

## CS 525 – QUIZ II

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### Question 1.1.1 Estimate Result Size (5 Points)

Estimate the number of result tuples for the query  $q = \sigma_{\text{instructor}=\text{Bob}}(\text{course})$  using the first assumption presented in class (values used in queries are uniformly distributed within the active domain).

$$T(q) = T(\text{course})/V(\text{course, instructor})$$

$$= 1000/500$$

$$= 2 \text{ tuples}$$

### Question 1.1.2 Estimate Result Size (5 Points)

Estimate the number of result tuples for the query  $q = \sigma_{\text{major}=\text{CS}} \vee \text{major}=\text{EE}}(\text{students})$  using the first assumption presented in class.

$$T(q) = T(\text{student})/V(\text{student, major})_{\text{CS}} + T(\text{student})/V(\text{student, major})_{\text{EE}}$$

$$= 40,000/50 + 40,000/50$$

$$= 800 + 800$$

$$= 1600 \text{ Tuples}$$

### Question 1.1.3 Estimate Result Size (7 Points)

Estimate the number of result tuples for the query  $q = \sigma_{\text{credits} \geq 29 \wedge \text{credits} \leq 40}(\text{course})$  using the first assumption presented in class.

$$T(q) = [(40 - 29 + 1) * T(\text{course})] / [\max(\text{course, credits}) - \min(\text{course, credits}) + 1]$$

$$= [12 * 1000] / [60 - 0 + 1]$$

$$= 197 \text{ Tuples}$$

#### Question 1.1.4 Estimate Result Size (8 Points)

Estimate the number of result tuples for the query  $q$  below using the first assumption presented in class.

$q = \sigma_{\text{credits} \leq 30}(\text{student}) \bowtie_{\text{CWID}=\text{student}} \text{registered} \bowtie_{\text{course}=\text{title}} \text{course}$

$$q_1 = \sigma_{\text{credits} \leq 30}(\text{student})$$

$$T(q_1) = [(30 - \min(\text{student}, \text{credits}) + 1) * T(\text{student})] / [\max(\text{student}, \text{credits}) - \min(\text{student}, \text{credits}) + 1]$$

$$= [(30 - 0 + 1) * 40,000] / [60 - 0 + 1]$$

$$= (31 * 40,000) / 61$$

$$= 20,328 \text{ Tuples}$$

$$q_2 = q_1 \bowtie_{\text{CWID}=\text{student}} \text{registered}$$

$$T(q_2) = T(q_1) * T(\text{registered}) / \max[V(q_1, \text{CWID}), V(\text{registered}, \text{student})]$$

$$= 20,328 * 20,000 / \max(20328, 5000)$$

$$= 20,000 \text{ Tuples}$$

$$q = q_2 \bowtie_{\text{course}=\text{title}} \text{course}$$

$$T(q) = T(q_2) * T(\text{course}) / \max[V(q_2, \text{course}), V(\text{course}, \text{title})]$$

$$= 20,000 * 1000 / \max(100, 1000)$$

$$= 20,000 \text{ Tuples}$$

### Question 1.2.1 External Sorting (5 Points)

You have a block of size of 128KB and 2GB memory. Compute the maximum size of file that can be sorted by 3-Pass Multiway Sort.

$$\text{Memory in KB} = 2 * 1024 * 1024 = 2,097,152 \text{ KB}$$

$$M = 2,097,152 / 128$$

$$= 16,384 \text{ buffers}$$

$$B(R) \leq M * (M-1)^2 \text{ blocks}$$

$$\leq 16,384 * (16,384 - 1)^2 * 128 \text{ KB}$$

$$\leq 32,766 \text{ GB}$$

### Question 1.2.2 I/O Cost Estimation (25 Points)

Consider two relations R and S with  $B(R) = 400,000$  and  $B(S) = 2,000$  blocks, and  $S(R) = 1/40$ , and  $S(S) = 1/10$ . You have  $M = 101$  memory pages available. Let R has an index on the joining attribute C. Compute the minimum number of I/O operations needed to join these two relations using tuple-based-nested-loop join (relations are not clustered), block-nested-loop join (relations are clustered), merge-join (the inputs are not sorted but clustered), and non-clustering index-join (read of clustered relation S with uniform distribution assumption and expected 20 matching tuples on R for S) and hash-join (relations are clustered but not sorted). You can assume that the hash function evenly distributes keys across buckets. Justify your result by showing the I/O cost estimation for each join method.

$$T(R) = 400,000 * 40 = 16,000,000$$

$$T(S) = 2000 * 10 = 20,000$$

- Tuple based nested loop join:

S is smaller, thus let's consider it as outer relation

$$\begin{aligned} \text{Total Cost} &= T(S) * (1 + T(R)) \\ &= 20,000 * (1 + 16,000,000) \\ &= 320,000,020,000 \text{ I/Os} \end{aligned}$$

- Block nested loop join:

S is smaller. Thus, let's consider 100 buffers for S and 1 buffer for R

Read chunk of S = 100 I/Os

Read of R = 400,000 I/Os

$$\# \text{ of S chunks} = B(S) / M - 1$$

$$= 2000/100$$

$$= 20$$

$$\text{Total cost} = 20 * (100 + 400,000)$$

$$= 8,002,000 \text{ I/Os}$$

- Merge Join:

Sort cost of R (3 Pass Multiway join algorithm) =  $6 * B(R)$

Sort cost of S (2 Pass Multiway join algorithm) =  $4 * B(S)$

Total cost = sort cost + join cost

$$= 7 * B(R) + 5 * B(S)$$

$$= 7 * 400,000 + 5 * 2000$$

$$= 2,800,000 + 10,000$$

$$= 2,810,000 \text{ I/Os}$$

- Non clustering Index join:

$$T(S) + [T(S) * T(R) / V(R,C)] = 20,000 + [20,000 * 20]$$

$$= 20,000 + 400,000$$

$$= 420,000 \text{ I/Os}$$

- Hash-Join (relations are clustered but not sorted):

$$3 (B(R) + B(S)) = 3 (400,000 + 2000)$$

$$= 1,206,000 \text{ I/Os}$$