



# Parallel Algorithms

EECS 221: Intro to High-Performance Computing

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April 20, 2017

# Inclusive Prefix-Sums (Scan)

- ▶ The prefix-sums operation takes a binary associative operator  $\oplus$ , and an array of  $n$  elements

$[x_0, x_1, \dots, x_{n-1}]$

and returns the array

$[x_0, (x_0 \oplus x_1), \dots, (x_0 \oplus x_1 \oplus \dots \oplus x_{n-1})]$

- ▶ Example: If  $\oplus$  is addition

3	1	7	0	4	1	6	3
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# Inclusive Prefix-Sums (Scan)

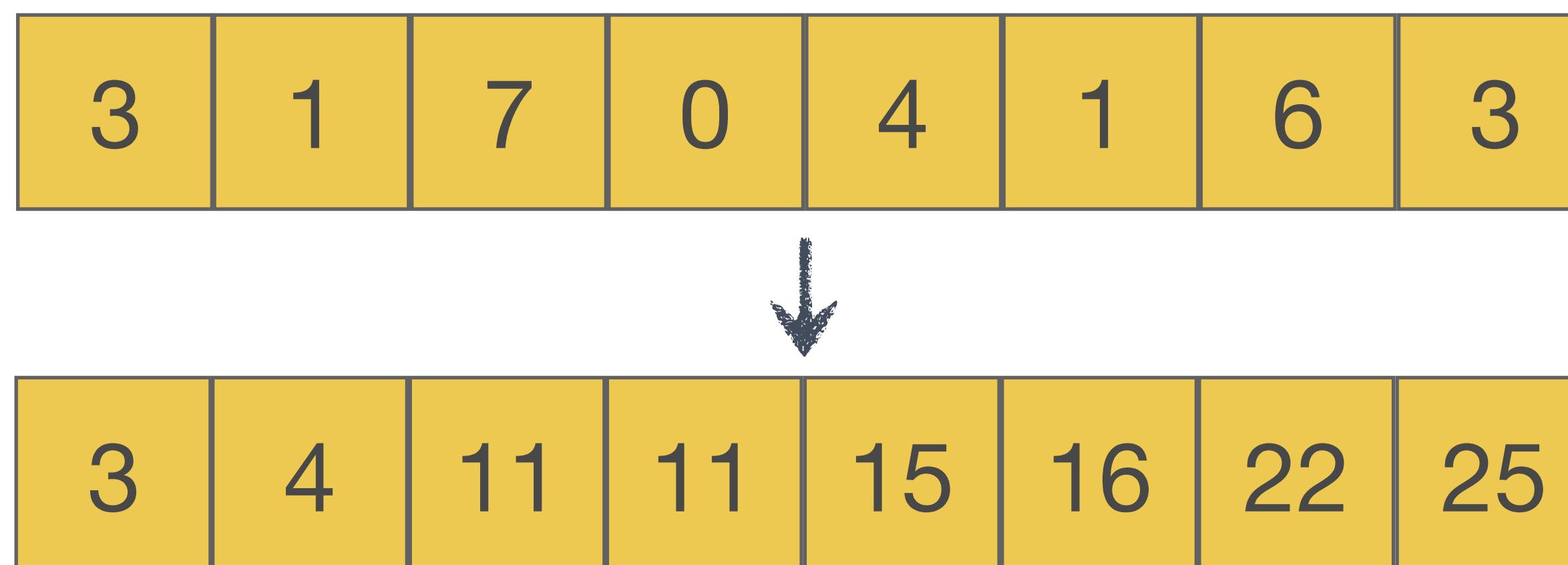
- ▶ The prefix-sums operation takes a binary associative operator  $\oplus$ , and an array of  $n$  elements

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# Exclusive Prefix-Sums (Scan)

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$[x_0, x_1, \dots, x_{n-1}]$

and returns the array

$[0, x_0, (x_0 \oplus x_1), \dots, (x_0 \oplus x_1 \oplus \dots \oplus x_{n-2})]$

- ▶ Example: If  $\oplus$  is addition

3	1	7	0	4	1	6	3
---	---	---	---	---	---	---	---

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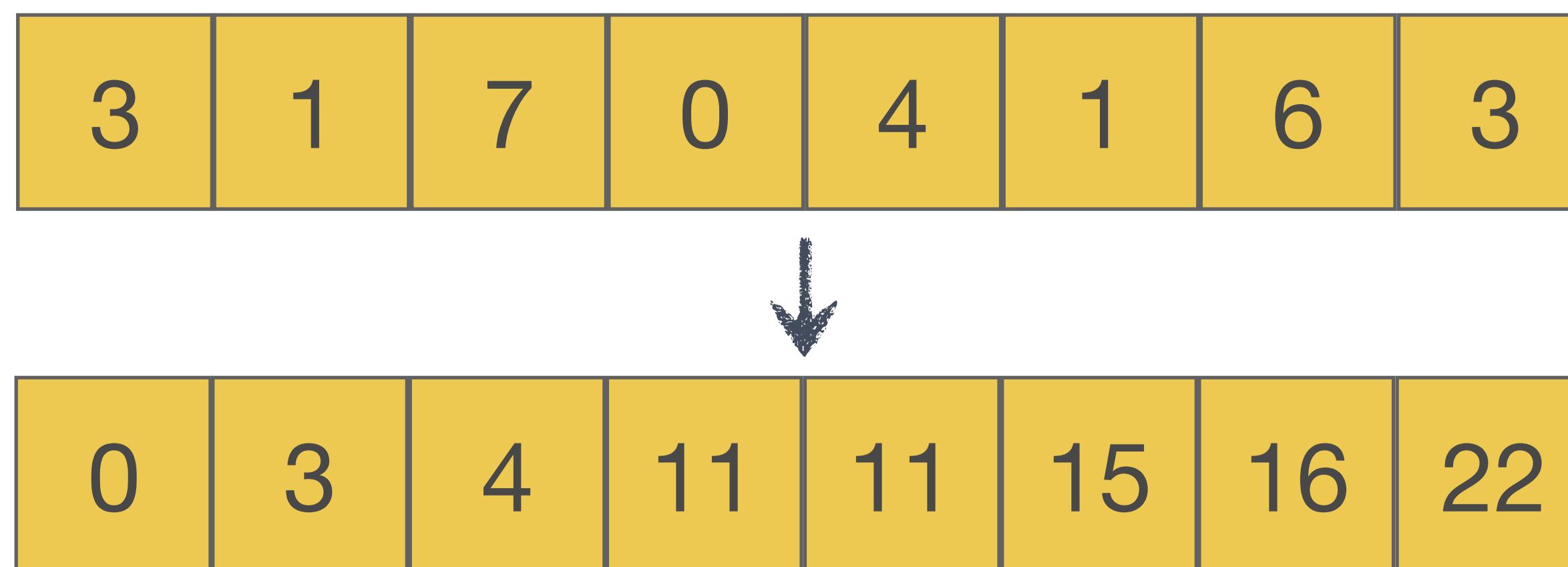
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- ▶ Example: If  $\oplus$  is addition



# Work Efficient Sequential Scan

```
for  $i \leftarrow 1$  to  $n - 1$  do  
     $A[i] \leftarrow A[i - 1] + A[i]$   
end for
```

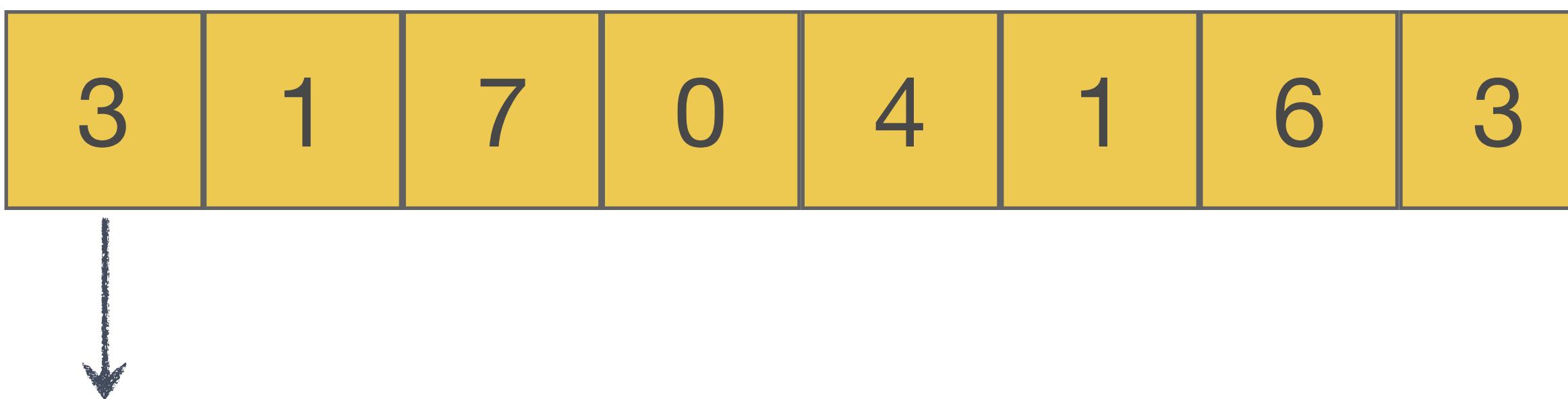
$$W_{\text{scan}}(n) = \Theta(n)$$

$$D_{\text{scan}}(n) = \Theta(n)$$

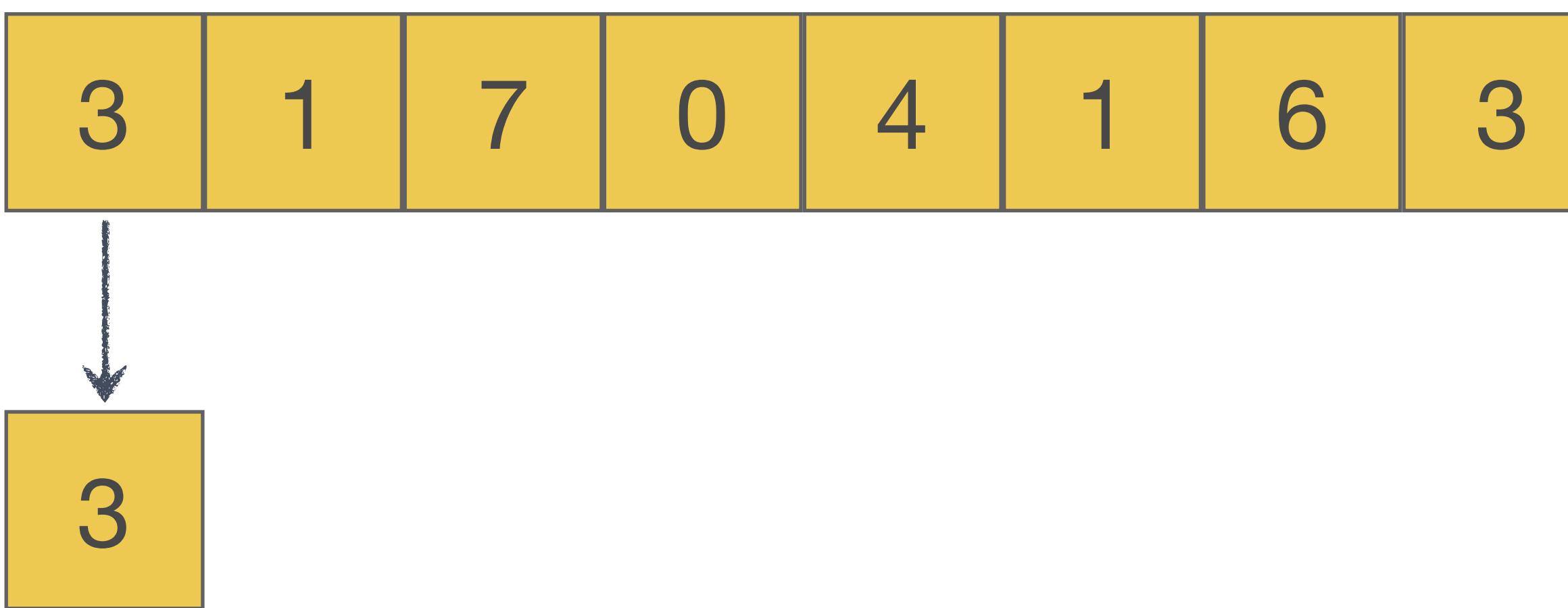
# Inclusive Parallel Scan

3	1	7	0	4	1	6	3
---	---	---	---	---	---	---	---

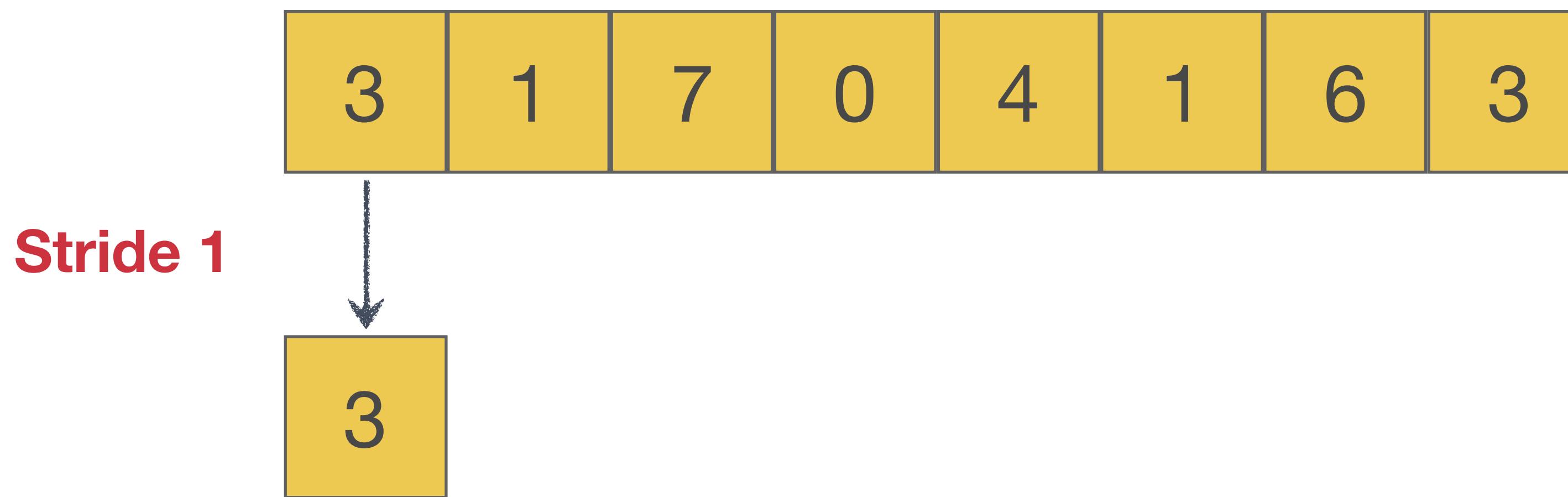
# Inclusive Parallel Scan



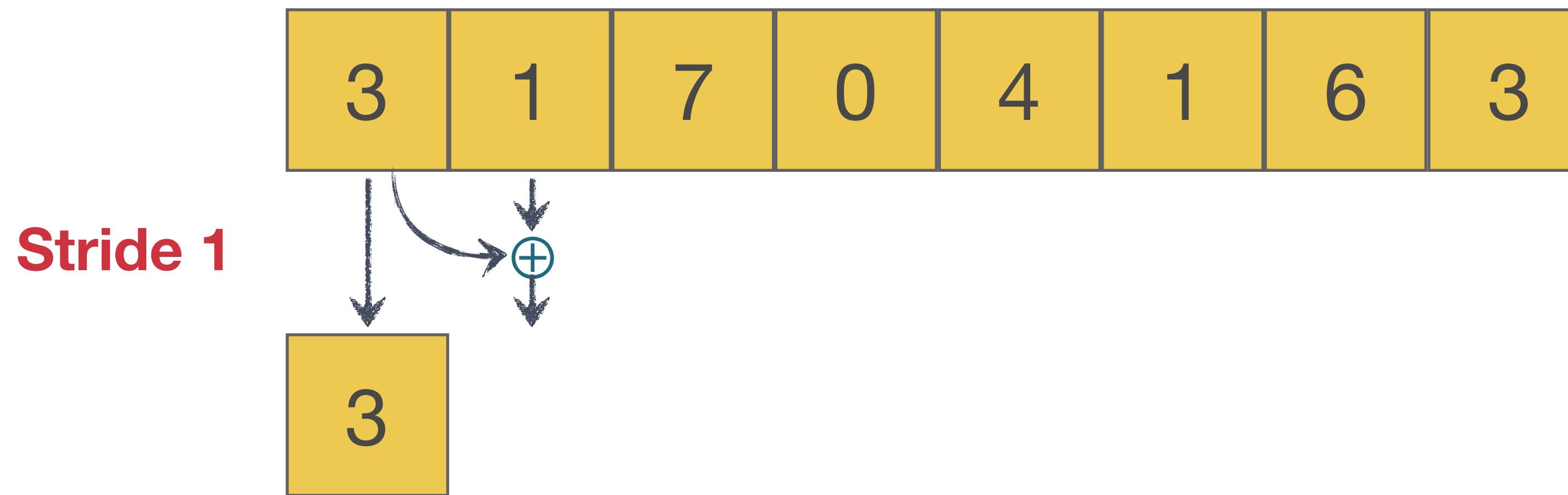
# Inclusive Parallel Scan



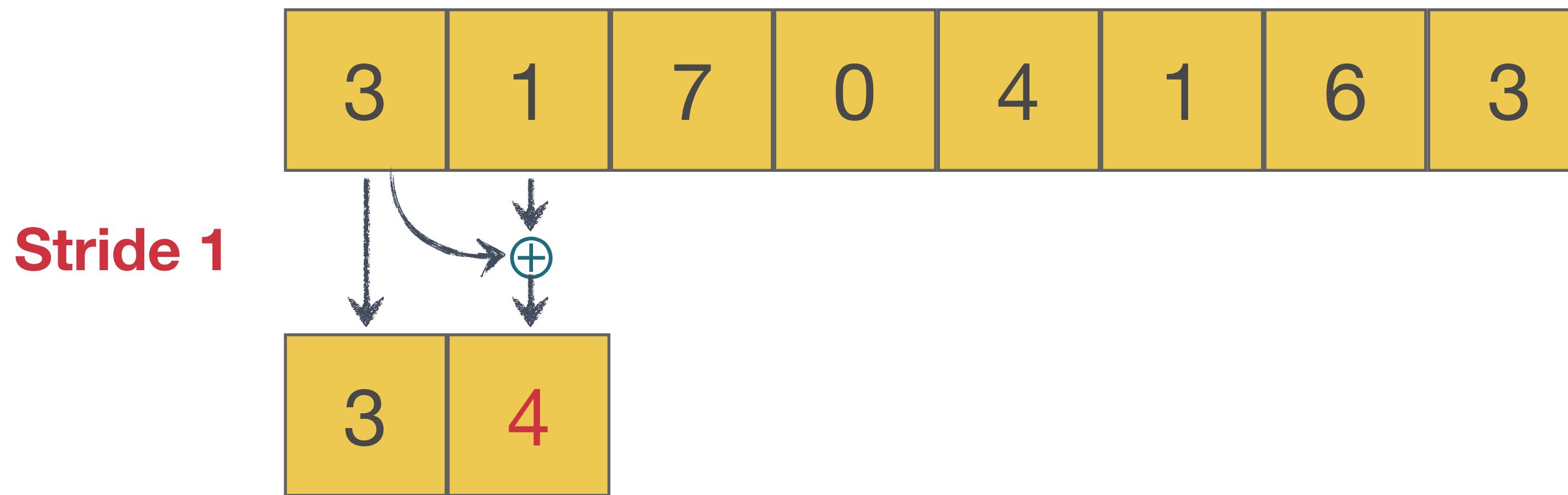
# Inclusive Parallel Scan



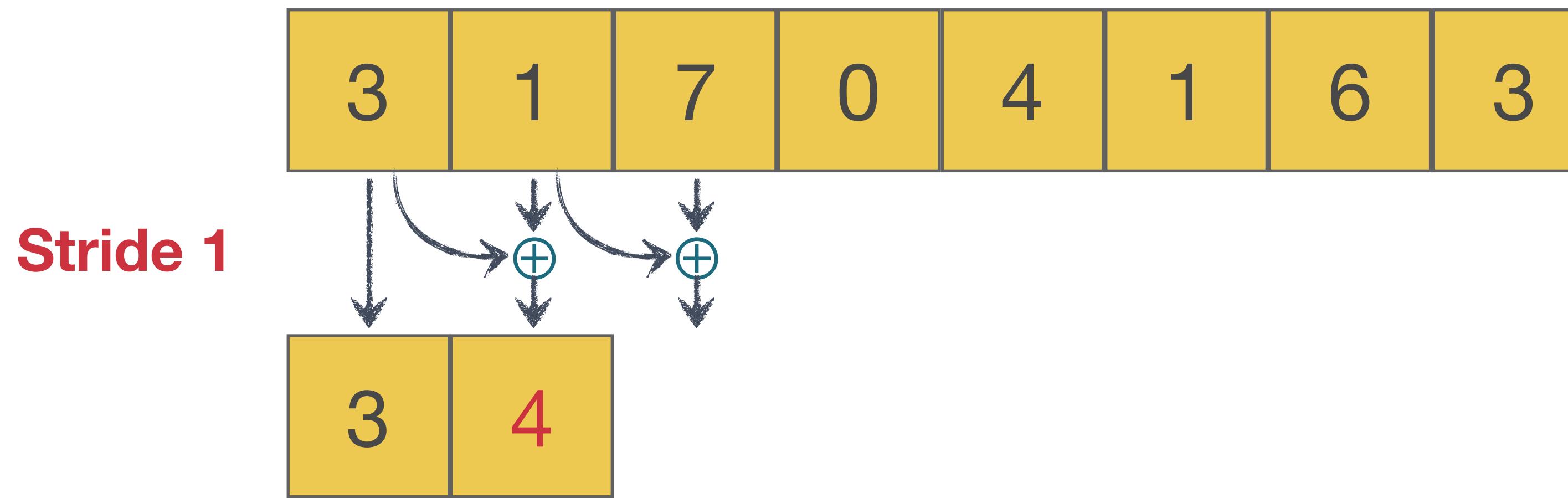
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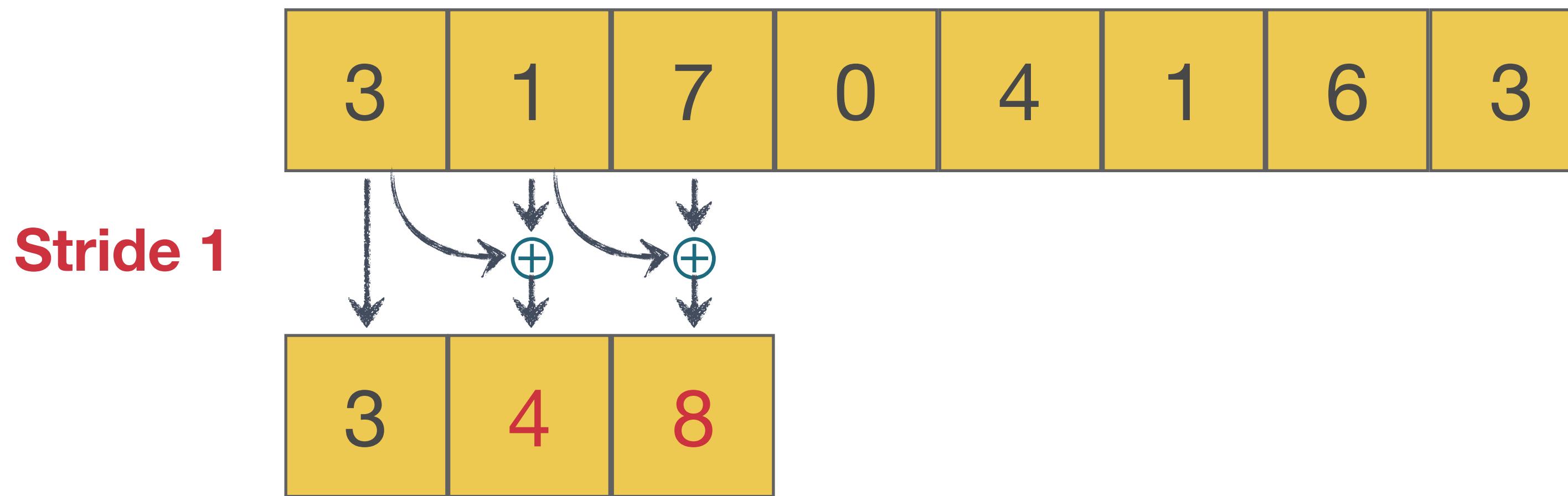
# Inclusive Parallel Scan



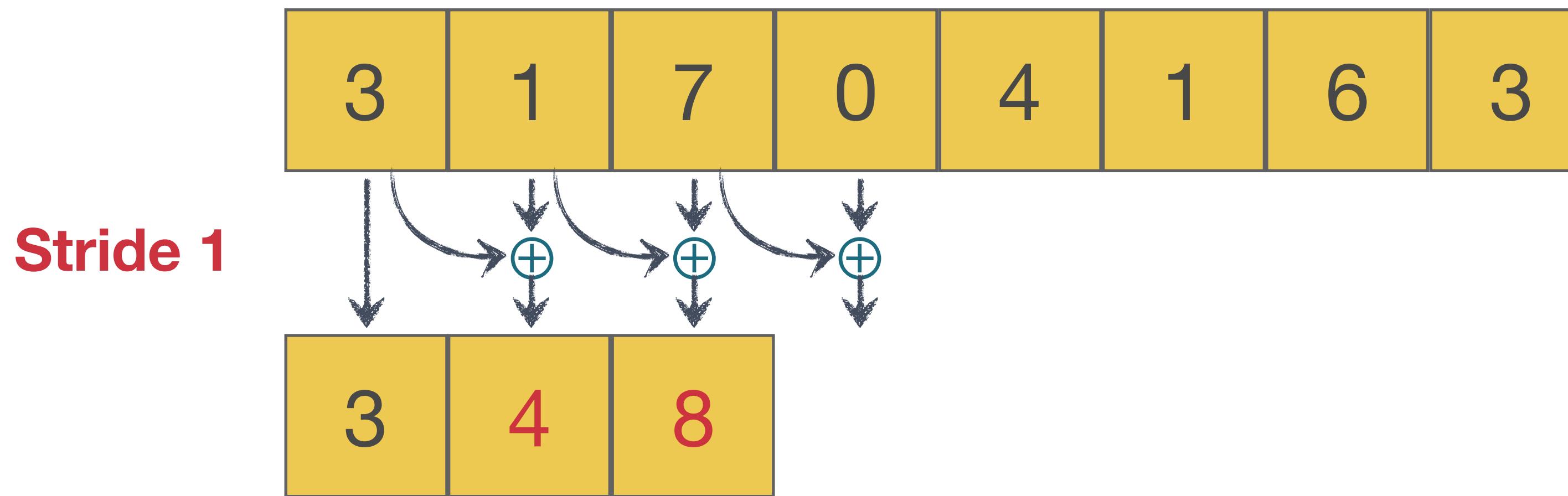
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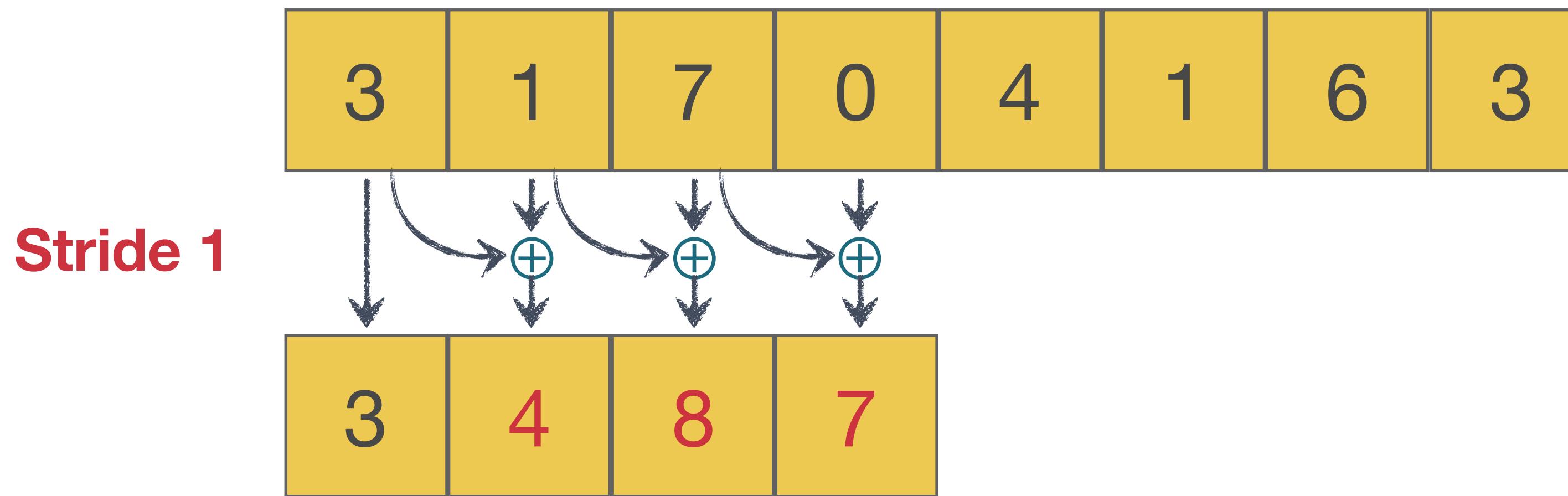
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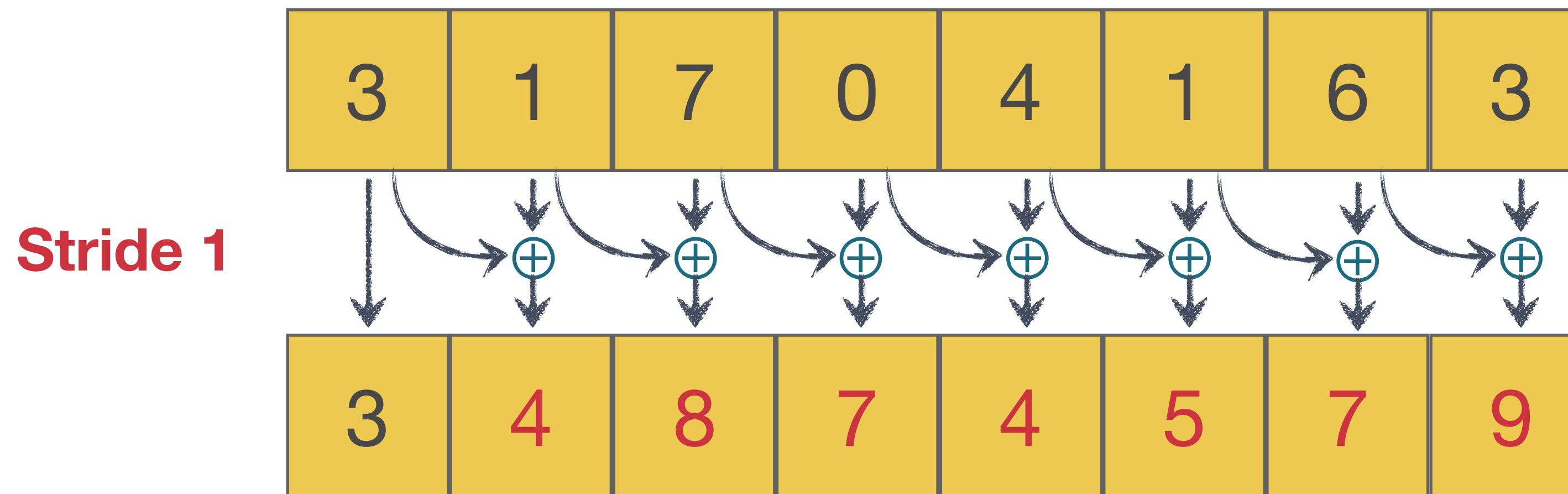
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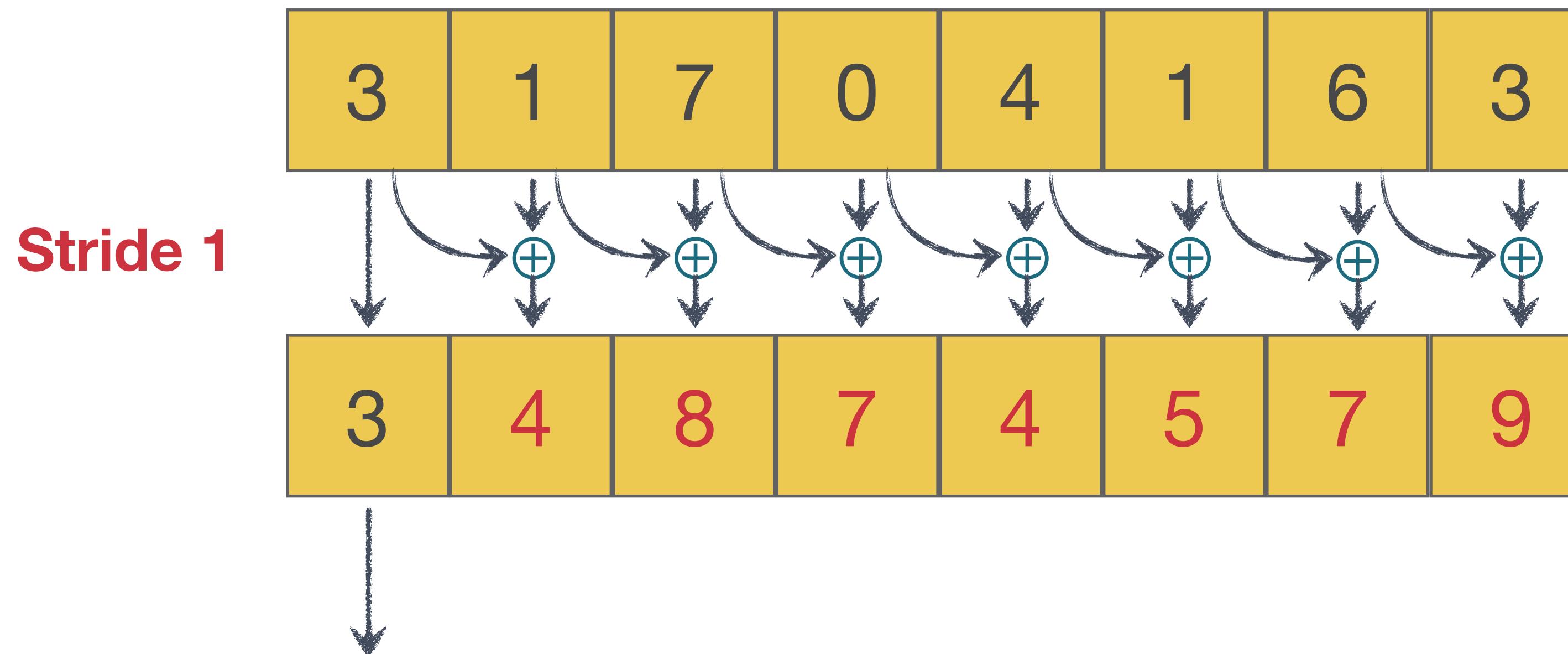
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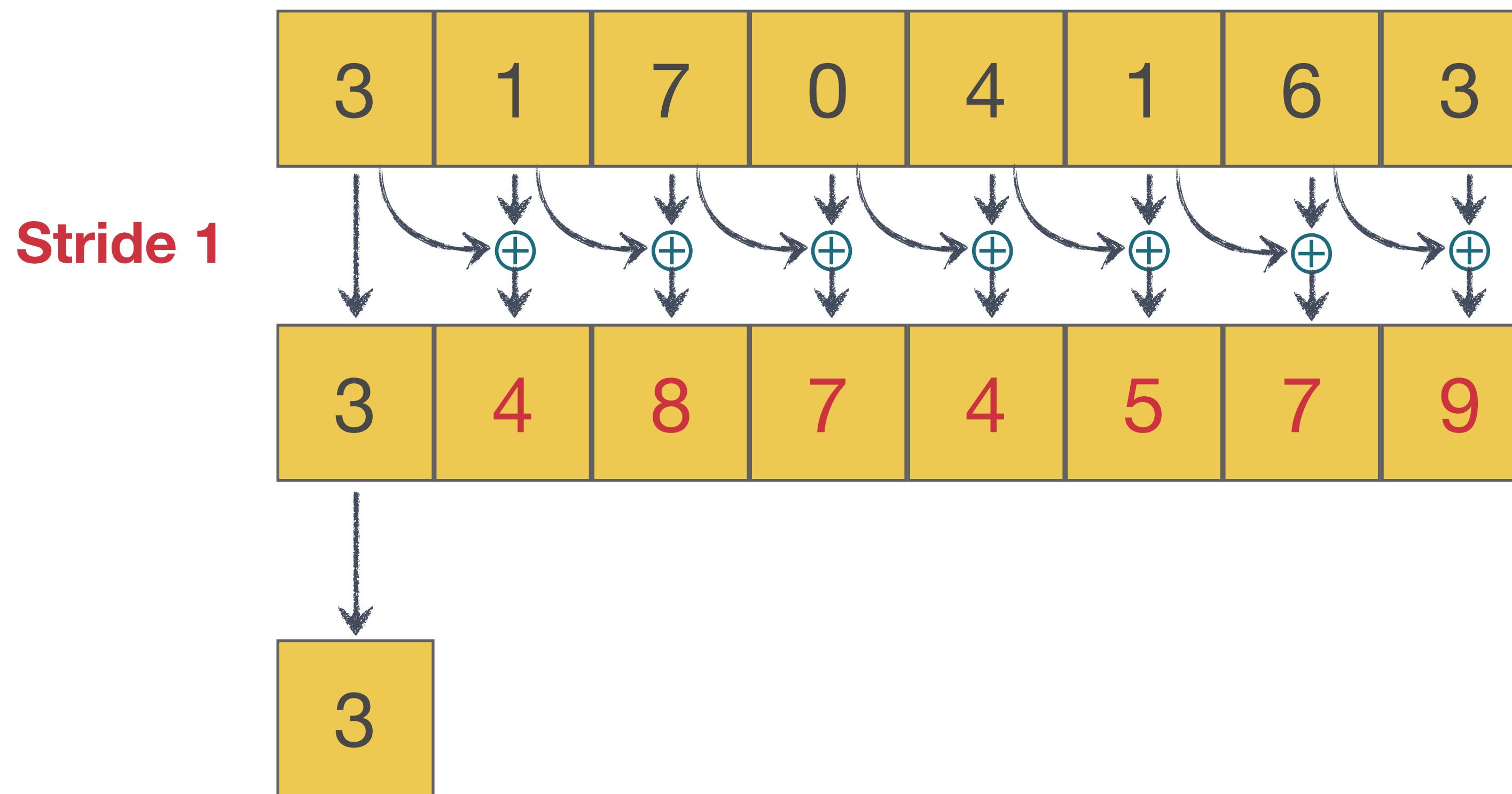
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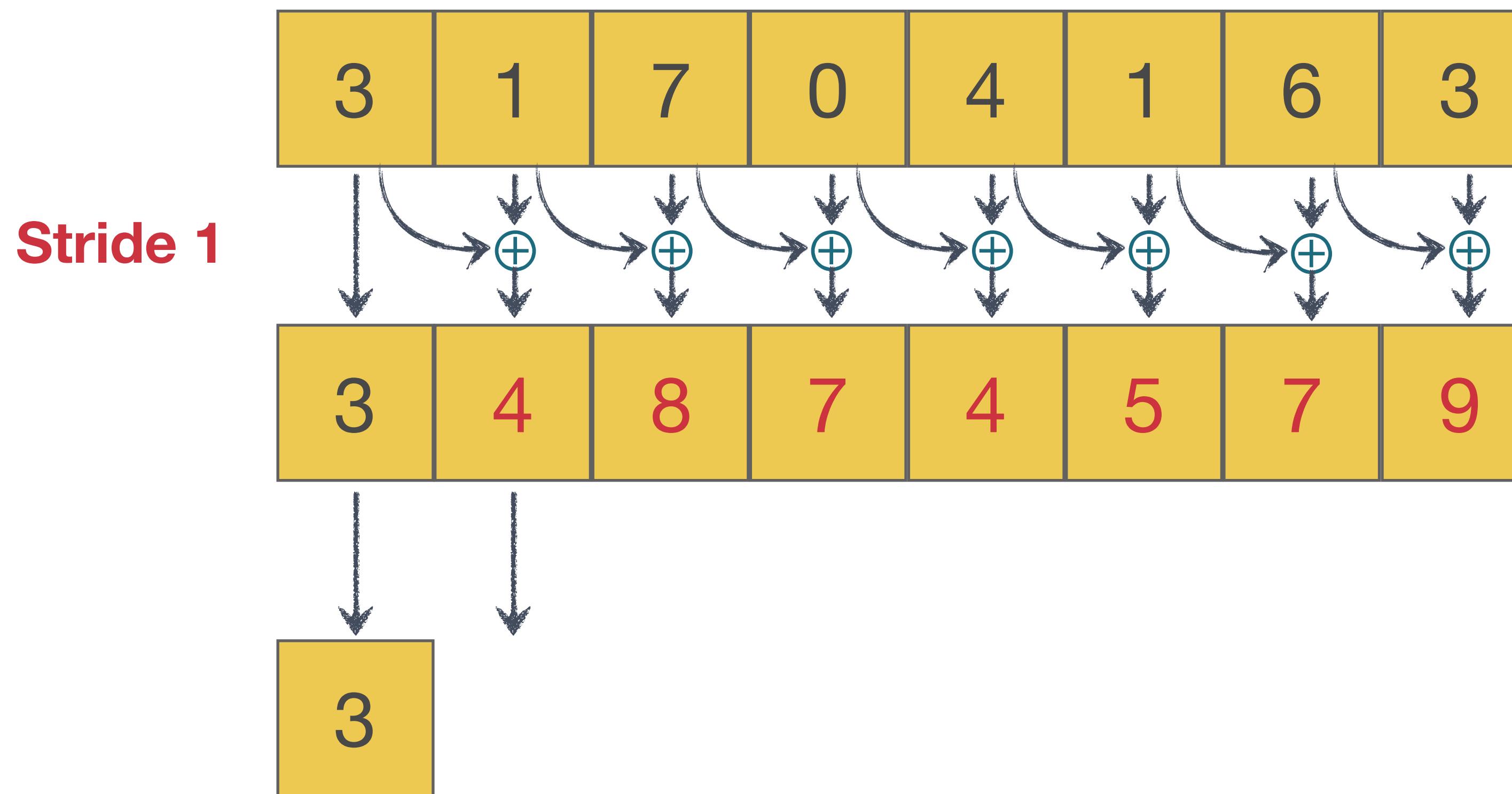
# Inclusive Parallel Scan



