1. Hoose Triples

Task 1.1

let s = while izndo x:= x + 1; i:= i+1 od

when we apply this value to S,

- -> Horse the pre-condition satisfied since i < n
- -> The loop(s)will goes on with out tour minating
- > The post condition will not be satisfied which is

Herea the above total correctness is unsatisfied.

b. fiz-1, n=53 to fizonnéog when we apply this to s,

- -> The post don pre-condition is satisfied since icn
- > Statement terminates with 11=0, 712-5}
- And the post condition of 120 An 50% is satisfied

Hence the above Bs peoitial correctness triple is Satisfied

-> Hose the pre-condition is feiled

I program is torminated since it did not enter the loop

And the post Condition is also failed, since if M

Hence the poortial Correctness triple is satisfied since
in the given state since pre-condition is failed and

it doesn't matter if the program is executed.

when we apply this value to 5, $\mathcal{E}_{0} = \{i=1, n=2, k=2\}$ $\mathcal{E}_{0} = \{i=1, n=2, k=2\}$ $\mathcal{E}_{0} = \{i=2, n=2, k=2\}$

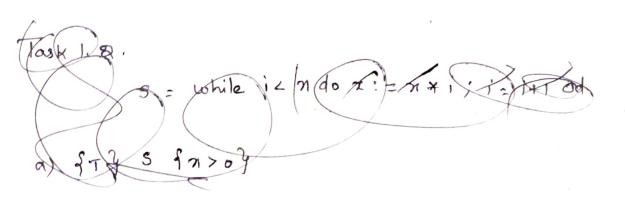
- Here the precondition is satisfied.

> Program is terminated since i is not less than or at

Si=2, m=2 y

And the post-Condition is also satisfied since sm=k?

Hence the paritial Correctness triple is satisfied



when we apply this value to s,

- -> Hore the pre-condition is satisfied
- -> The program doesn't tominate. so post wondition
- -> Hence the Partial Correctness triple is satisfied

Task 1.2

a) {TY S {n>0}

> Not Valid

counter example !

lets take fx=2, i=0}

$$z_0 = \{i = 0, x_2 = 2\}, \quad z_1 = \{x = 0, i = 1\}$$

Here precondition is satisfied, program gets terminated but the post condition fails since at 100 11=0.

Fix

To make this Valid late fix the post condition

Post condition \Rightarrow fix yso the triple will be fty y y

4

b, fx= k 3 s fx= k 3 -> Not Valid lets take frezz, i=-17 and after loop x=-2 and Consumpler example: Post condition Fails, Fix -> fixing the program s = while izy do x != x; i=i+1 od example giz1, x=2, K=2) after loop gi=2, x=2, x=2, x=2, ... > Pro-condition satisfied, Program and and Post condition satisfied so, fr=ky while izndox:=x; i=i+1 od fr=1g [i=1 An=KAn>o] S [i=KAx=k!] -> Not Valid Counter exampla fiz1, K=5,7=5} after loop fizz, K=5, x=5} -> Precondition satisfied, program executes but post condition failed, Fix the post condition, to [x= K] and priple will be fizi 1x2k 1x>0] S [x=k] example fi=1, K=5, 7=57 after loop fi=2, K=5, 7=53 Task 1.3 Precondition > m= 2K A JK EZ.K EO AV=1 So triple would be,

example:

lets take fm= -2, n= 2, 7=1}

- Love the procuondition is satisfied with given value.

+ Program terminates with &m=-2, n=0, r=93

-> post condition is satisfied since 3" = r(3 = 9)

Task 1.4

-> with this pre-condition the statement always executes without error and or with be upstated.

-> Also the post Condition is always time.

Task 1.5

a) #[P2] 82 [T]

Program:

- -> Here the pre condition is Satisfied
- I side the Statement will bevon terminate since x will be always growter than O
- -> Post condition is also always true.

-> Hence the total correctness triple is invalid.

Ь.

Since I have used while in previous ans, I am worting

- > Here the precondition is 21>0 and 4=0
- -) The Statement will be executed and and in error,
- -> Hence the total Correctness triple is invalid.

Task 1.6

Or.

S = while n > q do if even(n) then n := 5x+1 Else x := x/2 find.

triple valid > [x>8] S [n 41]
example.

lets take 7:4

 $\Sigma_0 = \frac{1}{3} \times \frac{1}{3} + \frac{1}{3}$, $\Sigma_1 = \frac{1}{3} \times \frac{1}{3} = \frac$

- -> Here the pre-condition is satisfied.
- > Program exminates when n=0.095

-> And the post condition is salisfied.

-> Hence the triple valid for total Correctors triple.

6.

