

# User's Guide to ddsip – A C Package for the Dual Decomposition of Two-Stage Stochastic Programs with Mixed-Integer Recourse

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## 1 Introduction

`ddsip` is a C-implementation of a number of scenario decomposition algorithms for two-stage stochastic linear programs with mixed-integer recourse. It expects a MIP, stochastic LPs are not supported.

The program is based on a previous Fortran 90-implementation of C.C. Carøe. Main idea of the decomposition algorithms is the Lagrangian relaxation of the nonanticipativity constraints and a branch-and-bound algorithm to reestablish nonanticipativity. The original scenarios decomposition algorithm has been developed in Carøe and Schultz (1999). Extensions including the treatment of mean-risk models have been made in Märkert (2004).

For the dual optimization we use `ConicBundle` – a C<sup>++</sup>-implementation kindly provided by C. Helmberg, see Helmberg (2012). We use the CPLEX callable library to solve the mixed-integer subproblems in the branch-and-bound tree, see CPLEX (2016). The current version of `ddsip` is ready for CPLEX 12.6.3 and requires at least version 11.2.

The implementation features risk minimization, too. This version supports mean-risk models involving the risk measure *expected excess of a target* (see Ogryczak and Ruszczyński (1999), here: *expected shortfall below target*), *excess probabilities* (as investigated in Schultz and Tiedemann (2003)), *absolute semideviation*, *worst-case-costs*, *tail value-at-risk*, *value-at-risk*, and *standard deviation*.

The code is of research quality, i.e. no production quality should be expected with respect to stability and efficiency. We kindly ask the user to support us fixing bugs by reporting them to us as they occur.

We did not include support for the SMPS format, but a rather basic input format is implemented, which is not very handy with stochastic matrix entries. Any contributions in this respect are welcome. Our code only supports the scenario formulation.

This manual describes the format of the input files and the data contained in the output files of `ddsip`. We try to provide all necessary information on the input parameters.

## 2 Stochastic programs with mixed-integer recourse

This implementation is appropriate to solve *recourse models*. Such problems were first investigated by Dantzig (1955) and Beale (1955). The conceptual idea behind recourse models is the following; assume the decisions are two-stage in the sense that some of them, say  $x$ , have to be taken immediately whilst others, say  $y$ , may be delayed to a time when uncertainty has revealed. We can write a random linear program of this type as

$$\inf_{x \in X, y(\omega) \in \mathbb{R}_+^m} \{c(\omega)x + q(\omega)y(\omega) : T(\omega)x + W(\omega)y(\omega) = h(\omega)\}. \quad (1)$$

The random parameter  $\xi := (c, q, T, W, h) : \Omega \rightarrow \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^{s \times n} \times \mathbb{R}^{s \times m} \times \mathbb{R}^s$  is defined on the probability space  $(\Omega, \mathbb{P}, \mathcal{A})$ . The set  $X \subset \{x \in \mathbb{R}_+^n : Ax = b\}$  contains all deterministic constraints on  $x$ . Inequality constraints can be handled in (1) by the introduction of appropriate slack variables. The model (1) is also referred to as *two-stage model*.

We note that the problem (1) is not yet well-defined. As the constraints include random parameters, the meaning of feasibility and thus of optimality is not clear. We complete the recourse model by adding an objective function criterion. Before we do so, we rewrite problem (1) as

$$\inf_{x \in X} \{c(\omega)x + \phi(x, \omega)\} \quad (2)$$

where

$$\phi(x, \omega) = \inf_{y \in \mathbb{R}_+^m} \{q(\omega)y : T(\omega)x + W(\omega)y = h(\omega)\}. \quad (3)$$

We note that, provided  $\phi : \mathbb{R}^n \times \Omega$  is measurable, we can regard  $\mathcal{Z} := \{c(\omega)x + \tilde{\phi}(x, \omega) : x \in X\}$  as a family of random variables. Now, each function  $\mathcal{R} : \mathcal{Z} \rightarrow \mathbb{R}$ , e.g. the expected value, some measure of risk, or a weighted sum of both, can serve as objective criterion

$$\inf_{x \in X} \mathcal{R}[c(\omega)x + \tilde{\phi}(x, \omega)]. \quad (4)$$

Our focus is on integer models, i.e. in addition to the constraints employed in problem (1) we may have integrality requirements on the variables  $x$  and  $y$ . Note that models with a linear second stage should be dealt with a version of the *L-shaped algorithm*, see Birge and Louveaux (1997).

We briefly discuss the implemented algorithm for the expected value case, i.e.

$$\mathcal{R}[c(\omega)x + \tilde{\phi}(x, \omega)] = \mathbb{E}[c(\omega)x + \tilde{\phi}(x, \omega)] := \int_{\Omega} c(\omega)x + \tilde{\phi}(x, \omega) \mathbb{P} d\omega.$$

Further algorithms suitable for the treatment of mean-risk models are described in Märkert (2004).

