

Message Passing Interface (MPI); Collective Communication

EECS 221: Intro to High-Performance Computing

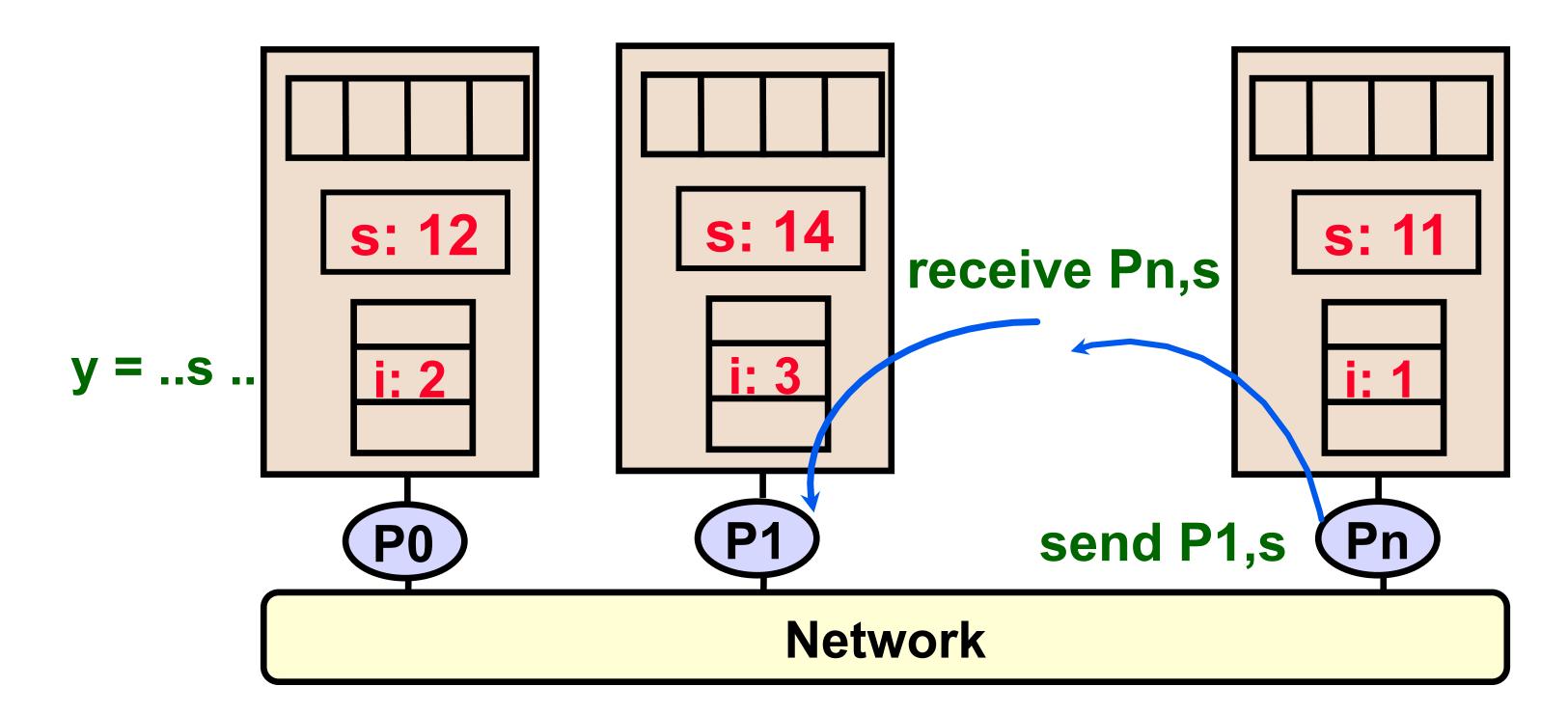
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Administrivia

- ► Homework 2: Hands-on lab on Tuesday May 2nd (next class)
- ► Readings posted on the class website

Message Passing Model

- Program consists of a collection of named processes
- ► Processes communicate by explicit send/receive messages
- ► Coordination is implicit in every communication
- ►MPI has become the de facto standard



Message Passing Interface

- All communication, synchronization require subroutine calls
 - No shared variables
- ► Communication primitives
 - Pairwise, or point-to-point: send & receive ([non]blocking, [a]synchronous)
 - Collectives
 - Move data: Broadcast, Scatter/gather
 - Compute and move: sum, product, max, prefix sum, etc,..
- ► Barrier synchronization
 - No locks because there are no shared variables to protect
- See: http://www.mpi-forum.org and http://www.mcs.anl.gov/research/projects/mpi

MPI environment

Two important questions arise in a parallel program

How many processes are participating in this computation?

► Which one am I?

