

Advanced Programming

Finite Automata and Regular Expressions

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① Regular Expression

| ② Finite Automata

OPERATIONS ON STRINGS

```
# Enclose letters, special characters, spaces,  
# digits using single quote or double quotation  
# marks  
  
first_name = "Arun"  
last_name = "Dravid"  
  
# concatenate strings  
greet = "hi " + ' ' + first_name + ' ' + last_name  
print(greet)  
  
#repetition of strings  
repeat_3_times = greet*3
```

STRING AS ARRAY

```
my_string = "hello"

#length of a string
print(len(my_string))

#locate the character(s) using indexes
print(my_string[1]) #output = 'e'

#slicing - extracting parts of a sequence
print(my_string[2:4]) # output = 'll' ? What are the start
# and end indexes used?
print(my_string[:4]) # output = 'hell'
print(my_string[2:]) # output = ?

#what does the following code do?
print(my_string([-1]))
```

METHODS/FUNCTIONS TO OPERATE ON STRINGS

```
message = "HELLO WORLD"
lower_case = message.lower()
# Output: "hello world"
upper_case = lower_case.upper()
#output = 'HELLO WORLD'

# split() function is used to split two words using a separator
words = "apple,banana"
fruit_list = words.split(",")
# Output: ["apple", "banana"]

# find whether a string exists in a sequence
# finds the first occurrence.
# If found, returns the position,
# else return -1

text = "Python is awesome"
position = text.find("awesome")
# Output: 12
```

STRINGS ...

```
name = "Roy"  
greeting = f"Hello, {name}!"  
print(greeting)  # Output: Hello, Roy!
```

STRINGS ...

```
input_string = 'Hello, World'
upper_string = input_string.upper()
# Output: HELLO, WORLD!
lower_string = input_string.lower()
# Output: hello, world!
```

```
input_string = 'Hello, World'  
index = input_string.find("world")  
# Output: 7 (index of the first occurrence)
```


STRINGS ...

```
# Checks if all characters are letters
is_alpha = input_string.isalpha()
#is_apha = True, if input_string contains
#alphabets [a-z]+, else False
# Checks if all characters are numbers
is_numeric = input_string.isnumeric()
```

STRINGS ...

```
text = 'Did you find Python in that corner?'
index = text.index("Python") # Returns 13
#starts with a character or string
if text.startswith("Did"):
    print("String starts with 'Did'")
if 'that' in text:
    print("Substring 'that' found in the string")
#output: index of the last 'o' = 29
does_text_ends_with = text.endswith('that corner?')
for i in text:
    print(i)
    #prints all the characters in the text, including
    #the white space
```

REGULAR EXPRESSIONS

In a text environment, it is often required to search a string from a huge collection of documents. If we want a specific string, we use a simple search. There are situations not a specific string but a pattern requires to be found in the entire document. In such cases, *Regular Expressions* are useful.

- ▶ A language for specifying text search strings
- ▶ An algebraic notation for characterizing a set of strings a language for specifying text search strings
- ▶ A tool for pattern matching within text
- ▶ Precise description of a set of strings using a combination of alphabets/symbols
- ▶ Case sensitive search - String 'Hello' is different from 'hello'

EXAMPLES

- ▶ Search for dates in the corpus - not a specific date but a set of dates with varying values

Find all strings that match "MM-DD-YYYY" 12-12-2024, 01-01-2022, 06-15-2021...

- ▶ Get all the currency values found in a document -\$12389.23, \$56789.12,...
- ▶ Find strings Hello and hello - [Hh]ello
- ▶ Find any digit in the text - It costs anywhere between 5 to 8 dollars