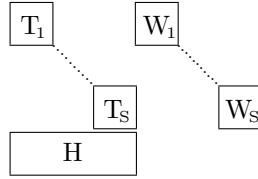


Abbildung 1: Constraints of (6)

Assume we are given a finite number of scenarios  $\xi_j$ ,  $j = 1, \dots, S$ , with corresponding probabilities  $\pi_j$ . Then, problem (4) turns into

$$\min_{x \in X, y_j \in \mathbb{Z}_+^m \times \mathbb{R}_+^{m'}} \left\{ c_j x + \sum_{j=1}^S \pi_j q_j y_j : T_j x + W_j y_j = h_j, \quad \forall j \right\}, \quad (5)$$

cf. Birge and Louveaux (1997), Kall and Wallace (1994), and Prekopa (1995). The so-called expected recourse function  $Q_E$  reads

$$Q_E(x) = c_j x + \min_{y_j \in \mathbb{Z}_+^m \times \mathbb{R}_+^{m'}} \left\{ \sum_{j=1}^S \pi_j q_j y_j : T_j x + W_j y_j = h_j, \quad \forall j \right\},$$

for all  $x \in \mathbb{R}^n$ . The program (5) is a large-scale deterministic mixed-integer linear program (MILP) with a block-angular structure. A recent and comprehensive overview of existing algorithms for problem (5) is provided in Louveaux and Schultz (2003). To mention some of the algorithmic approaches we refer to van der Vlerk (1995) for simple recourse models ( $W = (I, -I)$ ), to Laporte and Louveaux (1993) for two-stage models with a binary first stage, and to Ahmed et al. (2000) for models with an integer second stage and a fixed technology matrix  $T$ . An algorithm for problem (5) in its general form has been proposed in Carøe and Schultz (1999). It works on the expense of a branching on continuous first-stage variables. The latter algorithm is implemented in **ddsip** and described in what follows.

By introducing copies of the first-stage variables, an equivalent formulation of (5) is given by

$$\min_{x_j, y_j} \left\{ \sum_{j=1}^S \pi_j (c_j x_j + q_j y_j) : x_1 = \dots = x_S, (x_j, y_j) \in M_j, \forall j \right\} \quad (6)$$

where  $M_j = \{(x_j, y_j) : T_j x_j + W_j y_j = h_j, x_j \in X, y_j \in Y\}$ ,  $j = 1, \dots, S$ .

Considering the constraint matrix of (6) (cf. Figure 1), we can identify  $S$  single-scenario subproblems solely coupled by the equality (*nonanticipativity*) constraints on the copies of the first-stage variables and written as  $\sum_{j=1}^S H_j x_j = 0$ , where  $H = (H_1, \dots, H_S)$ . The problem decomposes when we relax the non-anticipativity constraints.

Upper bounds on the optimal value can be obtained by heuristics based on the solutions for the subproblems. We get a lower bound by solving the

Lagrangian dual, which is a nonlinear concave maximization

$$z_{\text{LD}} := \max_{\lambda \in \mathbb{R}^l} \min_{x_j, y_j} \left\{ \sum_{j=1}^S \pi_j (c_j x_j + q_j y_j) + \lambda \sum_{j=1}^S H_j x_j : (x_j, y_j) \in M_j, \forall j \right\}. \quad (7)$$

In general, the involved integrality restrictions lead to an optimality gap. If we are not satisfied with the bounds given by the above method, we can elaborate a branch-and-bound algorithm that successively reestablishes the equality of the components of the first-stage vector. Let  $\mathcal{P}$  denote a list of problems.

**Algorithm SD:** Scenario decomposition (Carøe and Schultz (1999))

*STEP 1* Initialization: Set  $z^* = \infty$  and let  $\mathcal{P}$  consist of problem (5).

*STEP 2* Termination: If  $\mathcal{P} = \emptyset$  then  $x^*$  with  $z^* = Q_E(x^*)$  is optimal.

*STEP 3* Node selection: Select and delete a problem  $P$  from  $\mathcal{P}$  and solve its Lagrangian dual. If the associated optimal value  $z_{\text{LD}}(P)$  equals infinity (infeasibility of a subproblem) go to *STEP 2*.

*STEP 4* Bounding: If  $z_{\text{LD}}(P)$  is greater than  $z^*$  go to *STEP 2*. Otherwise proceed as follows; if the first-stage solutions  $x_j, j = 1, \dots, S$ , of the subproblems are

- identical, then set  $z^* := \min\{z^*, Q_E(x_j)\}$ , delete all  $P' \in \mathcal{P}$  with  $z_{\text{LD}}(P') \geq z^*$  and go to *STEP 2*.
- not identical, then compute a suggestion  $\hat{x} := \text{Heu}(x_1, \dots, x_S)$  using some heuristic. Set  $z^* := \min\{z^*, Q_E(\hat{x})\}$  and delete all  $P' \in \mathcal{P}$  with  $z_{\text{LD}}(P') \geq z^*$ .

*STEP 5* Branching: Select a component  $x_{(k)}$  of  $x$  and add two new problems to  $\mathcal{P}$  that differ from  $P$  by the additional constraint  $x_{(k)} \leq \lfloor x_{(k)} \rfloor$  and  $x_{(k)} \geq \lfloor x_{(k)} \rfloor + 1$ , respectively, if  $x_{(k)}$  is integer, or  $x_{(k)} \leq x_{(k)} - \varepsilon$  and  $x_{(k)} \geq x_{(k)} + \varepsilon$ , respectively, if  $x_{(k)}$  is continuous.  $\varepsilon > 0$  has to be chosen such that the two new problems have disjoint subdomains. Go to *STEP 3*.

The algorithm is finite if  $X$  is bounded and if some stopping criterion is employed that prevents the algorithm from endless branching on the continuous components of  $x$ , see Carøe and Schultz (1999).

The function  $Q_E$  is evaluated at  $x$  by fixing the first stage to  $x$ , solving the scenario many subproblems, and calculating the expected value of the corresponding optimal values. Thus, infeasible suggestion are identified immediately.

We remark that problems related to random recourse have to be cared about by the users, cf. Walkup and Wets (1967).

### 3 Input files

ddsip requires 3 input files, a number of other files are optional.