MPI environment

Two important questions arise in a parallel program

► How many processes are participating in this computation? MPI_Comm_size returns the number of processors, *size*

► Which one am I? MPI_Comm_rank reports the *rank*, a number between **0** and **size-1**, identifying the calling process

```
int main(int argc, char *argv[])
 int rank, size;
 printf ("I am %d of %d\n", rank, size); // Execute in parallel
 return 0;
```

```
#include <mpi.h>
int main(int argc, char *argv[])
 int rank, size;
 printf ("I am %d of %d\n", rank, size);
 return 0;
```

```
#include <mpi.h>
int main(int argc, char *argv[])
 int rank, size;
 MPI_Init (&argc, &argv);
 printf ("I am %d of %d\n", rank, size);
 MPI_Finalize ();
 return 0;
```

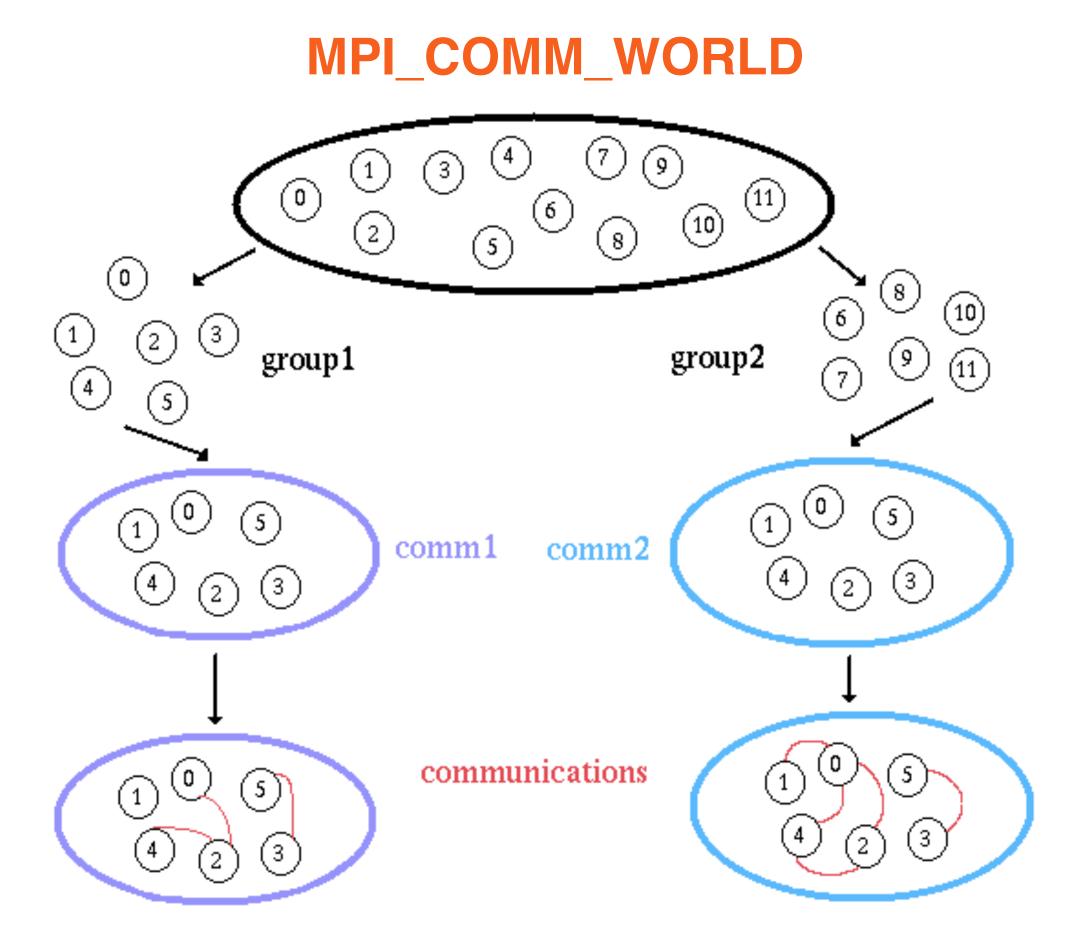
```
#include <mpi.h>
int main(int argc, char *argv[])
 int rank, size;
 MPI_Init (&argc, &argv);
 MPI_Comm_rank (MPI_COMM_WORLD, &rank);
 MPI_Comm_size (MPI_COMM_WORLD, &size);
 printf ("I am %d of %d\n", rank, size);
 MPI_Finalize ();
 return 0;
```

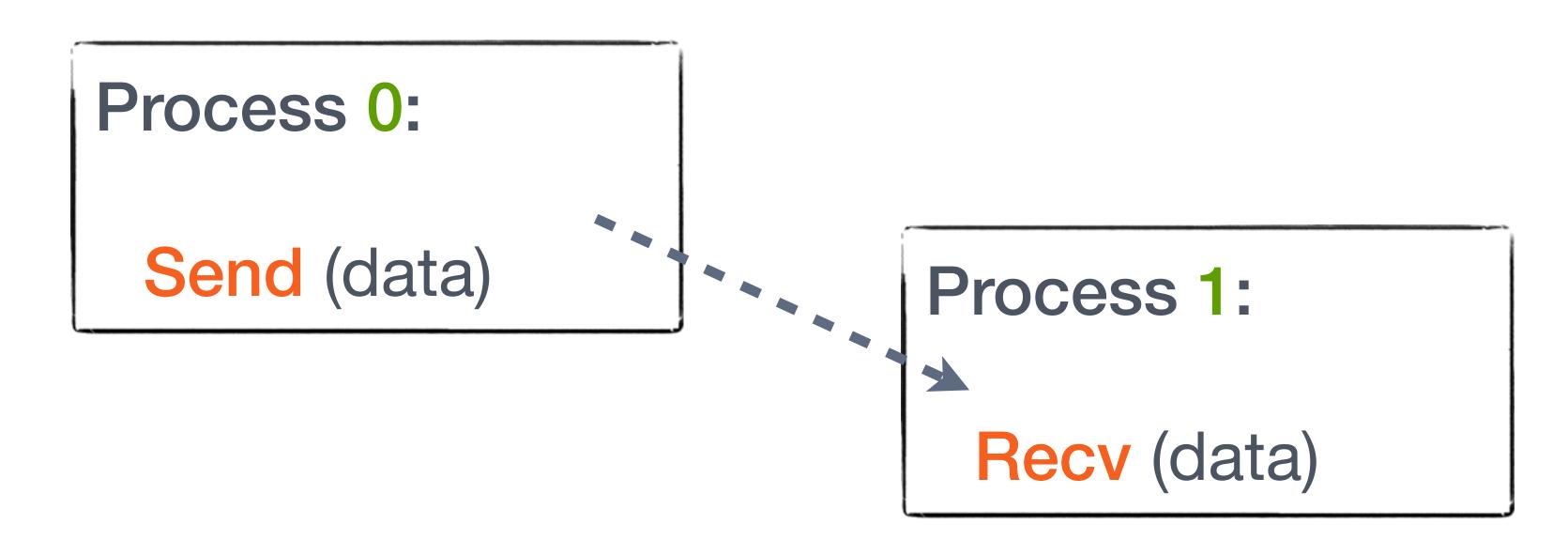
```
#include <mpi.h>
                                          Compiling:
int main(int argc, char *argv[])
                                           $ mpicc -o hello hello.c
 int rank, size;
 MPI_Init (&argc, &argv);
 MPI_Comm_rank (MPI_COMM_WORLD, &rank);
 MPI_Comm_size (MPI_COMM_WORLD, &size);
 printf ("I am %d of %d\n", rank, size);
 MPI_Finalize ();
 return 0;
```

