

Report

1-)

decreaseKey(index, key)

heap[index]-key <- key	O(1)
while heap[parent]-key > heap[child] – key do	O(lgn)
swap(heap[child], heap[parent])	O(5)
index <- (index -1)/2 //assigning parents index to the “index”	O(1)

While loop will operate as many as depth of the heap so asymptotic upper bound is $O(\lg n)$

Total bound is $O(\lg n) + O(7)$

Since $O(\lg n)$ is dominant term, running time of this function is $O(\lg n)$

extract(index)

if heap-size < 1	O(1)
then error “heap underflow”	O(1)
if index != 0	O(1)
decreaseKey(index, $-\infty$)	$O(\lg n)$
temporary <- heap[0]	O(1)
swap (heap[0], heap[heap-size - 1])	O(5)
heap.pop-back(0)	O(5)
return temporary	O(1)

Since there is no loop, sum is $O(\lg n) + O(15)$

Since $O(\lg n)$ is dominant term, running time of this function is $O(\lg n)$

insertVehicle(input)

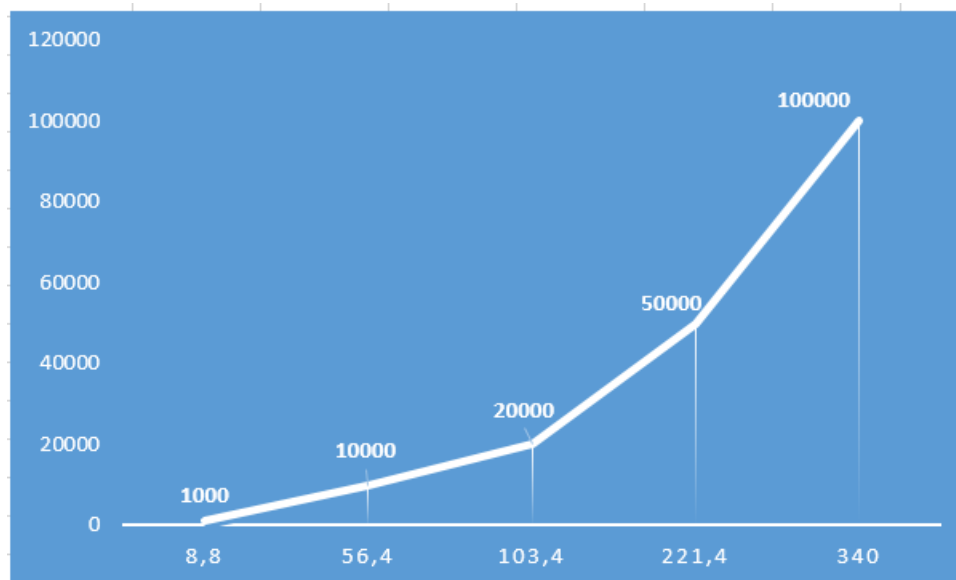
temporary_key <- input – key O(1)
input-key <- +∞ O(1)
heap-push_back(input) O(5)
decreaseKey(heap-size - 1, temporary_key) O(lgn)

Since there is no loop, sum is $O(lgn) + O(7)$

Since $O(lgn)$ is dominant term, running time of this function is $O(lgn)$

2-) Table is in milliseconds

	1	2	3	4	5	Avg.
1000	10	8	6	10	10	8.8
10K	44	58	64	56	60	56.4
20K	67	110	119	106	115	103.4
50K	254	263	271	163	156	221.4
100K	487	301	311	301	300	340



decreaseKey function has $O(lgn)$ asymptotic upper bound as well as extract and insert functions.

In total $O(lgn) + O(lgn) + O(lgn) = O(3lgn)$. While lgn is dominant term, it will be $O(lgn)$ anyway. When we observe the plot table above, we can see $n lgn$ function. The reason for that

both decrease, insert and extract functions are repeated n times due to n number of data in the heap. So we must multiply it for $\lg n$ to get the final runtime for overall algorithm.

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