- Specification file: a file containing the specifications of the stochastic program, some CPLEX- and b&b-parameters
- Model file: a file readable by CPLEX (e.g. lp- or mps-format, possibly in gzipped form) that specifies the single-scenario model
- Priority order file for subproblems (optional): a CPLEX order file corresponding to the model file
- RHS scenario file: a file containing the stochastic right-hand sides and the probabilities of the scenarios
- Cost scenario file (optional): a file containing stochastic cost coefficients
- Matrix scenario file (optional): a file containing stochastic matrix entries
- Priority order file for master (optional): a file containing a branching order for the master branch-and-bound procedure
- Start information file (optional): a file containing start informations such as a feasible solution and/or a lower bound

A comfortable way to invoke the program on a Unix system is the command

```
ddsip < files2sip
```

where the file *files2sip* may contain the lines displayed in Figure 2.

```
sip.in  
model.lp  
model.ord  
rhs.sc  
cost.sc  
matrix.sc  
order.dat  
start.in
```

Abbildung 2: Input file

### 4 The specification file

Comments could be included anywhere in this file apart from the CPLEX parameter section. An Asterisk identifies the beginning of a comment until the end of the line. This way explanations could be included and parameter lines

temporarily inactivated Comments or commenting out lines could not be used within the CPLEX parameter section.

Keywords are accepted if their start matches the strings given in the tables. Unknown keywords are simply ignored. Lines with the first occurrence of each keyword are evaluated.

A sample specification file is given in Appendix A.

#### 4.1 Parameters for the two-stage model

Hopefully, the identifiers in the first part of the file are self explaining. Otherwise, they should become clear when consulting the two-stage chapter of one of the standard text books on stochastic programming, see Birge and Louveaux (1997); Kall and Wallace (1994); Prekopa (1995). Some of the identifiers, as e.g. FIRSTVAR, are redundant. They serve to check consistency with the model file.

Either a prefix or a postfix of the variable names for identifying the first-stage variables can be specified. The default behaviour is to expect a postfix of the form "01" in their names.

Name	Type	Default/Description	
FIRSTC	Int	0	Number of First stage constraints
FIRSTV	Int	0	Number of First stage variables
SECCON	Int	0	Number of Second stage constraints
SECVAR	Int	0	Number of Second stage variables
POSTFIX	String	<i>none</i>	postfix for first-stage variables
PREFIX	String	<i>none</i>	prefix for first-stage variables
SCENAR	Int	0	Number of scenarios
STOCRHS	Int	0	Number of stochastic rhs elements
STOCCOST	Int	0	Number of stochastic cost coefficients
STOCMAT	Int	0	Number of stochastic matrix entries

Table 1: Problem specification parameters

#### 4.2 CPLEX parameters

CPLEX parameters to be set different from their defaults have to be specified following a line with the indicator *CPLEXBEGIN*.

The contained lines have to start with the CPLEX parameter number followed by the parameter value. This can be followed by a comment, starting with an asterisk. The parameter numbers are specified in the CPLEX callable library documentation, section parameter table.

The settings following directly the line with *CPLEXBEGIN* are applied to all problems solved.

Special settings for CPLEX parameters can be specified for the different stages of the program.

For the calculation of the EEV parameter values can be overwritten after the identifier *CPLEXEEV*,

for the evaluation of lower bounds after the identifier *CPLEXLB*,

for a continuation of the evaluation of lower bounds with different settings after the identifier *CPLEX2UB*,  
for the evaluation of upper bounds after the identifier *CPLEXUB*,  
for a continuation of the evaluation of upper bounds with different settings after the identifier *CPLEX2UB*,  
for the subproblems of the Lagrangian dual after the identifier *CPLEXDUAL*,  
for a continuation of the subproblems of the Lagrangian dual after the identifier *CPLEX2DUAL*.

For some problems it is beneficial to use two optimization calls with different parameter settings (e.g. mip emphasis, heuristic frequency) sequentially, the second one starting with the B&B tree and solutions obtained in the first call. This is enabled for lower and upper bound calculations and the Lagrangean dual if the sections beginning with *CPLEX2LB*, *CPLEX2UB*, or *CPLEX2DUAL* are specified.

The CPLEX parameter section has to end with the identifier *CPLEXEND*. Parameters not present in the mentioned chapters are set according to their CPLEX default values, cf. CPLEX (2016).

Note that in the section between *CPLEXBEGIN* and *CPLEXEND* there could be comments following the pair of numerical parameter identifier and value, but not whole lines could be comments, i.e. no line should start with an asterisk.

### 4.3 Parameters for the decomposition procedure

The parameters that effect the amount of output and the termination behavior are listed in the table below. The first column of the table contains the name of the parameter, the second one specifies whether it is an integer or a real (Dbl) parameter, and the last column explains the parameter's meaning.

The signal SIGTERM (as defined on LINUX/UNIX-environments) is handled by the program, other signals are not handled separately. SIGTERM causes the termination of *ddsip* after solving the current subproblem.

So far all termination parameters are static in the sense that they cannot be changed during the branch-and-bound algorithm. Therefore, a careful trade-off between these parameters and the termination parameters of the bundle method, see Section 4.6, is essential for the numerical performance, cf. Carøe and Schultz (1998).

Parameters that effect the behavior of the branch-and-bound algorithm are compiled in the Tables 2 and 3. The default values are marked with a star. The use of start values is described in Section 7.1, the use of priority order information in Section 7.2. The parameter RELAXL effects the relaxation levels during the evaluation of lower bounds. It relates to the scenario subproblems.