BASIC RULES OF REGEX

- ▶ `[0123456789]` specifies any one of digits - `[]` specifies **disjunction**
- ▶ `[ABCDEFGH]` specifies _____?
- ▶ `-` (hyphen) specifies the range
- ▶ `[0-9]` indicates any digit from 0 to 9
- ▶ `[A-Z]` means _____?
- ▶ `[a-z]` represents _____?
- ▶ `^[A-Z]` - indicates negation/NOT - do **NOT** match an upper case letter
- ▶ `?` can be used to match zero or one instance of the preceding character
`-colou?r` will match both `color` and `colour`
- ▶ `*` - zero or more instances of the preceding character or combinations -
`a*` - `a`, `aa`, `aaa`, `aaaa`, This is also known as **Kleene ***
- ▶ `[ab]*` represents a set of strings
`{a,aaa,b,bbb,ab,ababab,abbbbbb,...}`

BUILDING BLOCKS OF REGEX

Some more rules

- ▶ `[a-z][a-z]` - represents _____?
- ▶ `+` - positive closure or Kleene `+` - represents ONE or MORE
- ▶ `[a-z][0-9]+`
- ▶ `.` is a wildcard - `r.n` matches `run`, `ran`, `ron`, ...
- ▶ `[']{3}.*[']{3}` - matches any text between `" "` and `" "`
- ▶ **Alphabets** - Represent individual symbols like letters, numbers, and punctuation
- ▶ **Wild cards** - Match any single character (`"."`) or any character set (`"[]"`)
- ▶ **Quantifiers** - Specification for repetition of patterns - zero or more (`"*"`), one or more (`"+"`), or exactly n times (`"n"`)
- ▶ **Anchors** - Match positions within a string like beginning (`"^"`) or end (`"$"`) -

ANCHORS

Regex	Match	Examples
^	The beginning of the sentence	^Once - <u>Once</u> upon a time
\$	The end of a line	\w+\$ - How are <u>you</u> ?
\s	Any white space character	\s[a-z]\s I have <u>a</u> Python book \s[a-z]+\s I <u>have</u> <u>a</u> Python book \s[a-z]*\s I <u>have</u> <u>a</u> Python book
\S	Matches any visible character	
\b	word boundary	\b[a-z] - <u>I</u> <u>have</u> <u>a</u> Python <u>book</u>
\B	non word boundary	\B[a-z]\B - I <u>have</u> a <u>Python</u> <u>book</u>

Day (DD) Rules

1. Must be two digits
2. Must follow date semantics
3. It restricts the day part to valid values (01-30,31)
 - ☒ One of the many possible patterns $\wedge(0[1-9]|1[0-9]|2[0-9]|3[01])$

0[1-9] - accepts 0 as the first digit and any digits from 0...9

[12][0-9] - accepts either 1 or 2 as the first digit and any digit from 0...9

3[01] - accepts 3 as the first and 0 or 1 as the second digit

Identifying numbers (scientific notation of digits, float, etc)

$[+-]?[0-9]*[.]?[0-9]+([eE][+-]?[0-9]+)?$

DAY REGEX - PYTHON CODE

```
#include the module that contains all regex functions
import re
def check_day(input_pattern:str) -> bool:
    """
    Check whether the given pattern matches the dates between 01-31
    Anything not in this set should be considered as invalid
    """
    date_pattern = re.compile(r"^(0[1-9]|[12][0-9]|3[01])$")
    # If date_pattern matches input_pattern, return True;
    # else return False
    if ( date_pattern.match(input_pattern)):
        return True
    else:
        return False

if __name__ == '__main__':
    print(check_day(input("Input date: ")))
    # Test cases
    for date in range(40):
        print(f'{date} - {check_day(str(date))}')
```

EMAIL PATTERN

```
def check_email_address(input_pattern:str) -> str:
    """
    Check whether the given pattern matches
    the standard email address
    This may verify 75-85% of the email addresses
    """
    email_pattern = re.compile(r"^[a-zA-Z0-9._%+-]+@[a-zA-Z0-9.-]+\.[a-zA-Z]{2,}$")
    if ( email_pattern.match(input_pattern)):
        return "a valid email address"
    else:
        return "an invalid email address"
```

- ▶ **Text Search and Manipulation:** - All important applications where Find and replace specific patterns, extract data, validate formats are important
- ▶ **Compilers/Interpretors, Programs** - Define syntax, validate user input, perform data parsing
- ▶ **Web Technologies** - Build search engines, extract information from web pages, filter content

