

Ramaneek Gill 1000005754

## CSC263 Problem Set 5

An important thing to observe in this question is that DELETE will never reduce the size of the array. This means that DELETE will always run at  $O(1)$  time.

Amortized complexity =  $WCSC / m$

- Worst case run time for DELETE is  $O(1)$

- unknown for APPEND.

Let's compute APPEND's worst case run time by using the accounting method:

DELETE's charge and cost are 1. It needs no credits.

APPEND's actual cost is 1 but we will charge  $2+10k$ .  $k$  is the amount of times a new array is dynamically created

Credit Invariant/Breakdown:

- 1 is used right away to append

- 1 credit is used to copy itself over into the new array in the future

- 1 credit (from the  $10k$  credits part) is used to help copy over previous elements which don't have credits left to the new array

- this one might seem confusing, let  $n$  refer to the size of the array

Using induction:

Invariant:

- each element after being appended has  $1+10*k$  credits where  $k$  is the amount of times array has been dynamically increased

- if no more spaces left in array create a new array with 10 more spaces and copy all previous elements over

base case:

- $k = 0$

- //array hasn't been dynamically increased yet, so the first 10 elements have 1 extra credit each for the next move over

Second base case: //first dynamic increase occurs

- $k = k+1$

- elements use 1 credit to copy over

- all elements now have 0 credits //this is right before the 11th element is appended

Induction Step:

- suppose invariant holds and we need to dynamically increase array size by 10

- the next 10 elements have 1 credit stored for themselves for the future move

AND

- $10*k$  credits stored for all the previous elements to move over (since they have none)

- this process repeats every time the array size needs to be increased

- the last 10 elements in the array always give a credit to previous elements to make the move

So total cost  $\leq$  total charge  $= 2 + 10 \cdot k$ .

So for  $n$  insertions the amortised cost is  $O(n)$ . For  $n$  deletions the amortised cost is  $O(1)$ . So for any sequence of insertions and deletions the amortised cost will be  $O(n)$ .