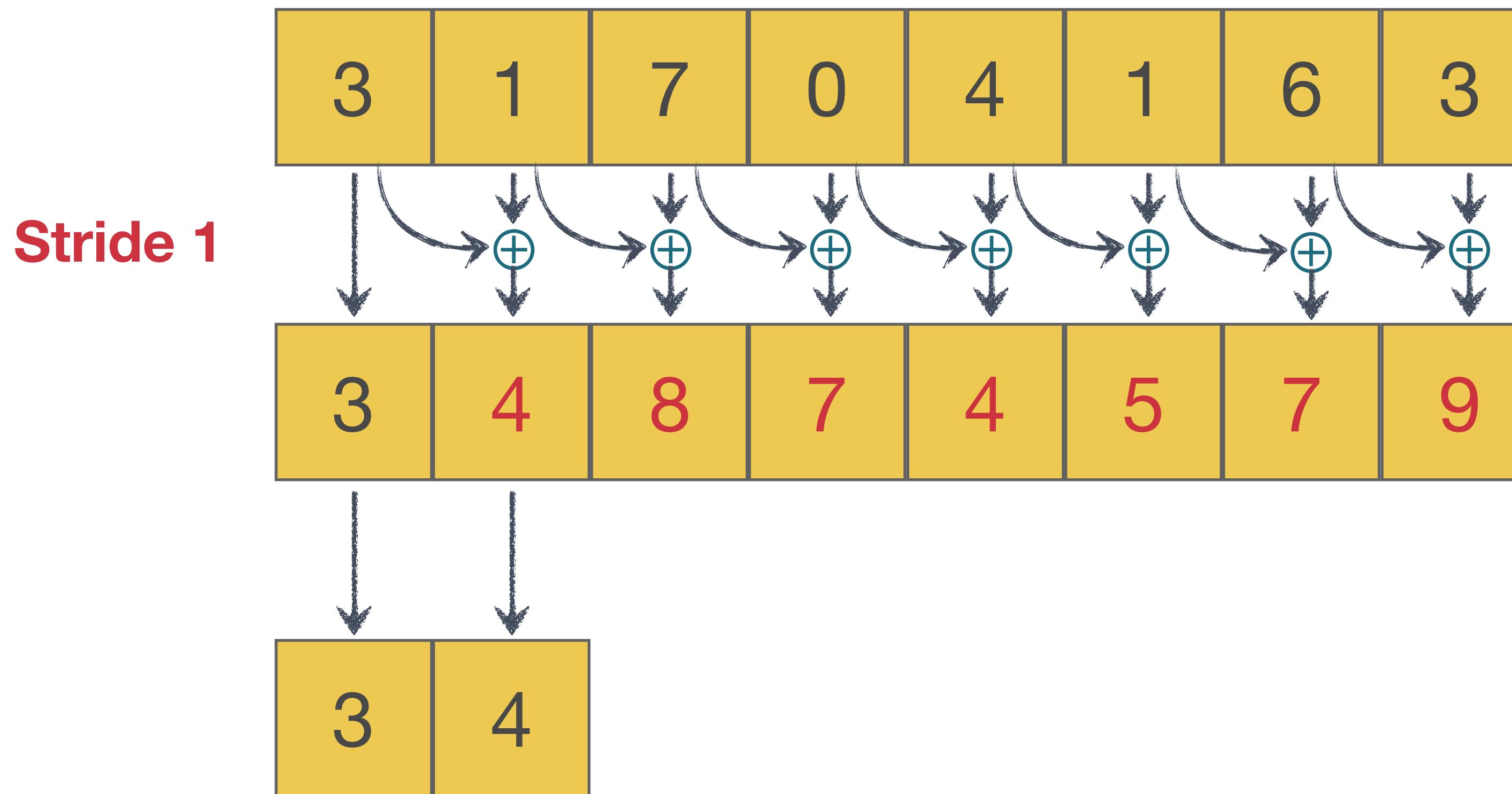
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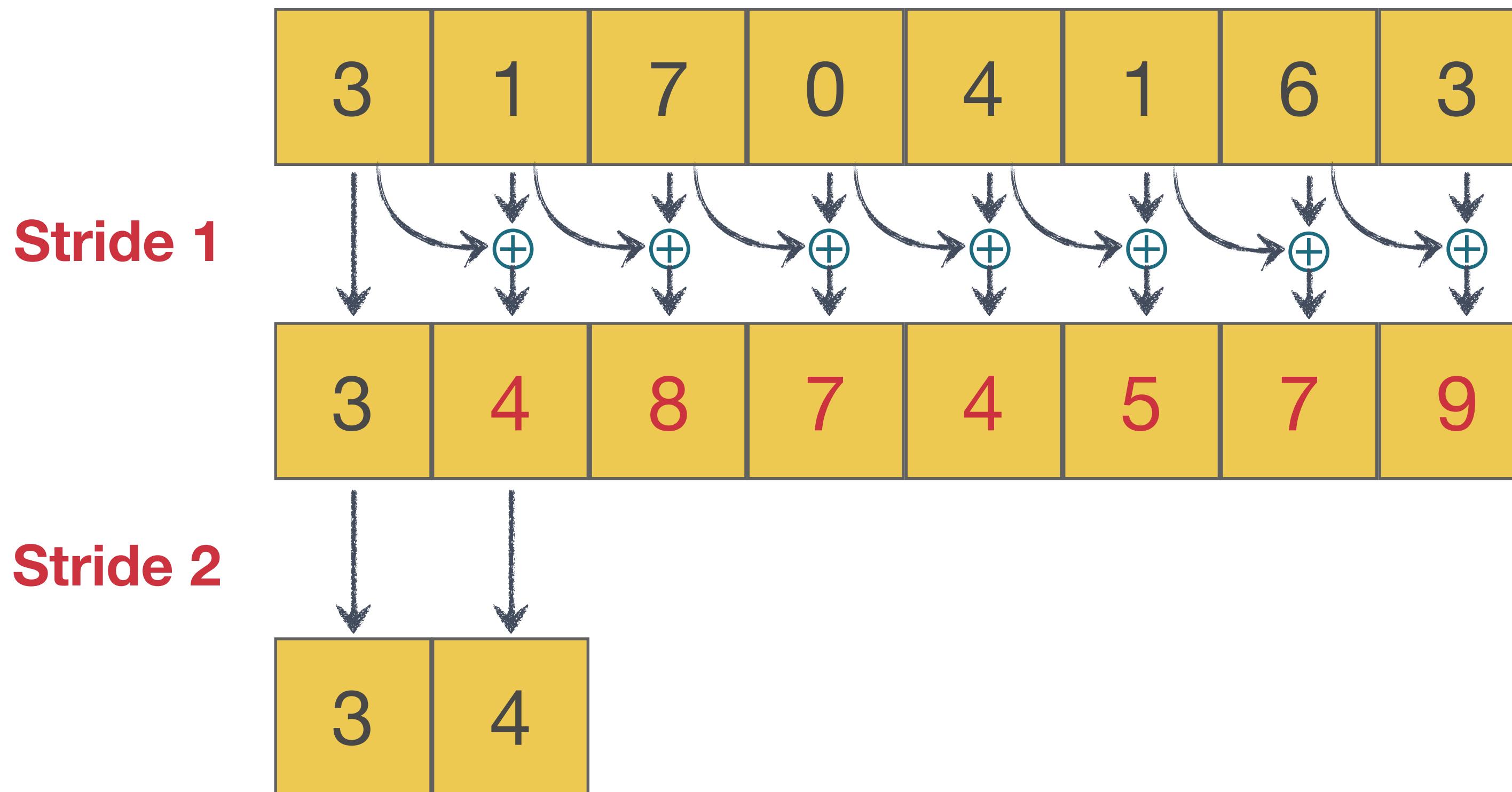
# Inclusive Parallel Scan



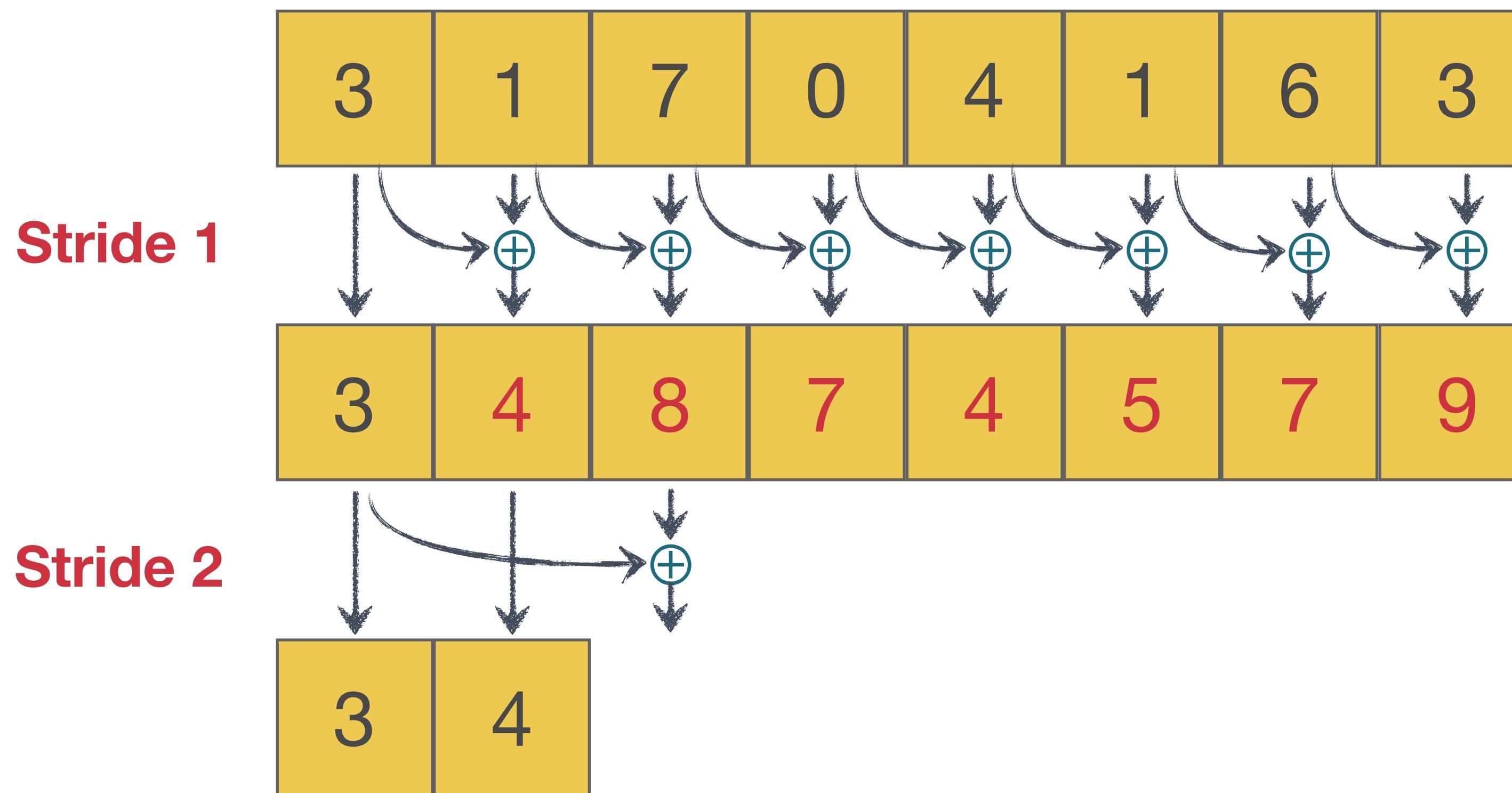
# Inclusive Parallel Scan



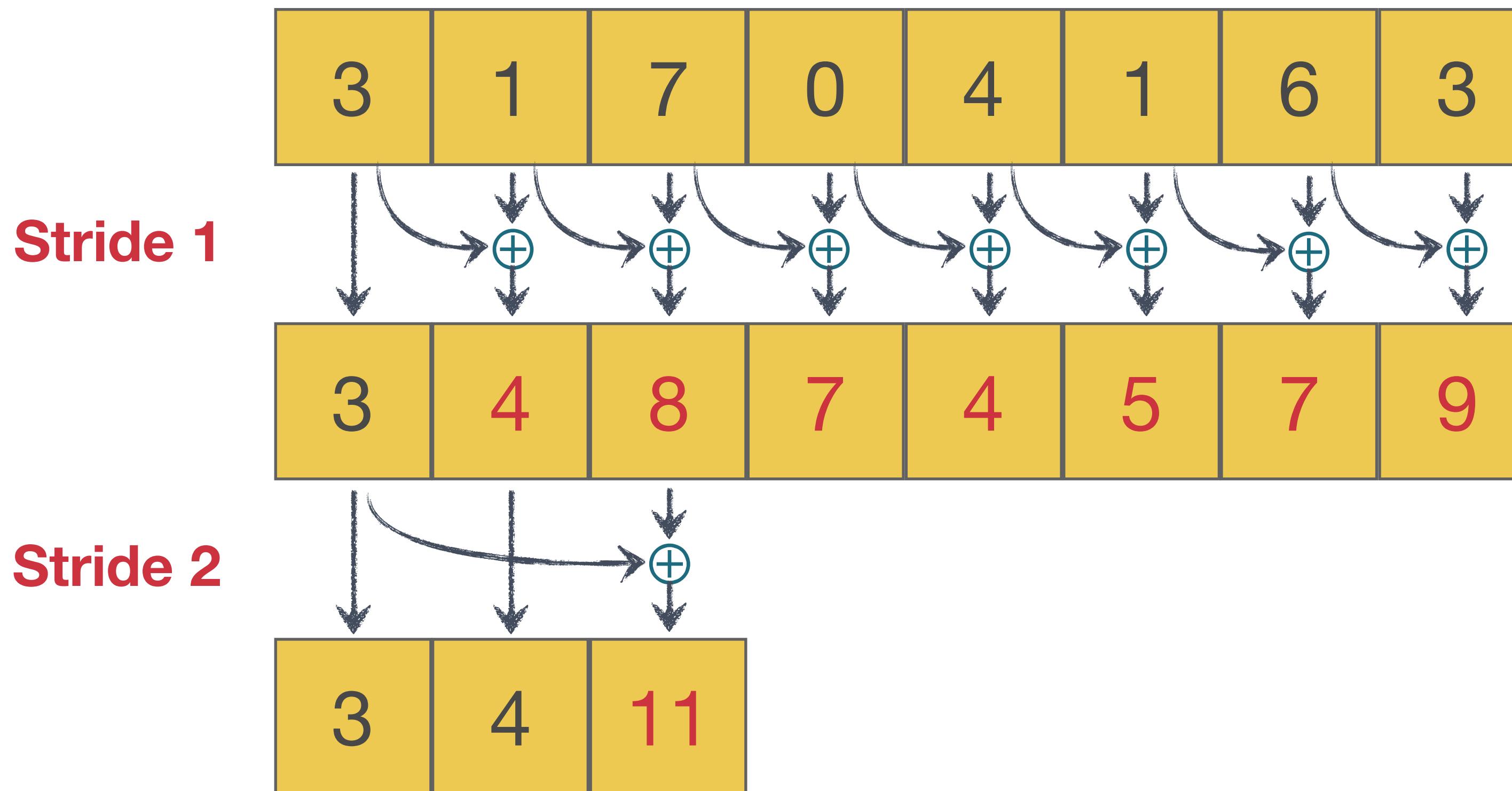
# Inclusive Parallel Scan



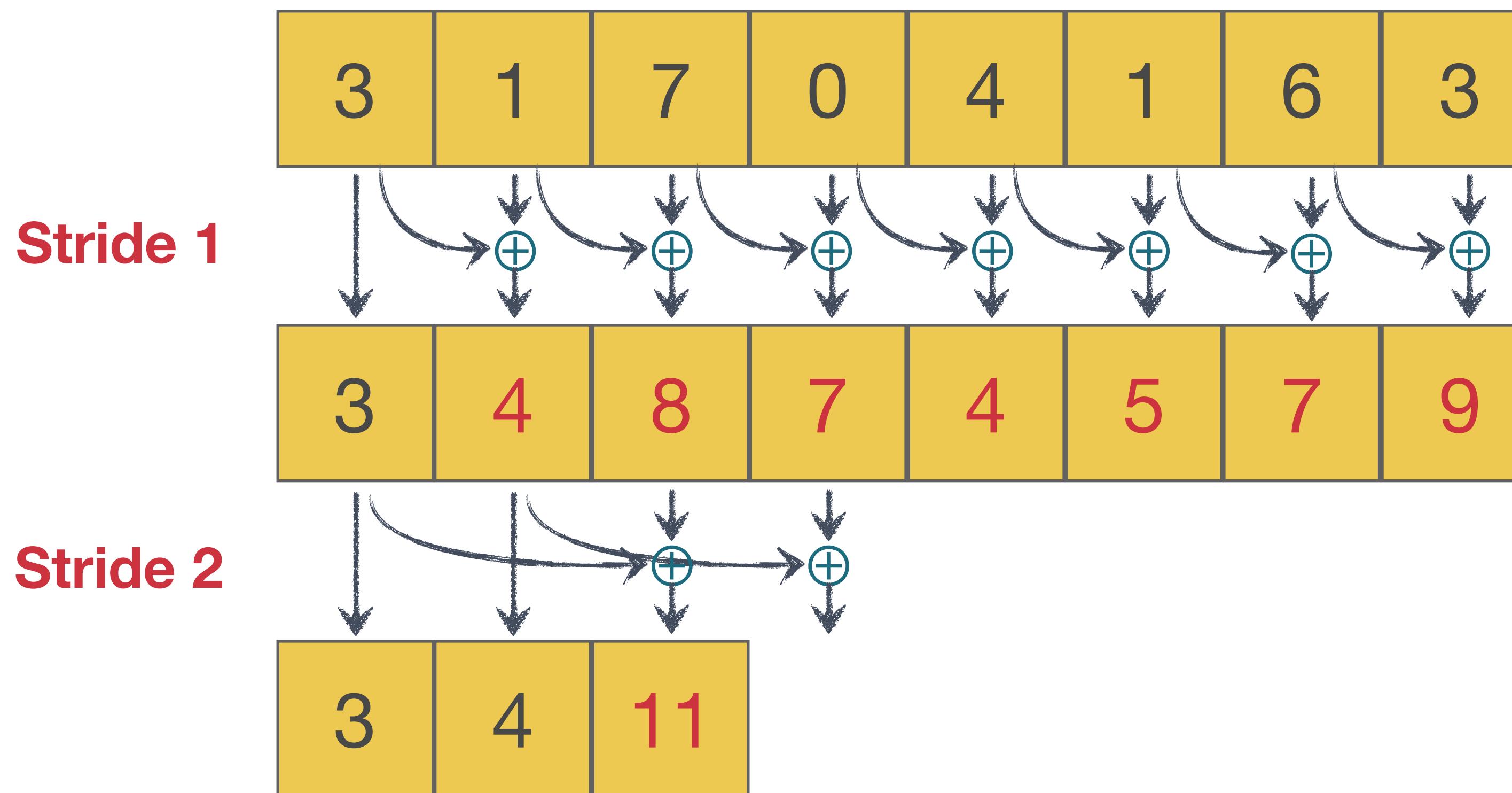
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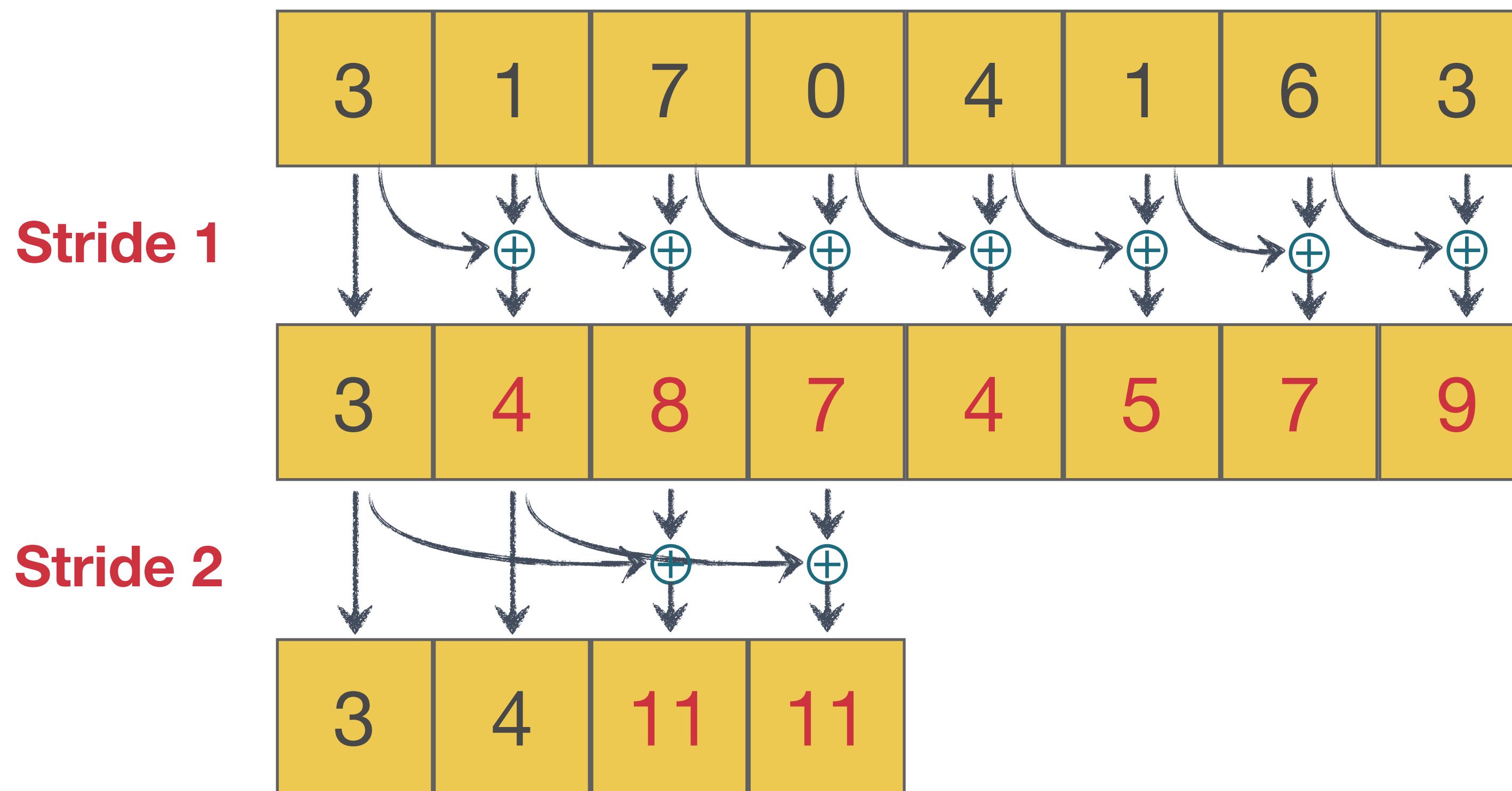
# Inclusive Parallel Scan



