Advanced Programming

Finite Automata and Regular Expressions

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DPG 1 / 30

1 Regular Expression



Finite Automata

ADPG 1 / 30

```
Enclose letters, special characters, spaces,
# digits using single quote or double quotation
# marks
first name = "Arun"
last name = "Dravid"
# concatenate strings
greet = "hi " + ' ' + first_name + ' ' + last name
print(greet)
#repetition of strings
repeat_3_times = greet*3
```

```
my string = "hello"
#length of a string
print(len(my string))
#locate the character(s) using indexes
print(my_string[1]) #output = 'e'
#slicing - extracting parts of a sequence
print(my string[2:4]) # output = 'll' ? What are the start
# and end indexes used?
print(my string[:4]) # output = 'hell'
print(my_string[2:]) # output = ?
#what does the following code do?
print(my string([-1])
```

```
message = "HELLO WORLD"
lower_case = message.lower()
# Output: "hello world"
upper_case = lower_case.upper()
#output = 'HELLO WORLD'
# split() function is used to split two words uing a separator
words = "apple, banana"
fruit list = words.split(",")
# Output: ["apple", "banana"]
# find whether a string exists in a sequence
# finds the first occurrence.
# If found, returns the position,
# else return -1
text = "Pvthon is awesome"
position = text.find("awesome")
# Output: 12
```

```
name = "Roy"
greeting = f"Hello, {name}!"
print(greeting) # Output: Hello, Roy!
```

```
input_string = 'Hello, World'
upper_string = input_string.upper()
# Output: HELLO, WORLD!
lower_string = input_string.lower()
# Output: hello, world!
```

```
input_string = 'Hello, World'
index = input_string.find("world")
# Output: 7 (index of the first occurrence)
```

```
# Checks if all characters are letters
is_alpha = input_string.isalpha()
#is_apha = True, if input_string contains
#alphabets [a-z]+, else False
# Checks if all characters are numbers
is_numeric = input_string.isnumeric()
```

```
text = 'Did you find Python in that corner?'
index = text.index("Python") # Returns 13
#starts with a character or string
if text.startswith("Did"):
    print("String starts with 'Did'")
if 'that' in text:
    print("Substring 'that' found in the string")
#output: index of the last 'o' = 29
does text ends with = text.endswith('that corner?')
for in text:
    print(i)
    #prints all the characters in the text, including
    #the white space
```

In a text environment, it is often required to search a string from a huge collection of documents. If we want a specific string, we use a simple search. There are situations not a specific string but a pattern requires to be found in the entire document. In such cases, *Regular Expressions* are useful.

- A language for specifying text search strings
- An algebraic notation for characterizing a set of strings a language for specifying text search strings
- ► A tool for pattern matching within text
- Precise description of a set of strings using a combination of alphabets/symbols
- Case sensitive search String 'Hello' is different from 'hello'

- Search for dates in the corpus not a specific date but a set of dates with varying values
 - Find all strings that match "MM-DD-YYYY" 12-12-2024, 01-01-2022, 06-15-2021...
- ► Get all the currency values found in a document -\$12389.23, \$56789.12,...
- Find strings Hello and hello [Hh]ello
- ► Find any digit in the text It costs anywhere between 5 to 8 dollars

BASIC RULES OF REGEX

- ► [0123456789] specifies any one of digits [] specifies **disjunction**
- ► [ABCDEF] specifies _____?
- (hyphen) specifies the range
- ▶ [0-9] indicates any digit from 0 to 9
- ► [A-Z] means _____?
- \triangleright [a-z] represents _____
- ► [`A-Z] indicates negation/NOT do **NOT** match an upper case letter

- ? can be used to match zero or one instance of the preceding character -colou?r will match both color and colour
- * zero or more instances of the preceding character or combinations a* a, aa, aaa, aaaa, This is also known as Kleene *
 - [ab]* represents a set of strings {a,aaa,b,bbb,ab,ababab,abbbbb,...}

Some more rules

- \triangleright [a-z][a-z] represents _____?
- ▶ + positive closure or Kleene + represents ONE or MORE
- \triangleright [a-z][0-9]+
- . is a wildcard r.n matches run, ran, ron, ...
- ► [']{3}.* [']{3} matches any text between "' and "'
- ▶ Alphabets Represent individual symbols like letters, numbers, and punctuation
- ▶ Wild cards Match any single character (".") or any character set ("[]")
- Quantifiers Specification for repetition of patterns zero or more ("*"), one or more ("+"), or exactly n times ("n")
- ► Anchors Match positions within a string like beginning ("^") or end ("\$") -

