Hopcroft-Karp Algorithm

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- Introduction
- 2 Definitions
- 3 Algorithm
- Time Complexity
- 6 Applications

Table of Contents

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- 2 Definitions
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Introduction

Hopcroft-Karp Algorithm

The Hopcroft-Karp algorithm is a graph algorithm that finds the maximum cardinality matching in a bipartite graph.

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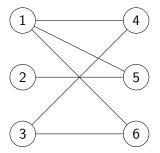


Figure: Bipartite graph

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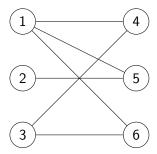


Figure: Bipartite graph

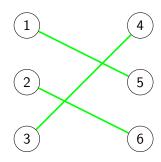


Figure: Maximum cardinality matching

Table of Contents

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- 4 Time Complexity
- 6 Applications

Definitions

Bipartite Graph

A graph is bipartite if its vertices can be divided into two disjoint sets such that every edge connects a vertex in one set to a vertex in the other set.

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Free vertex

a free vertex refers to a vertex in the left part of the bipartite graph that is not yet matched with any vertex in the right part of the graph.

Table of Contents

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Hopcroft-Karp(G)

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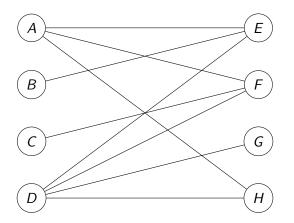
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- 2: **while** there exists an augmenting path P with respect to M **do**
- Use a breadth-first search to find an augmenting path P with respect to M.

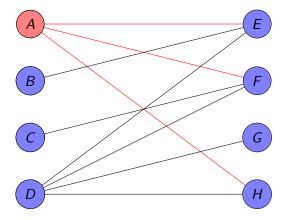
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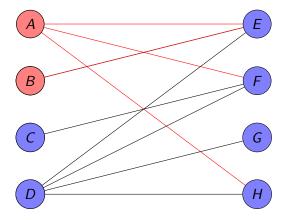
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- 3: Use a breadth-first search to find an augmenting path P with respect to M.
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- 5: end while
- 6: Output the matching M as the maximum cardinality matching.

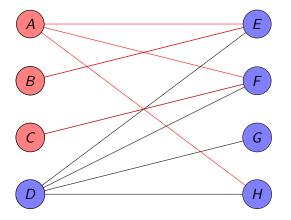
Example

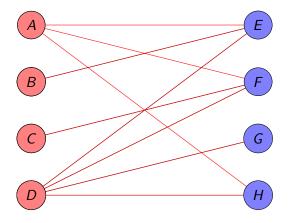
A bipartite graph is given below where we will find its maximum cardinality mathching:

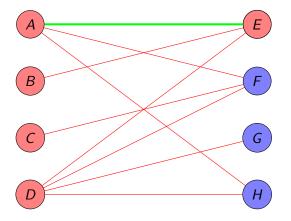


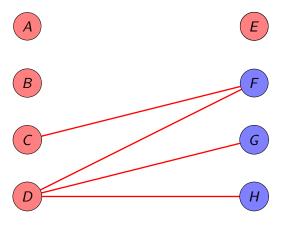


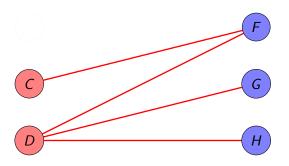


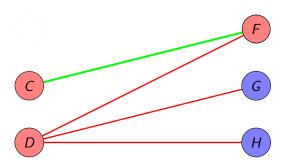


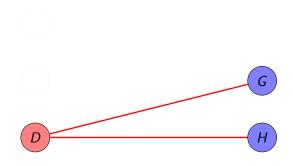






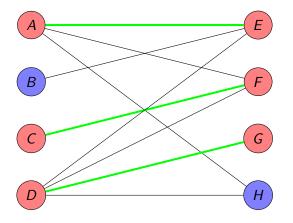


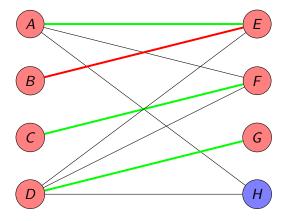


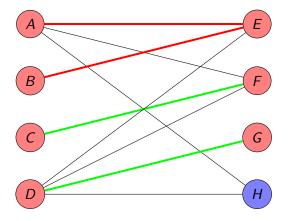


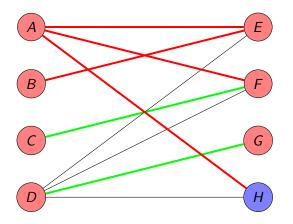
After first iteration

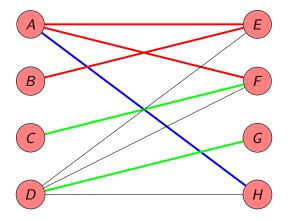
After running the DFS on the remaining part...

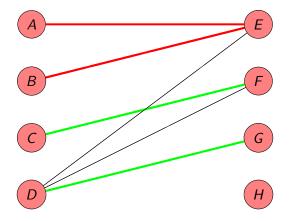


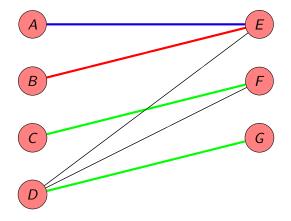


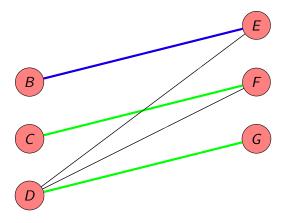


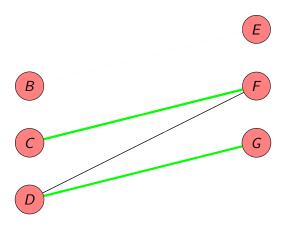


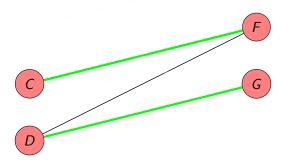












Algorithm termination..

As no more free vertex is available, the algorithm terminates and finally we get the following graph with maximum cardinaliy 4

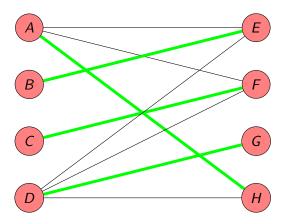


Table of Contents

- Introduction
- 2 Definitions
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- 4 Time Complexity
- 6 Applications

Time Complexity

The time complexity of the algorithm is O(E * sqrt(V)), where E is
the number of edges and V is the number of vertices in the graph.
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 linearly with the number of edges, but is also influenced by the square
 root of the number of vertices.
- The time complexity of the algorithm is considered efficient for most bipartite graphs, but may not be the best choice for extremely large graphs.

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Applications

• Image segmentation - finding matches between objects in an image and a pre-defined set of object templates.

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- Job scheduling matching workers with tasks based on their skills and availability.
- Online advertising matching ads with potential viewers based on demographic and behavioral data.