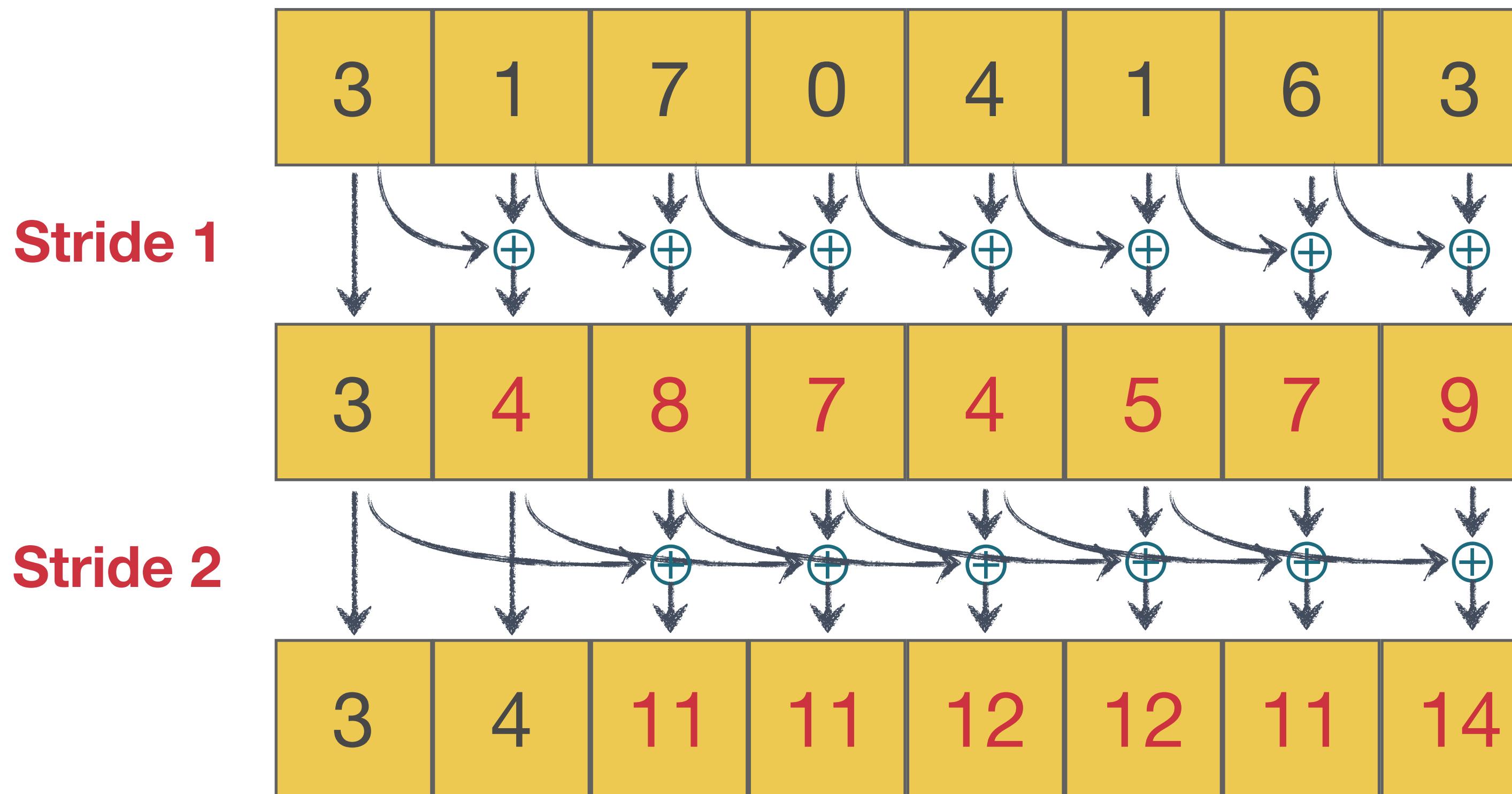
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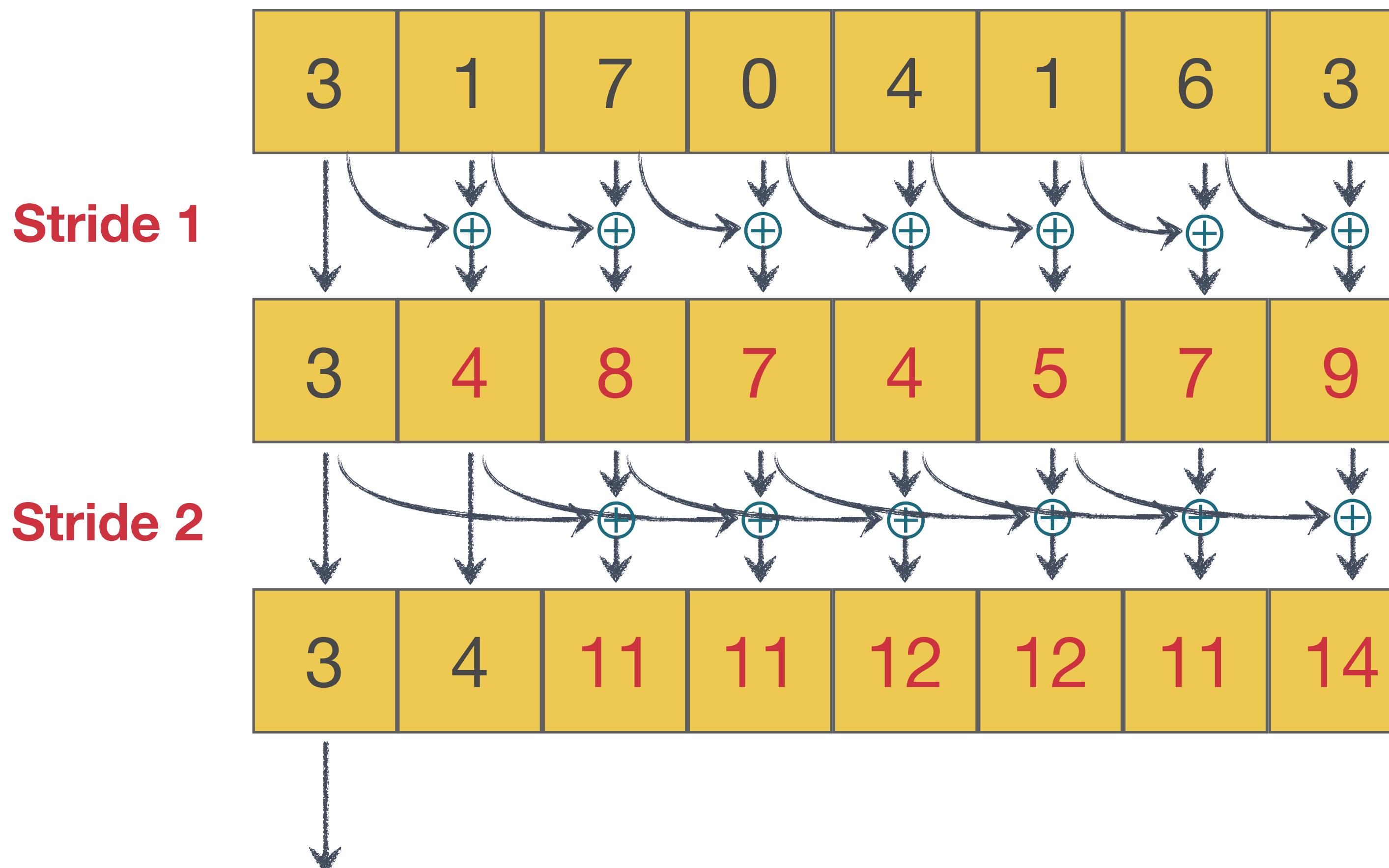
# Inclusive Parallel Scan



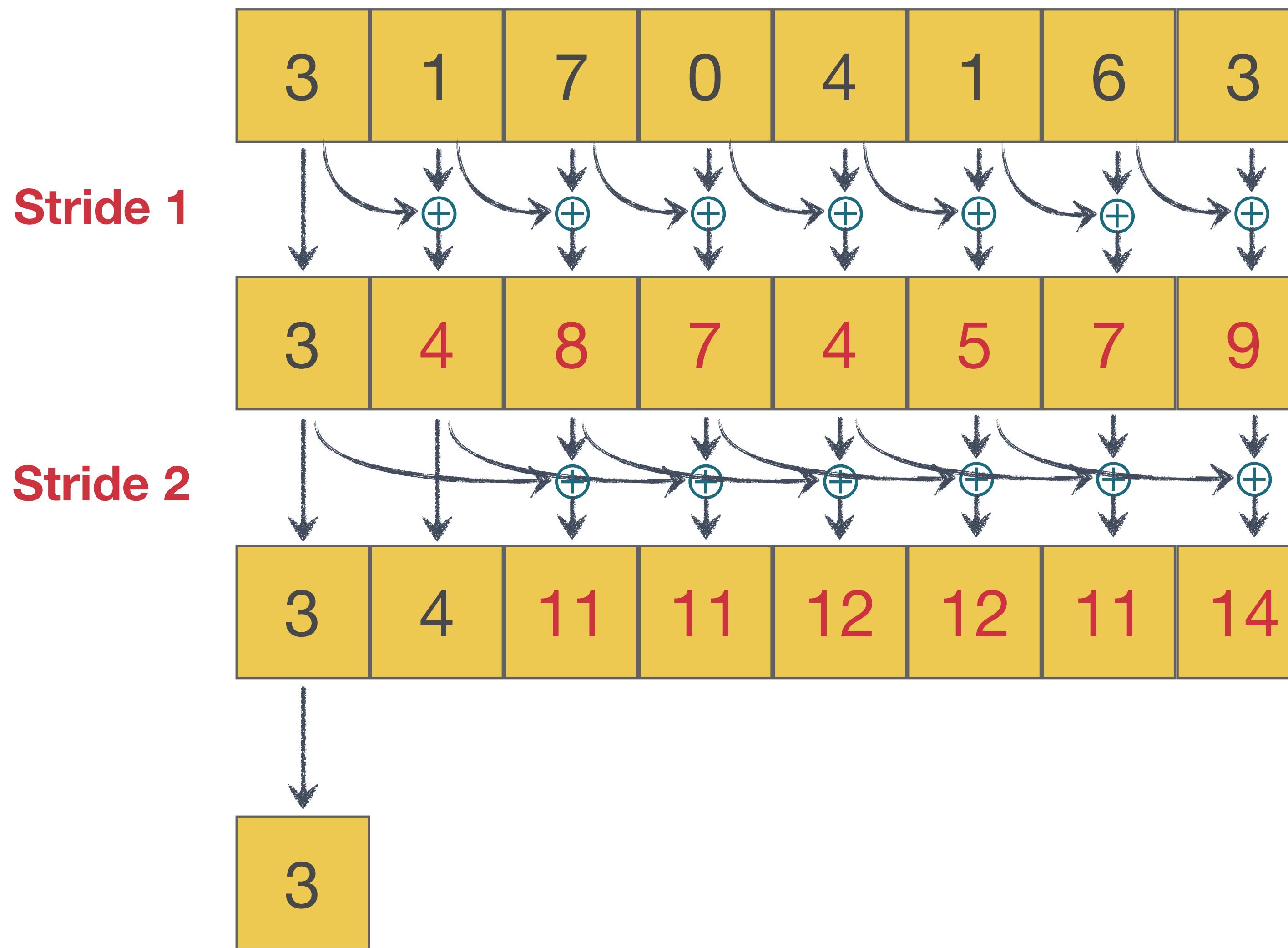
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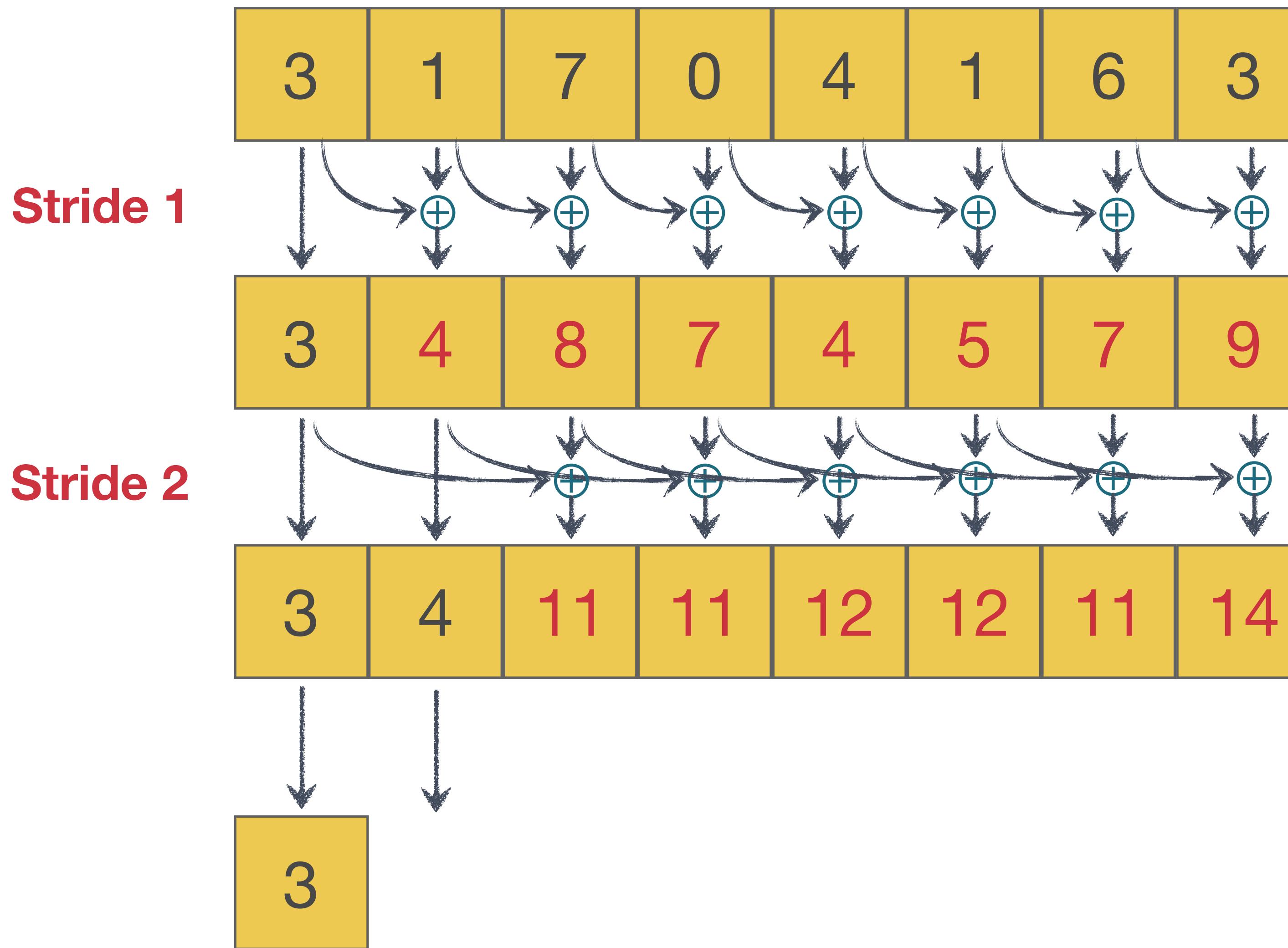
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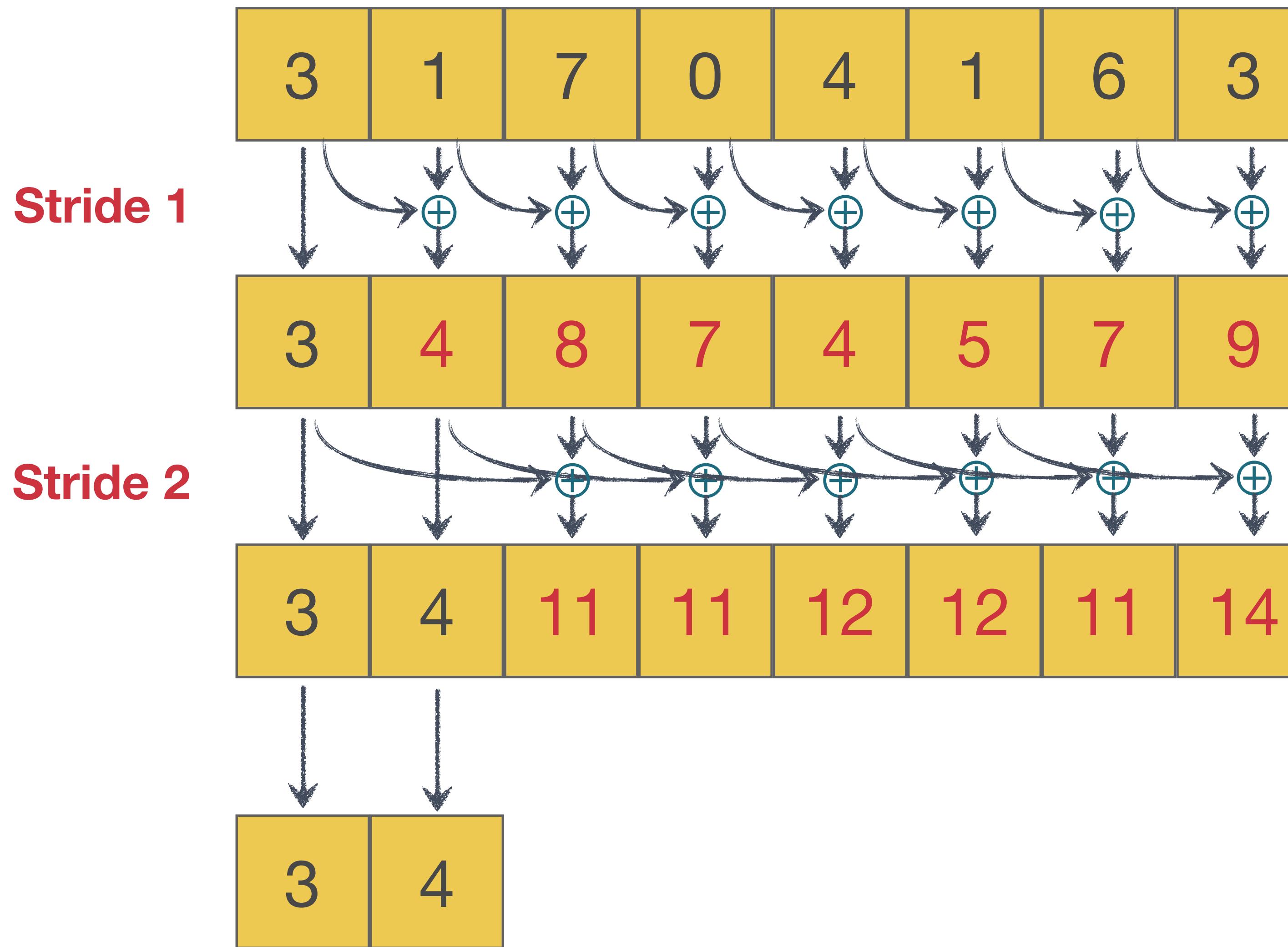
# Inclusive Parallel Scan