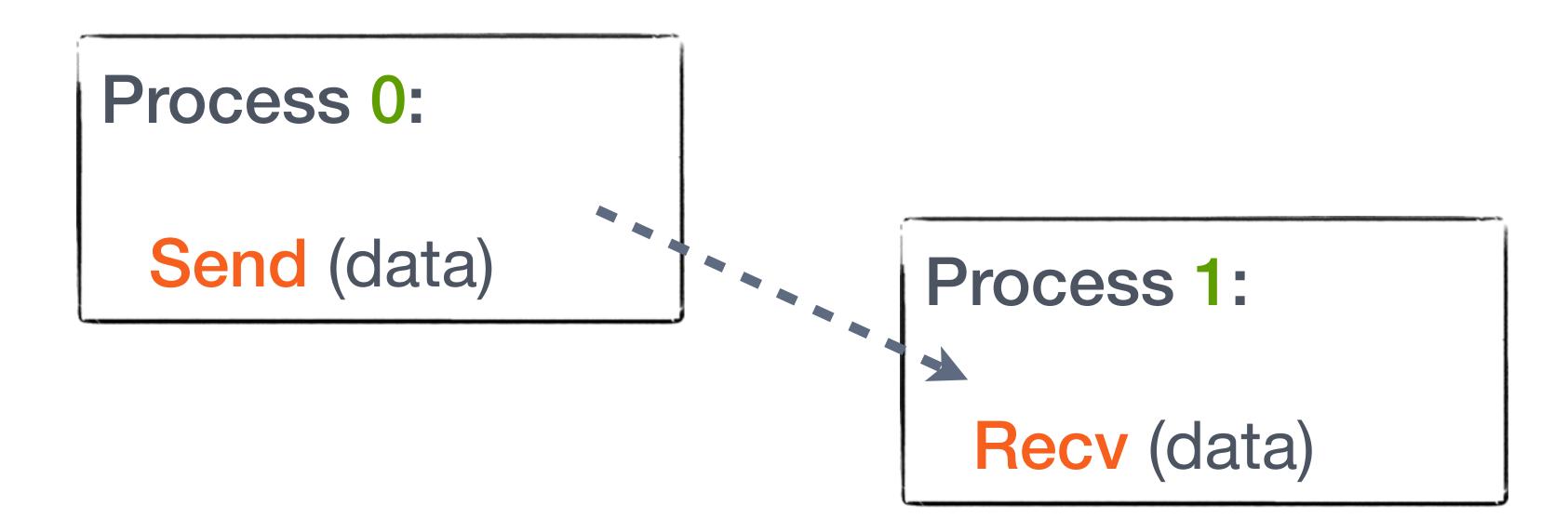
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#include <mpi.h>
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 int rank, size;
 MPI_Init (&argc, &argv);
 MPI_Comm_rank (MPI_COMM_WORLD, &rank);
 MPI_Comm_size (MPI_COMM_WORLD, &size);
 printf ("I am %d of %d\n", rank, size);
                                         Run:
 MPI_Finalize ();
                                          $ mpirun -np 4 ./hello
 return 0;
```

Basic concept: Communicators

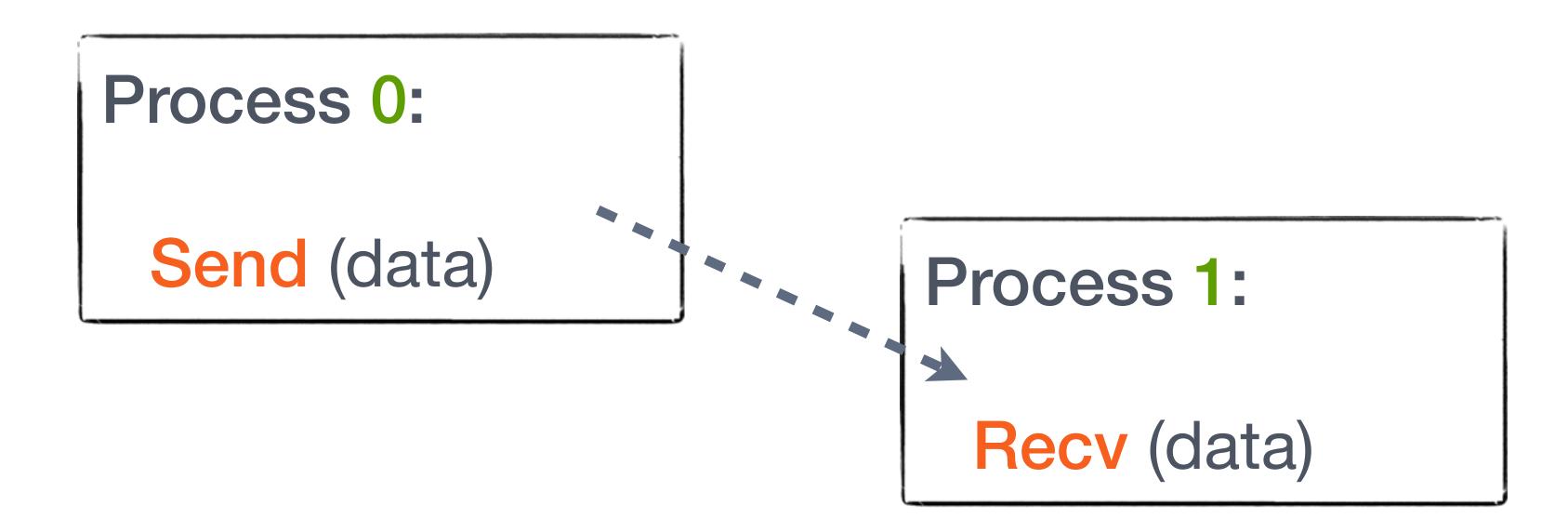
- ► Group: collection of processors
- Each message is sent in acontext, and must be received in the same context
- ► Communicator = group + context
- Rank: process ID in its communicator
- ►MPI_COMM_WORLD = Group consisting of all processes
- ►MPI_ANY_RANK = Wildcard rank (receive only)







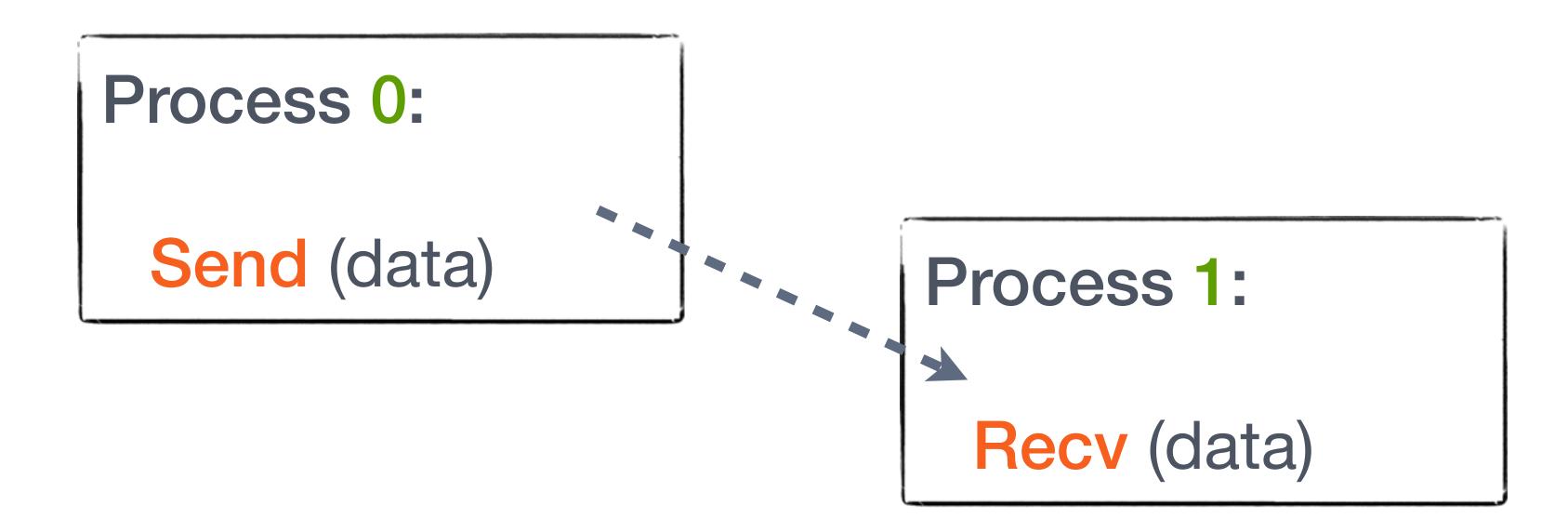
- ► How to identify tasks/processes?
- ► How to describe "data"?
- ► How will receiver recognize and screen messages?
- ► What does it mean for operations to complete?



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Basic concept: MPI datatypes

- In MPI call, "data" is described a triple: (address, count, type)
- ► Datatype = (recursively defined)
 - "Standard" predefined scalar types: MPI_INT, MPI_DOUBLE, MPI_CHAR, etc,.
 - A contiguous array of datatypes
 - A strided block of datatypes
 - An indexed array of blocks of datatypes
 - An arbitrary structure of datatypes
- Can construct custom datatypes, in particular ones for subarrays



- ► How to identify tasks/processes?
- ► How to describe "data"?
- ► How will receiver recognize and screen messages?
- ► What does it mean for operations to complete?

Basic concepts: MPI tags

Message tags

- Every message has a user-defined integer ID to assist the receiver in identifying the message
- Messages can be screened at the receiver by specifying a tag
- ► Wildcard: MPI_ANY_TAG

Basic concept: Status object

►Opaque structures that contains more information (e.g. size of the message, error)

