



中山大學
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Lecture 16

Recap

Algorithm Design

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Basic Knowledge

- Characteristics of algorithms
- Scientific notation for numbers
- Numerical errors
- Bitwise operations
- Asymptotic complexities
- Order of growth
- Best, average, worst case analysis
- Mathematical induction
- Recursion
- ...

Data Structures

- Stacks
- Queues
- Lists: Arrays and Linked lists
- Heap: Priority Queue
- Trees: Binary search trees, Tree traverse, Tree representation
- Hashing
- Disjoint-sets

Greedy

- Definition: A *greedy algorithm* is an algorithm in which at each stage a locally optimal choice is made.
- Characteristics:
 1. **Greedy-choice property:** A global optimum can be arrived at by selecting a local optimum.
 2. **Optimal substructure:** An optimal solution to the problem contains an optimal solution to subproblems.
- Greedy algorithms are usually extremely efficient, but they can only be applied to a small number of problems.

Greedy

- Examples:

- Activity selection problem
- Fractional Knapsack problem
- Spanning trees
- Huffman coding
- Change-making problem (找零问题)
- **Biggest number:** 设有 n 个正整数，将它们连接成一排，组成一个最大的多位整数。例如： $n=3$ 时，3个整数13，312，343，连成的最大整数为34331213。又如： $n=4$ 时，4个整数7，13，4，246，连成的最大整数为7424613。输入 N 个数，输出：连成的多位数

Divide-and-conquer

- A **divide-and-conquer** algorithm works by recursively breaking down a problem into two or more sub-problems of the same (or related) type, until these become simple enough to be solved directly. The solutions to the sub-problems are then combined to give a solution to the original problem.
- Example:
 - Mergesort / Quicksort
 - Strassen's Matrix Multiplication
 - Hanoi Tower

Dynamic Programming

- **Dynamic Programming (DP)** is a method for solving complex problems by breaking them down into simpler subproblems. It is applicable to problems exhibiting the properties of **overlapping subproblems** (子问题重叠) and **optimal substructure** (最优子结构).
- A problem is said to have **overlapping subproblems** if the problem can be broken down into subproblems which are reused several times or a recursive algorithm for the problem solves the same subproblem over and over rather than always generating new subproblems.
- A problem is said to have **optimal substructure** if an optimal solution can be constructed efficiently from optimal solutions of its subproblems.

Dynamic Programming

- You need to know the following concepts:
 - State (状态): A way to describe a situation, a sub-solution for the problem.
 - Top-down approach
 - Bottom-up approach
 - Memoization (记忆化)
 - Recurrence formulas / State transition equation (状态转移方程)
 - Traceback (for obtaining concrete solutions)

Dynamic Programming

- Examples:

- Street Walking: 一个城市的街道布局如下，从最左下方走到最右上方，每次只能往上或往右走，一共有多少种走法？
- Weighted Street Walking
- Longest Common Subsequence: $O(n^2)$
- Longest Non-Decreasing Subsequence: $O(n^2) \rightarrow O(n \log n)$
- 0-1 Knapsack Problem: $O(nW)$
- Matrix-Chain Multiplication: $O(n^3)$
- Traveling Salesman Problem $O(n!)$ $\rightarrow O(n^2 \cdot 2^n)$

Search

- Backtracking: recursively explore each node.
- Branch-and-bound: It is very similar to backtracking. It computes a number (**bound**) at a node to determine whether the node is promising. If that bound is no better than the value of the best solution found so far, the node is **non-promising**. Otherwise, it is **promising**.
- Branch-and-bound methods are adopted using the Depth-First Search (DFS) manner, the Breadth-First Search (BFS) manner, and the Best-First search manner.
- Use bounds to do pruning.
- A bound can be obtained by relaxing some constraints of the subproblem related to the node.

Search in Graphs

Uninformed Search

- Uninformed means that we only know the goal and the ***succs()*** function.
- Search strategies: BFS, UCS, DFS, IDS, Bi-directional search.

b: branching factor (assume finite) *d*: goal depth *m*: graph depth

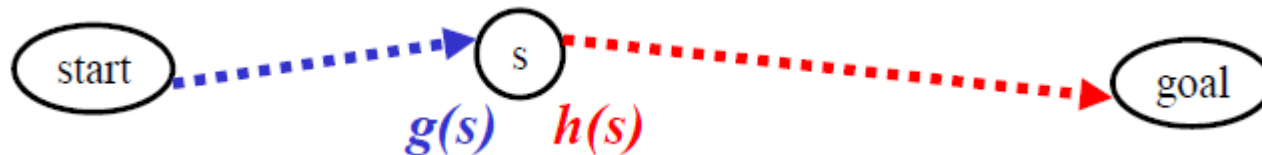
	Complete	optimal	time	space
Breadth-first search	Y	Y, if ¹	$O(b^d)$	$O(b^d)$
Uniform-cost search	Y	Y	$O(b^{C^*/\epsilon})$	$O(b^{C^*/\epsilon})$
Depth-first search	N	N	$O(b^m)$	$O(bm)$
Iterative deepening	Y	Y, if ¹	$O(b^d)$	$O(bd)$
Bidirectional search	Y	Y, if ¹	$O(b^{d/2})$	$O(b^{d/2})$

1. edge cost constant, or positive non-decreasing in depth

Search in Graphs

- Informed Search

- Informed means that we also know a heuristic $h(s)$ of the cost from node s to goal.



- Search strategies: A*, IDA*.
- For the heuristic function $h()$, if $h(s) \leq h^*(s)$, where $h^*(s)$ is the true cost from node s to the goal, then such heuristic function $h()$ is called **admissible**.

Search

- Examples:
 - Subset Tree Enumeration
 - Permutation Tree Enumeration
 - Sudoku
 - Traveling Salesman Problem
 - 8-Puzzle Problem
 - Arithmetic Calculation Problem
 - Graph Coloring Problem (图染色)
 - Maximum Clique Problem (最大团)
 - Many Path Finding Problems

Graph Algorithms

- Graph representations: adjacency matrix, adjacency list
- Graph traverse: DFS, BFS
- Topological sorting for DAGs: It is a linear ordering of nodes in which each node comes before all nodes to which it has outbound edges.
- Single source shortest path algorithms: Dijkstra algorithm ($O(n^2)$ or $O(E+V\log V)$), Bellman-ford algorithm
- All-pairs shortest path algorithms: Floyd-Warshall algorithm
- System of difference constraints
- Minimum spanning tree algorithms: Prim's algorithm, Kruskal's algorithm

Optimization Algorithms

- Approximation algorithms: providing provable solution quality and provable run-time bounds.
- Heuristic algorithms: the art of discovering new strategies (rules) to solve problems.
- Meta-heuristic algorithms: upper level general methodologies that can be used as guiding strategies in designing underlying heuristics to solve specific optimization problems
- Components in SA, TS, GA

Network flows

- Maximum flow
 - Ford-Fulkerson Algorithm
- Minimum cut
- Maximum matching in bipartite graph

Thank you!