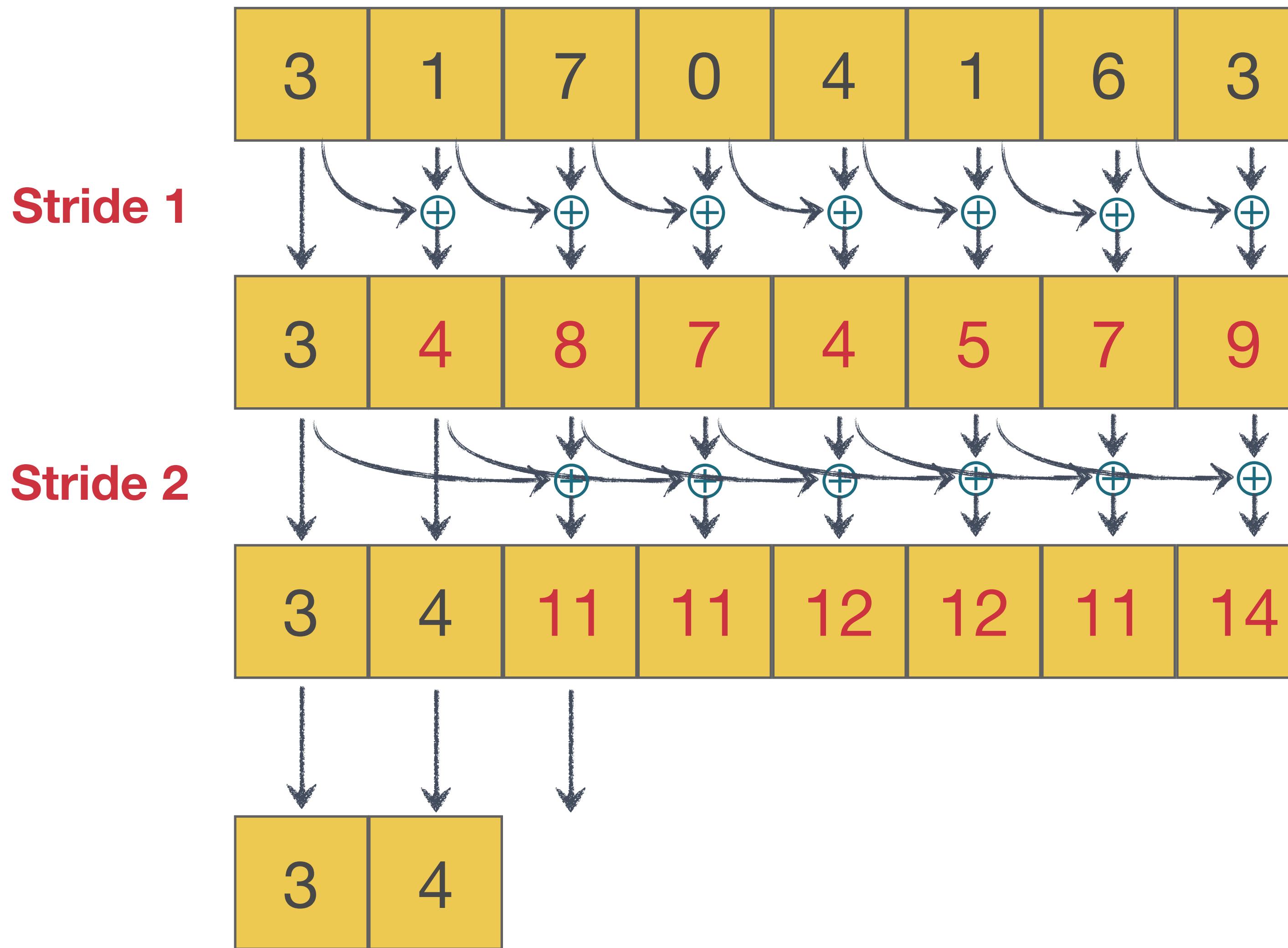
# Inclusive Parallel Scan



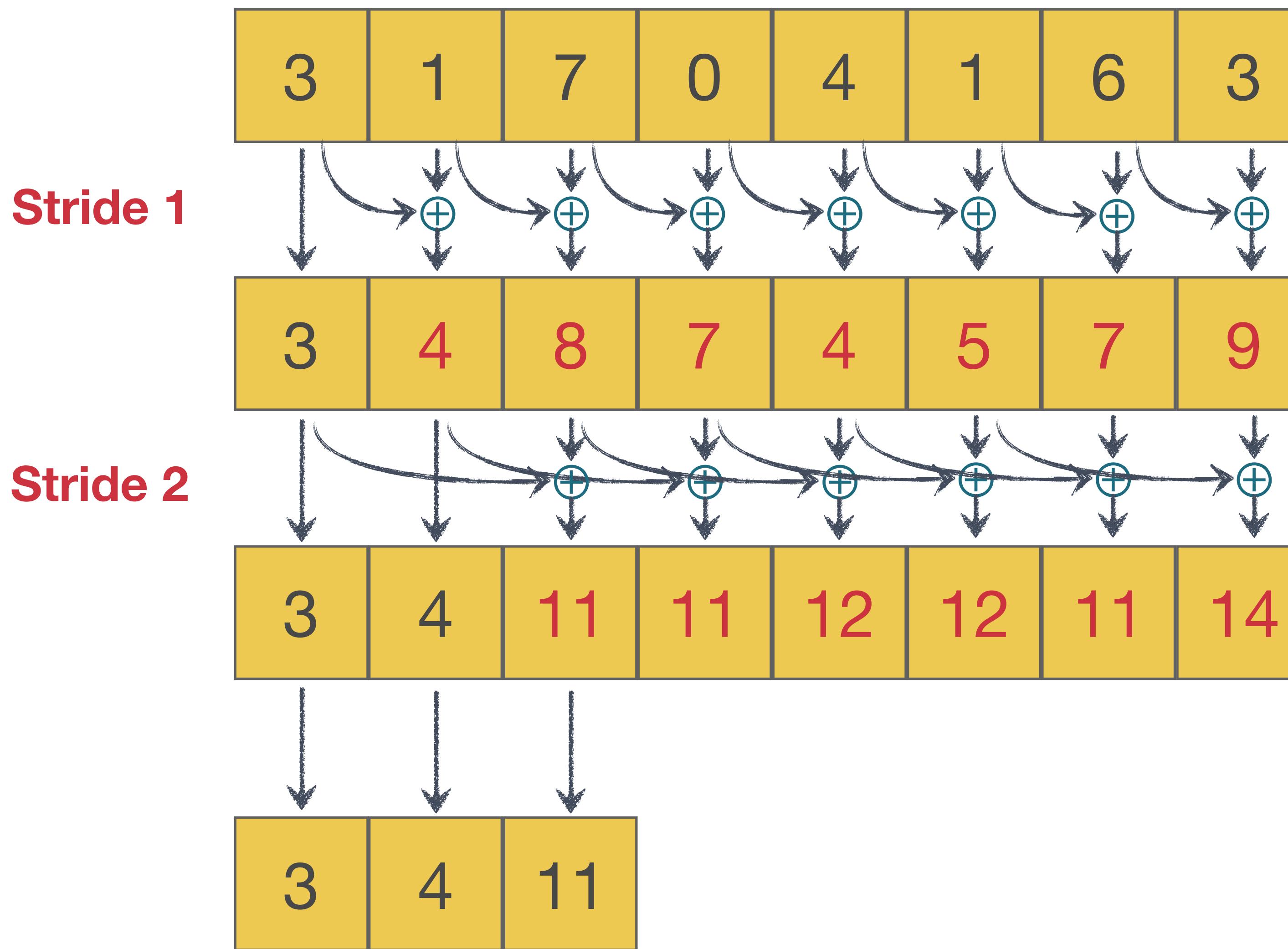
# Inclusive Parallel Scan



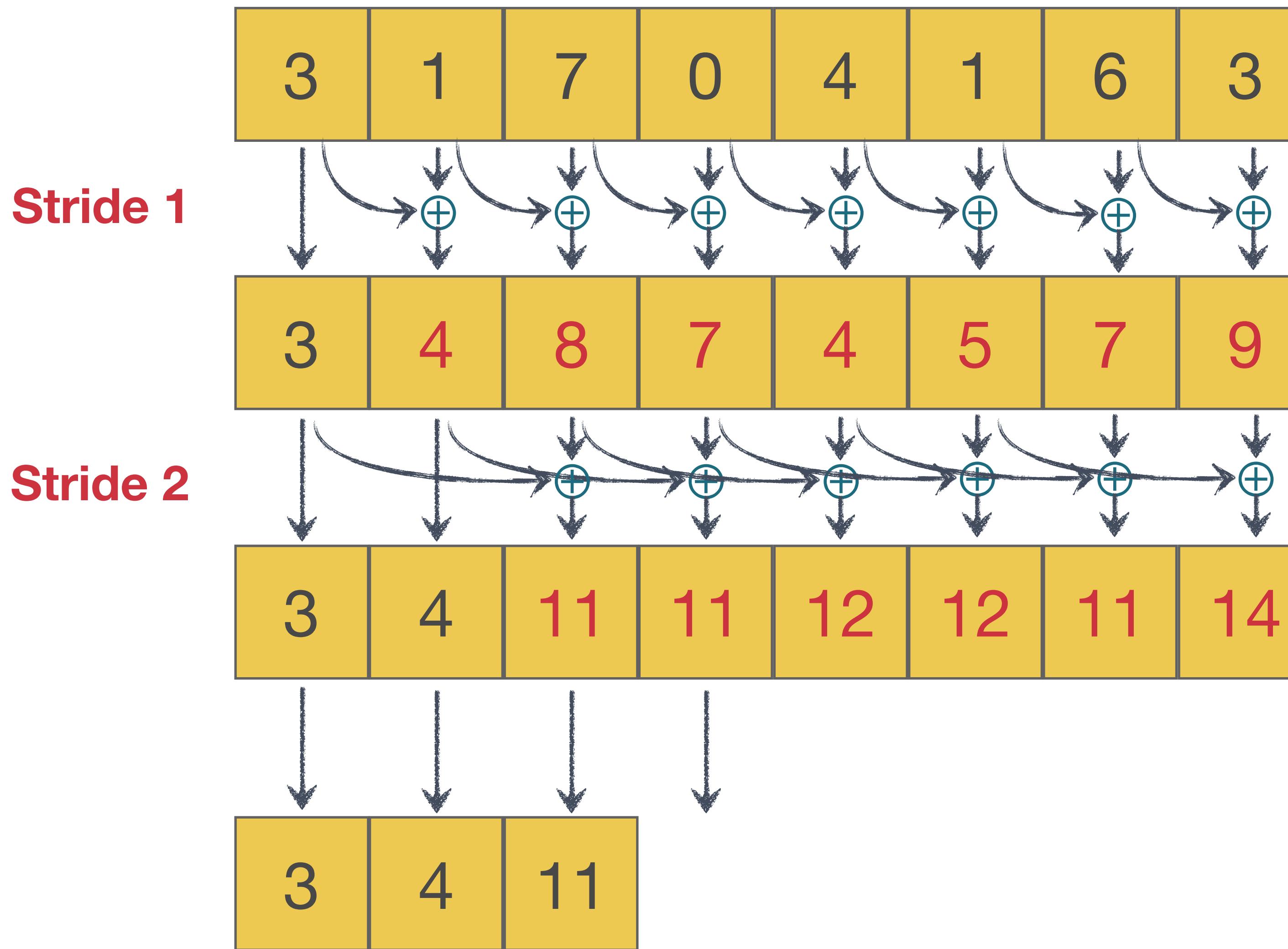
# Inclusive Parallel Scan



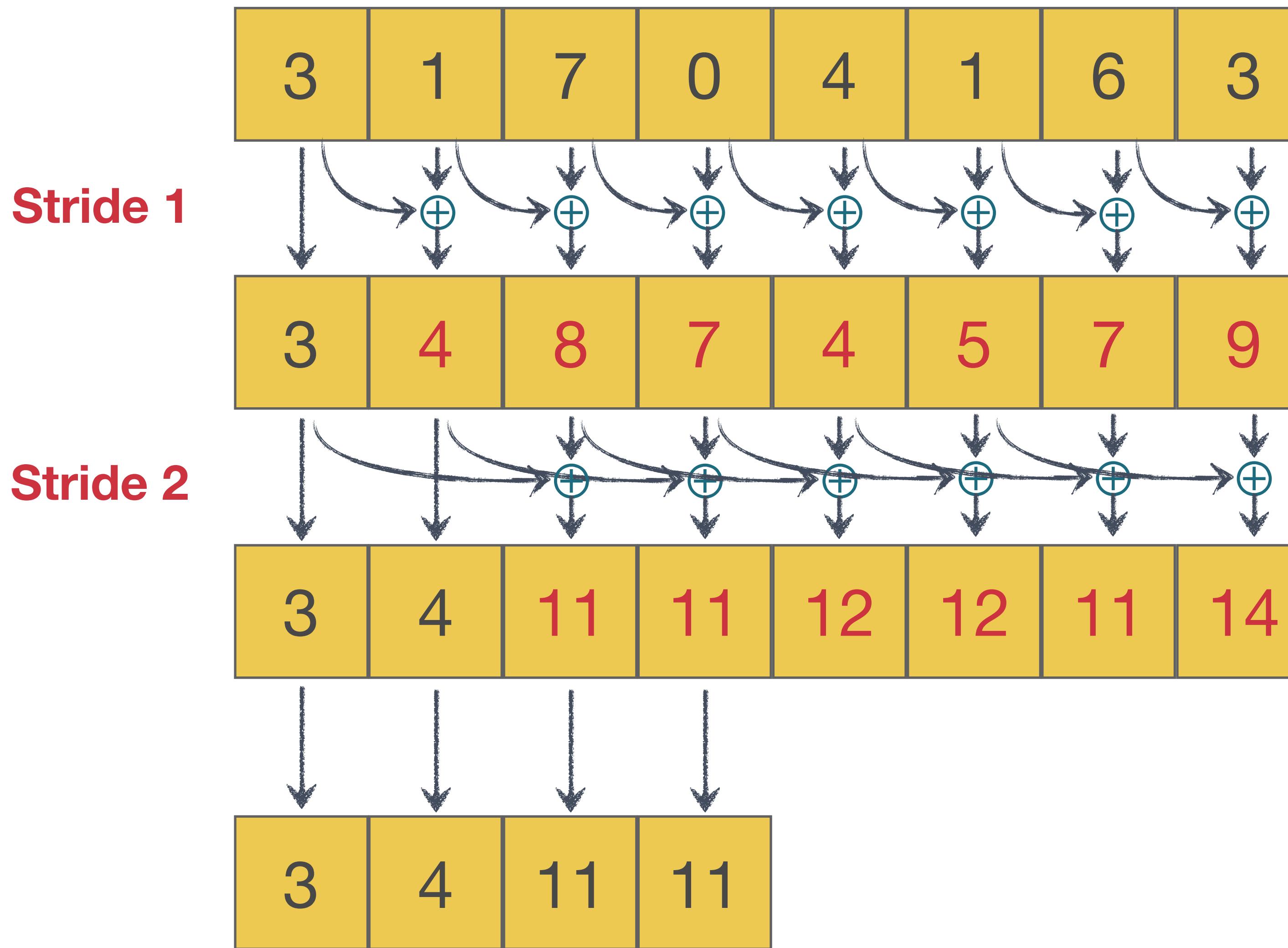
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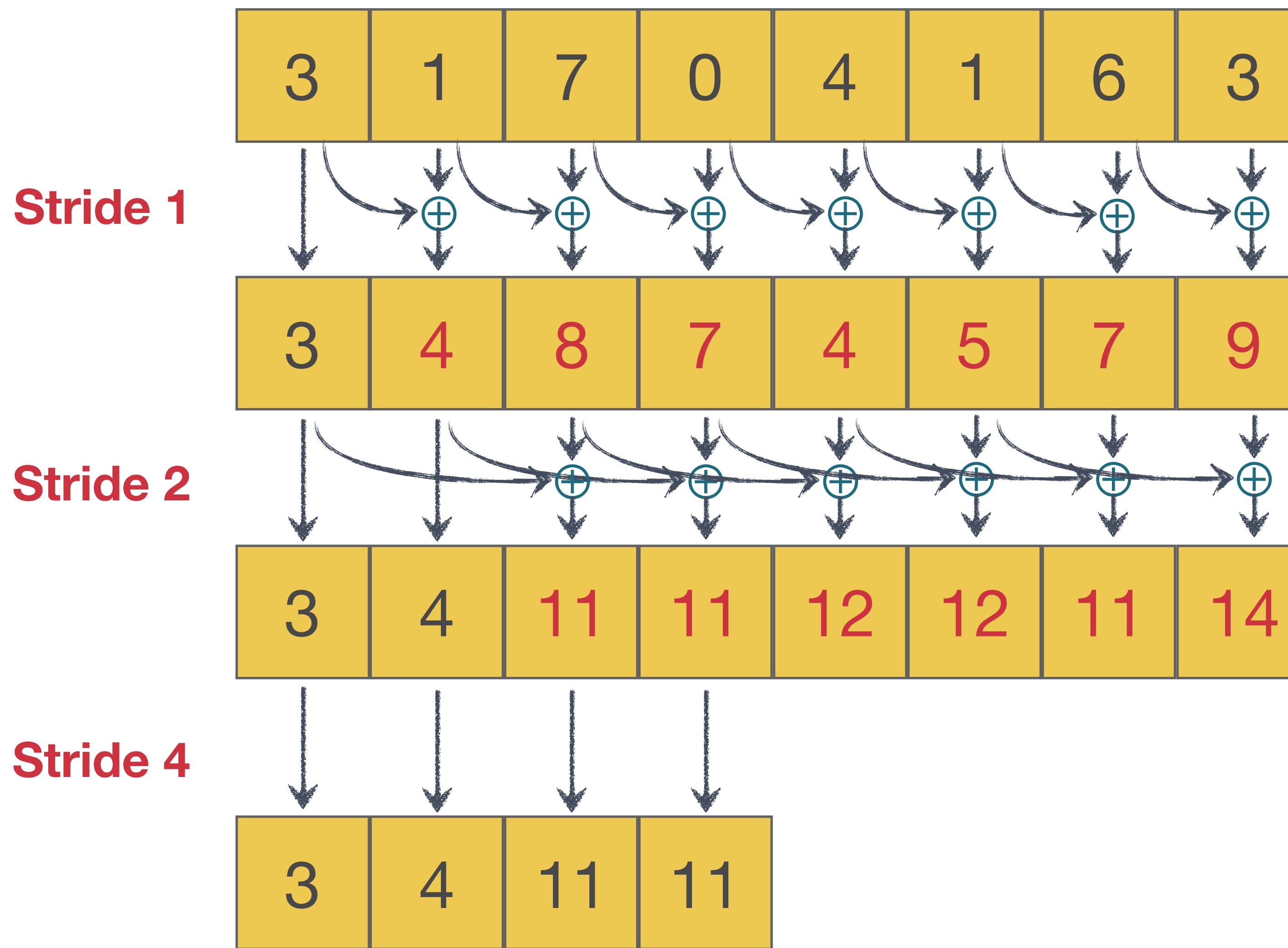
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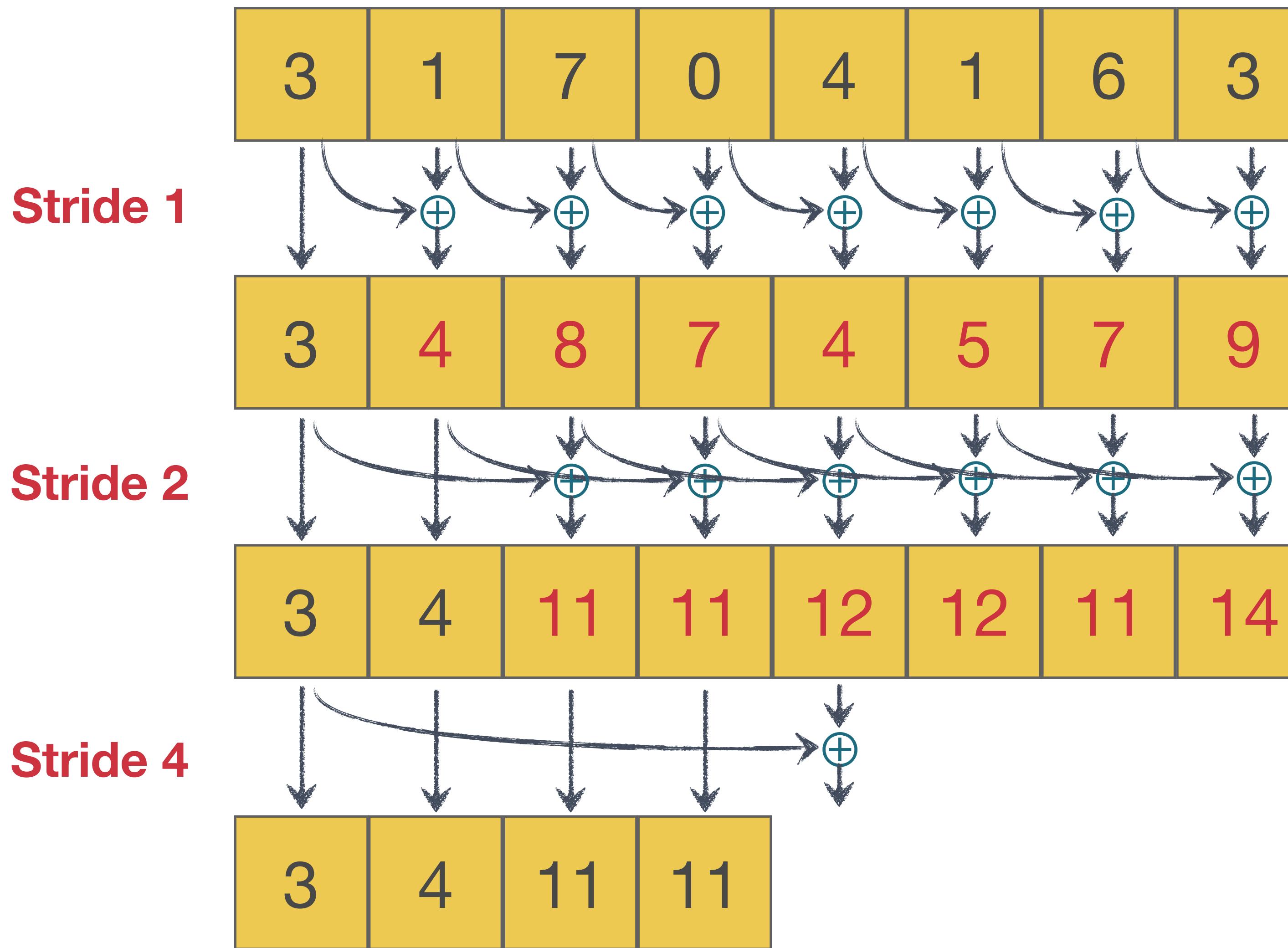
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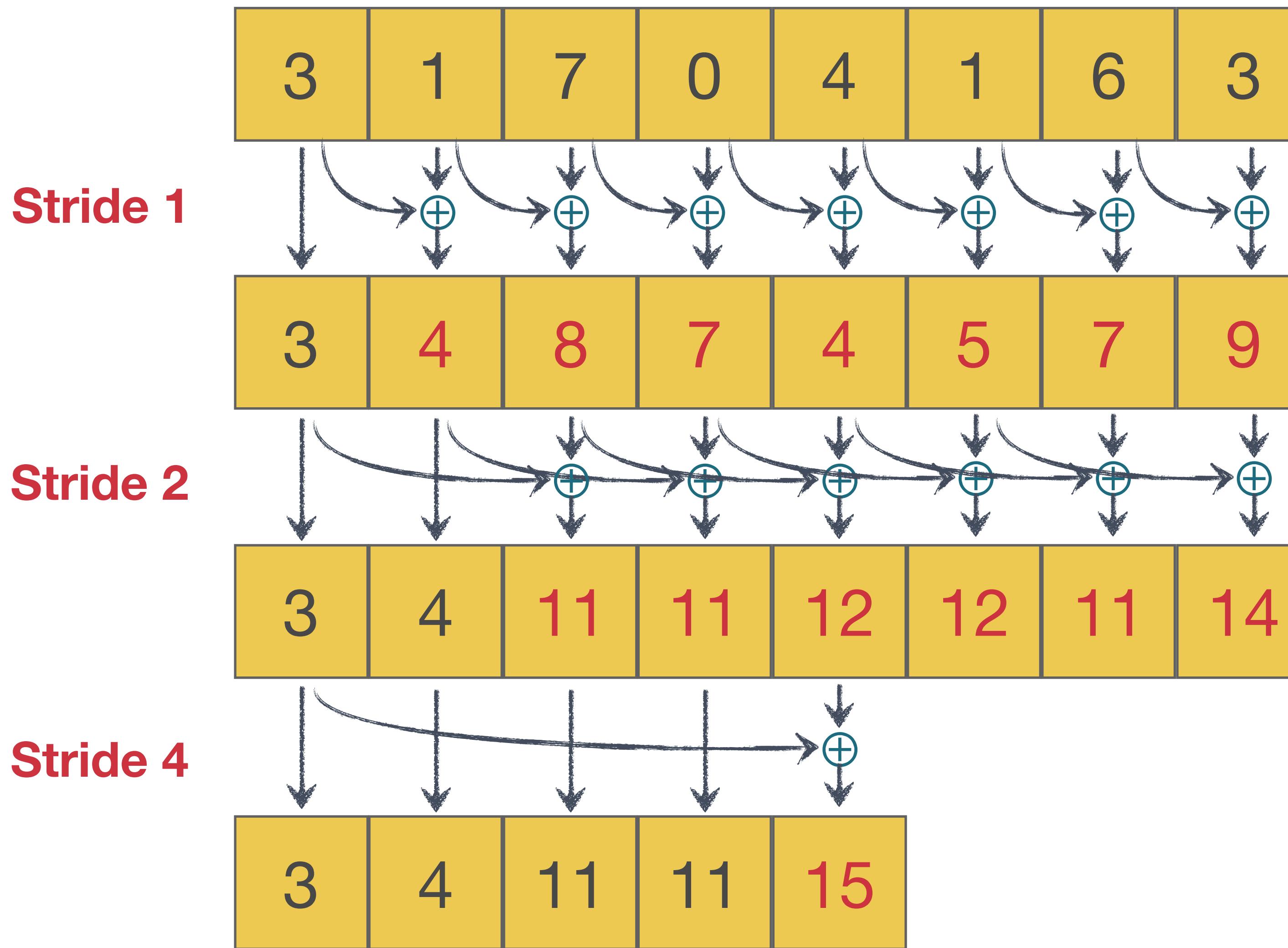
# Inclusive Parallel Scan



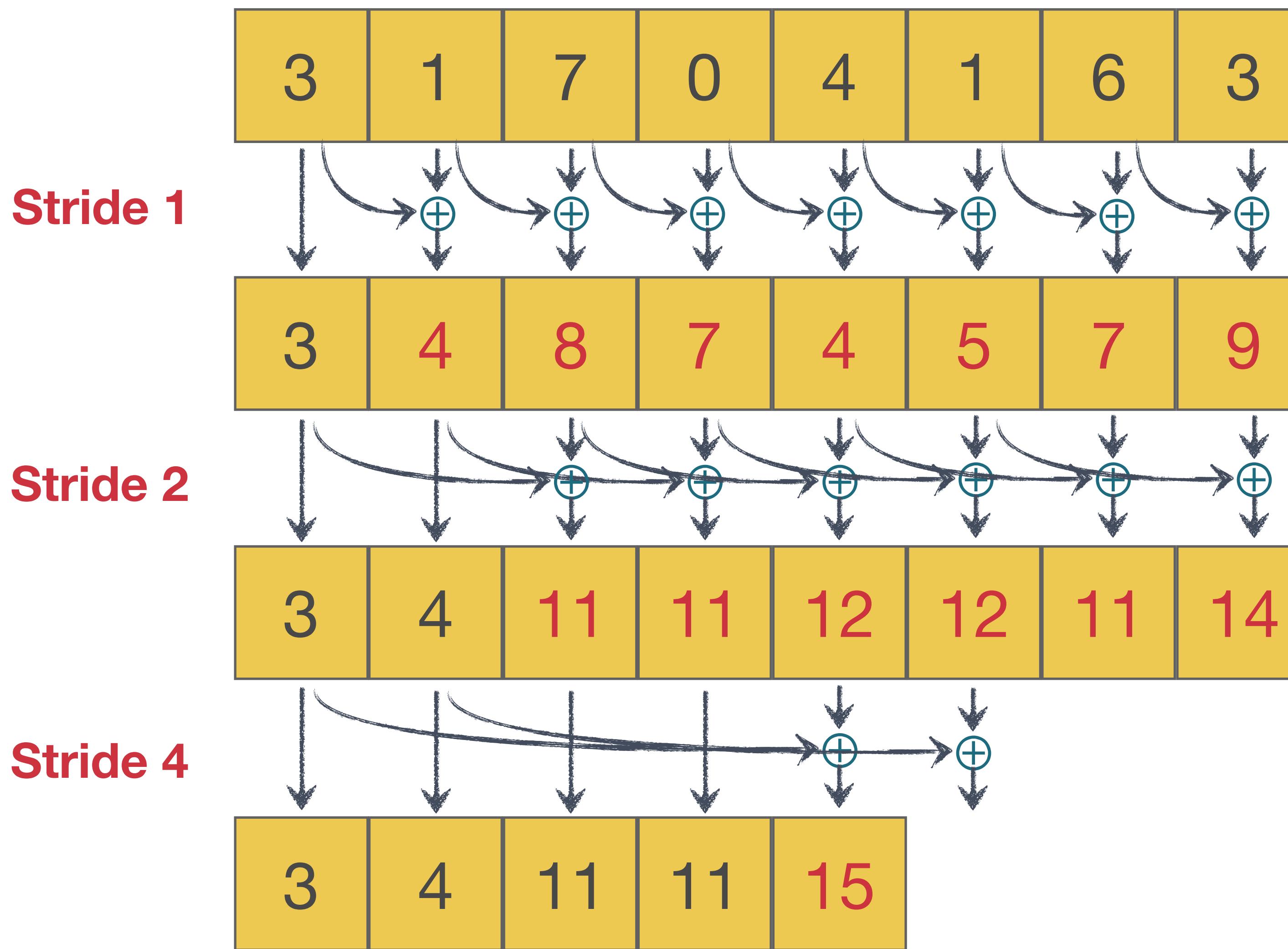
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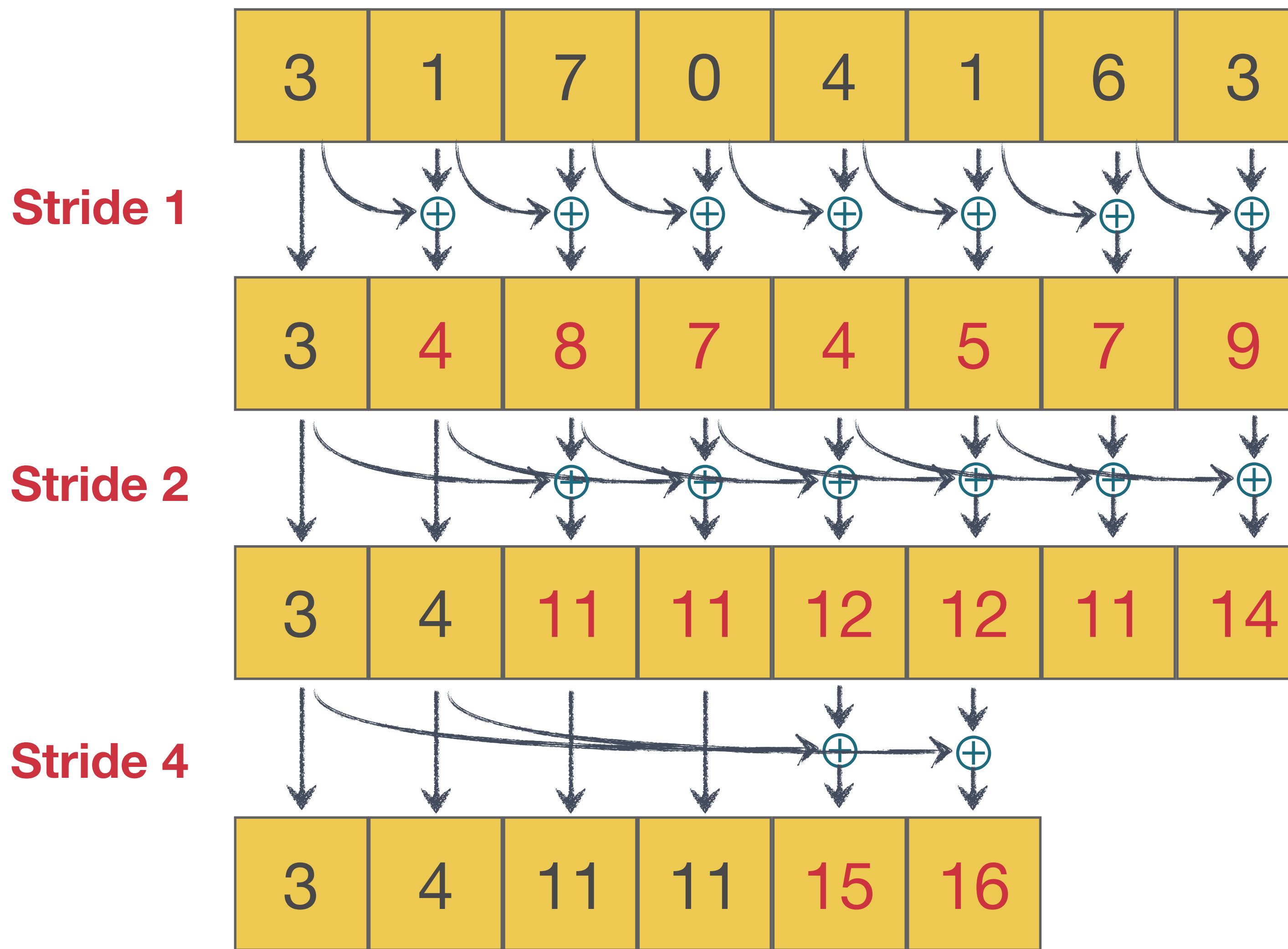
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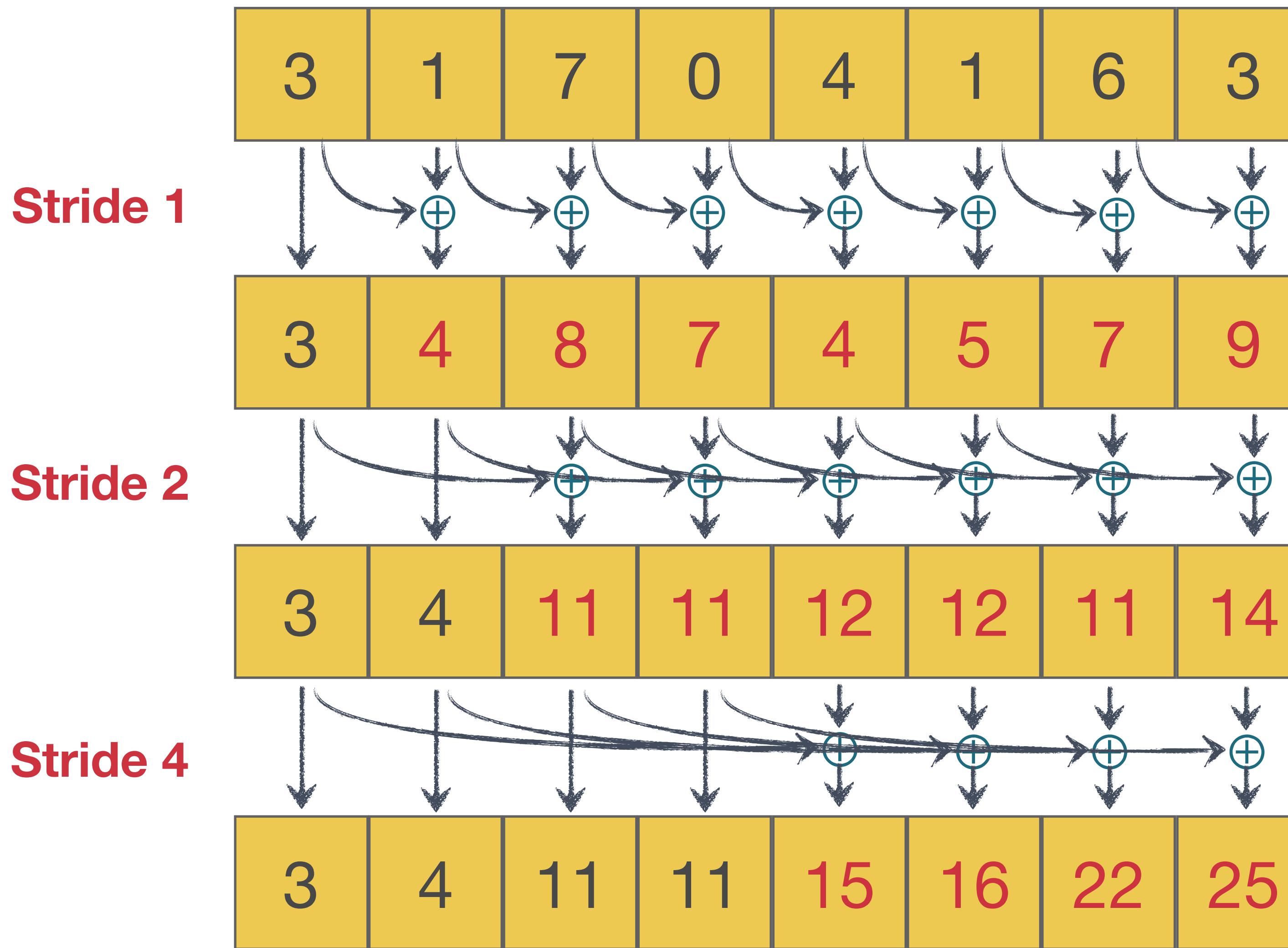
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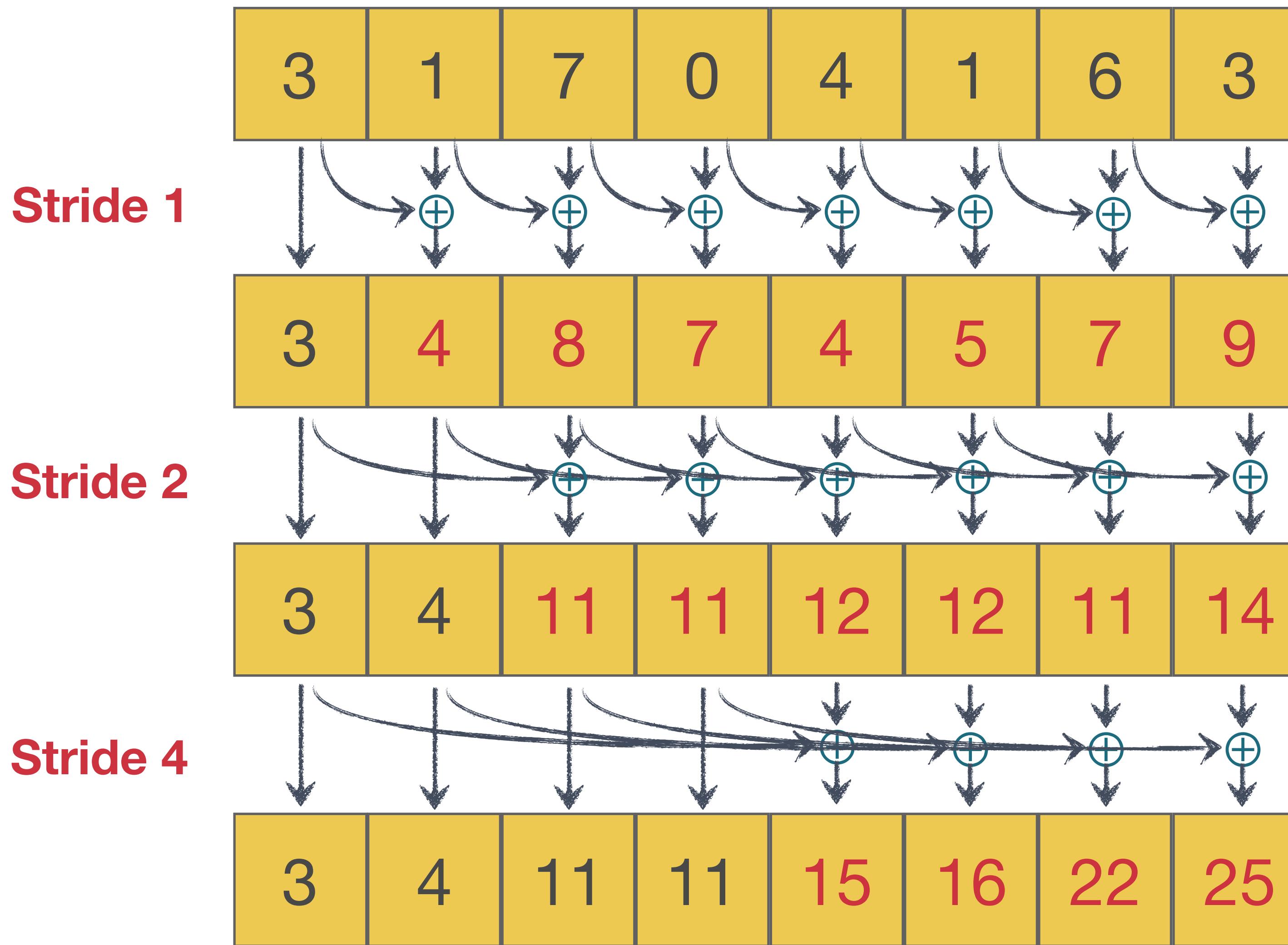
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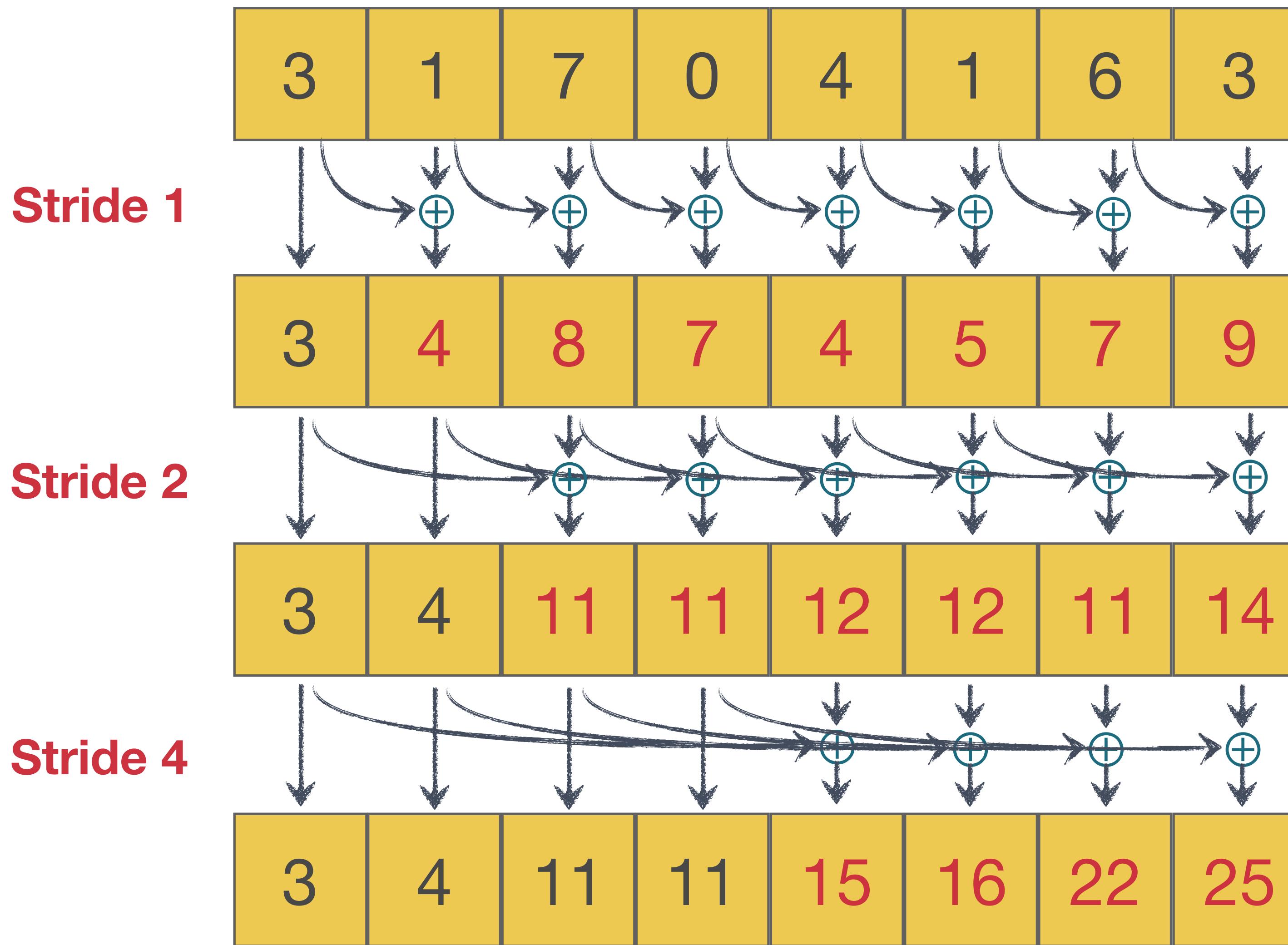


# Inclusive Parallel Scan



Iterate  $\log(n)$  times  
adding pairs of  
elements *stride*  
elements apart.

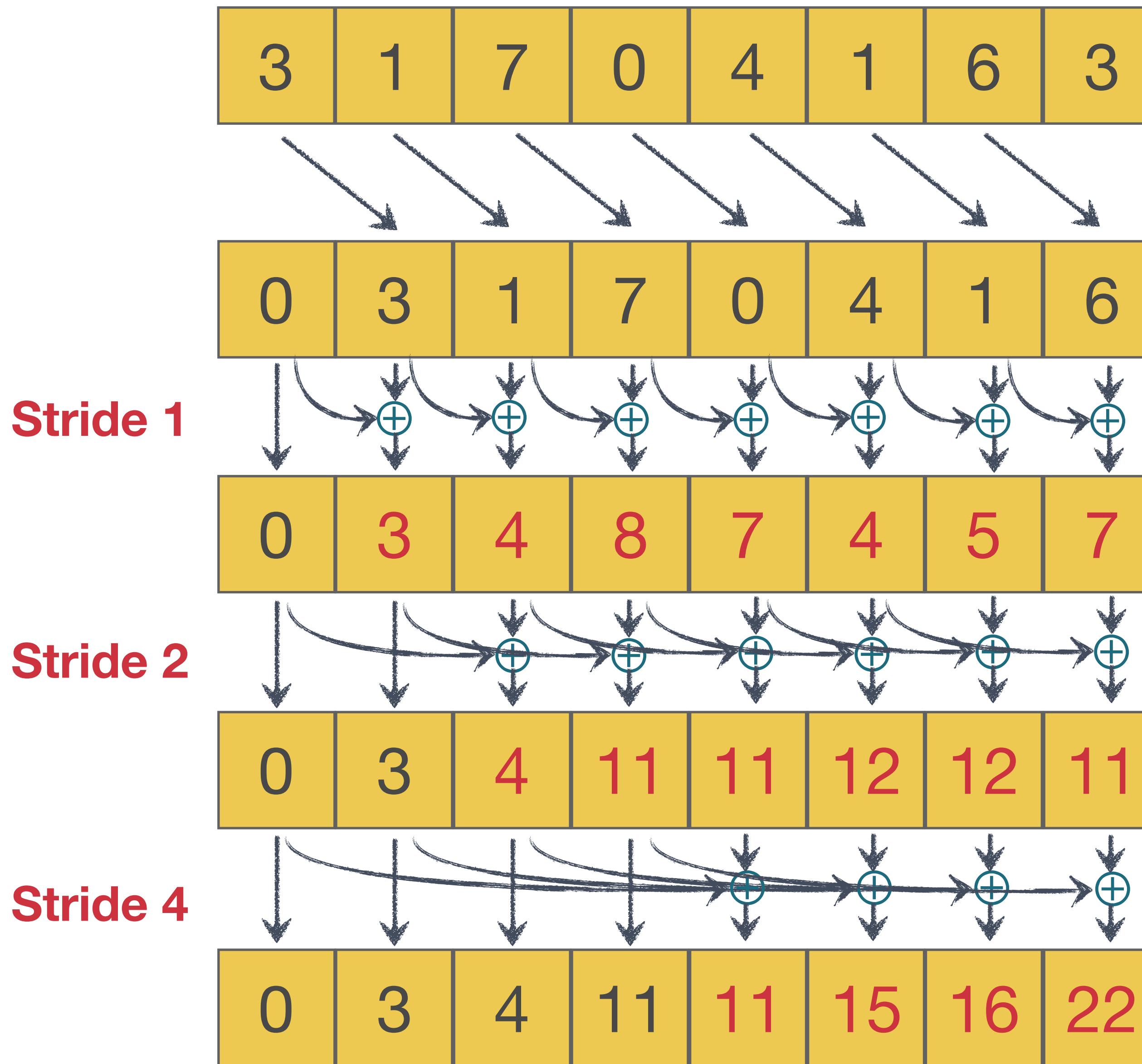
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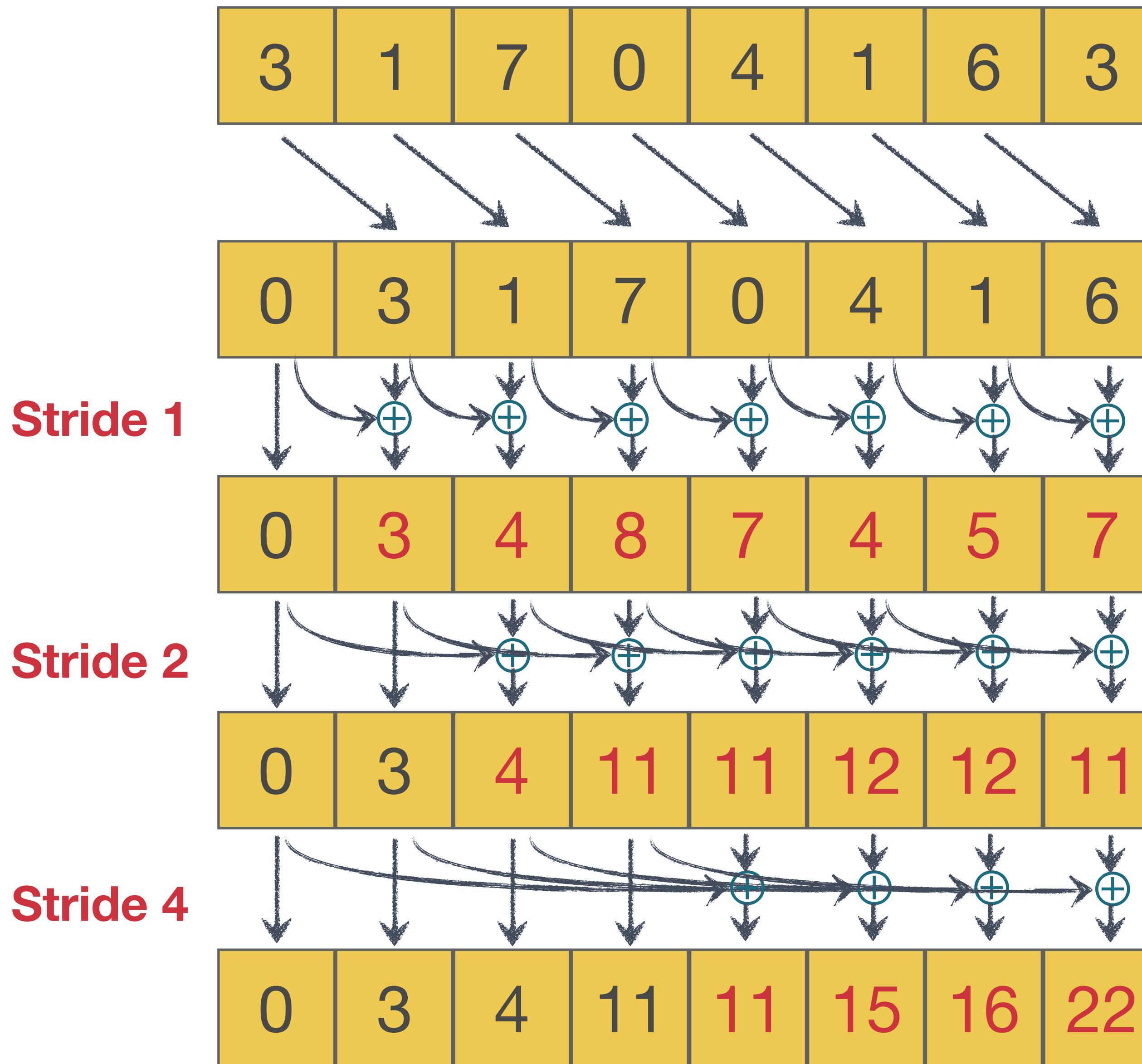
$$W(n) = \mathcal{O}(n \log n)$$
$$D(n) = \mathcal{O}(\log n)$$

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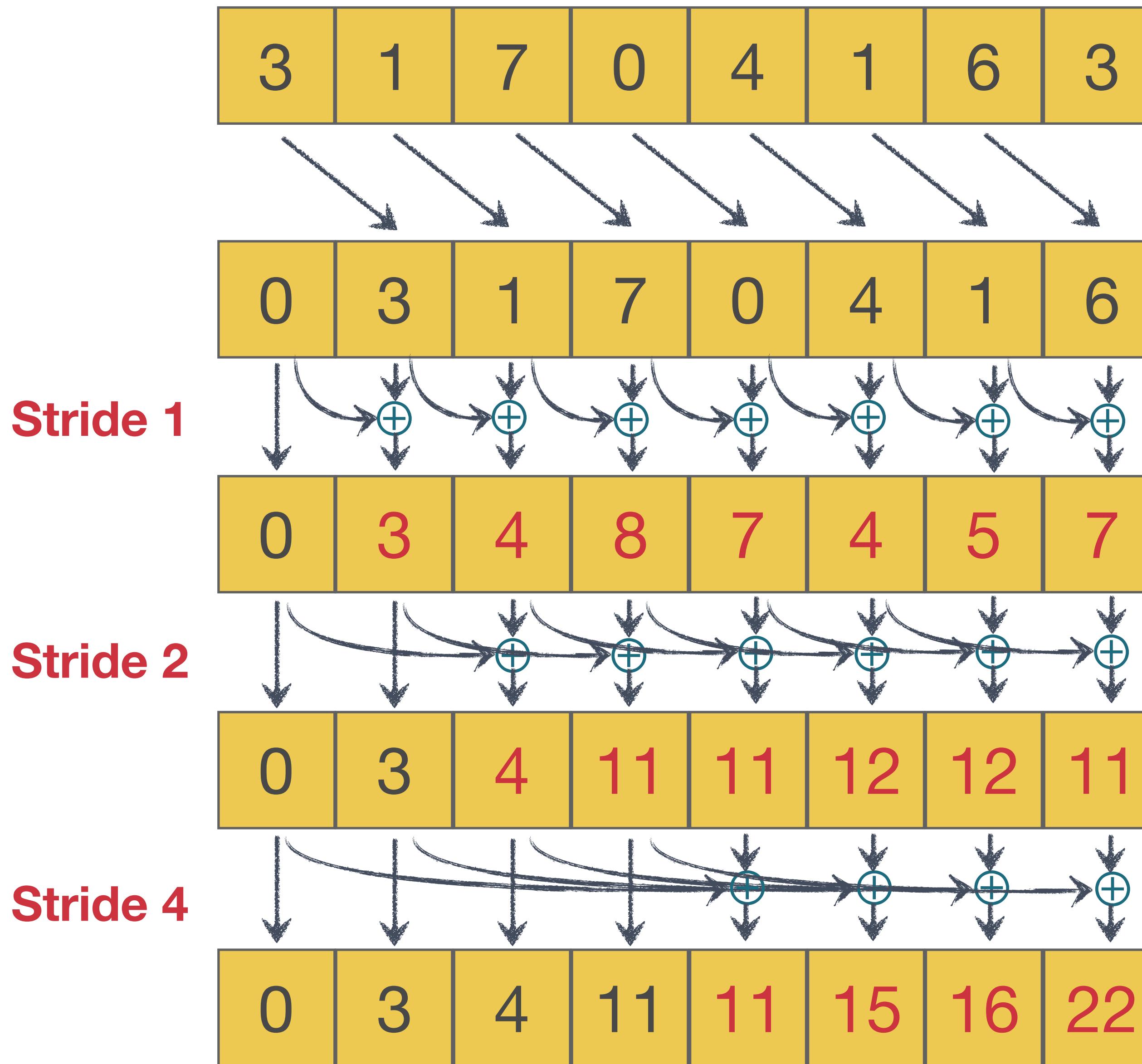
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Iterate  $\log(n)$  times  
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elements apart.

$$W(n) = \mathcal{O}(n \log n)$$
$$D(n) = \mathcal{O}(\log n)$$

Can we do better?

