the algorithm is O(E * sqrt(V)), where E is
the number of edges and V is the number of vertices in the graph.
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 linearly with the number of edges, but is also influenced by the square
 root of the number of vertices.
- The time complexity of the algorithm is considered efficient for most bipartite graphs, but may not be the best choice for extremely large graphs.

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Applications

• Image segmentation - finding matches between objects in an image and a pre-defined set of object templates.

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- Job scheduling matching workers with tasks based on their skills and availability.
- Online advertising matching ads with potential viewers based on demographic and behavioral data.