Regex	Match	Examples
\wedge	The beginning of the sentence	∧Once - <u>Once</u> upon a time
\$	The end of a line	\w+\$ - How are you?
\s	Any white space character	s[a-z] I have <u>a</u> Python book
		$\slash s[a-z] + \slash s$ I <u>have</u> <u>a</u> Python book
		$\slash s[a-z]*\slash $ I <u>have</u> <u>a</u> Python book
\S	Matches any visible character	
\b	word boundary	$ackslash b[a\!-\!z]$ - $ar{ t L}$ have $ar{ t a}$ Python $ar{ t b}$ ook
∖B	non word boundary	$ackslash \mathrm{B}[\mathfrak{a}\!-\!z]ackslash \mathrm{B}$ - I h <u>av</u> e a P <u>ytho</u> n b <u>oo</u> k

Day (DD) Rules

- 1. Must be two digits
- 2. Must follow date semantics
- 3. It restricts the day part to valid values (01-30,31)
 - oxdot One of the many possible patterns (0[1-9]|[12][0-9]|3[01])

0[1-9] - accepts 0 as the first digit and any digits from 0...9

[12][0-9] - accepts either 1 or 2 as the first digitand any digit from 0...9

3[01] - accepts 3 as the first and 0 or 1 as the second digit

Identifying numbers (scientific notation of digits, float, etc) [+-]?[0-9]*[.]?[0-9]+([eE][+-]?[0-9]+)?

```
#include the module that contains all regex functions
import re
def check day(input pattern:str) -> bool:
  Check whether the given pattern matches the dates beteeen 01-31
   Anything not in this set should be considered as invalid
   date_pattern = re.compile(r"^(0[1-9]|[12][0-9]|3[01])$")
   # If date pattern matches input pattern, return True;
  # else return False
  if ( date_pattern.match(input_pattern)):
      return True
  else:
       return False
if name == ' main ':
   print(check day(input("Input date: ")))
  # Test cases
  for date in range (40):
       print(f'{date} - {check day(str(date))}')
```

```
def check email address(input pattern:str) -> str:
    , , ,
    Check whether the given pattern matches
    the standard email address
    This may verify 75-85% of the email addresses
    , , ,
    email pattern = re.compile(r'' = a-zA-zO-9. %+-]+@[a-zA
   -Z0-9.-]+\.[a-zA-Z]\{2,\}")
    if ( email pattern.match(input pattern)):
        return "a valid email address"
    else:
        return "an invalid email address"
```

- ► **Text Search and Manipulation**: All important applications where Find and replace specific patterns, extract data, validate formats are important
- Compilers/Interpretors, Programs Define syntax, validate user input, perform data parsing
- ► Web Technologies Build search engines, extract information from web pages, filter content

FINITE AUTOMATA - ALPHABETS AND STRINGS

Definition

The term alphabet denotes any set of symbols

Example

```
English alphabets and numbers The set containing binary alphabet \{0,1\}
```

Definition

The term *Language* denotes any set of strings formed using the alphabet

Example

```
{000, 111, 0101, 1110000,...}
{a, b, aa, aaabbb, ababab,...}
```

OPERATIONS ON LANGUAGES

We can apply concatenation, union, and closure (Kleene¹ and positive²) operations on *Languages*. Let L be a set of alphabets $\{A,B,...,Z, a,b,...,z\}$ and D be set containing a set of 10 digts $\{0,1,...,9\}$

Name	Operation	Example
Union	$L \cup D = \{s \mid s \in L \text{ or } \in D\}$	s,5,2
Concatenation	$LM = \{sd \mid s \text{ is in and } d \in D\}$	12s,e4
Kleene Closure	$L^* = \bigcup_{i=0}^{\infty} L^i$	ε, 0000,1111,00110011
Positive Closure	$L^+ = \bigcup_{i=1}^{\infty} 4L^i$	0,1,010101
ϵ operations	$\epsilon s = s, \ s \epsilon = s, \ s^* = (s \mid \epsilon)^*$	

Finite Automata ADPG 20 /

¹L* denotes zero or more concatenation of L

²L⁺ denotes one or more concatenation of L

Let us consider a language whose identifiers constructed using the following alphabet