# Parallel Scan: Reduction

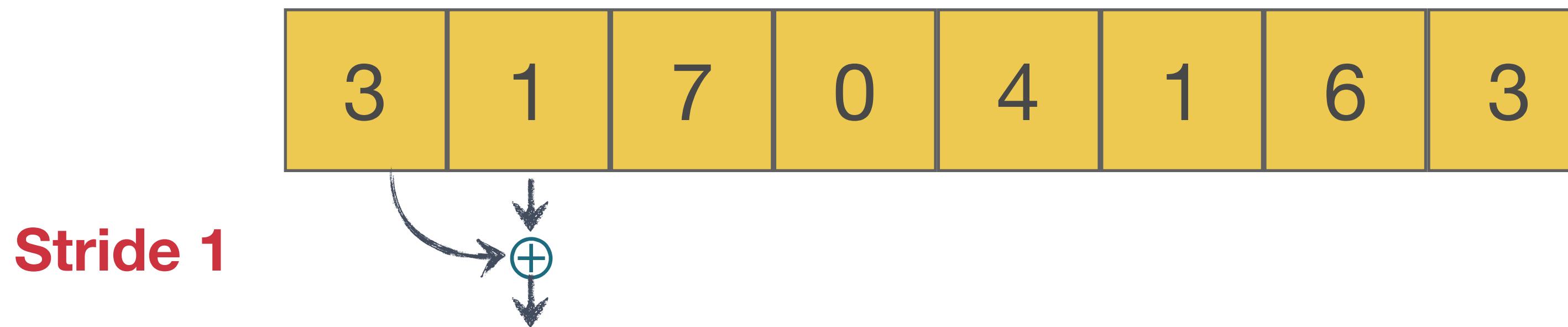
3	1	7	0	4	1	6	3
---	---	---	---	---	---	---	---

# Parallel Scan: Reduction

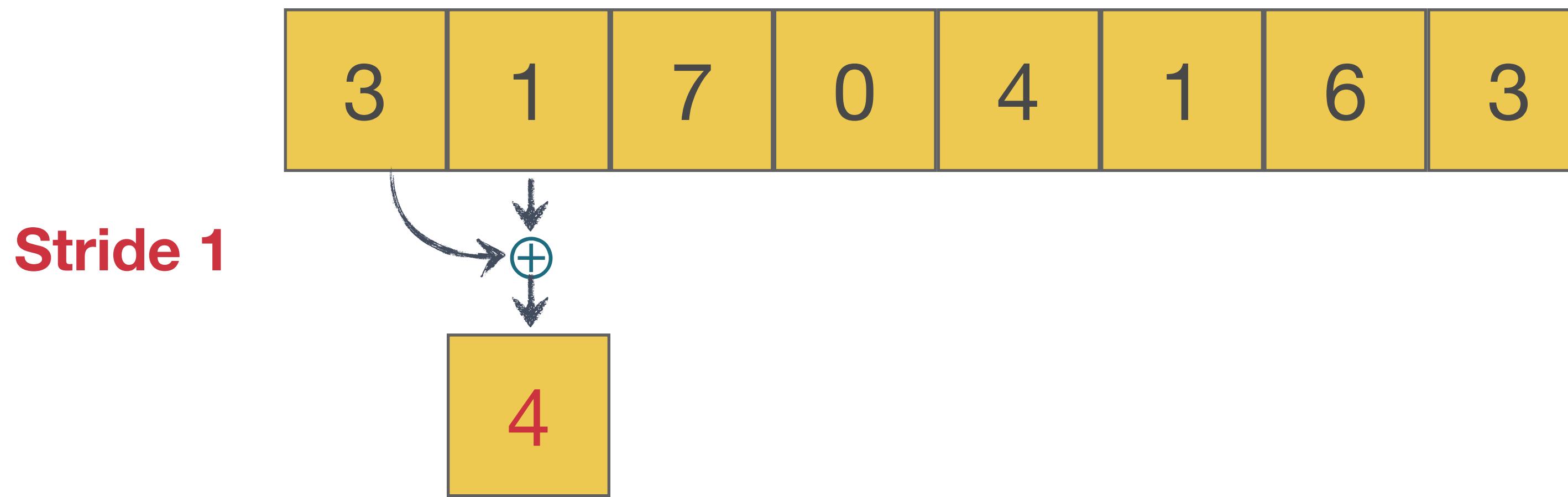


**Stride 1**

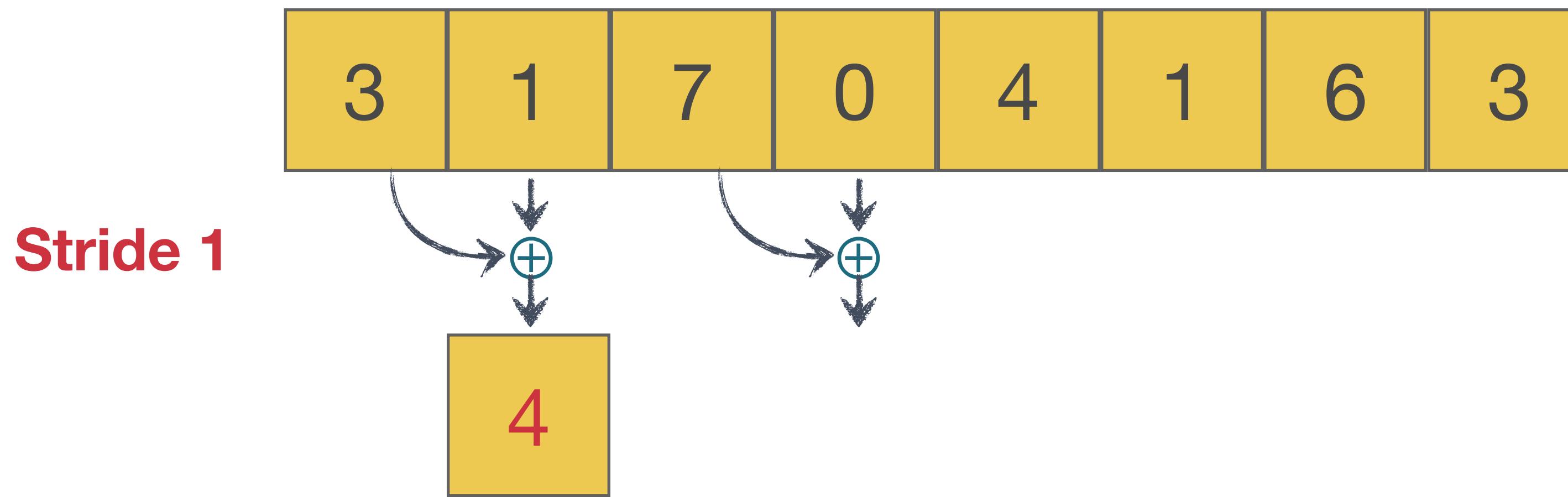
# Parallel Scan: Reduction



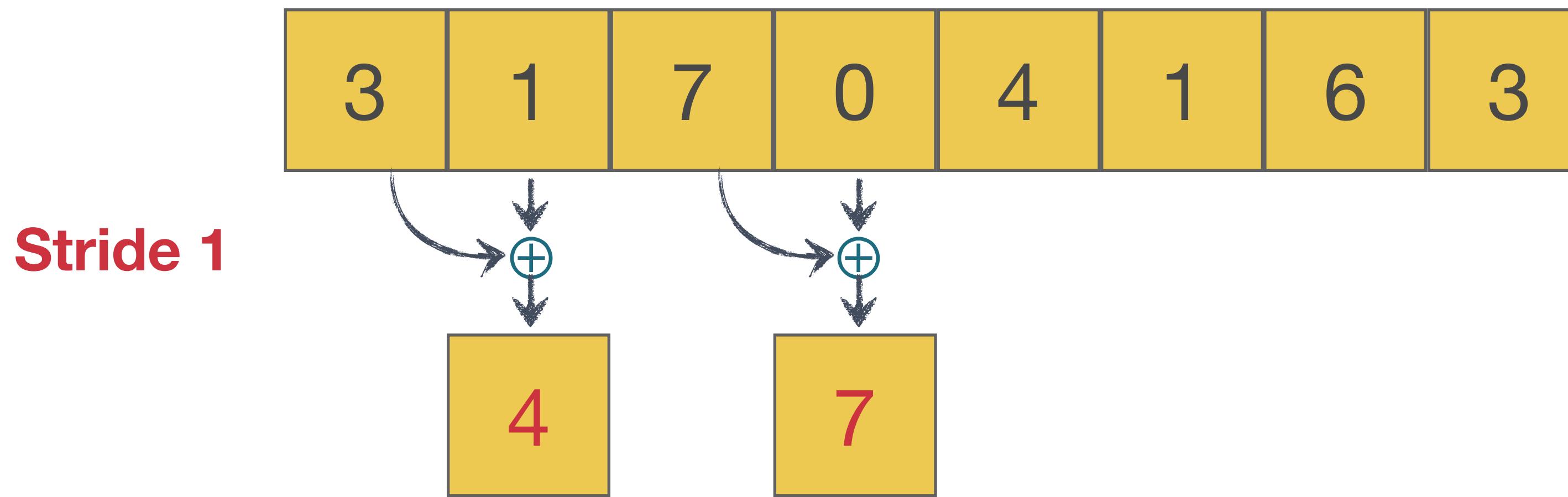
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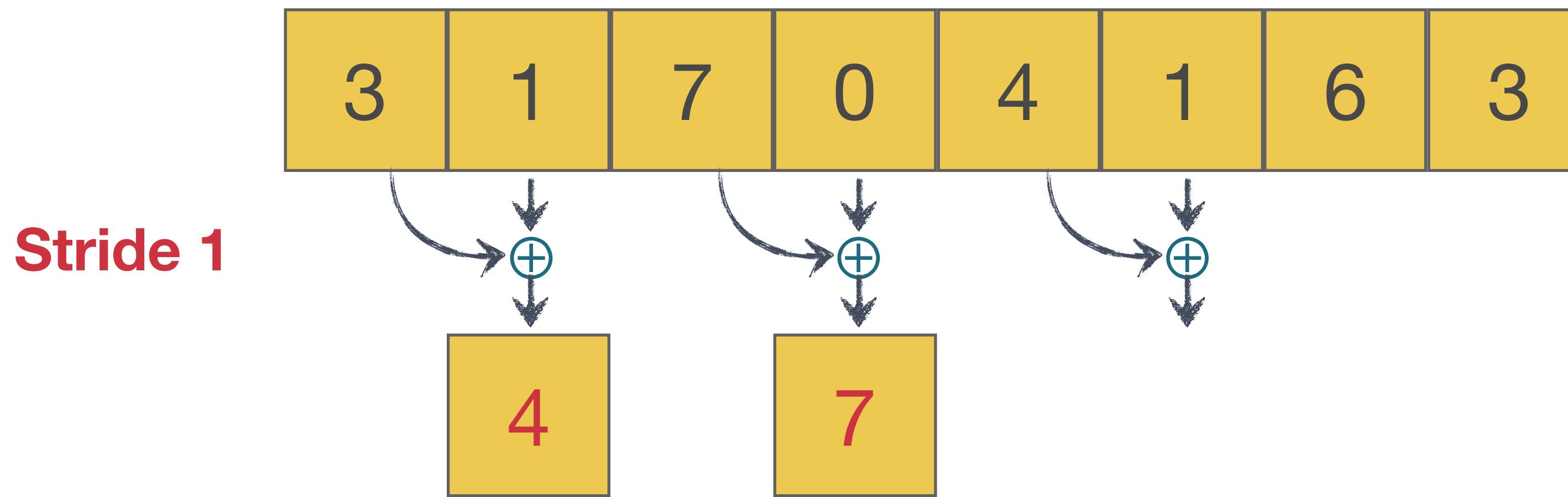
# Parallel Scan: Reduction



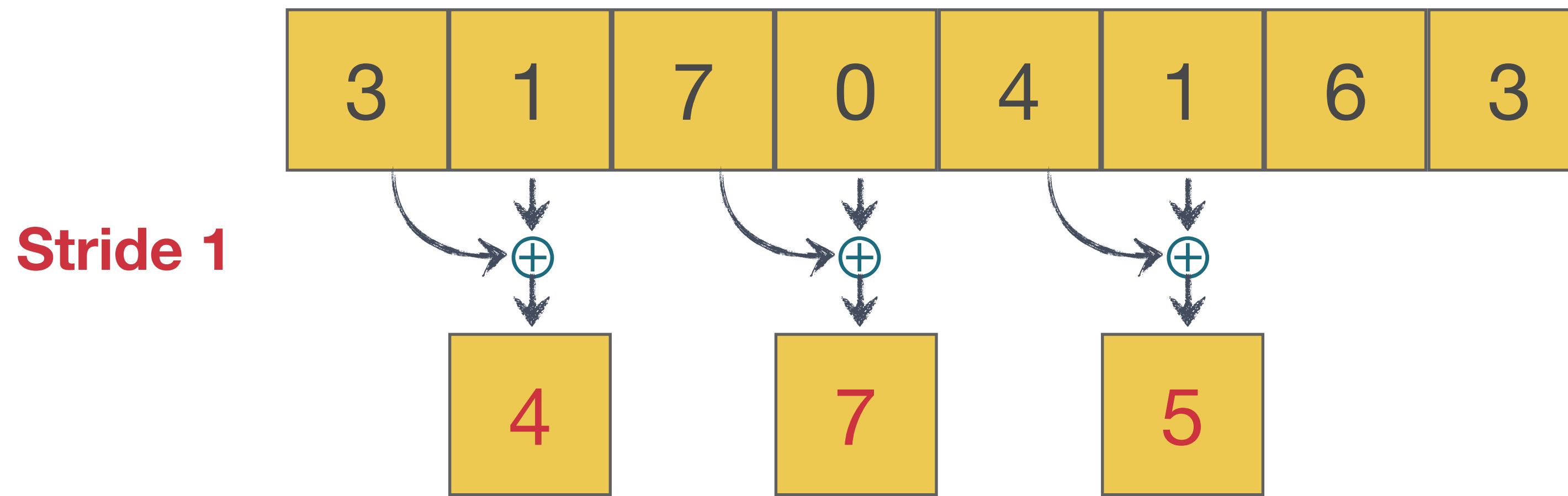
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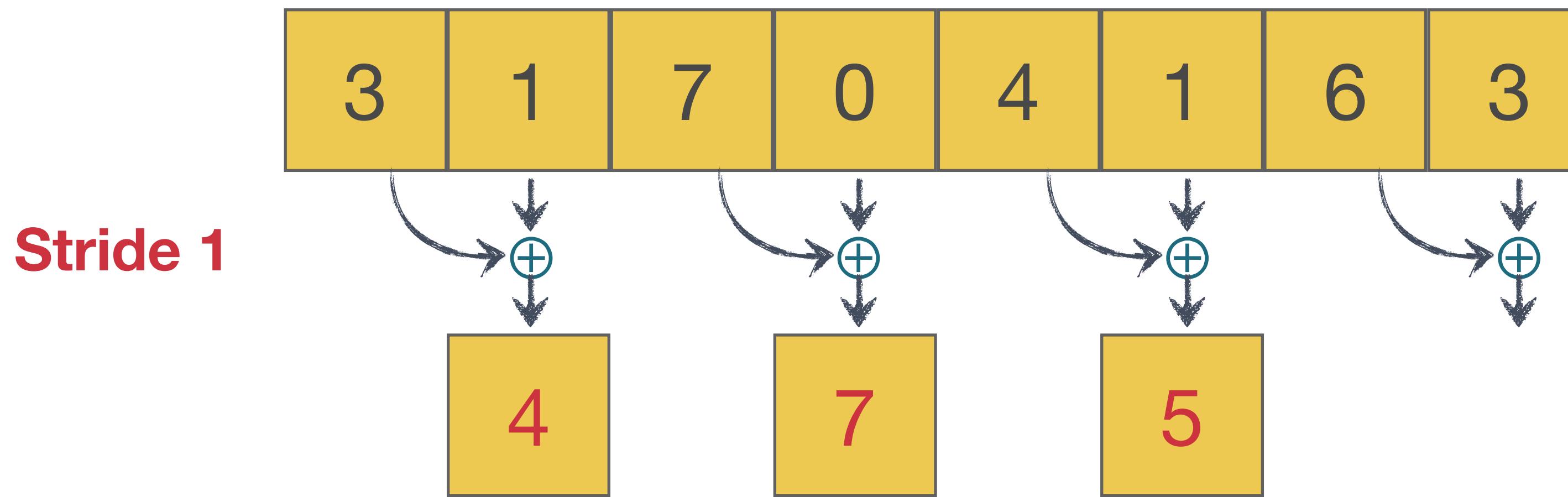
# Parallel Scan: Reduction



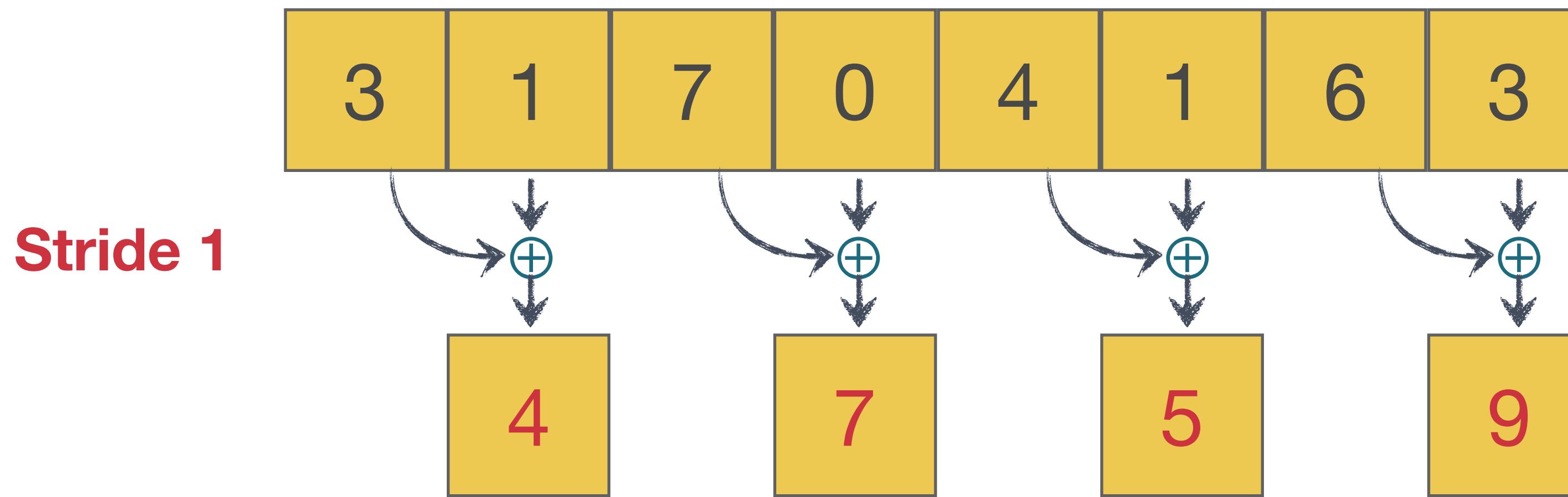
# Parallel Scan: Reduction



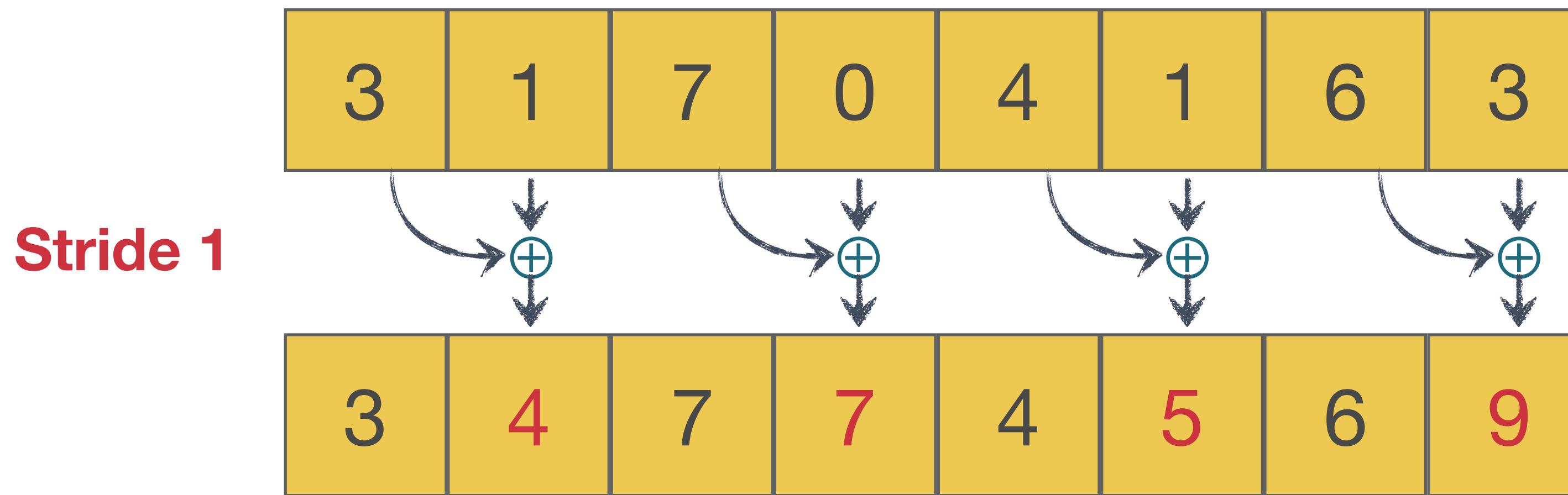
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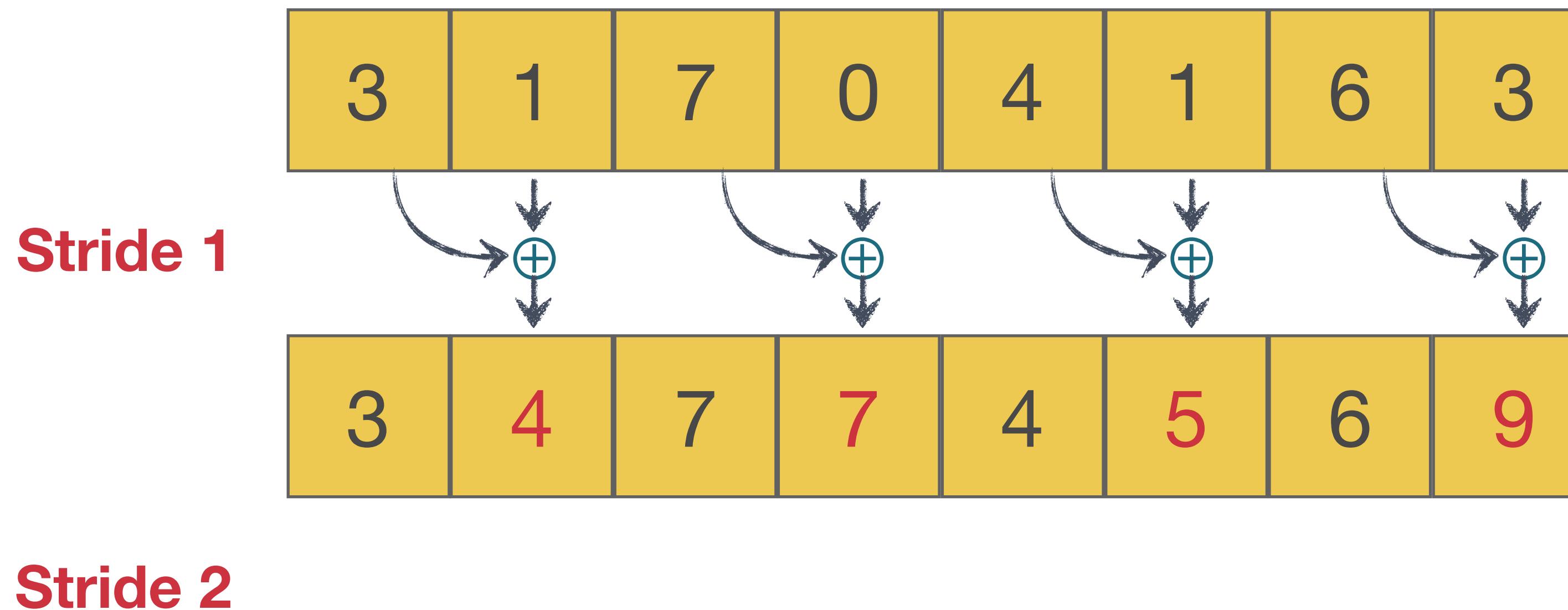
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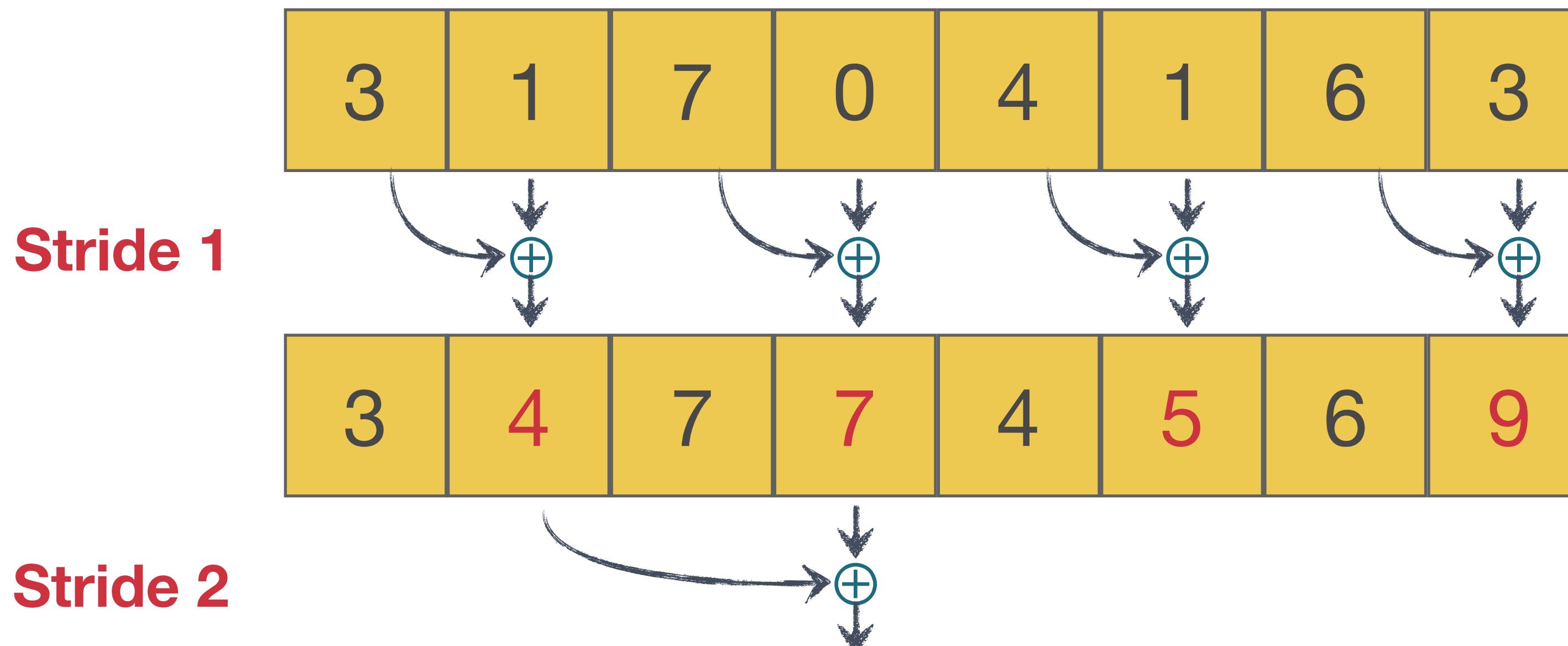
# Parallel Scan: Reduction



