

Chap 3.

part 1 : C^* -opera construction on infinite spin systems

quasi-local algebra, reps and translational symmetry

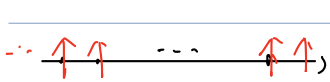
infinite lattice

naively. $\bigotimes_{i \in \mathbb{Z}} \mathcal{H}_i$ as the Hilbert space.

$\hookrightarrow \langle \psi, \xi \rangle = \prod_{i \in \mathbb{Z}} \langle \psi_i, \xi_i \rangle \neq 0$ does not converge

way out. 1. choose reference vector $\Omega_i \in \mathcal{H}_i \forall i$.

and consider only these sequences ψ s.t. $\psi_i \neq \Omega_i$ for finitely many i



$\Omega = |\uparrow \dots \uparrow\rangle$

$\psi = |\text{finitely many } \downarrow's\rangle$

\hookrightarrow completion $\mathcal{H}_\Omega \rightarrow \overline{\mathcal{A} = \mathcal{B}(\mathcal{H}_\Omega)}$

\downarrow
no local structure

\nearrow support on infinite space?

remark: different Ω 's often corresponds to inequiv. reps of opera

$$[\sigma_n^i, \sigma_m^j] = 2\delta_{nm} i \sum_{k=1}^3 \varepsilon_{ijk} \sigma_n^k$$

$\Omega_1 = |\uparrow \dots \uparrow\rangle \rightarrow \text{rep. 1}$

$\Omega_2 = |\downarrow \dots \downarrow\rangle \rightarrow \text{rep. 2}$

$\updownarrow \neq u$

2. quasi-local opera from local opera

What is the tensor product of two operas?

C^* -algebra = $*$ -algebra + completion w.r.t some norm

Def: suppose \mathcal{A}, \mathcal{B} are $*$ -alg.

then. their algebraic tensor product $\mathcal{A} \otimes \mathcal{B}$ consists of linear combinations

of elements $A \otimes B$. $A, B \in \mathcal{A}, \mathcal{B}$.

$\mathcal{A} \otimes \mathcal{B}$ is a $*$ -alg. by setting

$$(A_1 \otimes B_1)(A_2 \otimes B_2) := A_1 A_2 \otimes B_1 B_2$$

$$(A \otimes B)^* := A^* \otimes B^*$$

$\mathcal{A} \otimes \mathcal{B}$ is a C^* -alg. by some completion

T labels the sites of the system

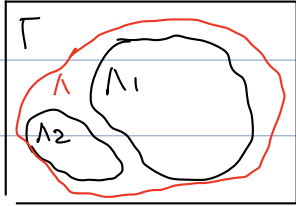
$\mathcal{P}(T)$ is the set of all subsets of T

$\mathcal{P}_f(T)$ is the subset of $\mathcal{P}(T)$ where all its elements are finite subsets

assume that $\forall x \in T$ there is an opera $\mathcal{A}(\{x\}) := M_d(\mathbb{C})$

if $\Lambda \in \mathcal{P}_f(T)$. $\mathcal{A}(\Lambda) = \bigotimes_{x \in \Lambda} \mathcal{A}(\{x\})$

which gives an assignment $\Lambda \mapsto \mathcal{A}(\Lambda) \quad \forall \Lambda \in \mathcal{P}_f(T)$



prop: this assignment preserves a local net structure

if $\Lambda_1, \Lambda_2 \in \mathcal{P}_f(T)$ and $\Lambda_1 \cap \Lambda_2 = \emptyset$.

then $[\mathcal{A}(\Lambda_1), \mathcal{A}(\Lambda_2)] = \{0\}$

here $\mathcal{A}(\Lambda_i)$ $i=1,2$ is embedded in to a sufficiently large Λ

local observables . $\mathcal{A}_{loc} = \bigcup_{\Lambda \in \mathcal{P}_f(T)} \mathcal{A}(\Lambda)$

\hookrightarrow completion w.r.t norms on each $\mathcal{A}(\Lambda) \rightarrow C^*$ algebra called \mathcal{A}
(quasi-local algebra)

additional construction

if $A \in \mathcal{A}(\Lambda)$ for some $\Lambda \in \mathcal{P}_f(T)$ A is localized in Λ

the smallest Λ is called $\text{supp}(A)$ we require $\mathcal{A}(\emptyset) = \mathbb{C}1$

$\Lambda \in \mathcal{P}(T)$'s complement is denoted as Λ^c

Q: can A be localized in infinite regions? . ($\Lambda \in \mathcal{P}(T) \wedge \Lambda \notin \mathcal{P}_f(T)$)

A: yes. $\mathcal{A}(\Lambda) = \overline{\bigcup_{\Lambda' \in \mathcal{P}_f(\Lambda)} \mathcal{A}(\Lambda')}^{|| \cdot ||}$ \nwarrow completion same as quasi-local algebra
 \downarrow
is $\mathcal{A}(T)$ quasi-local algebra?