Name	Type	Range	Default	Description
OUTLEV	Int	0..100	0	Amount of output to <i>more.out</i> . <b>Caution:</b> The file <i>more.out</i> may become large for high values of OUTLEV!
OUTFIL	Int	0..6	1	Amount of output files, see chapter 8. <b>Caution:</b> If OUTFIL is greater than 3, lp- or sav-files are written at each node and for each scenario!
LOGFRE	Int	0..	1	A line of output is printed every i-th iteration.
NODELI	Int	0..	10000	The node limit for the branch-and-bound procedure.
TIMELI	Dbl	0..	86400.	The total time limit in seconds (CPU-time) including the time needed to solve the EEV problem.
ABSOLU	Dbl	0..	0	The absolute duality gap.
RELATI	Dbl	1e-10..	1e-4	The relative duality gap.
EEVPRO	Int	0..1	0	If the parameter is 1, then solve the EEV-problem and report the VSS, cf. Birge and Louveaux (1997).

Table 2: Output and termination parameters

The *branching value* is used to create two new subproblems as described in *STEP 5* of the decomposition algorithm. For continuous components it depends on the parameter EPSILON.

The dispersion norm of a node is calculated as  $\max_j \{\max_i x_{ij} - \min_i x_{ij}\}$  where  $x_j$ ,  $j = 1, \dots, S$ , are the (first-stage part of the) solutions of the  $S$  subproblems and  $x_{ij}$ ,  $i = 1, \dots, n$ , are their  $i$ -th components. Nodes with a dispersion norm smaller than NULLDISP are considered as leaves of the B&B tree (i.e. no further branching on them).



Name	Type	Range	Default/Description
PORDER	Int	0..1	0 Use priority order for branching? If this parameter is 1, a priority order file has to be specified.
STARTI	Int	0..1	0 Use start information? If this is 1, a file containing the start values and/or lower bound has to be specified.
MAXINH	Int	0..	5 Maximal level of "inheritance" of scen. solutions in the B&B-tree.
HOTSTA	Int	0..6	0 No warm starts during B&B.
			1* Use solution pool of previous scenario and integer values of the same scenario in father as initial solutions.
			2 Use solution pool of previous scenario and lower bound from father node (not applicable for risk models).
			3 Use solutions of all previous scenarios in the same node.
			4 Use solutions of all previous scenarios of the same node as well as solutions of all scenarios in father node.
			5 Uses solution info as in 3 together with the lower bound from father node.
			6 Uses solution info as in 4 together with the lower bound from father node.
BRADIR	Int	-1*,1	-1 Branching down is done first.
			1 Branching up is done first.
BRASTR	Int	0..2	0 Use the middle between maximal and minimal value for the chosen variable as branching value.
			1 Use the expected value (weighted average) of the scenario solutions for the chosen variable as branching value.
			2* Use a point in between the weighted average of the scenario solutions for the chosen variable and the mean between maximal and minimal value for that variable as branching value.
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Name	Type	Range	Default/Description	
BRAEQU	Int	-1..1	1	"equal distribution" – Among the variables with maximal dispersion value choose one for branching which most equally divides feasibility of scenario solutions among both new nodes
			0*	"mixed" – in 3/5 of the iterations choose a variable such that most of the scenario solutions remain feasible for one of the new nodes, in the others equally distributed
			-1	"unequally" – in all iterations choose a variable such that most of the scenario solutions remain feasible for one of the new nodes
INTFIR	Int	0..1	0	Branching exclusively according to branching rules.
BOUSTR	Int	0..10	1*	Branching on integers first.
			0	choose node with best bound as next node
			1	heuristic: among nodes with sufficiently small bound choose node with biggest dispersion norm (unsolved nodes first)
			2	heuristic: among nodes with sufficiently small bound choose node with biggest dispersion norm
			3	heuristic: among nodes with sufficiently small bound choose node with least number of violations (unsolved nodes first)
			4	heuristic: among nodes with sufficiently small bound choose node with least number of violations
			5-9	heuristic: depth first until a incumbent is found, then switch to BOUSTRAT-5
			10*	heuristic: among last four generated nodes choose node with best bound (unsolved nodes first) as long as its objective value does not exceed a threshold. Choose the best bound node otherwise and also every 25th node.

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Name	Type	Range	Default/Description	
PERIOD	Int	1..	32	for BOUSTRAT>0: determines the period of bounding steps with changing tolerances for "sufficiently small"
TOLSMA	Int	1..PERIOD	16	for BOUSTRAT>0: determines the number of steps in each period with a small tolerance
ACCURA	Dbl	1e-13..1	1e-12	Accuracy – real values are considered equal if the absolute value of their difference is less than accuracy
EPSILO	Dbl	1e-10..	1e-10	Specifies the gap between disjoint subdomains if continuous variables are branched.
NULLDI	Dbl	5e-10..	5e-10	Branch nodes only if their dispersion norm is greater than NULLDI.
KAPPA	Int	0..2	0*	No gathering of kappa values by CPLEX
	Int		1	sample gathering of kappa values by CPLEX and reporting
	Int		2	full gathering of kappa values by CPLEX and reporting
RELAXL	Int	0..2	0*	No integrality relaxation.
			1	Relax first-stage variables.
			2	Relax first- and second-stage variables.
QUANTI	Int	0..	10	Number of Quantiles to be displayed at the end
PREMA	Int	0..1	1	Should premature stop in upper bounding be used?
MAXINH	Int	0..	5	maximal level of inheritance of solutions in lower bounding
HEURIS	Int	0..13,99,100	100	choice of heuristics, cf. Table 5

Table 3: Branch-and-bound parameters

The MAXINHERIT parameter specifies a maximal level of passing on solutions and bounds in the B&B-tree for the lower bound step. This facility is interrupted by an invocation of the conic bundle algorithm, which will change the Lagrangian multipliers. It saves half of the solutions of scenario problems, but in case these problems are not solved to optimality the lower bound might be inferior to the one possibly obtained for the descendants (with the additional bounds on the variables branched on). For this reason the level of inheritance can be bounded to the number of branching steps specified with the MAXINHERIT keyword.

