2.4 Assembly Programming Exercise - Computing Average

Note to students: Some important concepts to bring across are:

- 1. Understand how consecutive elements in an array can be accessed.
- 2. Understand how predetermine looping cycles can be implemented using a countdown counter.
- 3. Understand how unsigned and signed division by 2ⁿ can be implemented using the logical and arithmetic shift right operation.
- 4. Understand how the use of appropriate logical operators can mask out unwanted high-order bits.
- 5. Understand the efficient use of conditional execution instructions to avoid branching.

Write an ARM assembly language program to compute the average value of **eight** numbers in an array.

- (1) Your program should be written based on the following specifications:
 - a) These numbers are word-sized unsigned integers (i.e. **32-bit unsigned** numbers).
 - b) The eight numbers are in consecutive memory locations starting at address **0x0100**.
 - c) The average (quotient) and remainder are stored in registers **R0** and **R1** respectively.

Note: You are free to use the partial completed VisUAL ARM assembly program template *Tutorial 2_4-Template* to test out your answer.

Suggested solution:

Start	MOV	R1,#0x100	; initialize array pointer to start of array
	MOV	R0,#0	; initialize cumulating register R0 to zero
	MOV	R2,#8	; initialize loop counter for 8 loops
SumLoop	LDR	R3,[R1],#4	; get current element and increment array pointer
	ADD	R0,R0,R3	; add array element to current sum in R0
	SUBS	R2,R2,#1	; decrement loop counter R2 (check Z flag)
	BNE	SumLoop	; branch back until count is 0
GetQnR	MOV	R1,R0	; make copy of cumulated sum in R1
	AND	R1,R1,#0x07	; clear all bits except lowest 3 bits to get remainder in R1
	MOV	R0,R0,LSR #3	; logical shift right by 3 bits to divide by 8 to get quotient

Note the following:

- 1. The register R1 is used as an address pointer to the array of numbers. It is first initialized to 0x100, the start address of the number array. The LDR instruction uses R1 as an indirect register with index (and post increment) to retrieve each 32-bit number each time around the loop. 4 is added to the content in R1 because each number (4 bytes) is stored in 4 consecutive memory locations.
- 2. The register **R2** is the decrementing loop counter which is initialized to 8 at the start to do 8 loops. **BNE** will continue to loop back to the **SumLoop** label until the **SUBS** instruction sets the Z flag.
- 3. Once the sum of all 8 numbers have been computed in register R0, a copy of this is made in R1 in order to compute the remainder in R1. The remainder of any division by 8 (i.e. 2³) is the least significant 3 binary bits of original number before division. This is obtained by using the AND operation to mask all high order bits to 0, except bits 0 to 2.
- 4. The positive quotient is obtained in **RO** by doing a logical right shift (**LSR**) by 3 bits.

(2) How would your program be changed if these numbers are **32-bit signed** numbers instead?

Suggested solution:

For signed numbers, Arithmetic Shift Right (**ASR**) can be used. Arithmetic right shifts are equivalent to logical right shifts (**LSR**) for positive signed numbers. Arithmetic right shifts for negative numbers (2's complement) is similar in concept to division by a power of 2^n but the negative quotients with non-zero remainder will round towards negative infinity instead of towards 0 as is usually expected (e.g. -28/8 gives a quotient of -4 instead of -3, because -3.5 is rounded toward -4 and not -3. As such, additional steps need to be taken to handle this rounding behaviour when using **ASR** when a negative quotient is detected.

Method #1 – Convert to positive to compute quotient and then negate quotient again

```
GetQnR MOVS R1,R0 ; make copy of cumulated sum in R1 (and check N flag)

RSBMI R0,R0,#0 ; if negative, negate to positive

AND R1,R1,#0x07 ; clear all bits except lowest 3 bits to get remainder in R1

MOV R0,R0,LSR #3 ; logical shift right by 3 bits to divide by 8

RSBMI R0,R0,#0 ; if quotient was earlier negative, negate to negative again
```

Note: MOVS is used to check if the sum is negative. If the sum is negative, then the conditional execution version of RSBMI, which will execute only if the N flag is set (i.e. condition is MInus). Under this situation, RSB will negate the sum to its positive equivalent so that the quotient can be computed as in the unsigned example earlier. However, the positive quotient must be reverted to if negative equivalent with the second RSBMI. Remember RSB unlike SUB implements reverse subtraction, which is (0 - R0).

Method #2 – Use ASR but add #1 to negative quotient if there is a non-zero remainder

```
GetQnR MOV R1,R0 ; make copy of cumulated sum in R1

MOVS R0,R0,ASR #3 ; arithmetic shift right by 3 bits to divide by 8

ADDMI R0,R0,#1 ; add 1 if quotient is -ve to round quotient towards zero

ANDS R1,R1,#0x07 ; clear all bits except lowest 3 bits to get remainder in R1

SUBEQ R0,R0,#1 ; if remainder is zero, undo earlier add 1 operation quotient
```

Note: **ASR** is used to implement an arithmetic shift right, which will handle the sign of the quotient. However, if the quotient is negative, **ASR** will only give us a negative quotient that is rounded towards negative infinity. In this case, 1 is added to the negative quotient to get a rounding towards zero. This is done by the **ADDMI** instruction, which will only execute if the computed quotient is negative. However, this should only be done if there is a non-zero remainder. If there is no remainder, the addition of 1 will give us an incorrect answer. The **SUBEQ** instruction will subtract the negative quotient by 1 if the computed remainder done using the **ANDS** instruction yields a zero (i.e. Z flag set).