FINITE AUTOMATA - ALPHABETS AND STRINGS

Definition

The term *alphabet* denotes any set of symbols

Example

English alphabets and numbers

The set containing binary alphabet $\{0, 1\}$

Definition

The term *Language* denotes any set of strings formed using the alphabet

Example

$\{000, 111, 0101, 1110000, \dots\}$

$\{a, b, aa, aaabbb, ababab, \dots\}$

OPERATIONS ON LANGUAGES

We can apply concatenation, union, and closure (Kleene¹ and positive²) operations on *Languages*. Let L be a set of alphabets $\{A,B,\dots,Z, a,b,\dots,z\}$ and D be set containing a set of 10 digits $\{0,1,\dots,9\}$

Name	Operation	Example
Union	$L \cup D = \{s \mid s \in L \text{ or } \in D\}$	s,5,2
Concatenation	$LM = \{sd \mid s \text{ is in } L \text{ and } d \in D\}$	12s,e4
Kleene Closure	$L^* = \bigcup_{i=0}^{\infty} L^i$	ϵ , 0000,1111,00110011
Positive Closure	$L^+ = \bigcup_{i=1}^{\infty} L^i$	0,1,010101
ϵ operations	$\epsilon s = s, s\epsilon = s, s^* = (s \mid \epsilon)^*$	

¹ L^* denotes zero or more concatenation of L

² L^+ denotes one or more concatenation of L

LANGUAGE ANS STRINGS

Let us consider a language whose *identifiers* constructed using the following alphabet

Alphabet $\{A, B, \dots, Z, a, b, \dots, z, 0, 1, \dots, 9\}$

The rules for constructing the identifier is given below

identifier $\text{letter}(\text{letter} \mid \text{digit})^*$.

letter $a \mid b \mid c \mid \dots z$

digit $0 \mid 1 \mid 2 \mid \dots 9$

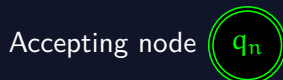
Explanation

- \mid is used in place of "or"
- The expression in $(\text{letter} \mid \text{digit})^*$ represents zero or more of either the letter or digit

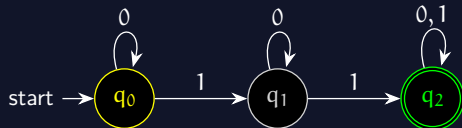
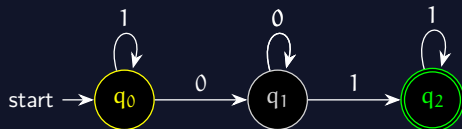
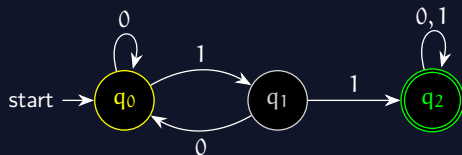
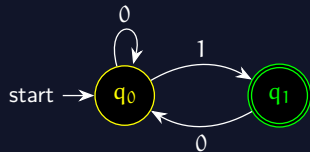
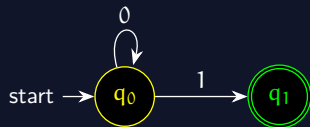
FINITE AUTOMATA

It is recognizer that determines whether a given string of alphabets adheres to the rules and patterns of a specific language. It functions as a gatekeeper, accepting only strings that conform to the language's grammar and structure, while rejecting those that do not. This is a mathematical model that consists of 5-tuple

1. Q , a set of states
2. Λ , a set of input symbols/alphabets
3. q_0 , a special start/initial state
4. F , a set of accepting states
5. $\delta : Q \times \Sigma$, a transition function that maps states-symbol pair to sets of states