S=while n>1 do if even(n) then x:=5x+1 else n:=n/2 di od

Malid triple fx < 33 9 f F3

Here the pre-condition will execute but never tour not

-> But the post- Londition is

always false. Correctness

the Statement will heren execute

2. Substitution

Task 2.1

a) $\begin{bmatrix} y+2/y \end{bmatrix} \exists_z \cdot \forall_n \cdot (n+y \geq x+y) \cdot d_n = \exists_z \cdot \begin{bmatrix} y+2/y \end{bmatrix} \cdot \forall_n = (n+y \geq x+y)$ $= \exists_z \cdot \forall_n \cdot \begin{bmatrix} y+2/y \end{bmatrix} \cdot (n+y \geq x+y)$ $= \exists_z \cdot \forall_n \cdot (n+(y+2)) \geq (x+(y+2))$ $= \exists_z \cdot \forall_n \cdot (n+(y+2)) \geq (x+(y+2))$ $= \exists_z \cdot \forall_n \cdot (n+(y+2)) \geq (x+(y+2))$

b)
$$\begin{bmatrix} y+2/m \end{bmatrix} \exists_z \cdot \forall_n \cdot (x+y \ge z+y)$$

 $= \exists_z \cdot \begin{bmatrix} y+2/m \end{bmatrix} \cdot \forall_n \quad (n+y \ge z+y)$
 $= \exists_z \cdot \forall_n \quad (x+y \ge z+y)$ (Since to free von1964)
 $= \exists_z \cdot \forall_n \quad (x+y \ge z+y)$ $\forall_n \in \{y \in Y\}$ $\{y \in Y$

c)
$$\left[x + 2/y \right] \exists_z \cdot \forall_n (n+y \ge 2+y)$$

$$= \exists_z . \left[\frac{\gamma + 2}{y} \right] . \forall_n \left(\frac{\gamma + y}{y} \ge \frac{\gamma + y}{y} \right)$$

In this to avoid the voniable capture, we need to do & conversion on bound variable x, and renaming

$$=\exists_{z}. \left[\eta + 2/y \right] \forall_{w} \left(w + y \geq z + y \right)$$

ofter applying of convension on bound variables

N. 2 with W. W respectively.

$$= (Z \geq Z \rightarrow (\exists \forall w, w \geq V) \land y > Z)$$

e. $[z/n](n \ge z' \to (\exists_{x,x+y} \ge z+y) \land y > z)$ after applying of conversion on bound variable xwith w.

$$= \left[\frac{z}{\pi} \right] \left(\frac{\pi}{2} \geq z \rightarrow \left(\frac{3}{\omega}, \frac{\omega+y}{2} \geq z+y \right) \wedge y > z \right)$$

$$= \left(\frac{z}{2} \geq z \rightarrow \left(\frac{3}{\omega}, \frac{\omega}{2} \geq z \right) \wedge y > z \right)$$

3. proofs and proof outlines.

Task 8.1

a. given size (a) = 2

 $S = if \alpha[0] > \alpha[1] then m:= \alpha[0]; else m:= \alpha[1] fi$

b.

IT Y S formax (a[o],a[i])}

to expand,

STY if a [o] > a[i] then m:= a[o]; else m=a[i] fi S(m = a[o] ∧ a[o] > a[i]) V(m=a[i] ∧ a[i]> a[o])

- C. Hilbert-Style for the above S=ifa(oJ)a[i] then m:=a[o]
 else m:=a[i] fi
 - 1. 9a[o] > a[i] 3 m: = a[o] 9 m= a[o] ∧ a[o] > a[i] 3
 - a. gt na[0] > a[i] g m := a[o] { m = a[o] na[o]>a[i] q
 - consequences of a cold a cold
- 4. 9 T 1 a [0] < a [1] 9 m = a [1] 9 m = a [1] 1 1 A a [1] > a [0] 3 Good (3)
- 5. FT 3 if a[07 > a[1] then m:=a[0] else m:=a[1] fi = a[0] \ a[0] \ a[1] \ V (m=a[1] \ \ a[1] \ if (1,2)

$$f T \wedge a [o] > a[i]$$

$$f m = a[o] \wedge m > a[i]$$

$$f T \wedge (a[i] > a[o])$$

$$f T \wedge (a[i] > a[o])$$

$$f a[i] = a[i] \wedge a[i]$$

$$f m = a[i] \wedge a[o] < a[i]$$

$$f (m = a[o] \wedge a[o] > a[i]) \vee a[o] \wedge a$$

Task 3.2

1. $\int_{1}^{1} z \ge 0 \wedge n y, 3 = 0 \int_{1}^{1} S = n \int_{1}^{1} S \ge 0 \wedge S y, 3 = 0 \int_{1}^{1} A \times n y, 3 = 0 \int_{1}^{1} S = n \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(1)$ 2. $\int_{1}^{1} n \ge 0 \wedge n y, 3 = 0 \int_{1}^{1} S = n \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(1)$ 3. $\int_{1}^{1} n y = 0 \wedge n y, 3 = 0 \int_{1}^{1} S = n - 1 \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(3)$ 4. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \wedge n y, 3 = 1 \int_{1}^{1} S = n - 1 \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(3)$ 5. $\int_{1}^{1} n - 2 \ge 0 \wedge n - 2 y, 3 = 0 \int_{1}^{1} S = n - 2 \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(5)$ 6. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 0 \int_{1}^{1} Conseq(5)$ 7. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 1 \int_{1}^{1} Conseq(5)$ 5. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 1 \int_{1}^{1} Conseq(5)$ 6. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 1 \int_{1}^{1} Conseq(5)$ 7. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 1 \int_{1}^{1} Conseq(5)$ 5. $\int_{1}^{1} n y \ge 0 \wedge n y, 3 \neq 0 \int_{1}^{1} S y, 3 = 1 \int_{1}^{1} Conseq(5)$

8. 77 ≥ 0 f if x11.3 = 0 then s:= 7 else if x1.3=1 thens:=x-1

else S:=x-2 fi fi \S'1,3=03 if (218)

4. It took me approximately 14 hours to complete the assignment.