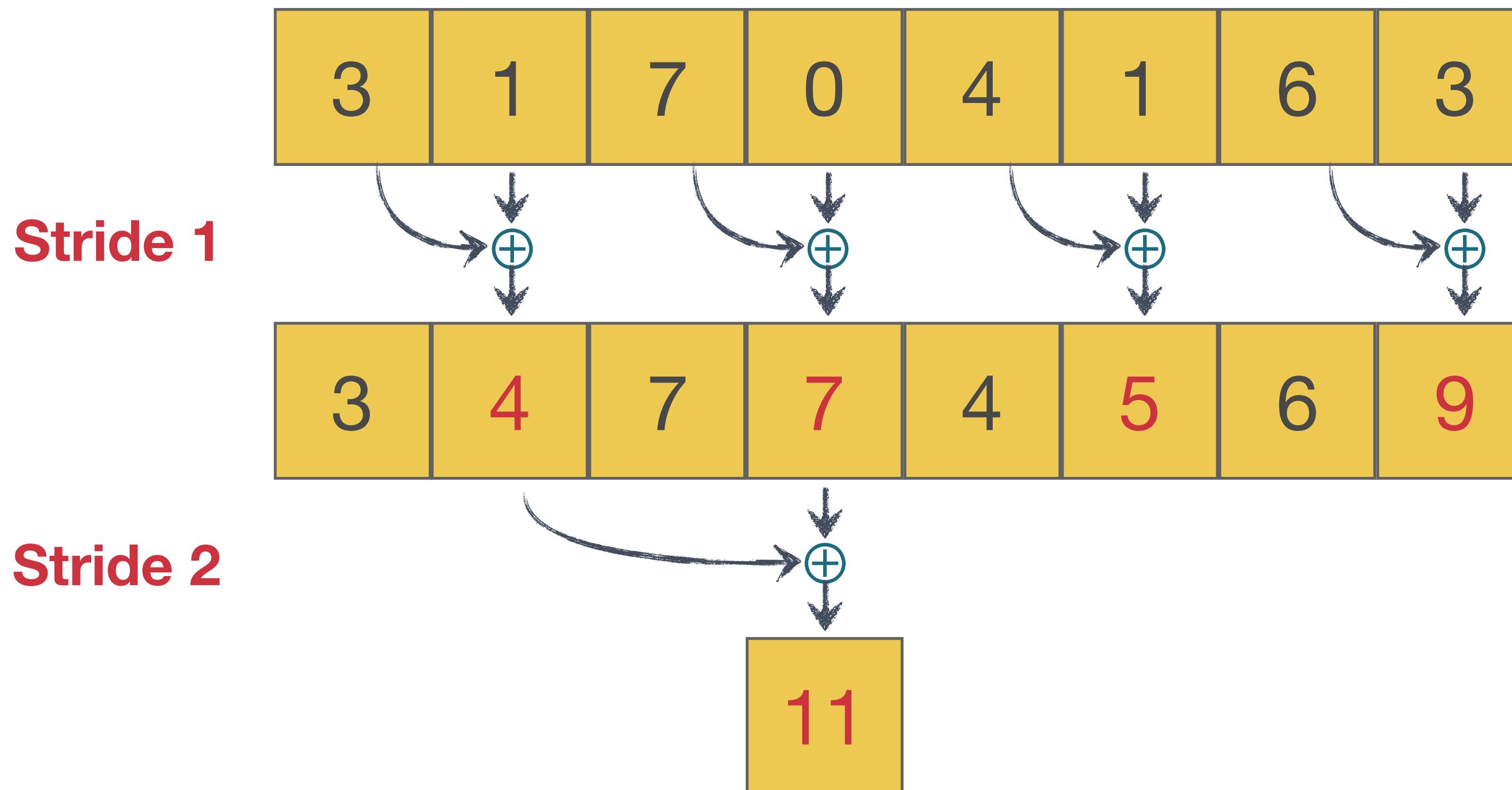
# Parallel Scan: Reduction



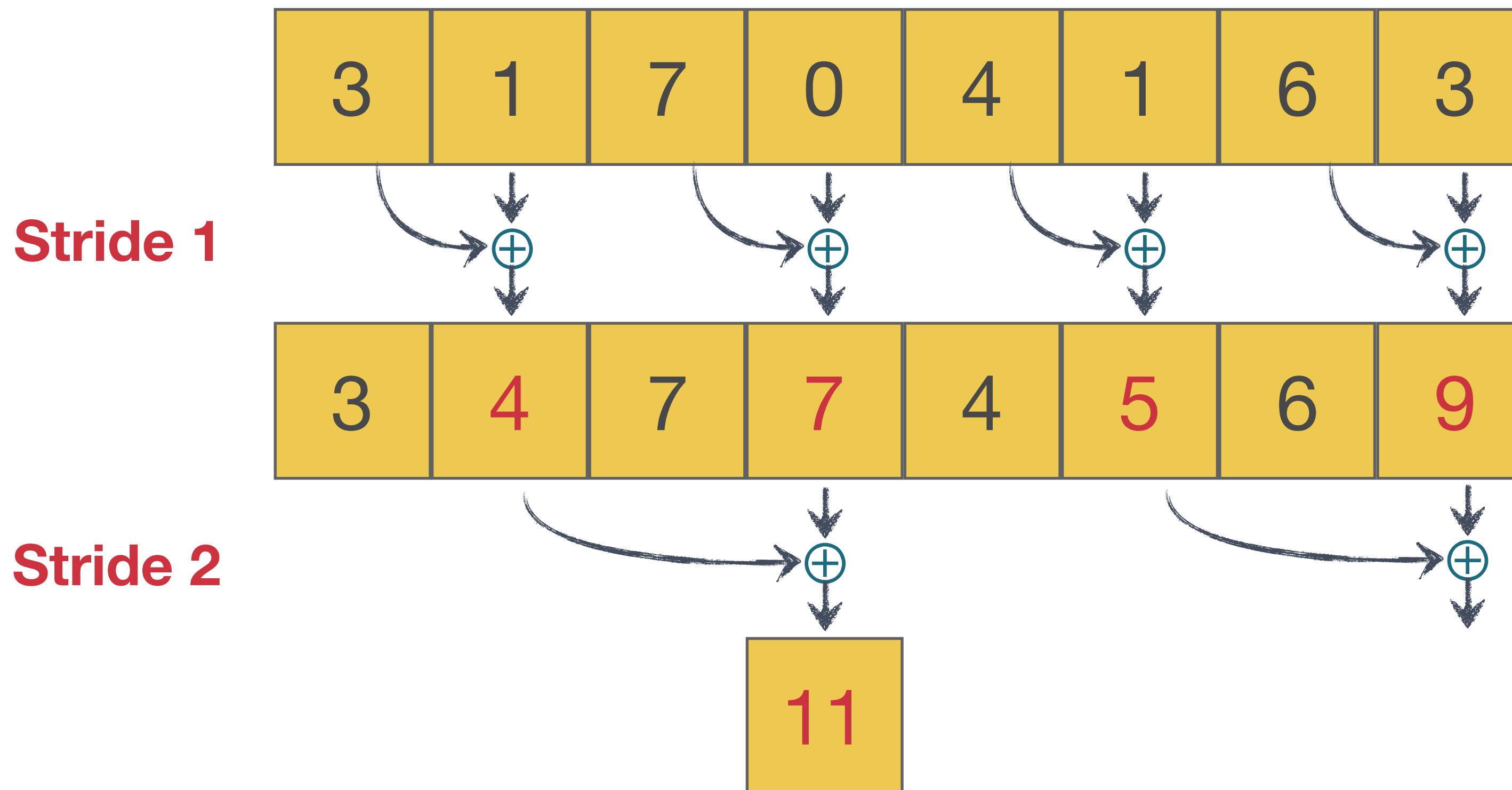
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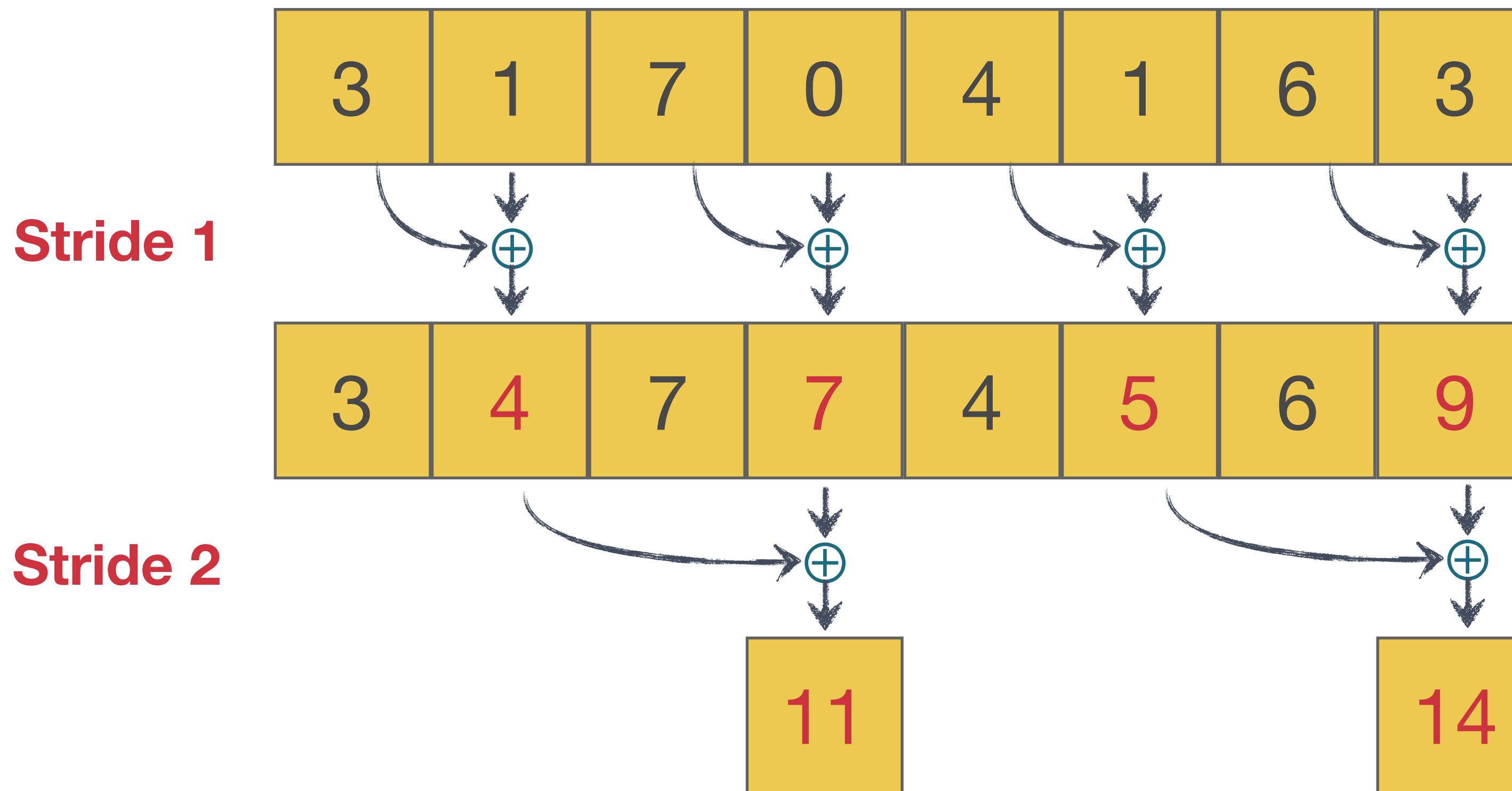
# Parallel Scan: Reduction



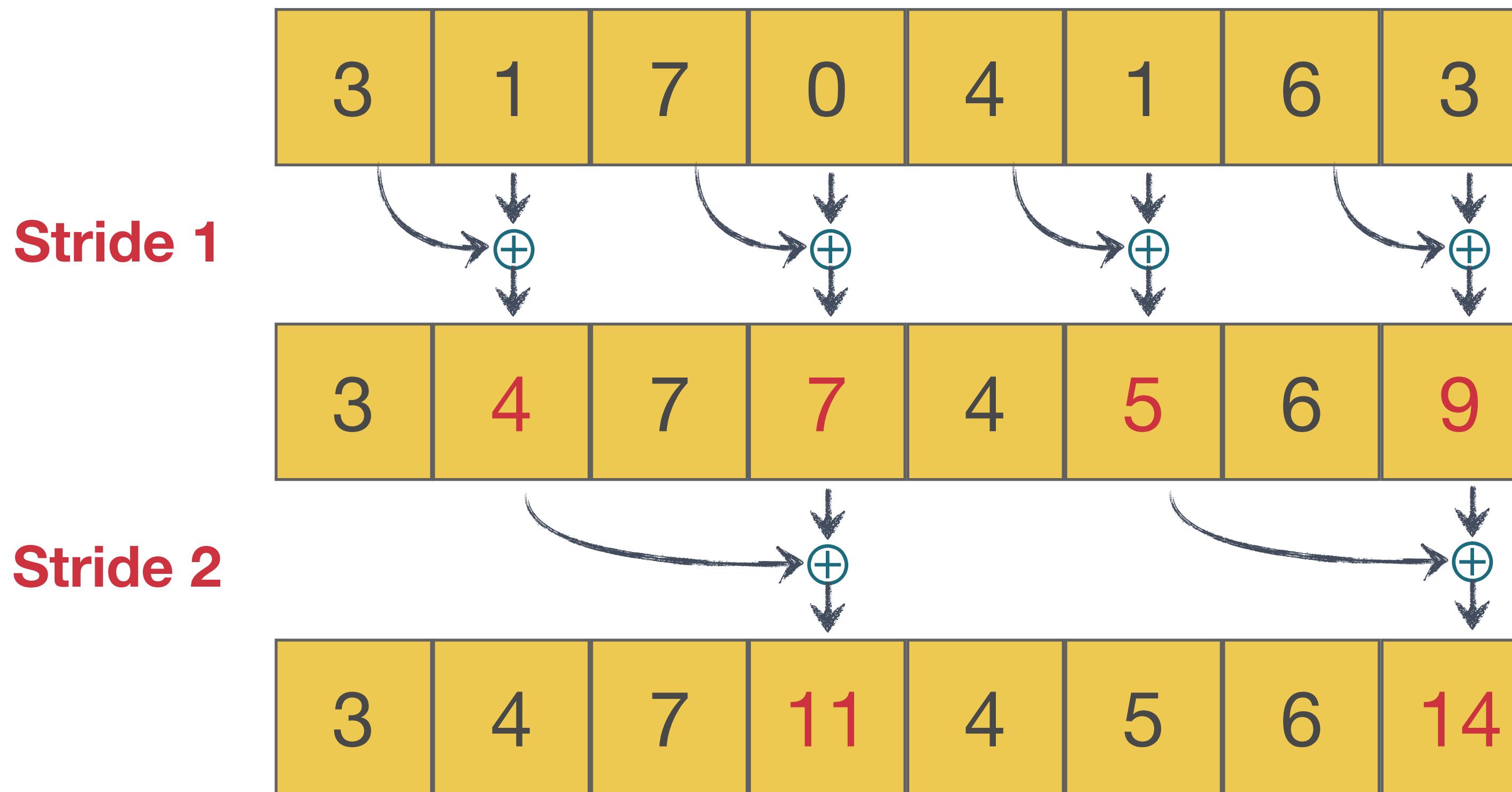
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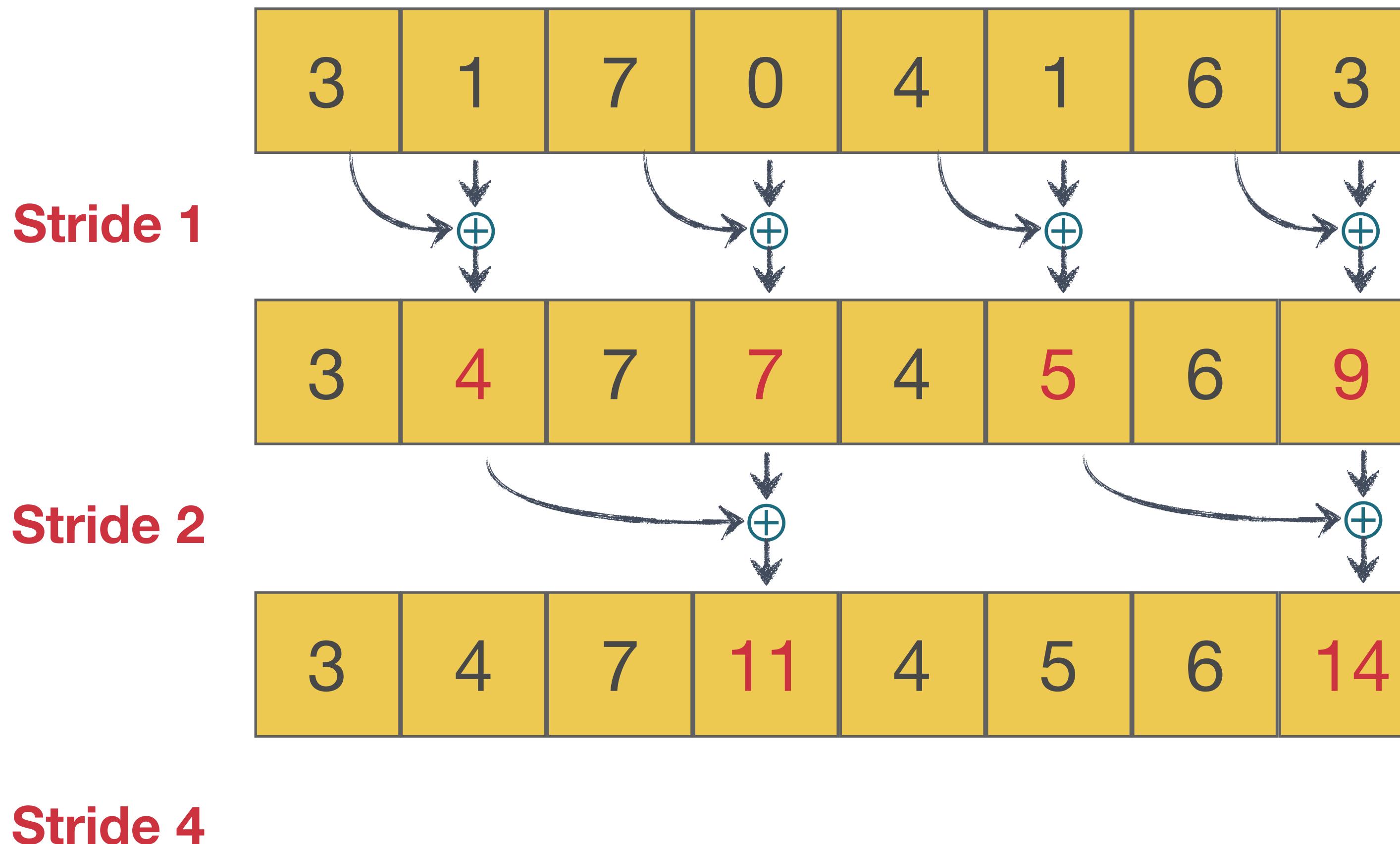
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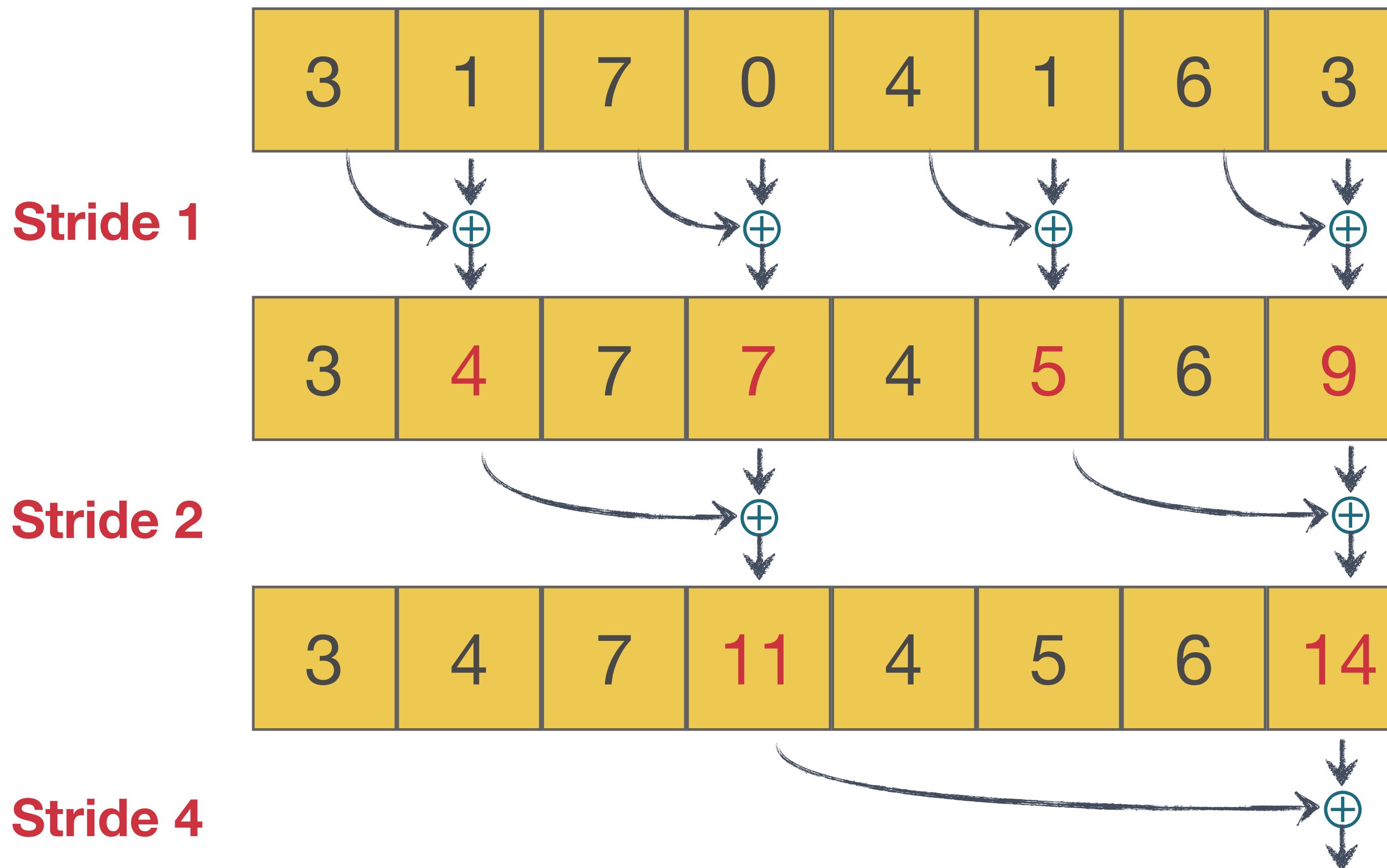
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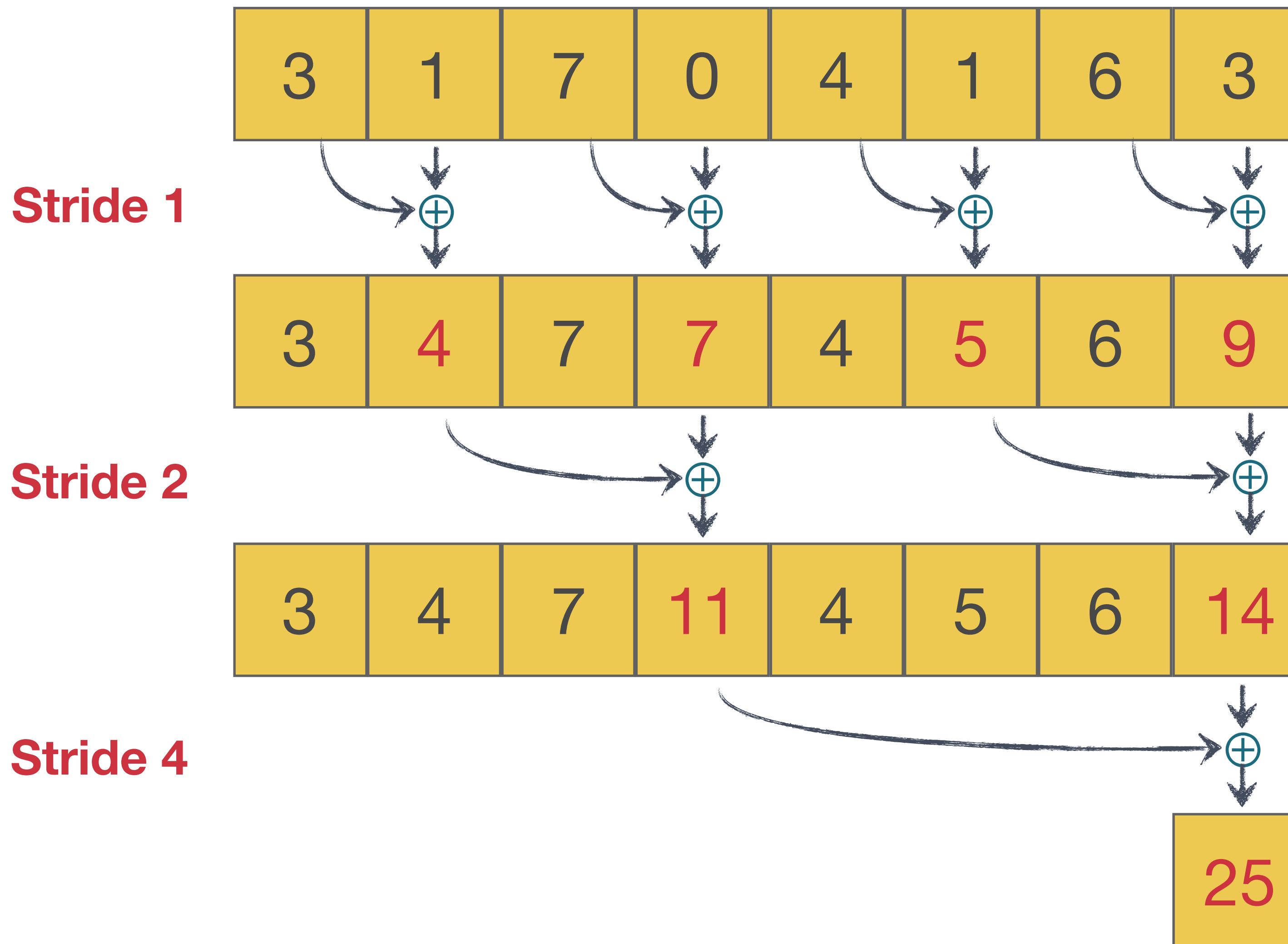
# Parallel Scan: Reduction



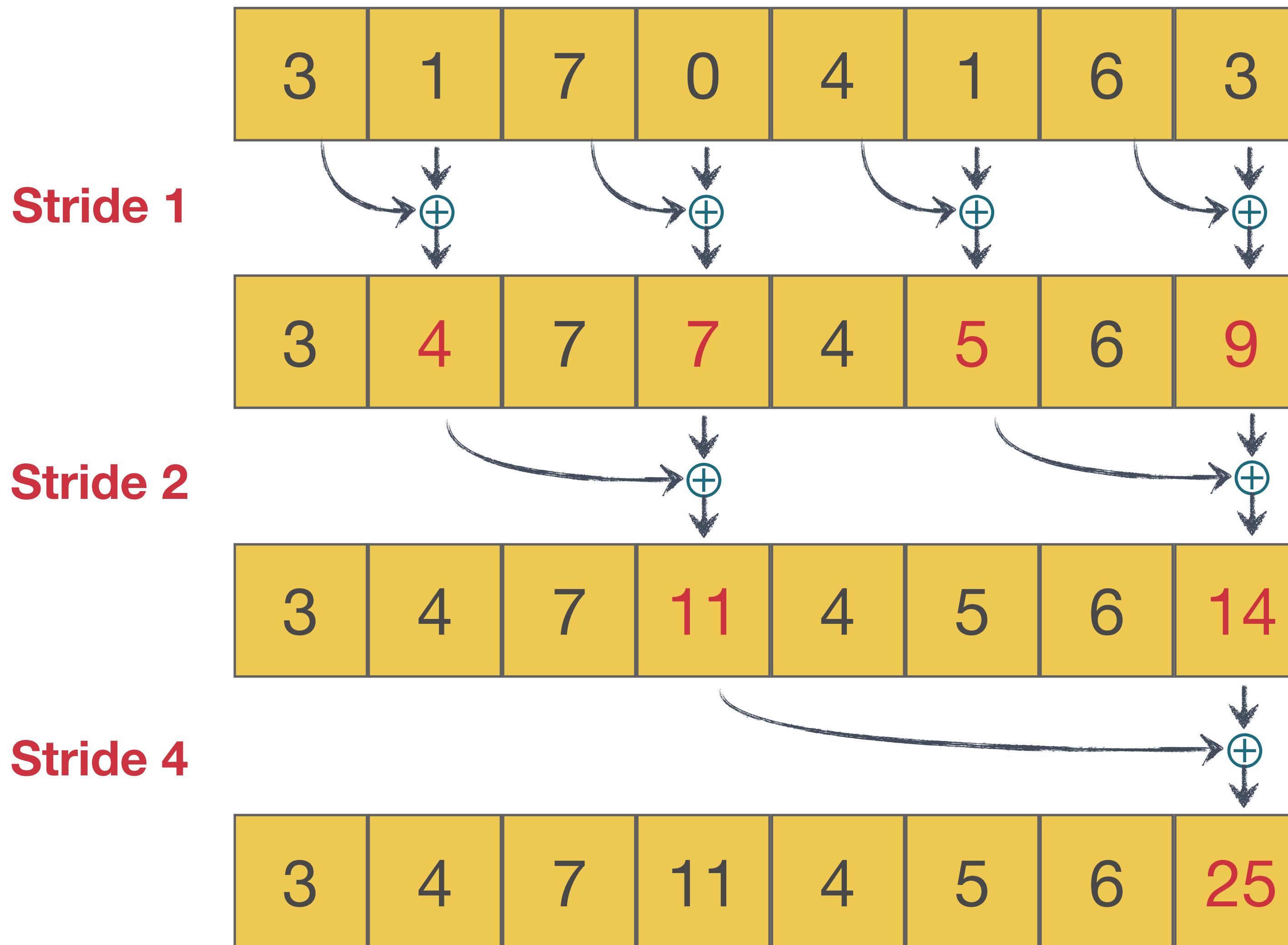
# Parallel Scan: Reduction



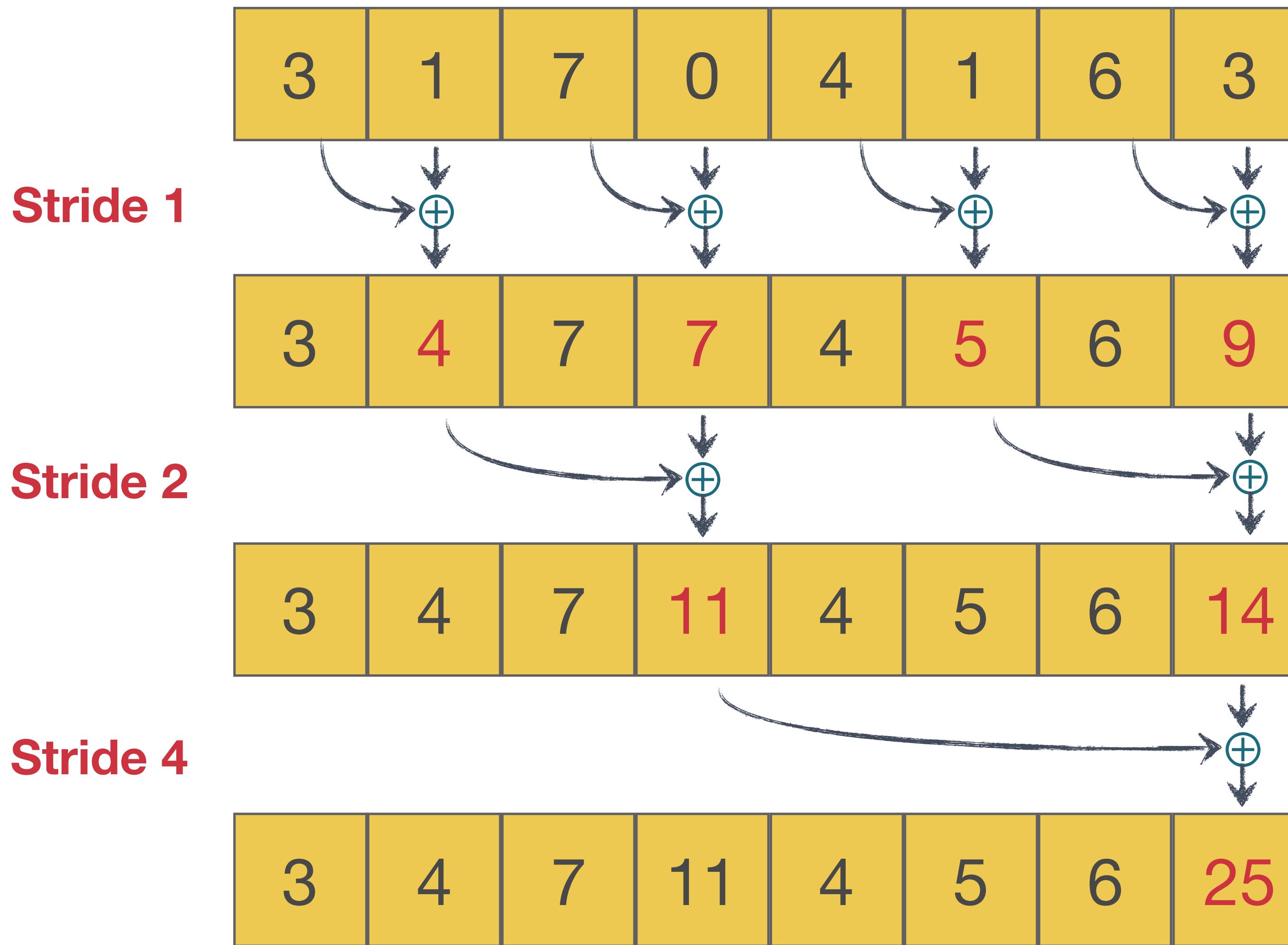
# Parallel Scan: Reduction



# Parallel Scan: Reduction



# Parallel Scan: Reduction



Iterate  $\log(n)$  times  
adding value *stride*  
elements away to  
its own value.

# Parallel Scan: Zero the last element

3	4	7	11	4	5	6	0
---	---	---	----	---	---	---	---

Since this is an **exclusive scan**, set the last element to zero. It will propagate back to the first element.

