

## 1. Linked List and Separation Logic

1.1

$$n_0 := !(\text{head} + \bar{1}) ;$$

$$n_1 := ! (n_0 + \bar{1}) ;$$

$$\text{head}_2 := ! (n_1 + \bar{1}) ;$$

$$!(n + \bar{1}) := \text{nil} ;$$

$$\text{head}_1 := \text{head} ;$$

1.2

$$\{ \text{list}(a_1, a_2, a_3, \text{head}, m+3) \}$$

$$\{ \text{head} \mapsto a_1, i_1 \times i_1 \mapsto a_2, i_2 \times i_2 \mapsto a_3, i_3 \times \text{list}(L_2, i_3, m) \}$$

$$n_0 := !(\text{head} + \bar{1}) ;$$

$$\{ (n_0 = i, \wedge \text{head} \mapsto a_1, n_0) \times i, \mapsto a_2, i_2 \times i_2 \mapsto a_3, i_3 \times$$

$$\text{list}(L_2, i_3, m) \} \Rightarrow \{ \text{head} \mapsto a_1, n_0 \times n_0 \mapsto a_2, i_2 \times i_2 \mapsto a_3, i_3 \times \text{list}(L_2, i_3, m) \}$$

$$n_1 := ! (n_0 + 1) ;$$

$$\{ \text{head} \mapsto a_1, n_0 * (n_1 = i_2 \wedge n_0 \mapsto a_2, n_1) * i_2 \mapsto a_3, i_3 * \text{list}(L_2, i_3, m) \} \Rightarrow$$

$$\{ \text{head} \mapsto a_1, n_0 * n_0 \mapsto a_2, n_1 * n_1 \mapsto a_3, i_3 * \text{list}(L_2, i_2, m) \}$$

$$\text{head2} := ! (n_1 + 1) ;$$

$$\{ \text{head} \mapsto a_1, n_0 * n_0 \mapsto a_2, n_1 * (n_1 \mapsto a_3, \text{head2} \wedge \text{head2} = i_3) * \text{list}(L_2, i_3, m) \} \Rightarrow \{ \text{head} \mapsto a_1, n_0 * n_0 \mapsto a_2, n_1 * n_1 \mapsto a_3, \text{head2} * \text{list}(L_2, \text{head2}, m) \}$$

$$!(n_1 + 1) = \text{nil} ;$$

$$\{ \text{head} \mapsto a_1, n_0 * n_0 \mapsto a_2, n_1 * n_1 \mapsto a_3, \text{nil} * \text{list}(L_2, \text{head2}, m) \}$$

$$\text{head1} := \text{head}$$

$$\{ \text{head1} \mapsto a_1, n_0 * n_0 \mapsto a_2, n_1 * n_1 \mapsto a_3, \text{nil} * \text{list}(L_2, \text{head2}, m) \} \\ \Rightarrow \{ \text{list}(a_1; a_2; a_3, \text{head1}, 3) * \text{list}(L_2, \text{head2}, m) \}$$

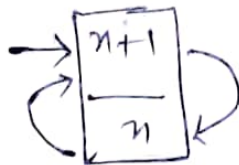
## 2. Resource logic

2.1

a)  $x \vdash x$



c)  $x \vdash x+1, x$

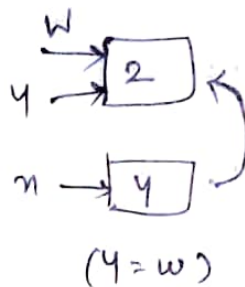


b)  $x \vdash 2 \text{ * } y \vdash 2$

$x \rightarrow [2]$

$y \rightarrow [2]$

d)  $w \vdash 2 \text{ * } (x \vdash y \wedge y = w)$



e)  $(\exists x. y \vdash x) \wedge y \vdash 2$

$y \rightarrow [2]$

2.2

a.  $x \leftrightarrow x$

(i)  $x \rightarrow [x] \Leftarrow$

$y \rightarrow [2]$

(ii)  $\rightarrow [x] \Leftarrow$

b.  $x \leftrightarrow 2 \wedge y \leftrightarrow 2$

(i)  $x \rightarrow [2]$   
 $y \rightarrow [2]$   
 $(x = y)$

(ii)  $x \rightarrow [2]$   
 $y \rightarrow [2]$

c.  $x \mapsto \_$

(i)  $x \mapsto \boxed{1}$

(ii)  ~~$x \mapsto \boxed{1}$~~   
 $x \mapsto \boxed{2}$

2.3

a.  $(x \mapsto 2 * y \mapsto 2) * z \mapsto 2 \Rightarrow (z \mapsto 2 * y \mapsto 2) * x \mapsto 2$

This is provable

$(x \mapsto 2 * y \mapsto 2) * z \mapsto 2$

$\Rightarrow z \mapsto 2 * (x \mapsto 2 * y \mapsto 2)$  ~~using rule~~ - ①.