```
int received_tag, received_from, received_count;
MPI_Status status;
MPI_Recv (..., MPI_ANY_SOURCE, MPI_ANY_TAG, ..., &status);
received_tag = status.MPI_TAG;
received_from = status.MPI_SOURCE;
MPI_Get_count (status, datatype, &received_count);
```

Blocking send

MPI_Send (buffer, count, datatype, dest-rank, tag, comm)

Process 0:

MPI_Send (A, 10, MPI_INT, 1, tag, comm);

Process 1:

MPI_Recv (B, 20, MPI_INT, 0, tag, comm, &stat);

- ► Message buffer = (buffer-start, count, datatype)
- ► Target = (dest-rank, tag, comm)
- On return:
 - ► Data delivered to the system
 - ► Buffer can be reused
 - ► Target may not yet have received the message

Blocking receive

MPI_Recv (buffer, count, datatype, source-rank, tag, comm, status)

```
Process 0:

MPI_Send (A, 10, MPI_INT, 1, tag, comm);
```

Process 1: MPI_Recv (B, 20, MPI_INT, 0, tag, comm, &stat);

- ► Message buffer = (buffer-start, count, datatype) Source = (source-rank, tag, comm)
- Returns when matching message (match on source triplet) is received
 - Source can be a wildcard, MPI_ANY_SOURCE
 - ► Tag can be a wildcard, MPI_ANY_TAG
 - Receiving fewer than *count* items is OK, but more is an error
 - May query status for more information (e.g. size of the message, source rank)

Simple MPI program

```
#include <mpi.h>
int main(int argc, char *argv[])
 int rank, buffer;
 MPI_Status status;
 MPI_Init (&argc, &argv);
 MPI_Comm_rank (MPI_COMM_WORLD, &rank);
 if (rank == 0) {
   buffer = 123;
   MPI_Send (&buffer, 1, MPI_INT, 1, 0, MPI_COMM_WORLD);
 else if (rank == 1) {
   MPI_Recv (&buffer, 1, MPI_INT, 0, 0, MPI_COMM_WORLD, &status);
   printf ("Received %d\n", buffer);
 MPI_Finalize ();
 return 0;
```

Process 0:

Recv (data ← 1)

Send (data → 1)

Process 1:

Recv (data ← 0)

Send (data → 0)

Process 0:

Recv (data ← 1)

Send (data → 1)

Process 1:

Recv (data ← 0)

Send (data → 0)

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Send (data → 1)

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Process 1:

Send (data → 0)

Recv (data ← 0)

Process 0:

Recv (data ← 1)

Send (data → 1)

Process 1:

Recv (data ← 0)

Send (data → 0)

Process 0:

Send (data → 1)

Recv (data ← 1)

Process 1:

Send (data → 0)

Recv (data ← 0)

- ► Unsafe orderings of send/recv
- ► How to avoid?

►Use safe orderings of send/recv

Process 0:

Send (data → 1)

Recv (data ← 1)

Process 1:

Recv (data ← 0)

Send (data → 0)

►Use simultaneous send/recv

Process 0:

SendRecv (data → 1)

Process 1:

SendRecv (data → 0)

MPI_SendRecv (Sbuffer, Scount, Stype, dest, Stag, Rbuffer, Rcount, Rtype, src, comm, &stat)

► Unsafe orderings of send/recv

Process 0: Send (data → 1) Recv (data ← 1)

```
Process 1:
Send (data → 0)
Recv (data ← 0)
```

►Use non-blocking send/recv

```
Process 0:
    Isend (data1 → 1)
    Irecv (data2 ← 1)
    Waitall
```

```
Process 1:

Isend (data1 → 0)

Irecv (data2 ← 0)