Representation .

\nearrow a C^* -alg \mathcal{A} is simple
iff only closed
two-sided ideals are $\{0\}$ and \mathcal{A}

prop: $\mathcal{A}(\mathbb{Z}^d)$ (quasi-local algebra on \mathbb{Z}^d) is simple.

cor: every non-zero rep of $\mathcal{A}(\mathbb{Z}^d)$ is faithful and hence isometric

$$\mathcal{A}(\mathbb{Z}^d) \subset \mathcal{B}(\mathcal{H})$$

we don't need to do the entire GNS construction

equivalence under GNS rep. ?

prop: $\mathcal{A} := \mathcal{A}(T)$ is a quasi-local algebra

and ω_1, ω_2 are pure states on \mathcal{A} .

the following criteria are equivalent:

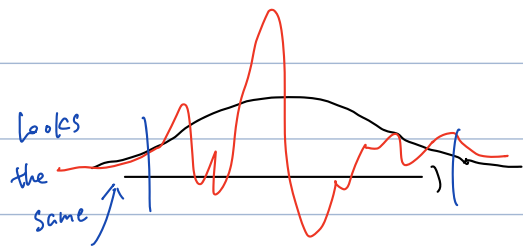
1. corresponding GNS reps of ω_1, ω_2 are equivalent

2. $\forall \varepsilon > 0 \exists \Lambda \in \mathcal{P}_f(T)$ s.t.

$$|\omega_1(A) - \omega_2(A)| < \varepsilon \|A\|$$

$$\forall A \in \mathcal{A}(\Lambda) \text{ with } \Lambda \in \mathcal{P}_f(\Lambda_\varepsilon^c)$$

if we restrict
to observables "far away"
the two states
look the same

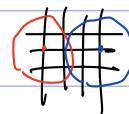


Translation symmetry

if translation symmetry exists on $\Gamma = \mathbb{Z}^d$

we can define a map $\tau_y: \mathcal{A}_{loc} \rightarrow \mathcal{A}_{loc} \quad \forall y \in \mathbb{Z}^d$

by defining $\tau_y(A(x)) = A(x+y)$



this map can be extended to an automorphism of $\mathcal{A}(\mathbb{Z}^d)$

$$\tau_y(\mathcal{A}(\Lambda)) = \mathcal{A}(\Lambda+y) \quad \forall \Lambda \in \mathcal{P}(\mathbb{Z}^d)$$

$\Rightarrow \tau: \mathbb{Z}^d \rightarrow \text{Aut}(\mathcal{A}(\mathbb{Z}^d))$ is a group homomorphism

translation sym $\left\{ \begin{array}{l} \text{site} \\ \text{opera} \\ \text{state} \end{array} \right.$

translation invariant state. can be defined as a state ω .

such that $\omega(\tau_x(A)) = \omega(A) \quad \forall x \in \mathbb{Z}^d$ and $A \in \mathcal{A}$

ω has the property that. \exists a unitary representation

$x \mapsto U(x)$. $U(x)$ is a unitary in GNS rep w.r.t ω

implementing the translation

$$\pi(\tau_x(A)) = U(x) \pi(A) U(x)^*$$

asymptotic Abelianess of quasi-local algebra,

if we move a local operator far enough.

it will commute with any other local operator

Thm: for a quasi-local $\mathcal{A}(\mathbb{Z}^d)$ with $x \mapsto T_x$ as the action of

translation group. then $\forall A, B \in \mathcal{A}(\mathbb{Z}^d)$,

$$\lim_{|x| \rightarrow \infty} \|[T_x(A), B]\| = 0$$

part 2. time translation, dynamics. \rightarrow Hamiltonian

\downarrow

GS, & KMS state