# **Practical-5**

Aim: Implement a program to validate that given grammar is LL(1) or not.

#### Code:

```
#include <iostream>
#include <map>
#include <set>
#include <vector>
#include <string>
using namespace std;
// Function to calculate the FIRST set for a given grammar
void calculateFirst(const map<char, vector<string>>& grammar, map<char,
set<char>>& firstSets) {
  bool changed = true;
  while (changed) {
     changed = false;
     for (auto& rule : grammar) {
       char nonTerminal = rule.first;
       for (const string& production : rule.second) {
          char firstSymbol = production[0];
          if (isupper(firstSymbol)) { // Non-terminal
            for (char f : firstSets[firstSymbol]) {
               if (firstSets[nonTerminal].insert(f).second) {
                 changed = true;
               }
            }
          } else { // Terminal or epsilon
            if (firstSets[nonTerminal].insert(firstSymbol).second) {
               changed = true;
            }
          }
       }
```

```
}
  }
}
// Function to calculate the FOLLOW set for a given grammar
void calculateFollow(const map<char, vector<string>>& grammar, map<char,
set<char>>& firstSets, map<char, set<char>>& followSets, char startSymbol) {
  followSets[startSymbol].insert('$'); // End of input marker
  bool changed = true;
  while (changed) {
     changed = false;
     for (auto& rule : grammar) {
        char nonTerminal = rule.first;
       for (const string& production : rule.second) {
          for (int i = 0; i < production.size(); i++) {
             char symbol = production[i];
             if (isupper(symbol)) {
               set<char> follow;
               if (i + 1 < production.size()) {
                  char nextSymbol = production[i + 1];
                  if (isupper(nextSymbol)) {
                     follow = firstSets[nextSymbol];
                     follow.erase('\sim'); // '\sim' represents epsilon (\epsilon)
                  } else {
                     follow.insert(nextSymbol);
               } else {
                  follow = followSets[nonTerminal];
               }
               for (char f : follow) {
                  if (followSets[symbol].insert(f).second) {
                     changed = true;
                  }
               }
             }
          }
       }
     }
```

```
}
}
// Function to check if the grammar is LL(1)
bool isLL1(const map<char, vector<string>>& grammar, map<char, set<char>>&
firstSets, map<char, set<char>>& followSets) {
  for (auto& rule : grammar) {
     char nonTerminal = rule.first;
     set<char> checked;
     for (const string& production : rule.second) {
        set<char> first = firstSets[production[0]];
        if (first.find('\sim') != first.end()) { // '\sim' represents epsilon (\epsilon)
           first.insert(followSets[nonTerminal].begin(), followSets[nonTerminal].end());
        for (char f : first) {
           if (checked.find(f) != checked.end()) {
              return false; // Conflict in predictive parsing table
           checked.insert(f);
        }
     }
  }
  return true;
}
int main() {
  // Example Grammar
  map<char, vector<string>> grammar = {
     {'E', {"TA"}},
     \{'A', \{"+TA", "\sim"\}\}, //'\sim' \text{ represents epsilon } (\epsilon)
     {'T', {"FB"}},
     \{'B', \{"*FB", "\sim"\}\}, //'\sim' \text{ represents epsilon } (\epsilon)
     {'F', {"(E)", "a"}}
  };
  char startSymbol = 'E';
  // Calculate FIRST and FOLLOW sets
  map<char, set<char>> firstSets;
  map<char, set<char>> followSets;
```

```
for (auto& rule : grammar) {
    firstSets[rule.first] = set<char>();
    followSets[rule.first] = set<char>();
}

// Calculate FIRST sets
    calculateFirst(grammar, firstSets);

// Calculate FOLLOW sets
    calculateFollow(grammar, firstSets, followSets, startSymbol);

// Check if the grammar is LL(1)
if (isLL1(grammar, firstSets, followSets)) {
    cout << "The grammar is LL(1)" << endl;
} else {
    cout << "The grammar is NOT LL(1)" << endl;
}
return 0;</pre>
```

# **Output:**

```
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The grammar is LL(1)

Process returned 0 (0x0) execution time : 0.135 s

Press any key to continue.
```

## **Questions**

### 1) How to create predictive parsing table?

Ans: **Step 1:** Calculate the FIRST set for each production.

**Step 2:** Calculate the FOLLOW set for each non-terminal.

**Step 3:** For each production, place it in the table under the columns for terminals in its

FIRST set. If it can derive  $\varepsilon$ , add it in the columns for terminals in the FOLLOW set.

Step 4: Ensure no two productions for the same non-terminal share a table cell.

#### 2) How to verify that given grammar is LL(1) or not?

Ans: Condition 1: The grammar should have no left recursion.

Condition 2: The FIRST sets of a non-terminal's productions must not overlap. If any production derives  $\epsilon$ , its FOLLOW set should be disjoint from the other FIRST sets. Verification: Build the predictive parsing table and check for conflicts (multiple entries in the same cell).