After the first node the scenarios are sorted descending wrt. their contribution to the expected value. This serves earlier detection of inferiority and allows stopping of evaluation of scenarios.

If conic bundle is used, the lower bound for each scenario is not fully reliable in upper bounding (which is done without Lagrangean parameters). The parameter PREMATURE (default 1) to allow/disallow premature stopping in upper bounding can be set to 0 when conic bundle is used.

#### 4.4 Heuristic

The decomposition algorithm uses a heuristic to guess feasible first-stage solutions based on the solutions of the subproblems in the current node. The heuristic is set by means of the parameter HEURIS. The possible values of HEURIS are listed in the following table.

More than one of the heuristics could be used sequentially in every node when specifying the value 99, followed by a list of heuristics to be used, the specs file line could look like this:

HEURISTIC 99 2 3 12 \* heuristics 2, 3, **and** 12 are specified

Value	Description
1	The average of the solutions is used. Integer components are rounded down.
2	The average of the solutions is used. Integer components are rounded up.
3	The average of the solutions is used. Integer components are rounded to the next integer.
4	The solution occurring most frequently is used.
5	Use the solution of a scenario that is closest ( $l_1$ norm) to average.
6	The first-stage solution with highest probability is chosen (adding probabilities in case of identical first-stage scenario solutions)
7	Solution with best objective value
8	Solution with worst objective value
9	Solution with minimal sum of first-stage variables
10	Solution with maximal sum of first-stage variables
12	Try solutions of all scenarios
13	Apply heuristic 3 and 5 alternating
99	Indicates that a list of heuristics is given in the sequel. See the example in the Appendix.
100*	Heuristics 1-10 are applied plus the more expensive heuristic 12, if one of the others found a new incumbent. Heuristic 12 is also applied in the first 11 nodes to facilitate dicoverly of an incumbent early in the B&B tree.

Table 5: Values of parameter HEURIS

In order to save time the order of evaluation of the scenario problems is changed if an infeasible one is encountered. Since often the same scenarios gives infeasibility these ones are evaluated first in the sequel.

#### 4.5 Parameters for the risk model

The parameters for the risk models are displayed below. Some parameter settings lead to ill-posed models, cf. Märkert (2004). `ddsip` provides only limited consistency checks with this respect.

The same holds true for the choice of an algorithm. Algorithms similar to the scenario decomposition as outlined above only apply to mean-risk models with *decomposable* linear risk measures. The algorithm FSD can be used with any risk measure that is consistent with first-order stochastic dominance. The weakest algorithm NFSD allows the minimization of any (nonlinear) risk measure. A detailed description of the different algorithms can be found in Märkert (2004).

Name	Type	Range	Default	Description
RISKMO	Int	-7..7	0	Risk model
WEIGHT	Dbl	0..	0.9	Weight on risk term in objective.
TARGET	Dbl	..	-1e+20	Target for target measures.
PROBLE	Dbl	0..1	0	Probability level for (T)VaR.
RISKAL	Int	0..2	0*	Scenario decomposition.
			1	FSD-algorithm (advised for TVaR, also for pure risk models except 4 (worst case costs) and 7 (standard deviation)).
			2	NFSD-algorithm.
BRAETA	Int	0..1	1	Branch order of additional first-stage variable in models 4 and 5. A value of 0 means that this variable will not be branched and its dispersion ignored.
RISKBM	Dbl	..	1e+20	Big M for risk model (-)2, also used as bound for $\eta$ with risk model 5.

Table 6: Parameters for the risk model

The program offers 7 different mean-risk models and the associated risk models, in which the expected value is not minimized. In the following table we have compiled the possible settings.

Name	Value	Description
RISKMO	1, −1	Expected excess above target.
	2, −2	Excess probabilities.
	3, −3	Absolute semideviation.
	4, −4	Worst-case-costs.
	5, −5	Tail value-at-risk.
	6, −6	Value-at-risk.
	7, −7	Standard deviation

Table 7: Risk models

Positive values of RISKMO implement the model  $\min_X EX + \alpha \mathbb{R}X$ , negative ones the risk model  $\min_X \mathbb{R}X$ . These pure risk models are often harder to solve.

For the (mean-)absolute semideviation model the TARGET has to be set less or equal the optimal expected value when using RISKALG 0 (cf. Märkert (2004)).

#### 4.6 Parameters for the dual method

There is an online manual of `ConicBundle` available Helmberg (2012). We refer also to Helmberg (2000) where the method is described in the context of semidefinite programming. Below, we only explain the parameters defined in our implementation.

The dual method cannot be used in combination with the algorithm based on the FSD-consistency of the risk measure, cf. Märkert (2004). The user has to take care about this detail.

The use of stochastic cost coefficient in combination with the Lagrangian dual is not supported, yet. We recommend to reformulate the problem by an additional variable and an additional constraint representing the objective function. This leads to a stochastic matrix.

The `ConicBundle` weight is changed between single dual steps under certain conditions by `ConicBundle` itself. `DDSIP` offers the possibility to heuristically change (increase or decrease) the weight between descent steps. By default this feature is turned on and could be deactivated by parameters.