# Parallel Scan: Post Scan

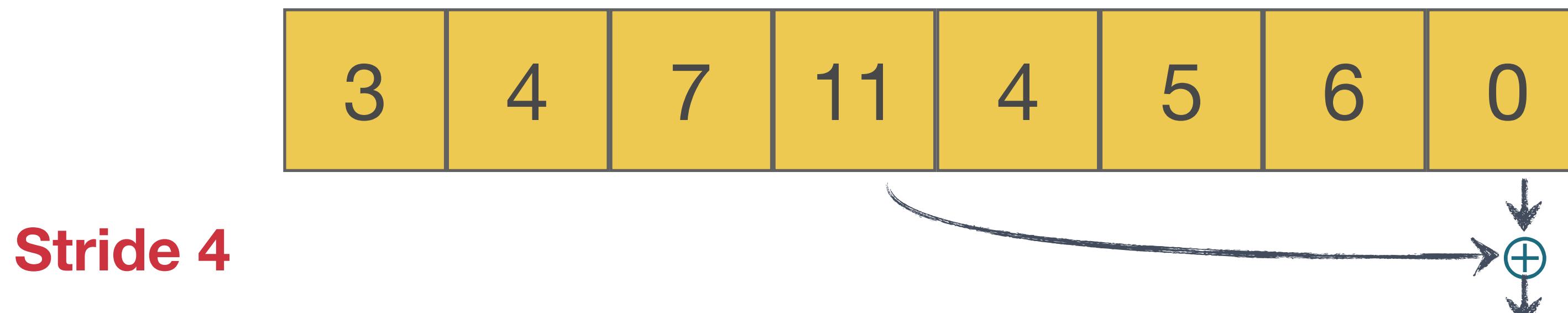
3	4	7	11	4	5	6	0
---	---	---	----	---	---	---	---

# Parallel Scan: Post Scan

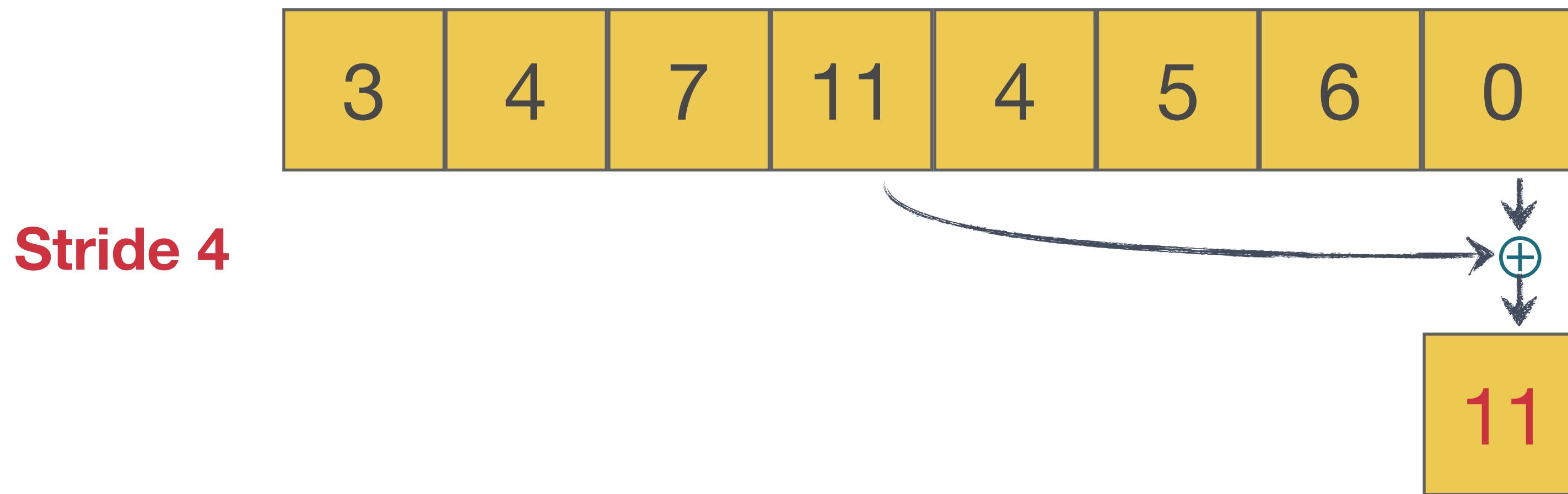
3	4	7	11	4	5	6	0
---	---	---	----	---	---	---	---

**Stride 4**

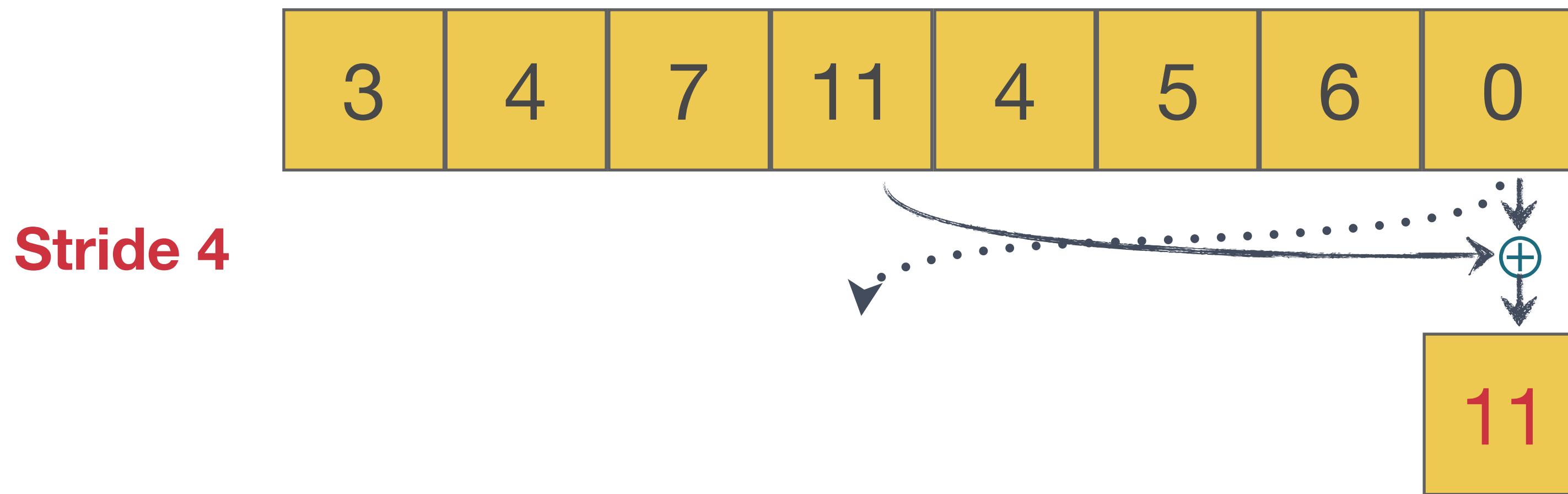
# Parallel Scan: Post Scan



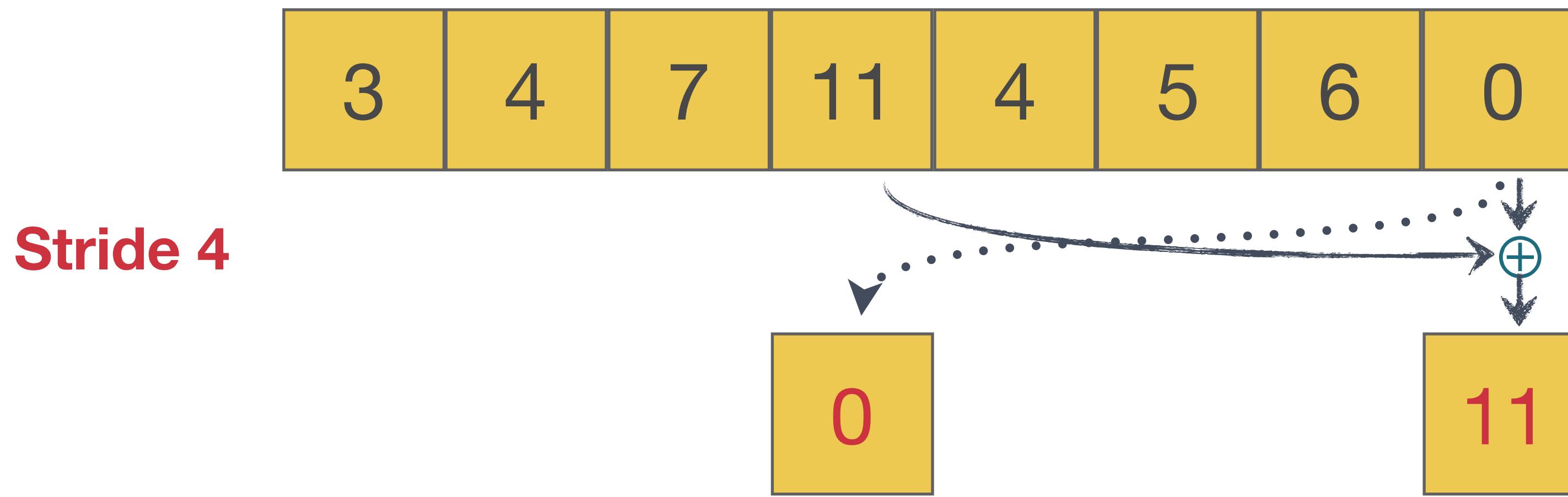
# Parallel Scan: Post Scan



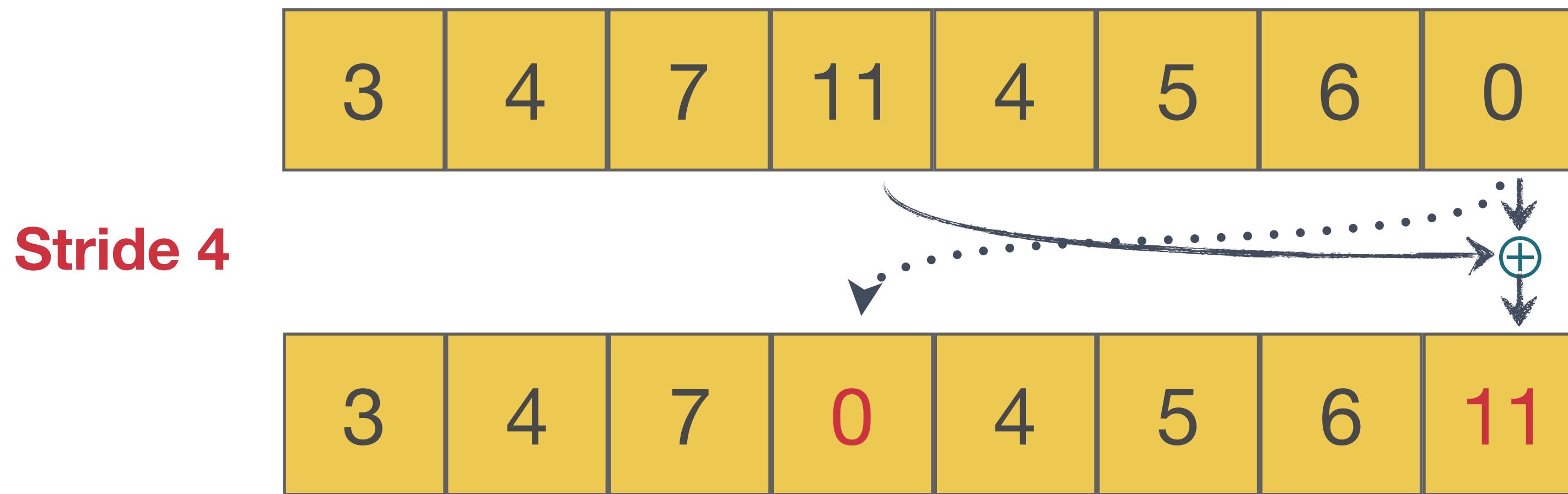
# Parallel Scan: Post Scan



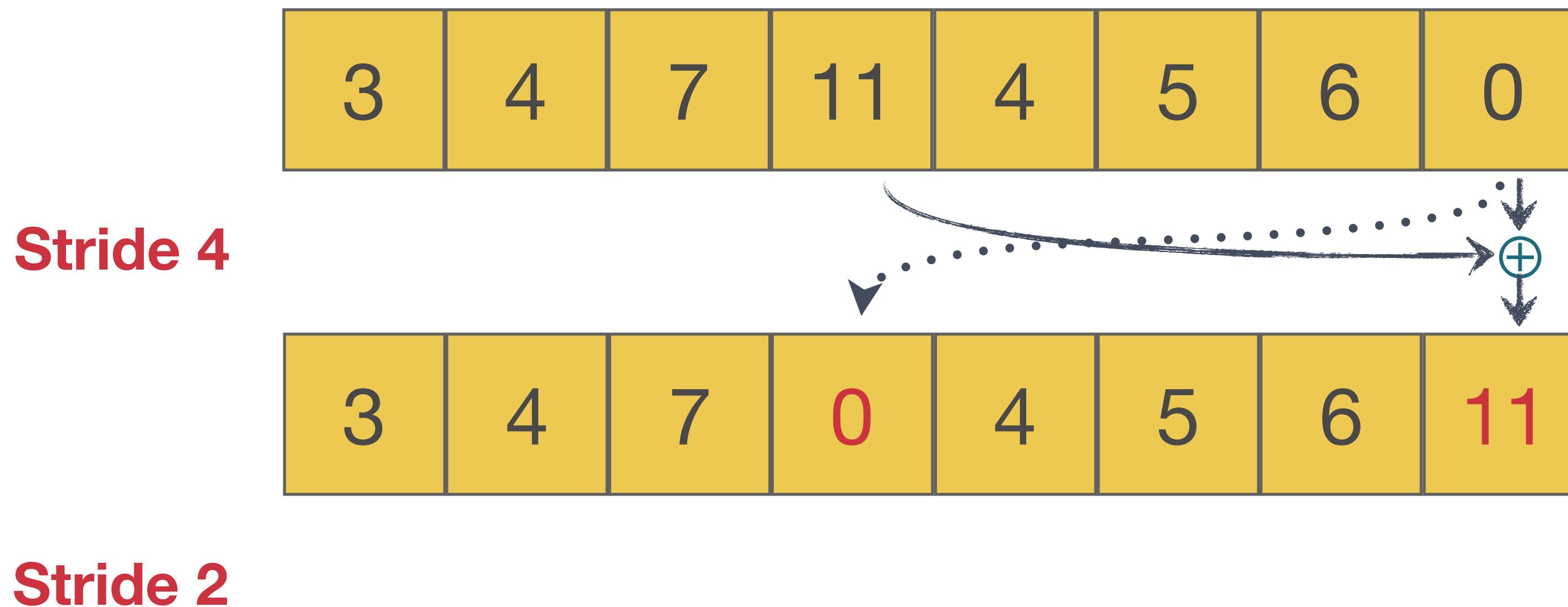
# Parallel Scan: Post Scan



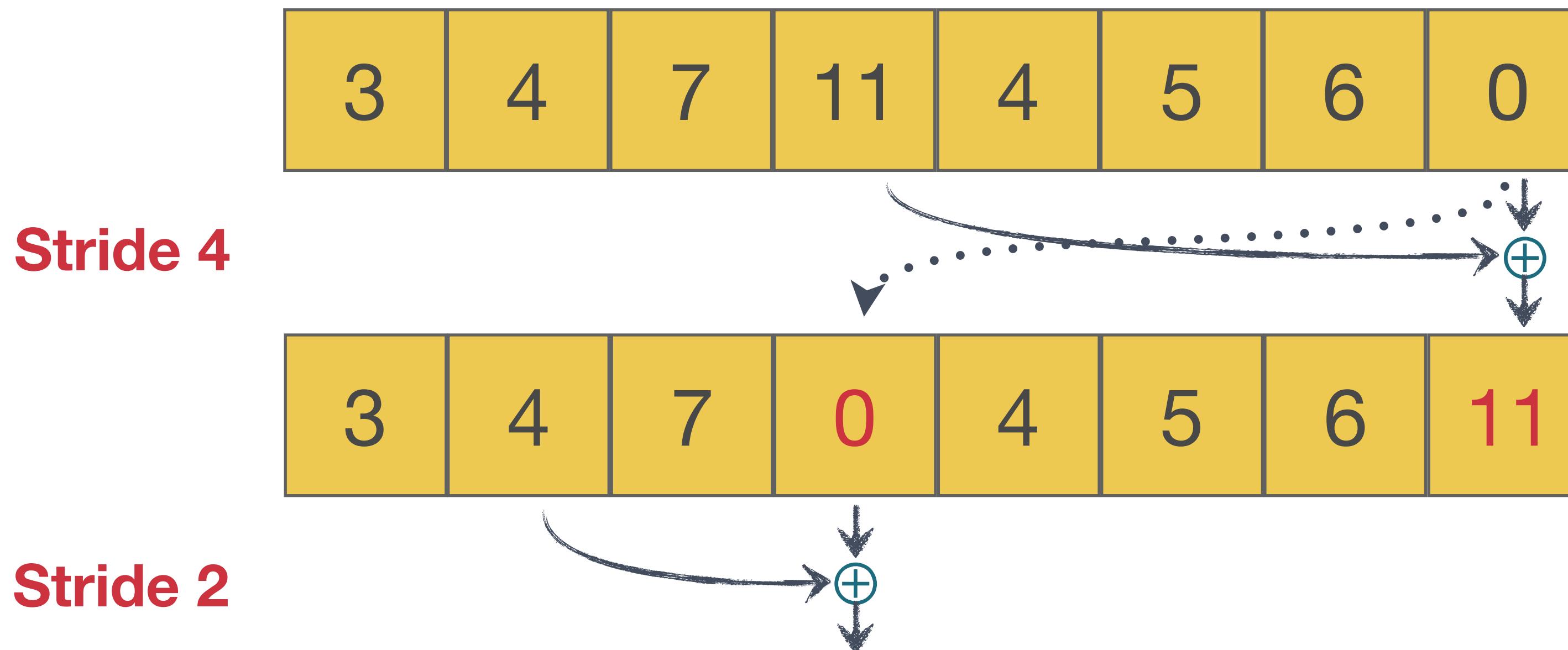
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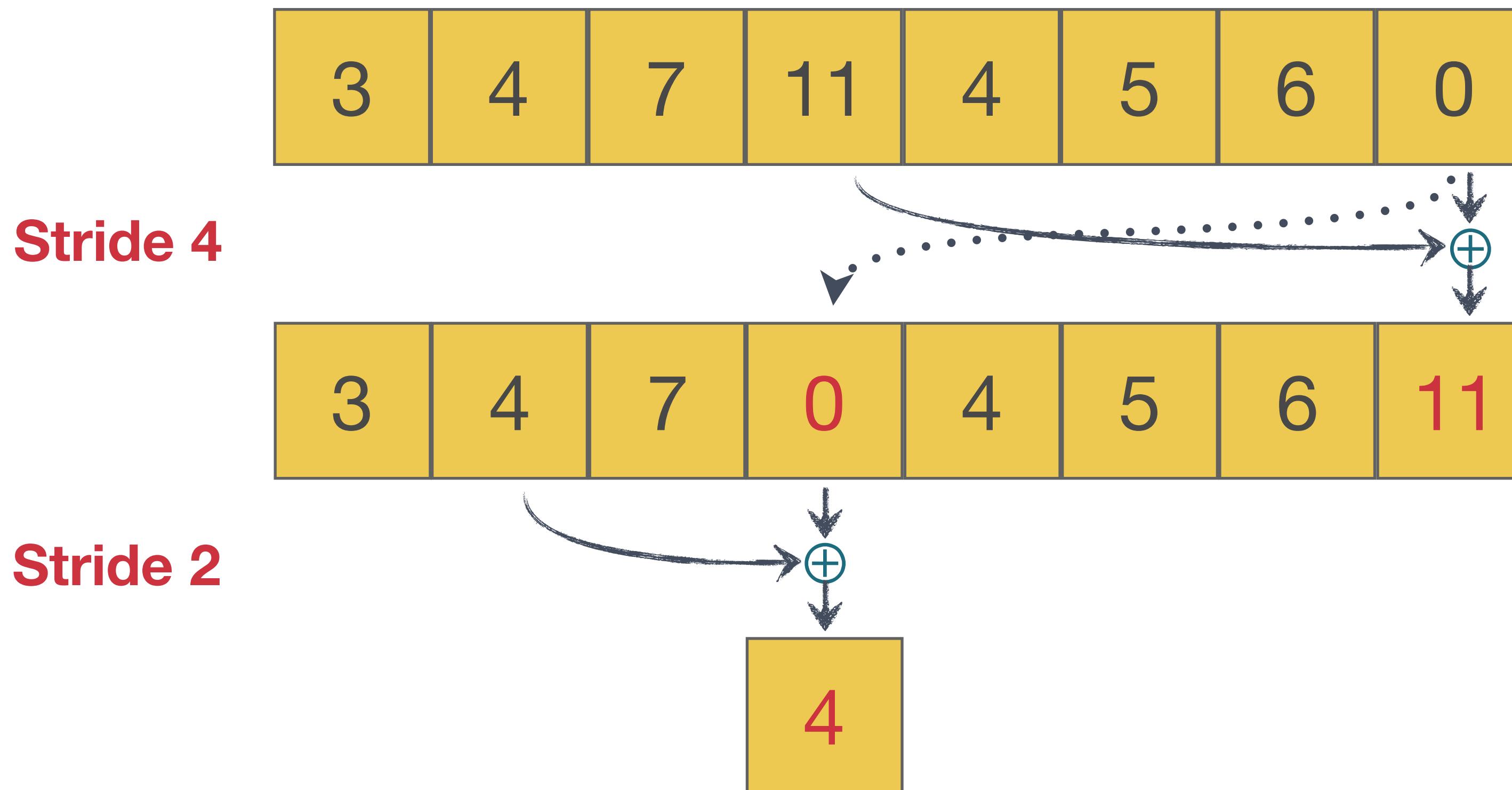
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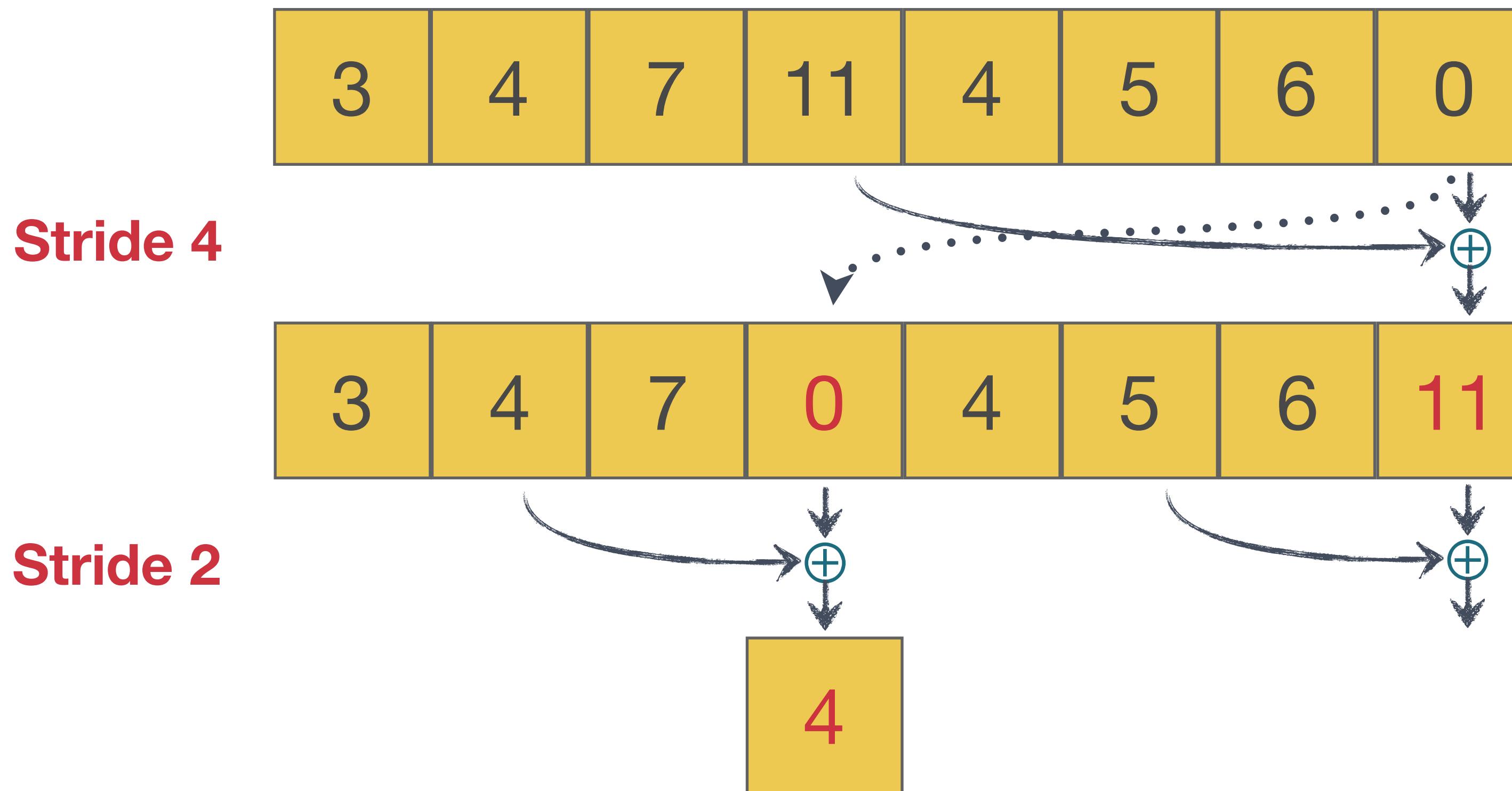
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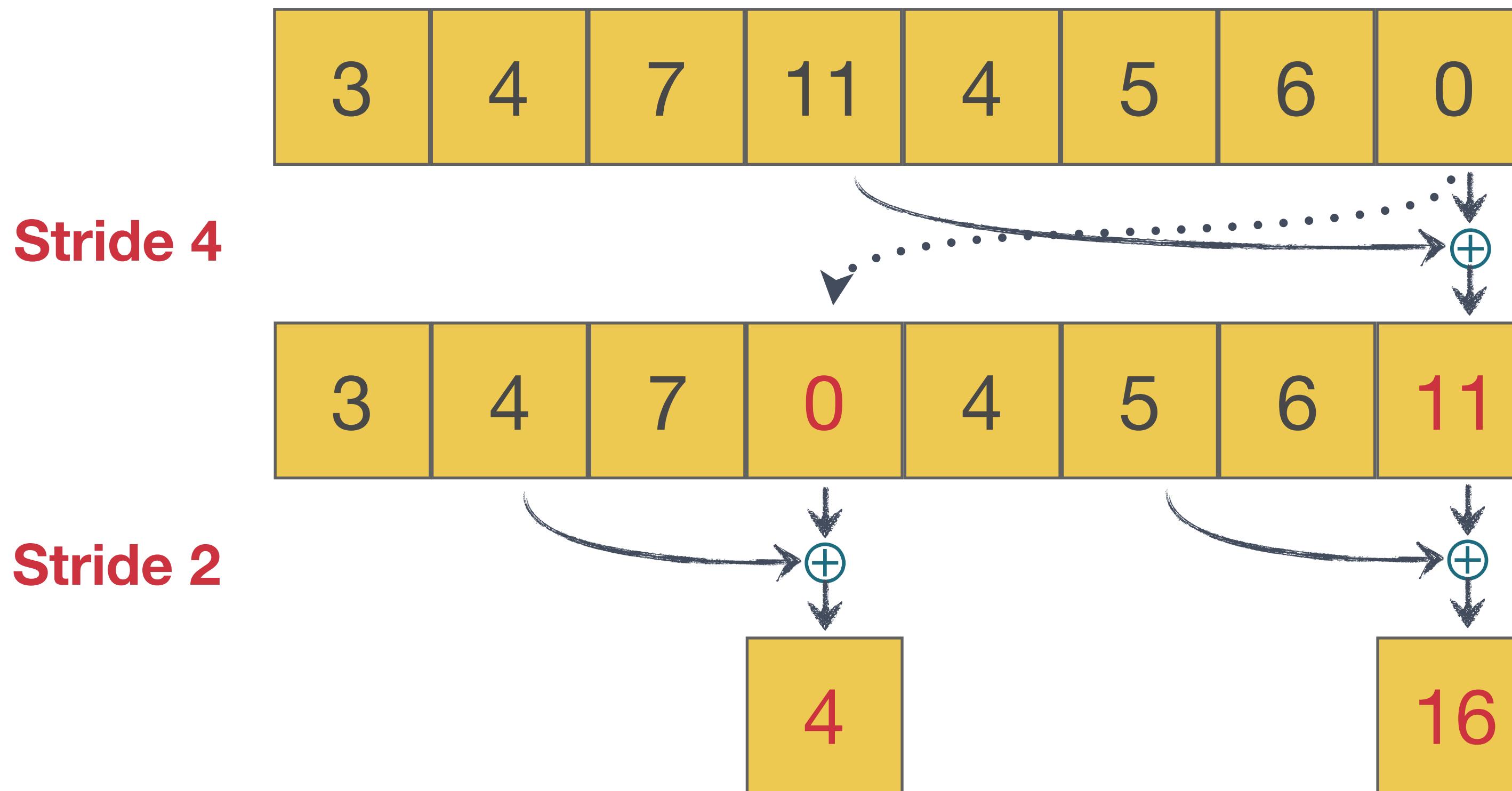
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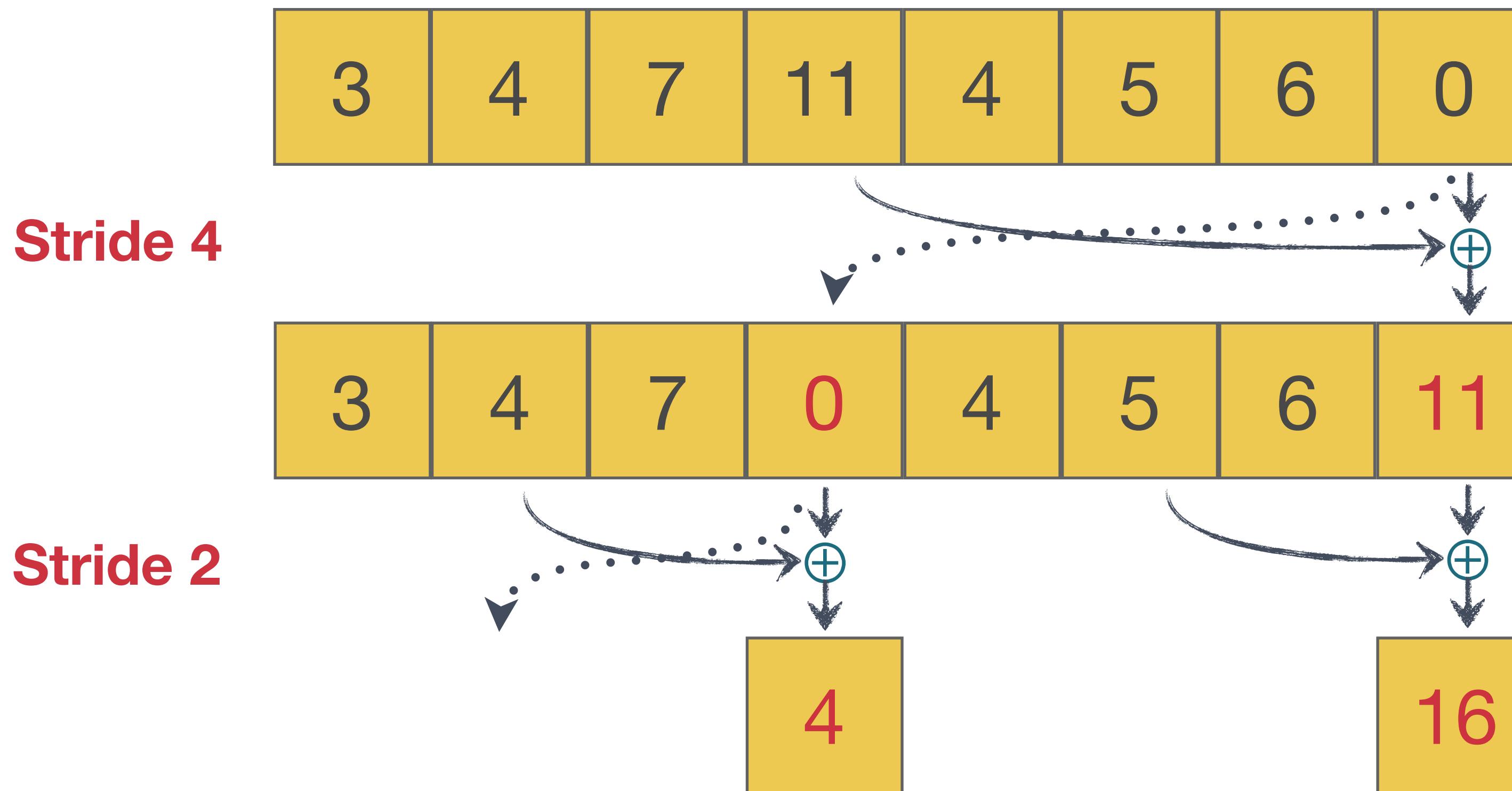
# Parallel Scan: Post Scan



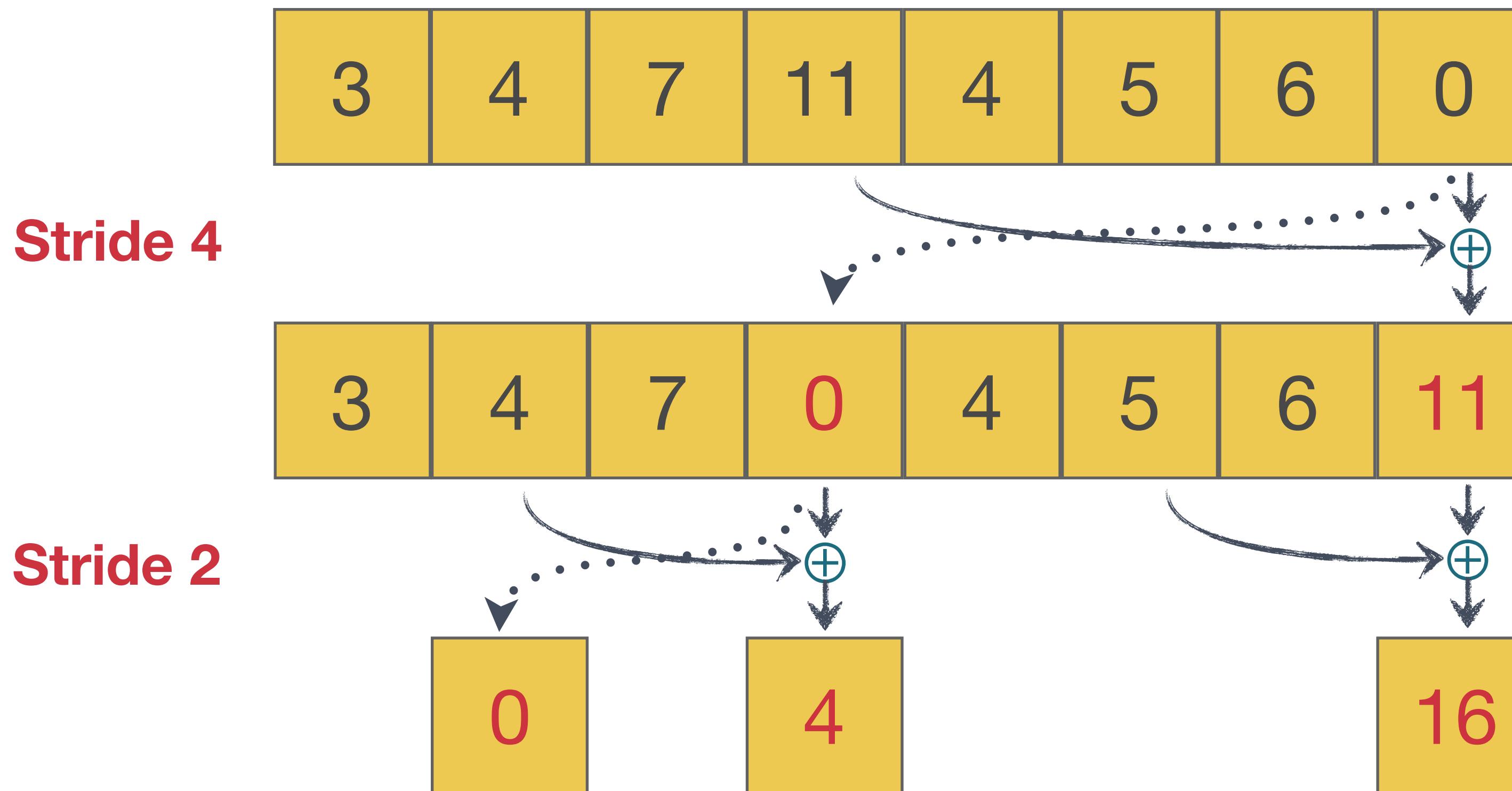
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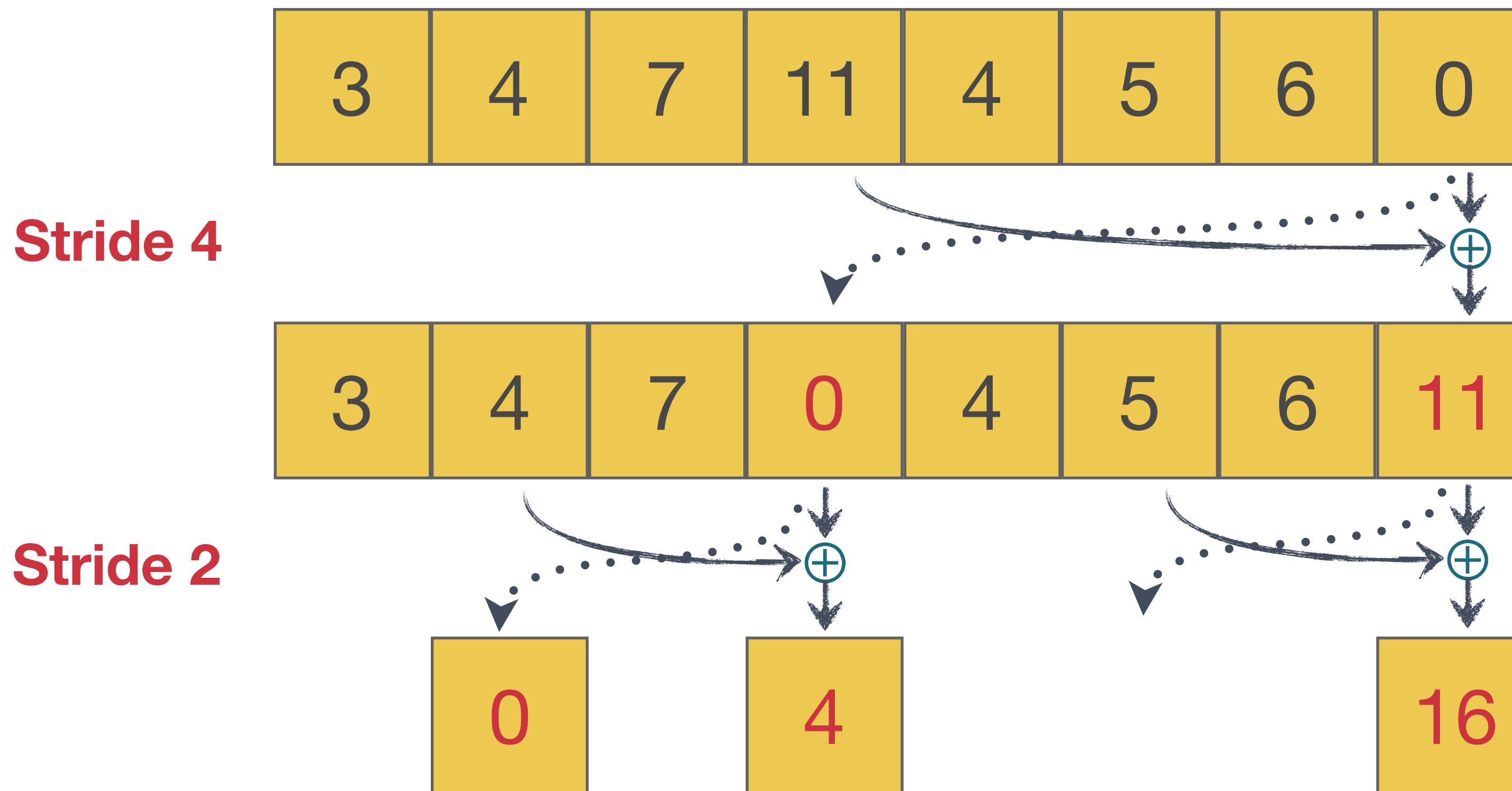
# Parallel Scan: Post Scan