$\Rightarrow z \mapsto 2 * (y \mapsto 2 * x \mapsto 2)$  - ②

$\Rightarrow (z \mapsto 2 * y \mapsto 2) * x \mapsto 2$  - ③

1. Using rule  $P_1 * P_2 \Leftrightarrow P_2 * P_1$

2. Using rule  $P_1 * P_2 \Leftrightarrow P_2 * P_1$

3. Using rule  $(P_1 * P_2) * P_3 \Leftrightarrow P_1 * (P_2 * P_3)$

b)  $(x \mapsto 3 \wedge y \mapsto 3) \Rightarrow x = y$

This is not provable

Since the above is imprecise assertion, it can be ~~the~~ below 2,

(i)  $x \mapsto \boxed{3}$   
 $y \mapsto \boxed{3}$   
 $(x=y)$

(ii)  $x \mapsto \boxed{3}$   
 $y \mapsto \boxed{3}$

} By this figure, we can see  
 $x \neq y$

Hence it is not provable

$$c, \text{ emp} * ((\exists n. y \mapsto n) * w \mapsto 2) \Rightarrow \exists n. (y \mapsto n * w \mapsto 2)$$

This is provable

$$\Rightarrow \text{emp} * ((\exists n. y \mapsto n) * w \mapsto 2)$$

$$\Rightarrow ((\exists n. y \mapsto n) * w \mapsto 2) * \text{emp} \rightarrow \textcircled{1}$$

$$\Rightarrow (\exists n. y \mapsto n) * w \mapsto 2 \rightarrow \textcircled{2}$$

$$\Rightarrow \exists n. (y \mapsto n * w \mapsto 2) \rightarrow \textcircled{3}$$

1. Using rule  $P_1 * P_2 \Leftrightarrow P_2 * P_1$

2. Using rule  $P * \text{emp} \Leftrightarrow P$

3. Using rule  $(\exists n. P_1) * P_2 \Leftrightarrow \exists n. (P_1 * P_2)$  when  $n$  is not free in  $P_2$

$$d. (\exists n. y \mapsto n) * x \mapsto 2 \Rightarrow \exists n. (y \mapsto n * x \mapsto 2)$$

This is not provable

We can use the rule  $(\exists n. P_1) * P_2 \Leftrightarrow \exists n. (P_1 * P_2)$  only when  $n$  is not free in  $P_2$ .

However, in the given question  $x$  is ~~not~~ been present as  $P_2$  is  $x \mapsto 2$  hence it can not be used.

## 9. Non determinism

3.1

a)

$S_1 \triangleq \text{while } \{n \geq 0 \rightarrow y := n/y \mid n \leq 0 \rightarrow y := y \times n\}$

on the state  $\sigma \triangleq \{n=0, y=1\}$

$$\langle S, \{n=0, y=1\} \rangle$$

$$\rightarrow \langle y := n/y ; S, \{n=0, y=1\} \rangle$$

$$\xrightarrow{2} \langle S, \{n=0, y=0\} \rangle$$

$$\rightarrow \langle y := y \times n ; S, \{n=0, y=0\} \rangle$$

$$\xrightarrow{2} \langle S, \{n=0, y=0\} \rangle$$

$$\rightarrow \langle y := y \times n ; S, \{n=0, y=0\} \rangle$$

$$\xrightarrow{2} \langle S, \{n=0, y=0\} \rangle$$

b)

$$M(S, \sigma) = \{ \} \text{ or } \perp$$

4.

~~Q2~~4.1

$$S_2 \triangleq [x := 4; y := y + 1 \parallel y := y + 3], \theta \triangleq \{x = 3, y = 3\}$$

$$\langle [x := 4, y := y + 1 \parallel y := y + 3], (\{x = 3, y = 2\}, h) \rangle$$



$$\langle [y := y + 1 \parallel y := y + 3], (\{x = 2, y = 2\}, h) \rangle$$

$$\langle [x := 4, y := y + 1 \parallel \text{skip}], (\{x = 3, y = 5\}, h) \rangle$$



$$\langle [\text{skip} \parallel y := y + 3], (\{x = 2, y = 3\}, h) \rangle$$

$$\langle [y := y + 1 \parallel \text{skip}], (\{x = 2, y = 5\}, h) \rangle$$

$$\langle [y := y + 1 \parallel \text{skip}], (\{x = 5, y = 5\}, h) \rangle$$



$$\langle [\text{skip} \parallel \text{skip}], (\{x = 2, y = 6\}, h) \rangle$$

$$\langle [\text{skip} \parallel \text{skip}], (\{x = 5, y = 6\}, h) \rangle$$

5.

5.1

It took me 4 hours to finish this.