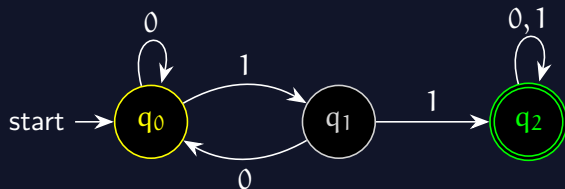


EXAMPLES OF FA



TRANSITION DIAGRAM

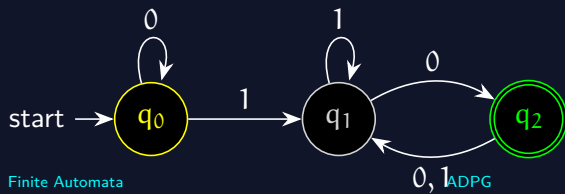
$$M = \{\{q_0, q_1, q_2\}, \{0, 1\}, \delta, q_0, \{q_2\}\}$$



	0	1
q_0	q_0	q_1
q_1	q_0	q_2
q_2	q_2	q_2

Transition Table

$$M = \{\{q_0, q_1, q_2\}, \{0, 1\}, \delta, q_0, \{q_1\}\}$$



	0	1
q_0	q_0	q_1
q_1	q_0	q_2
q_2	q_1	q_1

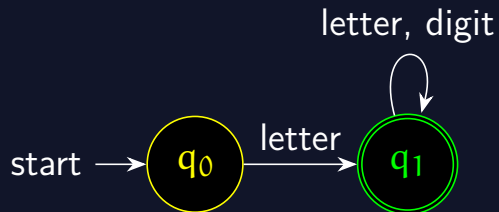
Transition Table

Grammar

identifier :: letter(letter|digit)*

letter :: [a-z]

digit :: [0-9]



	letter	digit
q_0	q_1	
q_1	q_1	q_1

identifier :: (`_letter`)(`_letter`|`digit`)*

letter :: `[a-z]`

digit :: `[0-9]`

FA SIMULATION

```
function get_next_state(current_state, symbol)
{
    // read the transition table to get the next state corresponding
    // to the current state and the symbol
}

current_state := initial_state

for symbol in input_string:
    next_state = get_new_state((current_state, symbol))

    if next_state is None then
        return 'Invalid string for the automaton'

    current_state = next_state

    if current_state == accepting_state then
        return 'Valid string'
    else
        return 'Invalid string'
end
```

FA SIMULATION

```
def is_string_accepted(transition_table:dict, initial_state:str,
    accepting_states, string) -> bool:
    current_state = initial_state

    for symbol in string:
        next_state = transition_table.get((current_state, symbol))
        if next_state is None: # No transition for the given input
            return False

    current_state = next_state

    return current_state in accepting_states

if __name__ == '__main__':
    transition_table = {
        ('q0', '0'): 'q0', ('q0', '1'): 'q1',
        ('q1', '0'): 'q2', ('q1', '1'): 'q1',
        ('q2', '0'): 'q2', ('q2', '1'): 'q0'
    }

    initial_state = 'q0'
    # set of accepting states
    accepting_states = {'q2'}

    string = input('Input a binary string: ')

    if is_string_accepted(transition_table, initial_state, accepting_states,
        string):
        print('The string is accepted.')
    else:
        print('The string is not accepted')
```

EQUIVALENCE OF REGULAR EXPRESSION AND FINITE AUTOMATA

Definition (Deterministic Finite Automata - DFA)

Every step of a computation follows in a unique way from the preceding step

Definition (Nondeterministic Finite Automata - NFA)

- ▶ Several choices may exist for the next state at any point. Nondeterminism is a generalization of determinism
- ▶ Every nondeterministic finite automaton has an equivalent deterministic finite automaton
- ▶ A language is regular if and only if some regular expression describes it
- ▶ The Languages accepted by the finite automata are precisely the language denoted by the regular expressions.
- ▶ For every regular expression, there is precisely a NFA with ϵ -transition

FIND IDENTIFIERS IN A PYTHON PROGRAM

```
def find_identifier(line:str) ->list:
    identifier:list = []
    # Use regex rules to find the identifier(s)
    return identifier

if __name__ == '__main__':
    with open("tmp.py", "r") as f:
        # Read the entire text file
        # text = f.read()
        # Read the entire text file and
        # store all the lines into list
        # lines:list = f.readlines()

        #process the lines
        for line in f.readlines():
            find_identifier(line)
```