Waitall
```

```
MPI_Isend (buffer, count, datatype, dest-rank, tag, comm, request)
MPI_Waitall (count, request[], status[])
```

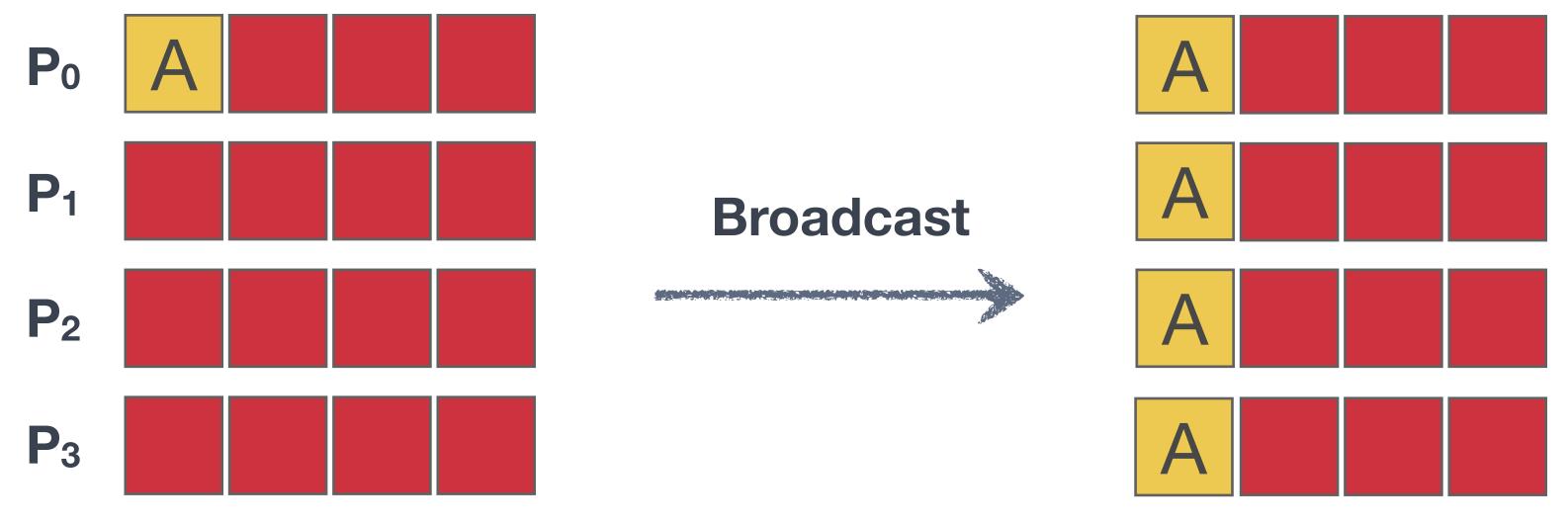
Summary: MPI

- Many parallel programs can be written using six commonly used MPI primitives
 - ► MPI_Init
 - ► MPI_Finalize
 - ► MPI_Comm_rank
 - ► MPI_Comm_size
 - ► MPI_Send
 - ►MPI_Recv
- Non-blocking primitives for correctness and performance

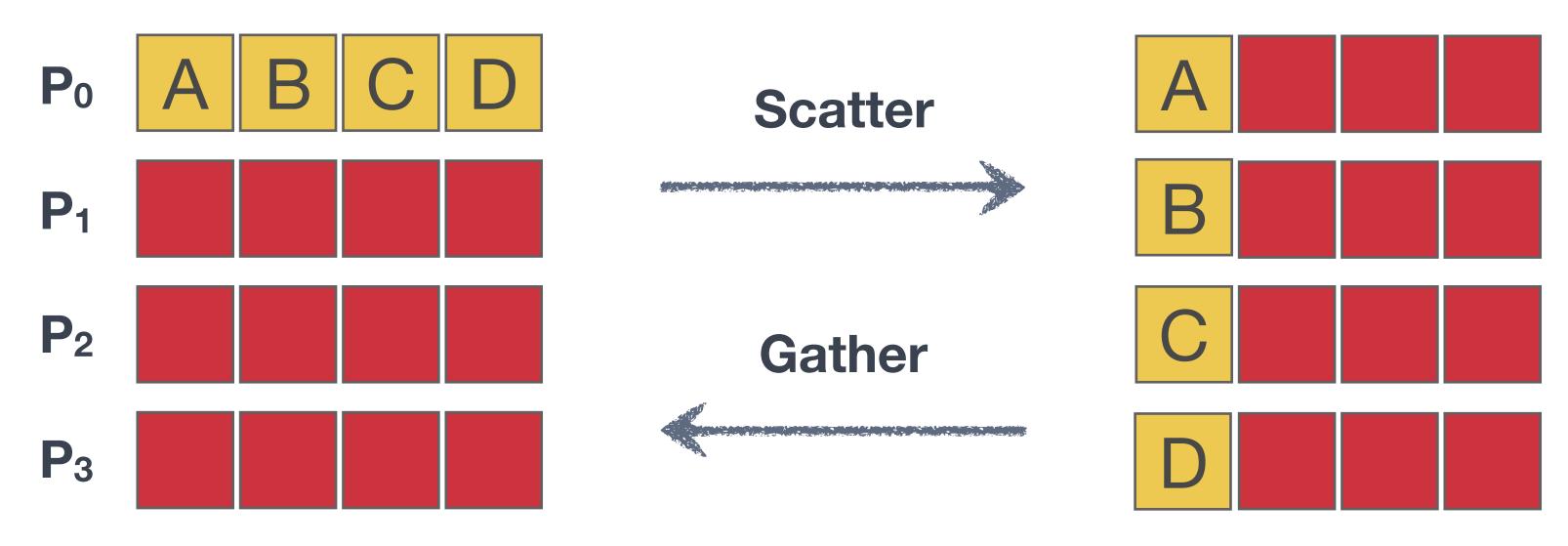
Alternate approach

▶ Collective communication

- ► Higher-level communication primitives
 - ► MPI_Bcast: Broadcast data to all processes
 - ► MPI_Reduce: Combine data from all processes to one process
 - ► MPI_Barrier
- Each process executes the same communication operation
- ►MPI provides a rich set of collective operations