Name	Type	Range	Default/Description
CBFREQ	Int	..	-16 Use <b>ConicBundle</b> in every $i$ th node. Negative numbers are interpreted as $-i$ —th node plus preceding one.
CBITLI	Int	0..	20 Iteration limit: descent steps.
CBRITLI	Int	0..	CBITLIM+10 Iteration limit in root node: descent steps.
CBSTEP	Int	0..	12 Maximal number of iterations within one descent step.
CBTOTI	Int	0..	5000 Iteration limit: total iterations, i.e. descent and null steps.
CBCONT	Int	1..	4 Used when CBFREQ<0: number of nodes with invocation of <b>ConicBundle</b> in every node after CBBREAK.
CBBREA	Int	1..	24 Used when CBFREQ<0: invocation of <b>ConicBundle</b> in nodes 1..CBBREAK only for nodes, which inherited more than half of the scenario solutions.
CBPREC	Dbl	0..	1e-14 Precision of bundle method.
CBPRIN	Int	0..	0 <b>ConicBundle</b> output level.
NONANT	Int	1..3	The identity of the first-stage vectors is represented by $1^* \quad x_1 = x_2, x_1 = x_3, \dots, x_1 = x_n$ $2 \quad x_1 = x_2, x_2 = x_3, \dots, x_{n-1} = x_n$ $3 \quad p_i x_i = \sum_{j=1, j \neq i}^n p_j x_j \quad \forall i$
CBBUNS	Int	1..	500 Maximal bundle size
CBWEIG	Dbl	0..	1. Initial weight (cf. CB manual)
CBFACT	Dbl	0..1	0.01 Initial weight in the next node is final (1-CBFACTOR)*(weight in the father) + CBFACTOR*(initial weight)
CBINHE	Int	0,1	0 Should the solutions of the father node be inherited in the initial evaluation?
CBCHAN	Int	0,1	0 Should the tolerances and time limits for CPLEX be temporarily changed after a number of unsuccessful steps?
CBREDU	Int	0,1	1 Should the weight be decreased after successful steps with no or few null steps?
CBINCR	Int	0,1	1 Should the weight be increased after steps with too many null steps?

Table 8: ConicBundle parameters

## 5 The model file

The model file contains the single-scenario model. Its format can be one of the formats readable by CPLEX, as e.g. the lp- or mps-format. The gzipped forms of these files (with extensions .lp.gz or .mps.gz) can be used for the sake of saving disk space. This is especially advisable in case OUTFILE is chosen greater than 2.

The first-stage variables and constraints have to be identified by a prefix or a postfix appended to their names as specified in the specs file.

## 6 The scenario files

Unfortunately, the program `ddsip` is not conform to the SMPS data format, yet. But in contrast to the SMPS format due to the requirements for the variable and constraint names and the input format for the stochastic coefficients no specific ordering is required in the model file. The proprietary data format of the scenario files is described below.

Each file has to start with the keyword 'Names' followed by a list of the names of constraints for the stochastic right-hand sides, variables for the stochastic objective coefficients, and pairs of names for constraint and variable for the stochastic matrix entries. These name lists determine the order of objects the coefficients in the different scenarios are assigned.

The identifier *sce* indicates the begin of the entries for the next scenario. Figure 3 displays the different scenario files.

right-hand sides	cost coefficients	matrix entries
Names	Names	Names
constr25	var16	cons5 var15
constr30	var18	cons5 var16
⋮	⋮	⋮
scenario1	scenario1	scenario1
0.1	23.1	0.125
12.12	33.33	34
32.4		
⋮	⋮	⋮
scenarioN	scenarioN	scenarioN
0.1	25.56	30.3
30.2	11.5	22.56
41.0		
⋮	⋮	⋮

Abbildung 3: Format of the scenario files



In the right-hand side scenario file, the first number after the identifier is the scenario probability. For problems without stochastic right-hand sides, this file has to contain only the probabilities of the individual scenarios.

In the cost coefficient scenario file the probability entries are left out. Otherwise the entries are the same as those in the right-hand side scenario file, cf. Figure 3. The number of stochastic cost coefficients  $c$  has to be specified in the specification file.

## 7 Optional files

### 7.1 Start information file

Start information can be provided as displayed in Figure 4. Hereby, BEST

BEST	2345
BOUND	2000
SOLUTION	
1	
2	
:	
MULTIPLIER	
3	
4	
:	

Abbildung 4: File with start values

should be an upper bound, BOUND a lower bound, SOLUTION a feasible first-stage solution, and MULTIPLIER a Lagrangian multiplier. Any item of these four different informations may be left out.

The user has to care for the consistency of the data.

### 7.2 The priority order files

If indicated by setting the CPLEX parameter CPX\_PARAM\_MIPORDIND, the name of a CPLEX branching order file to be used for the scenario problems has to be specified following the name of the model file (named model.ord in Figure 2). Its format is described in the CPLEX manual.

If the parameter PORDER is set, a file with branching order information for the master problem has to be given (named order.dat in Figure 2). The file has two columns whereby the first one contains **indices** (starting with 0 as in the output of **siphelp**) of first-stage variables and the second one integer values. High values lead to early branching.

