

Today's content:

- (i) Workers allocation
- (ii) Aggressive cows.

10. Given N tasks, k workers and time taken for each task, find min time in which we can complete all tasks.

- Note:
- i) A single worker can only do continuous set of tasks
 - ii) All workers start their assigned task at same time
 - iii) A task can only be assigned to 1 worker.

ex:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	Time taken.
$N=15,$	3	5	1	7	8	2	5	3	10	1	4	7	5	4	6	
$k=3.$	✱ = 9			✱ = 17			✱ = 45									45.
$k=3.$		t = 24m			t = 21m			t = 26m								26.
$k=3$		t = 24m			t = 25m			t = 22m								25

ans = 25

ex:

	0	1	2	3	4	5	Time taken.
$N=5,$							
$k=2.$	w1					w2	10

Ideal: Brute force.

Q: In how many ways can I choose 2 sticks out of 14 sticks?

$${}^{14}C_2.$$

Q: In how many ways can I choose $(k-1)$ sticks out of $(n-1)$ sticks?

$${}^{n-1}C_{k-1} * N. \leftarrow TC.$$

↳ prefix sum $\Rightarrow N + {}^{n-1}C_{k-1} * k \leftarrow TC.$

Idea 2: Binary search.

a) Target \longrightarrow min time.

b) search space \rightarrow $\left[\begin{array}{cc} \underline{\text{low}} & \underline{\text{high}} \\ \text{max of} & \text{sum of all} \\ \text{all times} & \text{times} \end{array} \right]$

Ex1: ar : [10 2 9 8].

$k=1$, Time taken = 29.

Ex 2: ar : $\begin{bmatrix} \underline{w_1} & \underline{w_2} & \underline{w_3} & \underline{w_4} \\ 10 & 15 & 11 & 13 \end{bmatrix}$.

$k=4$: Time taken.

Discard.



Case 1: If we can finish the task in mid time.

ans = mid

goto left.

Case 2: If we cannot finish task in mid time.

goto right.

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 | N=15, k=4.

3 5 1 7 8 2 || 5 3 10 1 4 7 5 4 6

[W₁=26min] [W₂=30min.] [W₃=15min] W₄=[] idle.

[W₁=9min] [W₂] [W₃] [W₄] . Tasks left.

Can I finish all the tasks in 10min?

7	8	9	10
F	F	F	F

Can I finish all the tasks in 30 min?

30, 31, 32, 33 - - - -

1 1 1 1 . .

Code:

```
def mintime(ar, k, n)
    |
    |   |   |
    |   |   |---> # workers
    |   |   |---> #Tasks.
    |   |   |---> Time for each task.
```

Tc: $N \times \log_2(\# \text{ elements})$

```
{
    l = max(ar) , h = sum(ar) , ans = sum(ar).
    while (l <= h)
    {
        m = (l+h)/2.
        // can we finish all the tasks in m time.
        if (check(n, k, ar, m))
        {
            ans = m
            h = m-1 // goto left.
        }
        else // goto right
            l = m+1
    }
    return ans
}
```

elements = $h-l+1$

= $\text{sum(ar)} - \text{max(ar)}$.

```
def check(n, k, ar, t)
```

```
{
```

```
    c = 0, s = 0
```

```
    for i in range(0, n)
```

```
        s = s + ar[i] // adding task time.
```

```
        if (s > t)
```

```
            c = c + 1, s = ar[i]
```

```
            if (c == k) // utilized k workers, but tasks
                        |   are left.
                        |   return false
```

```
    return true
```

```
}
```

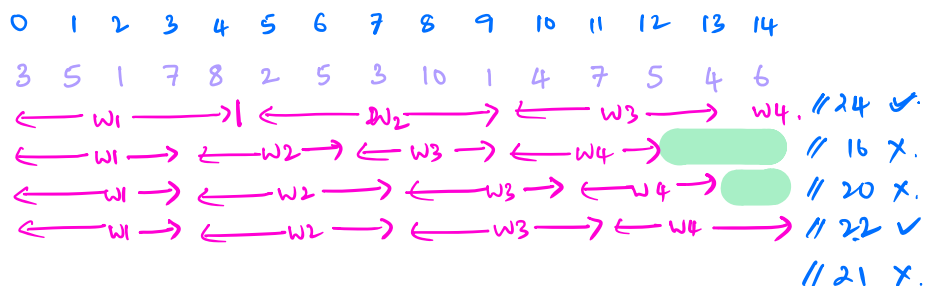
Idea for check fun.