- Presumably optimized/tuned for hardware



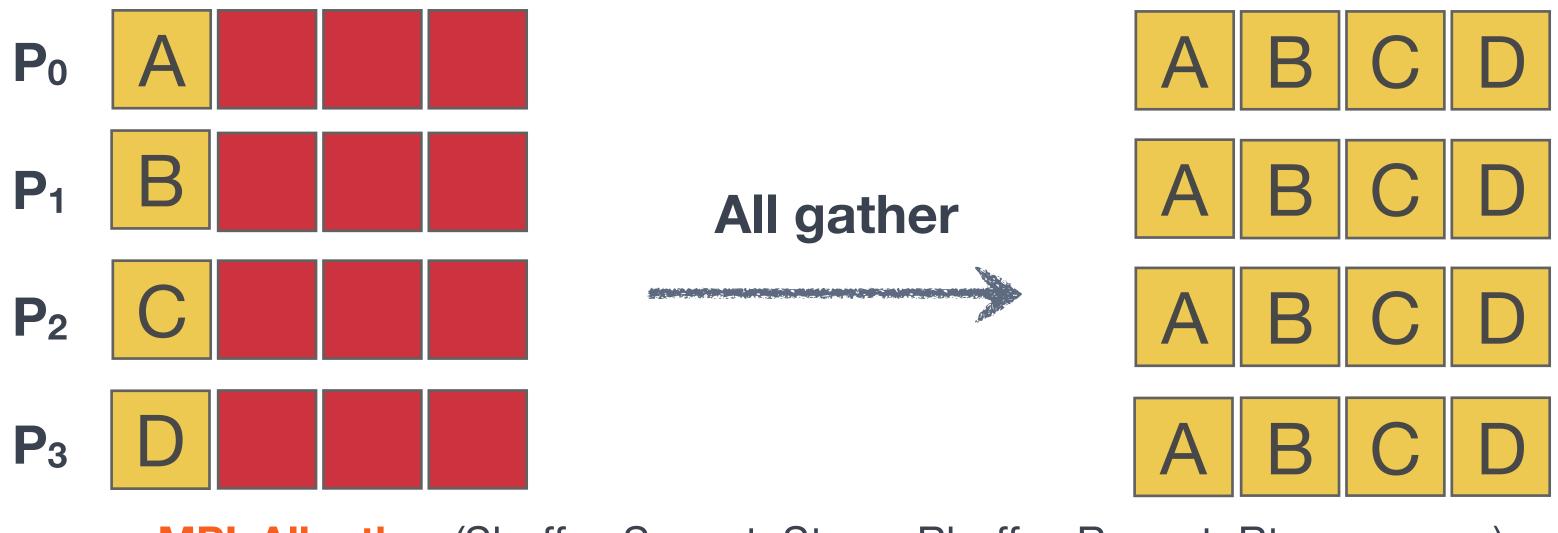
MPI_Bcast (buffer, count, datatype, root, comm)



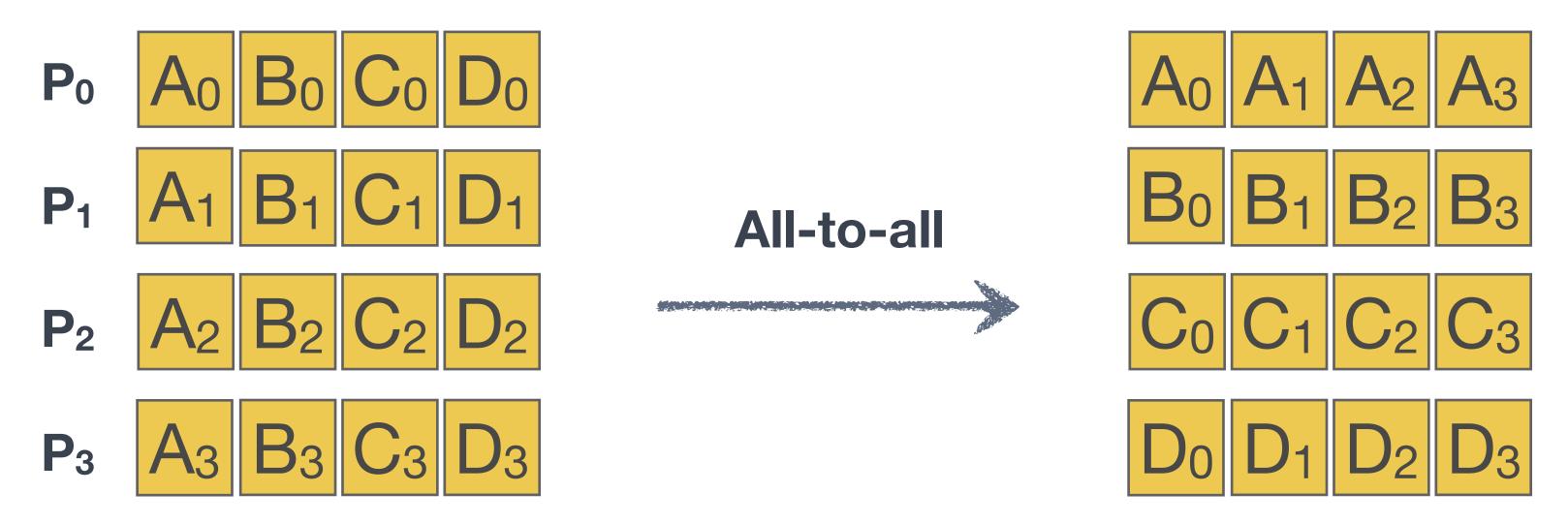
MPI_Scatter (Sbuffer, Scount, Stype, Rbuffer, Rcount, Rtype, root, comm)
MPI_Gather (Sbuffer, Scount, Stype, Rbuffer, Rcount, Rtype, root, comm)

Case study: Broadcast

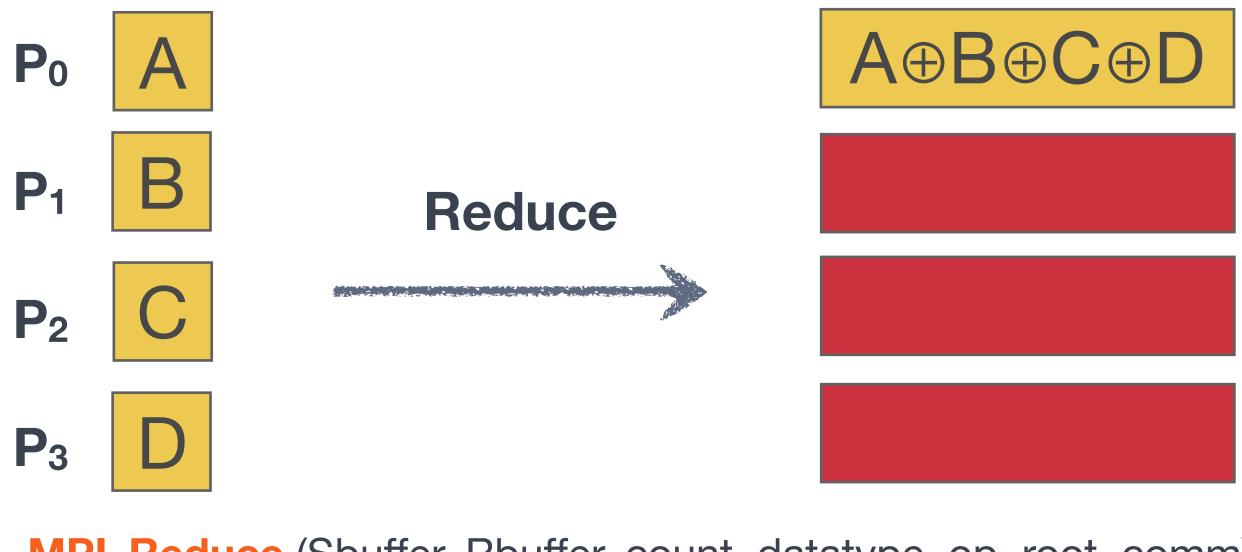
- ► All collective operations must be called by all processes in the communicator
- ▶Eg,. MPI_Bcast is called by both the root process and all the processes that receive the broadcast
- Demo example on HPC cluster!



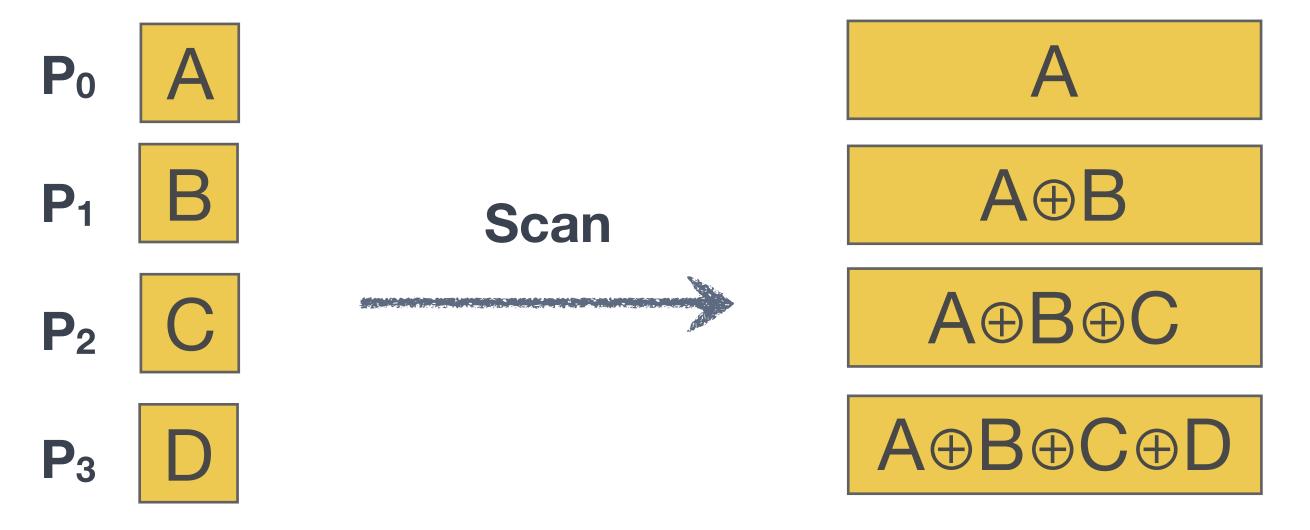
MPI_Allgather (Sbuffer, Scount, Stype, Rbuffer, Rcount, Rtype, comm)



MPI_Alltoall (Sbuffer, Scount, Stype, Rbuffer, Rcount, Rtype, comm)



MPI_Reduce (Sbuffer, Rbuffer, count, datatype, op, root, comm)



MPI_Scan (Sbuffer, Rbuffer, count, datatype, op, comm)

MPI Built-in Operations

MPI Reduction Operation	
MPI_MAX	maximum
MPI_MIN	minimum
MPI_SUM	sum
MPI_PROD	product
MPI_LAND	logical AND
MPI_BAND	bit-wise AND
MPI_LOR	logical OR
MPI_BOR	bit-wise OR
MPI_LXOR	logical XOR
MPI_BXOR	bit-wise XOR
MPI_MAXLOC	max value and location
MPI_MINLOC	min value and location

Other collective routines

- Many routines: Gather, Gatherv, Allgather, Allgatherv, Reduce_scatter, Reduce, Allreduce, Scatter, Scatterv, Bcast, Scan
- ► All = data to all processes
- V = allows data chunks to have variable sizes
- ►MPI-2 added Alltoallw, Exscan, inter-communicator variants for most routines