## 8 Output

### 8.1 Output on screen

The following shows typical lines of output as produced by `ddsip` (accuracy of the values cut due to line length). The meaning of the single entries is rather

Node	Nodes	Left	Objective	Heuristic	Best Value	Bound	Viol./Dispersion	Gap	Wall Time	CPU Time	Father	
0	1	1	44841.3565	infeasible		44841.356	5 3		0h 00:03	0h 00:03	-1	
1	3	2	47944.6273	infeasible		44841.356	7 2.7888		0h 00:06	0h 00:06	0	
2	3	2	44844.5722	infeasible		44844.572	5 2.7888		0h 00:07	0h 00:07	0	
3	5	3	44953.0125	infeasible		44844.572	7 2.7888		0h 00:09	0h 00:08	2	
4	5	3	44848.1886	multiple		44848.188	4 2.7888		0h 00:09	0h 00:09	2	
5	7	4	45549.6459	infeasible		44848.188	5 2.7888		0h 00:10	0h 00:10	4	
6	7	4	44854.2998	infeasible		44854.299	3 2.7888		0h 00:12	0h 00:12	4	
*	7	9	5 Heuristic	2	66346.70463	66346.70463			0h 00:14	0h 00:13		
	7	9	5	45712.4035	66346.70463	66346.70463	44854.299	6 2.7888	32.39%	0h 00:14	0h 00:13	6
* 8	9	5 Heuristic	2	44864.77749	44864.77749				0h 00:16	0h 00:16		
	8	9	1	44859.2107	44864.77749	44864.77749	44859.210	2 2.7888	0.0124%	0h 00:16	0h 00:16	6
	9	11	1	cutoff	44864.77749	44859.210			0.0124%	0h 00:17	0h 00:17	8
	10	11	1	44863.2847	44865.25242	44864.77749	44863.284	1 2.7888	0.00332%	0h 00:20	0h 00:19	8
	11	13	2	44864.4325	44865.05453	44864.77749	44863.284	1 1.162	0.00332%	0h 00:22	0h 00:21	10
* 12	13	2 Heuristic	12	44864.65876	44864.65876				0h 00:24	0h 00:23		
	12	13	2 Heuristic	3	44864.38172	44864.38172			0h 00:25	0h 00:24		
	12	13	1	44863.5109	44864.38172	44864.38172	44863.510	1 1.6268	0.00194%	0h 00:25	0h 00:24	10
	13	15	2	44864.1805	44864.54332	44864.38172	44863.510	1 0.67783	0.00194%	0h 00:27	0h 00:26	12

straightforward. Here is a short explanation.

Node	gives the number of the currently solved node.
Nodes	counts the number of nodes generated in the tree.
Left	counts the number of nodes in the front tree.
Objective	is the lower bound of the current node.
Heuristic	is the upper bound returned by the heuristic.
Best Value	is the overall upper bound.
Bound	is the overall lower bound.
Viol.	is the number of variables violating the nonanticipativity.
Dispersion	is the greatest difference of first-stage variable values.
Gap	is the relative gap between <i>Best Value</i> and <i>Bound</i> .
Wall Time	gives the wall time passed since invoking the program.
CPU Time	gives the CPU time passed since invoking the program.
Father	gives the number of the father node.

The pruning of nodes is indicated by the entry *cutoff* in the row ‘Objective’. The row ‘Heuristic’ may also contain the entries *n-stop* (inferiority, evaluation stopped at n-th scenario) or *multiple* (a solution has been evaluated previously). When more than one solution is evaluated in a branch-and-bound step, the entry indicates the status of the last solution.

An asterisk in the first position indicates that a new best value was found. In the example a list of heuristics is used in every node, the lines beginning with an asterisk indicate which heuristic found a new best solution. If the scenario solutions of a node fulfil the nonanticipativity and yield a new best upper bound, in the output line for the node an asterisk is prepended, too.

## 8.2 Output in the file *sip.out*

All output files will be placed in the subdirectory *sipout* of the current directory. This directory will be created if it does not exist. A further run of **ddsip** overwrites the existing output files. The output on screen is also directed to the file *more.out* in case `OUTLEV > 51`. In addition, this file contains the parameters read and some information on the solution:

- The value of *Status* indicates the solution status:
  - 1 the process was terminated by the user,
  - 1 the node limit has been reached,
  - 2 the gap (absolute or relative) has been reached,
  - 3 the time limit has been reached,
  - 4 the maximal dispersion, i.e. the maximal difference of the first-stage components within all remaining front nodes, was less then the parameter `NULLDISP` (null dispersion),
  - 5 the whole branching tree was backtracked.
- *Time* is the total CPU time needed by **ddsip**.
- *Upper bounds* is the number of evaluated upper bounds.
- *Tree depth* is the depth of the branch-and-bound tree.
- *Nodes* is the total number of nodes.
- *EEV* is the solution of the EEV problem according to Birge and Louveaux (1997).
- *VSS* is the value of stochastic programming according to Birge and Louveaux (1997).
- *Expectedvalue* is the expected value.
- *Riskmeasure* is the value of the used risk measure.

## 8.3 Other output files

The number of additional output files and their content is ruled via the parameters `OUTFIL` and `OUTLEV`. Values of `OUTFIL` greater than 1 are intended for debugging purposes.

- The file *more.out* contains more or less information depending on the parameter `OUTLEV`, e.g. the subproblem solutions, the branch-and-bound tree, and the objective function composition.
- The file *solution.out* contains information on the solution with the optimal first-stage solution among them. It is not written when `OUTFILES` is set to 0.
- If `OUTFILES` is greater than 1, then in addition the following files are written to *sipout*:

- The files *rhs.out*, *cost.out*, *matrix.out*, and *model.lp.gz* offer a check for a correct reading of the scenario files and the model file, respectively.
- The files *risk.lp.gz* and *eev.lp.gz* document the changes made during defining the risk model and solving the expected value problem, respectively.
- If OUTFILES equals 3, then in addition the following files are written to sipout:
  - The single scenario files *lb\_sc\_\*.n0\_gc\*.lp.gz* and *ub\_sc\_\*.n0\_gc\*.lp.gz* solved in the root node for lower and upper bounds, respectively.
- If OUTFILES equals 4, then in addition the following files are written to sipout:
  - The single scenario files *lb\_sc\_\*.n\*\_gc\*.lp.gz* and *ub\_sc\_\*.n\*\_gc\*.lp.gz* solved in all nodes for lower and upper bounds.
- If OUTFILES equals 5, then the files *lb\_sc\_\*.n\*\_gc\*.lp.gz* and *ub\_sc\_\*.n\*\_gc\*.lp.gz* are not written, but instead
  - The single scenario files *lb\_sc\_\*.n\*.sav.gz* solved in all nodes for lower bounds plus the used MIP starts as *lb\_sc\_\*.n\*.mst.gz* and the CPLEX settings as *lb\_sc\_\*.n\*\_1.prm* and *lb\_sc\_\*.n\*\_2.prm*. In the case of conic bundle use these are overwritten in every dual step, just the ones of the final dual step of each node (or the one causing a program crash) remain.
- If OUTFILES equals 6, as in the case 5 the single scenario problem files are not written as lp.gz, but instead
  - The single scenario files *lb\_sc\_\*.n\*\_gc\*.sav.gz* solved in all nodes for lower bounds plus the used MIP starts as *lb\_sc\_\*.n\*\_gc\*.mst.gz* and the CPLEX settings as *lb\_sc\_\*.n\*\_1.prm* and *lb\_sc\_\*.n\*\_2.prm*. This is the same amount of output files as with the setting 5 in case no conic bundle is used. In the case of conic bundle use these are written in every dual step, increasing the value after gc (global counter) with each step.

## 9 License and bugs

### 9.1 License

The program is free software. It is distributed under the GNU General Public License as published by the Free Software Foundation, see GNU Project (2004).

### 9.2 Bug report

Please report all bugs via email to [ralf.gollmer@uni-due.de](mailto:ralf.gollmer@uni-due.de). If possible, provide the input files to facilitate reproduction of the error.

## A Sample specs file

\* Parameters to specify the two-stage model

FIRSTCON	9	* First stage constraints
FIRSTVAR	30	* First stage variables
SECCON	280	* Second stage constraints
SECVAR	326	* Second stage variables
STOCCRHS	70	* Number of stochastic rhs elements
SCENARIOS	5	* Number of scenarios
STOCCOST	0	* Number of stochastic cost coefficients * (stoch. obj. doesn't work in combination with risk!)
STOCMAT	0	* Number of stochastic matrix entries

\* Parameters for dual decomposition procedure

OUTLEVEL	5	* Print info to <i>more.out</i>
OUTFILES	1	* Print files (models and output)
STARTINFO	0	* Use start solution/bound
PREPRO	0	* Sort scenarios
NODELIM	90	* <b>ddsip</b> node limit
TIMELIM	800	* <b>ddsip</b> time limit
HEURISTIC	99 3 4	* Heuristic to produce a feasible solution
EEVPROB	1	* Solve EEV, cf. Birge and Louveaux (1997)
ABSOLUTE	0	* Absolute duality gap
RELATIVE	0.001	* Relative duality gap
PORDER	0	* Use branching priority order
BRADIR	-1	* Branching direction
BRASTRAT	0	* Branching strategy
BOUSTRAT	1	* Bounding strategy
EPSILON	1e-4	* Epsilon used for branching on real variables
RELAXL	1	* Relaxation level (1= relax first-stage variables * in scenario problems)
LOGFREQ	1	* Output log frequency
INTFIRST	1	* Branching on integers first
ACCURACY	1e-8	* Accuracy used to compare real values
NULLDISP	0	* Tolerance level for null-dispersion * (used together with ACCURACY!)
QUANTILES	10	* Number of Quantiles

\* Risk model

RISKMOD	1	* Risk model
TARGET	1e+3	* Target for target measures
WEIGHT	1e+0	* Weight of the risk term
RISKBM	1e+4	* Big M used in modeling risk models 2 and 5
BRAETA	0	* Branch on auxiliary variable eta (models 4 and 5)?

\* CPLEX Parameters (See CPLEX manual CPLEX (2016))

CPLEXBEGIN

1035 1 \* Output on screen indicator

2009 1e-7 \* Relative gap

CPLEXEEV

2009 0.1 \* Relative gap for EEV

CPLEXLB

2058 3 \* MIP emphasis

2031 10 \* heuristic frequency

2009 0.002 \* Relative Gap

1039 900 \* Time limit

CPLEX2LB

2058 4 \* MIP emphasis

2031 5 \* heuristic frequency

2061 20 \* RINS neighbourhood search heuristic frequency

2009 1e-6 \* Relative Gap

1039 350 \* Time limit

CPLEXUB

2058 3 \* MIP emphasis

2031 10 \* heuristic frequency

1039 750 \* Time limit

2009 0.001 \* Relative Gap

CPLEX2UB

2058 4 \* MIP emphasis

2031 15 \* heuristic frequency

2061 60 \* RINS neighbourhood search heuristic frequency

1039 300 \* Time limit

CPLEXDUAL

2009 1e-4 \* Relative gap for Lagrangian Dual

1035 0 \* Output on screen indicator for Lagrangian dual

CPLEXEND

\* Parameters specifying the use of ConicBundle

CBFREQ 1000 \* Conic Bundle in every i-th node

CBITLIM 20 \* Limit for number of descent steps in Conic Bundle

CBTOTIT 1000 \* Limit for total number of CB iterations (incl. nullsteps)

NONANT 2 \* Nonanticipativity representation

CBPRINT 1 \* Output level of Conic Bundle

CBBUNS 200 \* Maximal bundle size

CBMAXS 10 \* Maximal number of subgradients

CBWEIGHT 5. \* Conic Bundle initial weight

Specification file

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