Alphabet
$$\{A, B, ..., Z, a, b, ..., z, 0, 1, ..., 9\}$$

The rules for constructing the identifier is given below

```
\begin{array}{ll} \textbf{identifier} & letter(letter \mid digit)^*. \end{array}
```

digit 0 | 1 | 2 | ... 9

Explanation

- | is used in place of "or"
- The expression in (letter|digit)* represents zero or more of either the letter or digit

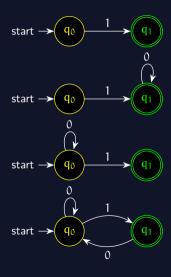
It is recognizer that determines whether a given string of alphabets adheres to the rules and patterns of a specific language. It functions as a gatekeeper, accepting only strings that conform to the language's grammar and structure, while rejecting those that do not. This is a mathematical model that consists of 5-tuple

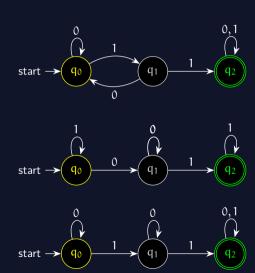
- 1. Q, a set of states
- 2. Λ , a set of input symbols/alphabets
- 3. q₀, a special start/initial state
- 4. F, a set of accepting states
- 5. $\delta: Q \times \Sigma$, a transition function that maps states-symbol pair to sets of states



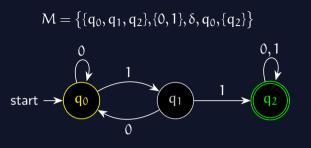
Accepting node



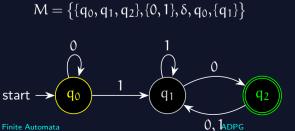




TRANSITION DIAGRAM



	0	1
qo	qo	q ₁
q ₁	qo	q ₂
q ₂	q ₂	q ₂
Transition ⁻		Table





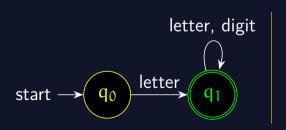
Transition Table

Grammar

identifier::letter(letter|digit)*

letter :: [a-z]

digit :: [0-9]



	letter	digit
qo	q ₁	
q_1	q ₁	q ₁

```
identifier :: (_letter)(_|letter|digit)*
    letter :: [a-z]
    digit :: [0-9]
```

```
function get next state(current state, symbol)
    // read the transition table to get the next state corresponding
    // to the current state and the symbol
current state := initial state
for symbol in input_string:
    next state = get new state((current state, symbol))
    if next state is None then
        return 'Invalid string for the automaton'
    current_state = next_state
    if current_state == accepting_state then
        return 'Valid string'
   else
       return 'Invalid string'
end
```

```
def is string accepted(transition table:dict, initial state:str,
   accepting states, string) -> bool:
   current_state = initial_state
   for symbol in string:
        next_state = transition_table.get((current_state, symbol))
        if next state is None: # No transition for the given input
           return False
   current_state = next_state
   return current_state in accepting\_states
if name == ' main ':
   transition \ table = {
        ('q2', '0'): 'q2', ('q2', '1'): 'q0'
   initial state = 'q0'
   # set of accepting states
   accepting states = {'q2'}
   string = input('Input a binary string: ')
   if is string accepted (transition table, initial state, accepting states,
   string):
        print('The string is accepted.')
        print('The string is not accepted')
```

Definition (Deterministic Finite Automata - DFA)

Every step of a computation follows in a unique way from the preceding step

Definition (Nondeterministic Finite Automata - NFA)

- Several choices may exist for the next state at any point. Nondeterminism is a generalization of determinism
- Every nondeterministic finite automaton has an equivalent deterministic finite automaton
- A language is regular if and only if some regular expression describes it
- ► The Languages accepted by the finite automata are precisely the language denoted by the regular expressions.
- \blacktriangleright For every regular expression, there is precisely a NFA with ϵ -transition

```
def find identifier(line:str) ->list:
    identifier:list = []
    # Use regex rules to find the identifier(s)
    return identifier
if __name__ == '__main__':
    with open("tmp.py", "r") as f:
    # Read the entire text file
    # text = f.read()
    # Read the entire text file and
    # store all the lines into list
    # lines:list = f.readlines()
    #process the lines
    for line in f.readlines():
        find identifier(line)
```