$N=15$
 $K=4$
 $T=20$
 $C=0$



Even after utilizing 4 people, tasks are left out, return false

$N=15$
 $K=4$



l h m can I finish all tasks in m time ans update.

10	71	40	✓	ans = 40, goto left, h = m-1.
10	29	24	✓	ans = 24, goto left, h = m-1
10	23	16	X	goto right, l = m+1.
17	23	20	X	goto right, l = m+1
21	23	22	✓	ans = 22, goto left, h = m-1
21	21	21	X	goto right, l = m+1
22	21	(l > h) X	break return 22.	

Break.

10:28:00

10:38:00

10:40:00

20. Given k cows and N stalls, all stalls are on x -axis at different locations, Place all cows in such a way that minimum distance between the 2 cows is maximized. (Keep the aggressive cows as far as possible)

Note: (i) Only one cow can be present at one stall.
(ii) All cows need to be placed, ($N > k$). All stalls positions are sorted.

Ex1:

Stalls = 5

cows = 3

0 1 2 3 4
1 2 4 8 9
 $C_1 \overset{1}{-} C_2 \overset{2}{-} C_3$
 $C_1 \overset{2}{-} C_2 \overset{5}{-} C_3$
 $C_1 \overset{3}{-} C_2 \overset{5}{-} C_3$

min Distance.

1
2
3

ans = 3.

Ex2:

Stalls = 9

cows = 4

0 1 2 3 4 5 6 7 8
2 6 11 14 19 25 30 39 43.
 $C_1 \overset{4}{-} C_2 \overset{5}{-} C_3 \overset{3}{-} C_4 \dots$
 $C_1 \overset{9}{-} C_2 \overset{8}{-} C_3 \overset{11}{-} C_4$
 $C_1 \overset{12}{-} C_2 \overset{16}{-} C_3 \overset{13}{-} C_4$

Min distance

3

8

12

Idea1: Given N stalls, ways to choose k stalls for k cows.

$\underbrace{N C_k}_{\downarrow \text{No. of ways}} * k \rightarrow$ To calculate min distance for a selection.

Idea 2: Binary Search.

a) Target : maximum min distance between 2 cows.

b) Search space : $\left[\begin{array}{cc} \text{low} & \text{high} \\ \text{min of} & \text{stalls}[N-1] \\ \text{adjacent elements} & - \text{stalls}[0]. \\ \text{difference.} & \end{array} \right]$

Ex 1:

stalls[4] : [3 8 14 20].

k = 2 cows : c1 ————— c2
17

stalls[4] : [3 8 14 20].

k = 4 cows : c1 c2 c3 c4.
5 6 6.

BS discard.

Case 1: say we can place all cows at a distance of "mid".

ans = mid
move right.



Case 2: say we cannot place all cows at a distance of "mid".

move left:



stalls = 9

cows = 4

0 1 2 3 4 5 6 7 8
2 6 11 14 19 25 30 39 43.
c1 ————— 23 ————— c2
9 8 6
c1 ————— c2 ————— c3 ————— c4

c3, c4, x. → 20.

✓ for 5

Can I place cows atleast at 20 distance apart?

Can I place cows atleast at 5 distance apart?

20 21 22 23
F F F F

--- 2 3 4 5
T T T T

Code:

no of stalls.
no of cows.

```
def maxDistance(n, k, stalls)
```

```
l = min adjacent diff. in stalls[], h = stalls[m-1] - stalls[0], ans =
```

```
while (l ≤ h)
```

```
    m = (l+h)/2
```

```
    // If you can place all cows at 'm' distance apart.
```

```
    if (check(n, k, stalls, m))
```

```
        ans = mid
```

```
        l = m+1
```

```
    else
```

```
        h = m-1
```

```
return ans
```

(h-l+1).

↑

Tc: $N * O(\log_2(\text{#elements}))$

m=20.

```
def check(n, k, stalls, m)
```

```
{
```

```
    c=1, j=0
```

```
    for i in range(1, n)
```

```
        if (stalls[i] - stalls[j] ≥ m)
```

```
            c = c+1
```

```
            j = i
```

```
return c ≥ k
```

```
}
```

0	1	2	3	4	5	6	7	8
2	6	11	14	19	25	30	39	43.

c1

c2

Doubts:

0	1	2	3	4	5	6	7	8
2	6	11	14	19	25	30	39	43.
c1			c2			c3		c4

l	h	m	Can I place all cows	update ans.
3	41	22	X	goto left, $h = m - 1$.
3	21	12	✓	ans = 12, goto right, $l = m + 1$
13	21			