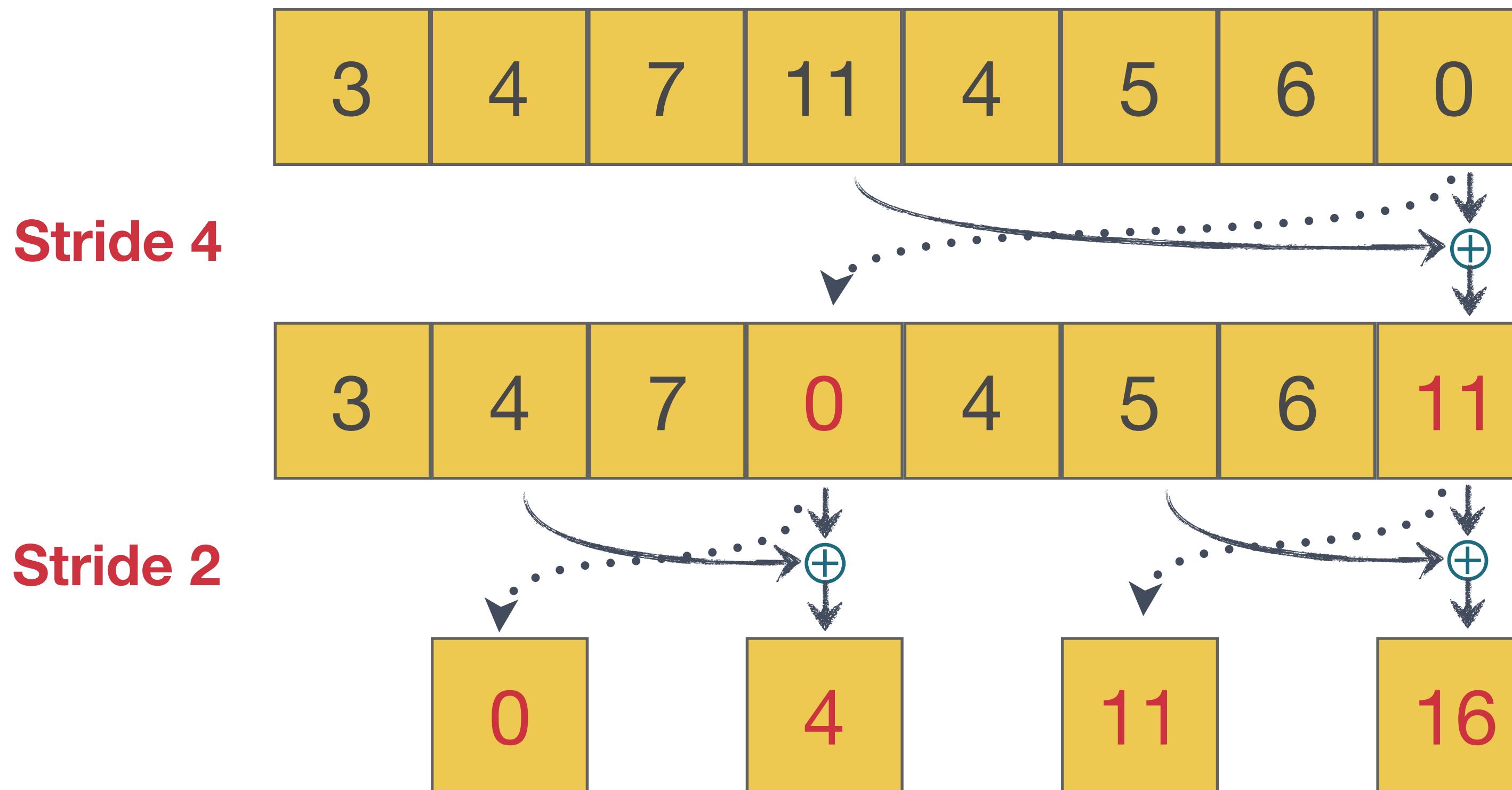
# Parallel Scan: Post Scan



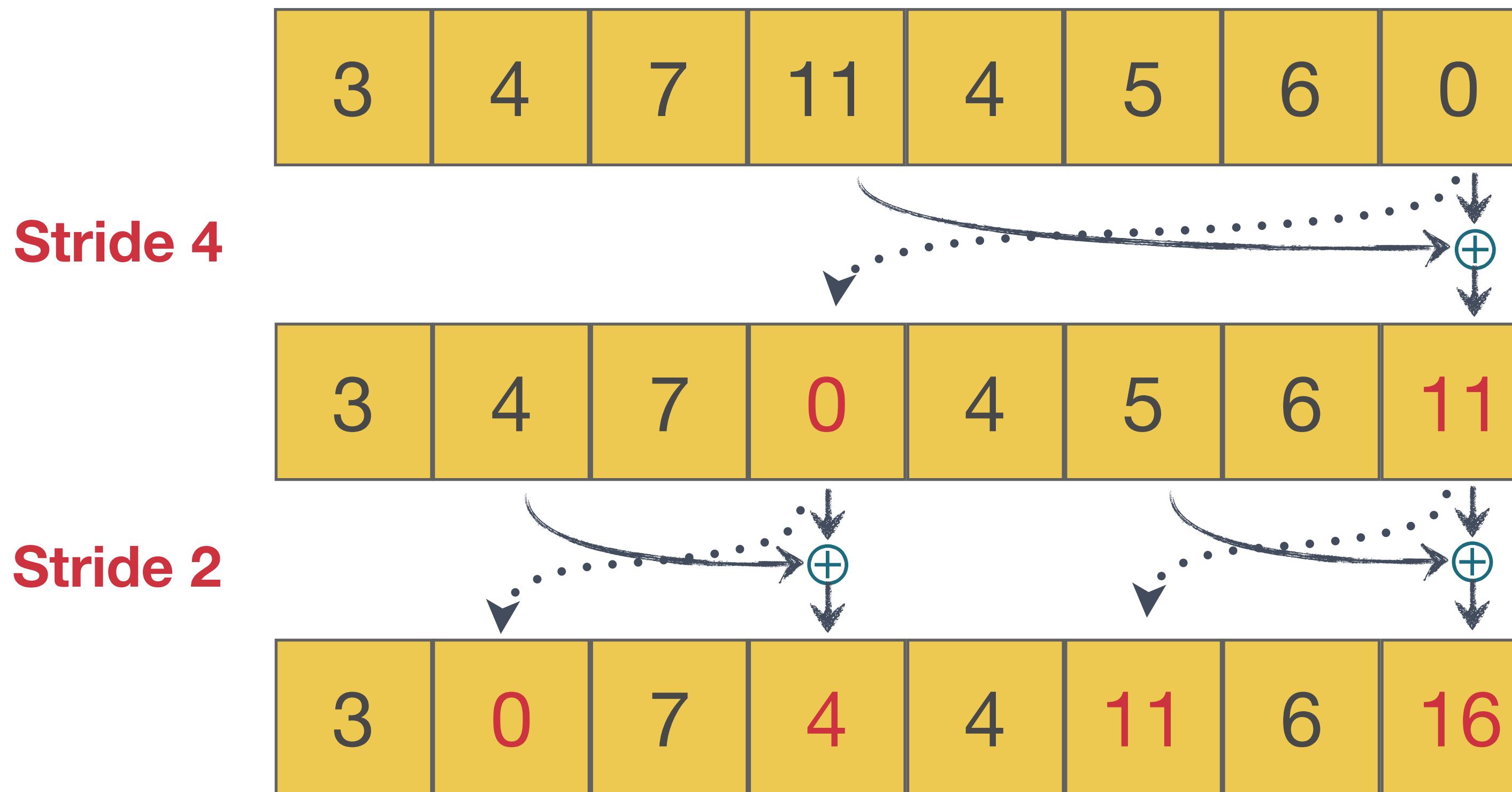
# Parallel Scan: Post Scan



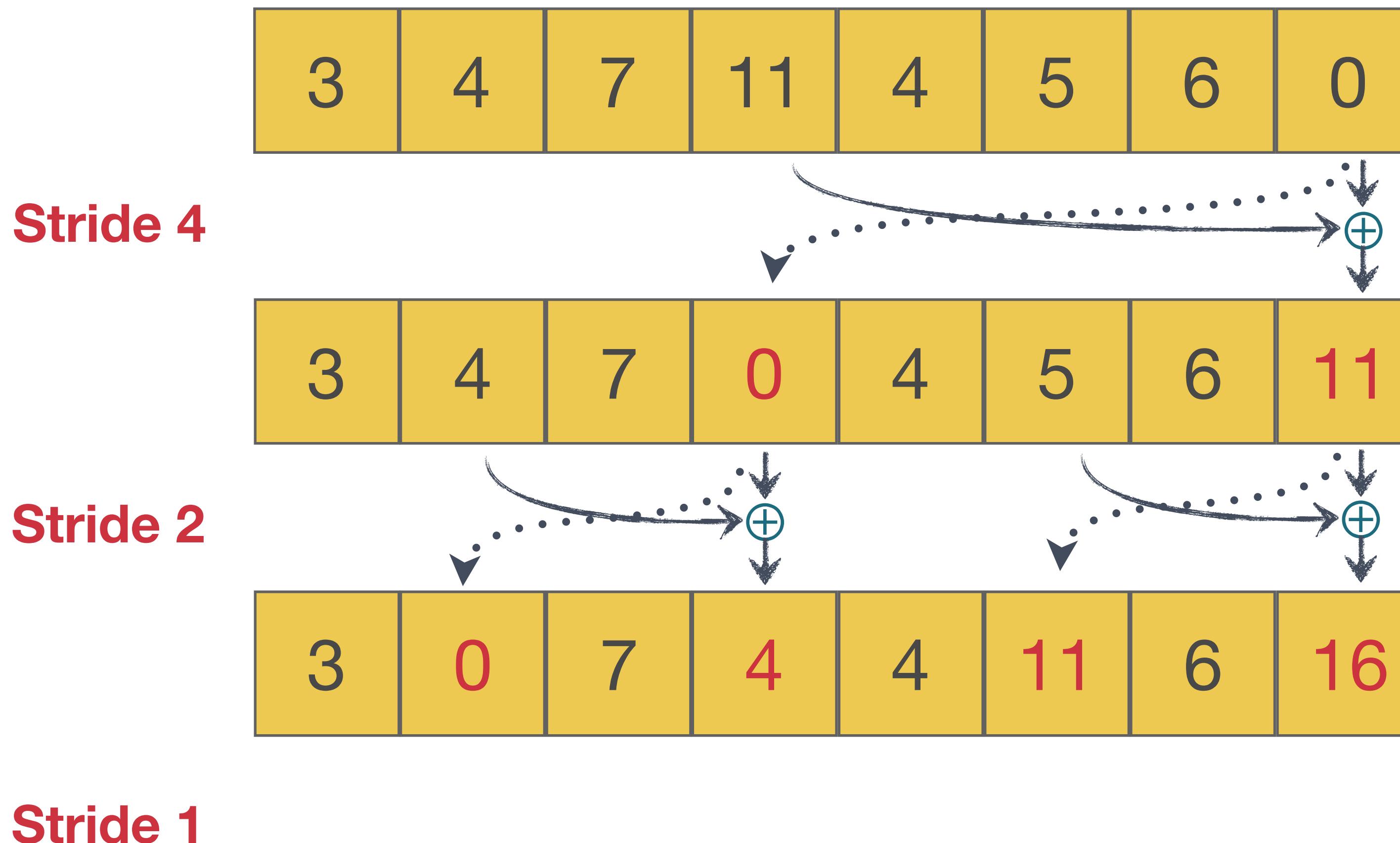
# Parallel Scan: Post Scan



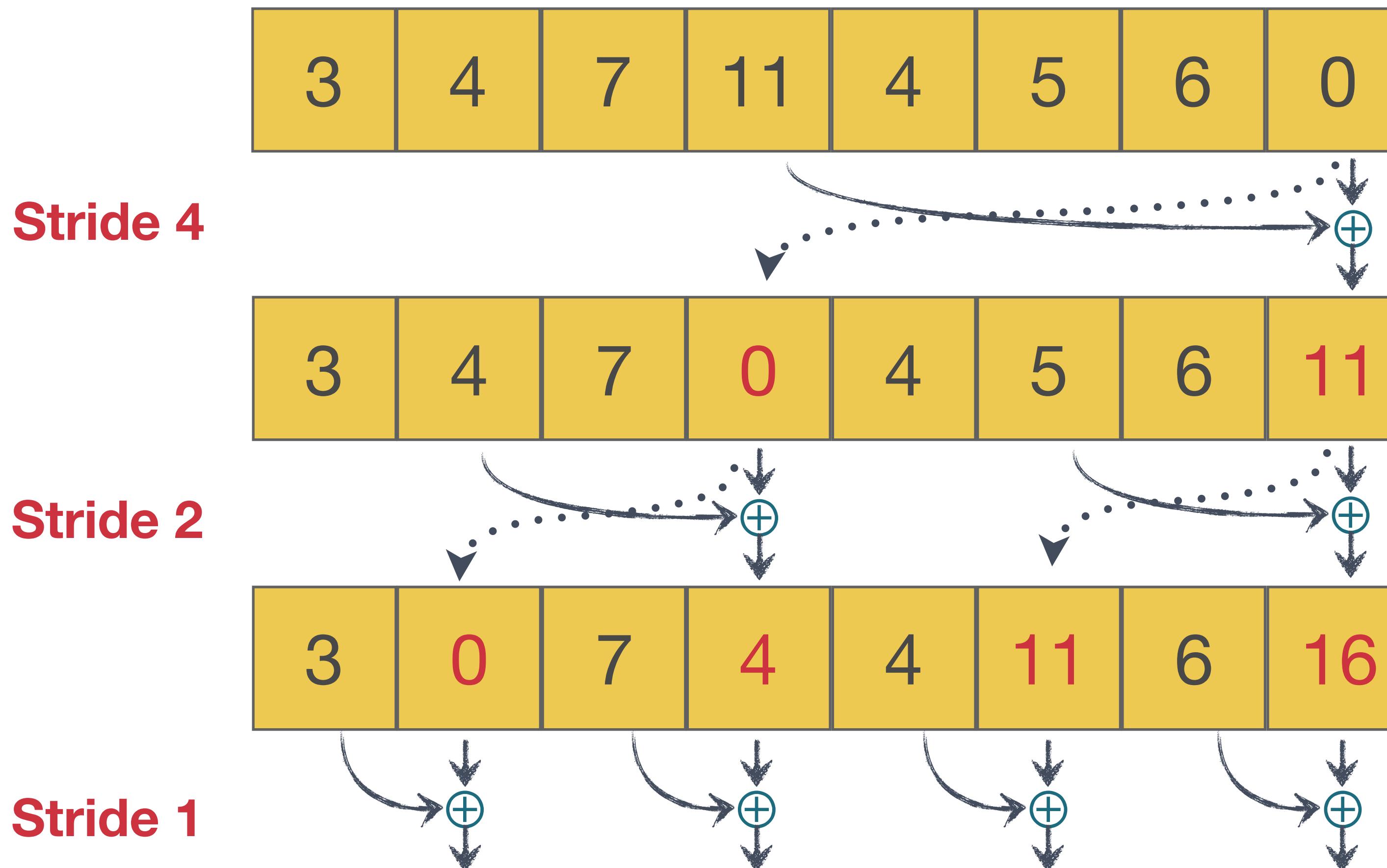
# Parallel Scan: Post Scan



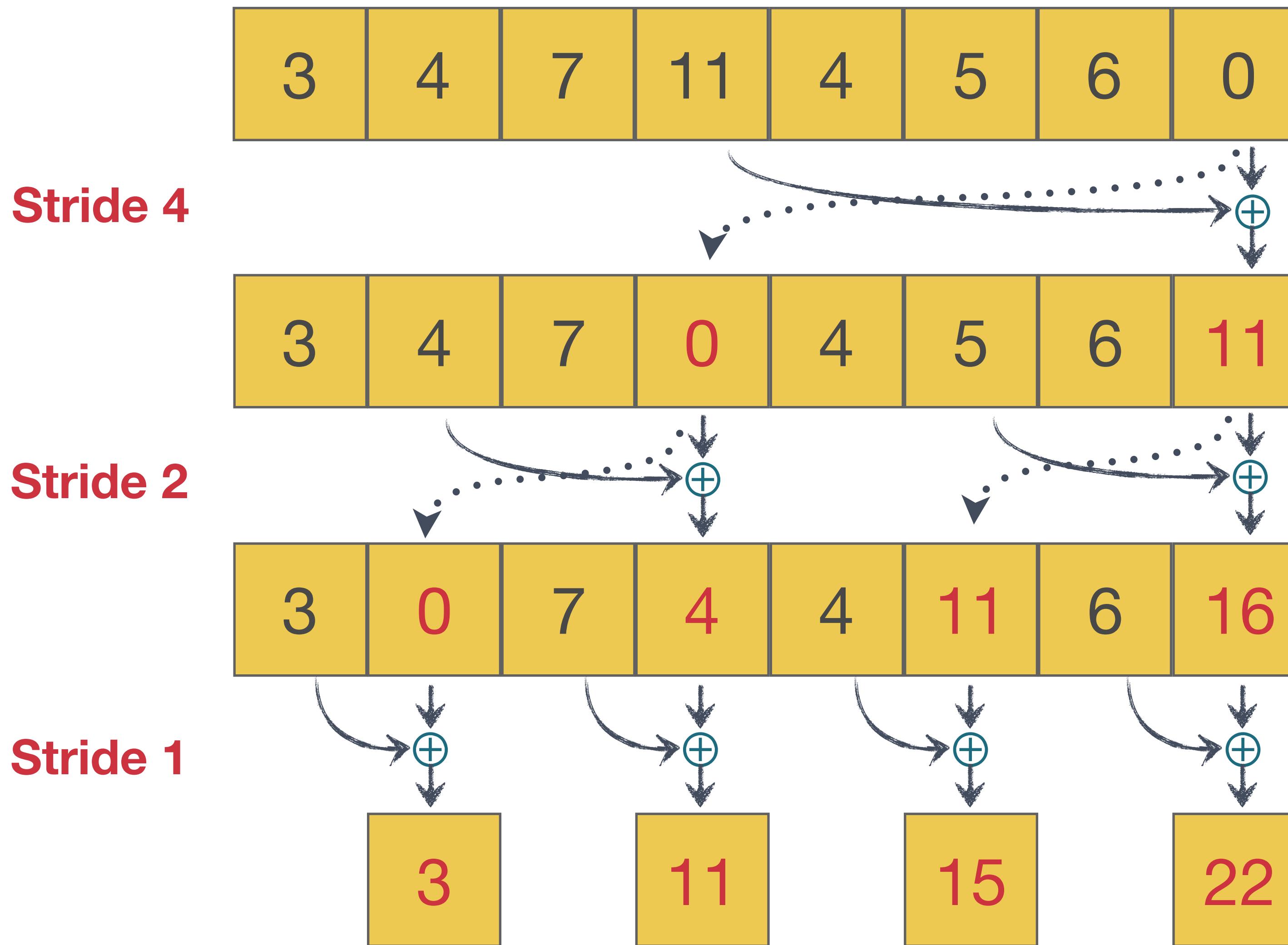
# Parallel Scan: Post Scan



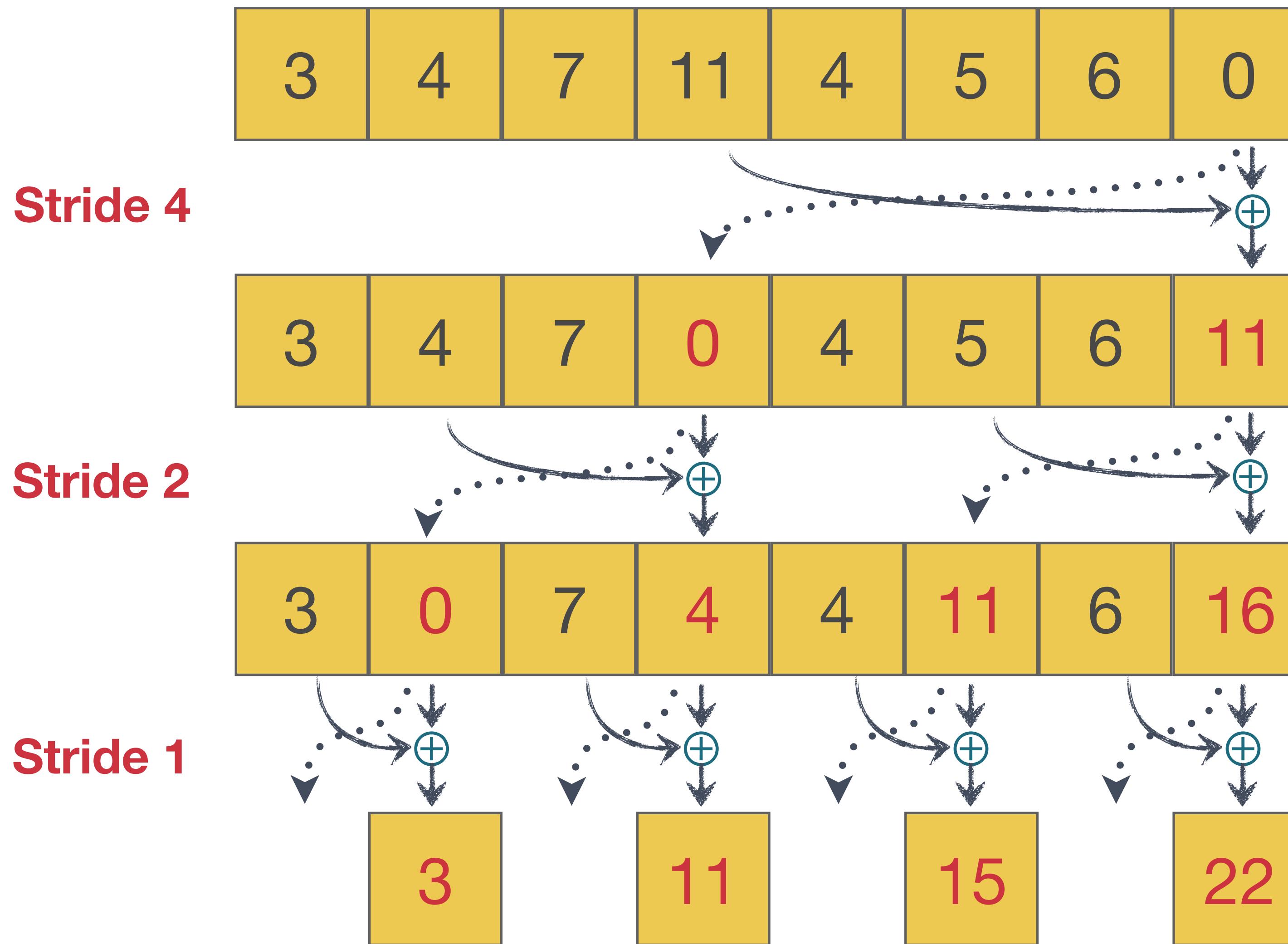
# Parallel Scan: Post Scan



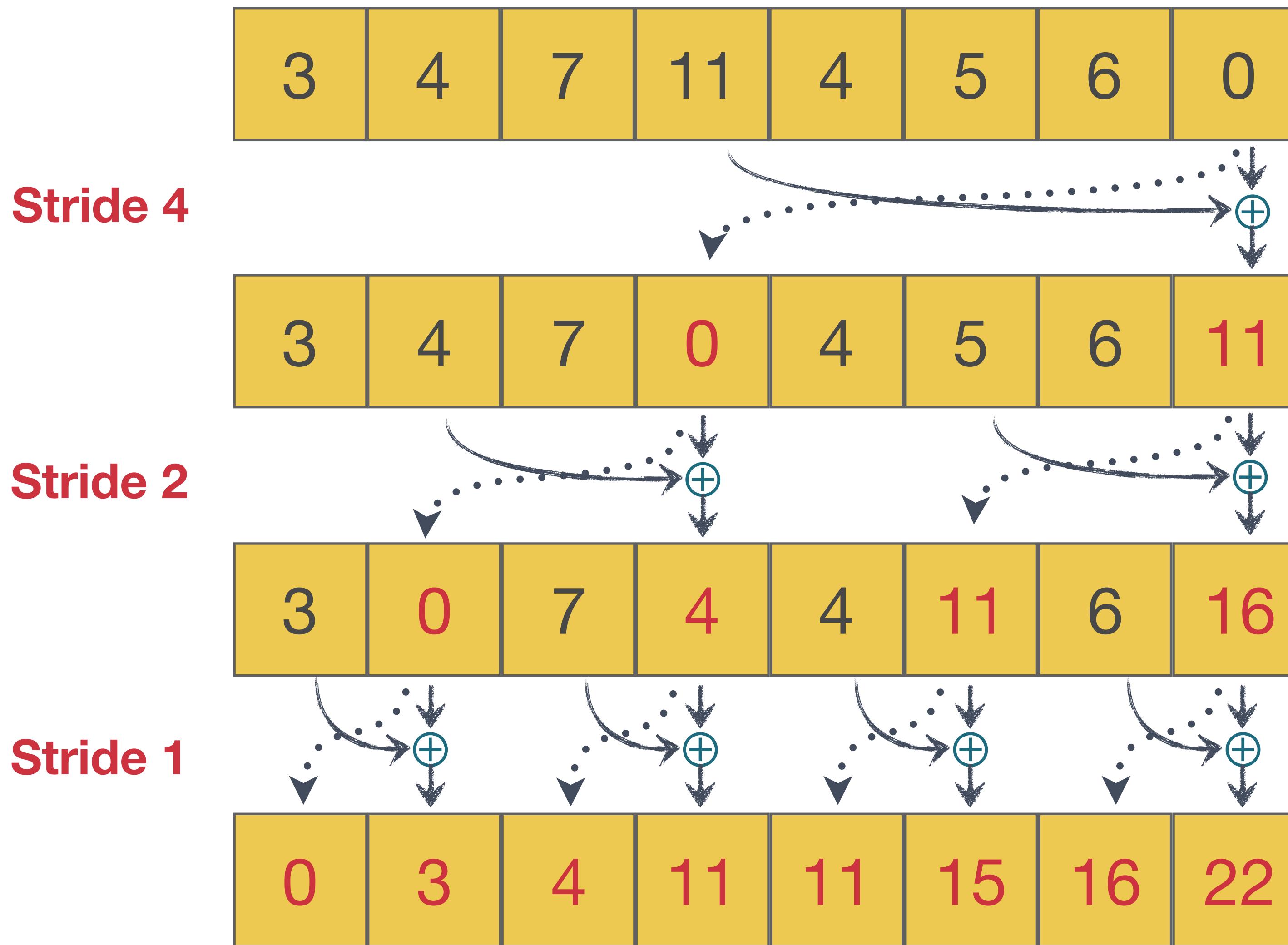
# Parallel Scan: Post Scan



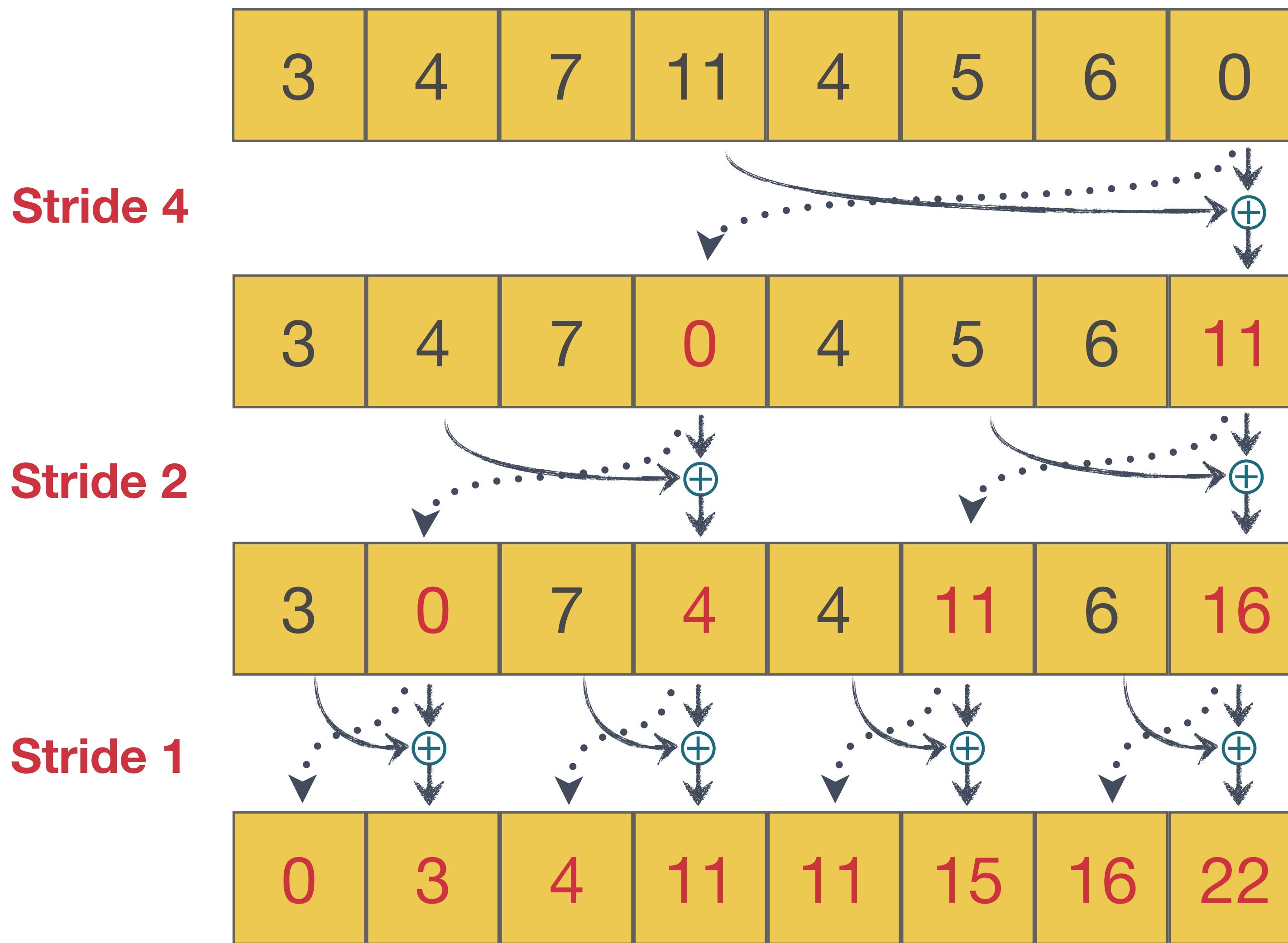
# Parallel Scan: Post Scan



# Parallel Scan: Post Scan



# Parallel Scan: Post Scan



Iterate  $\log(n)$  times adding value *stride* elements away to its own value, and set the value *stride* elements away to its previous value.

$$W_{scan}(n) = \Theta(n)$$

$$D_{scan}(n) = \Theta(\log n)$$



$$p \equiv \frac{W}{D} = \Theta\left(\frac{n}{\log n}\right)$$