

Beautiful Arrangement II

Given two integers n and k , you need to construct a list which contains n different positive integers ranging from 1 to n and obeys the following requirement:

Suppose this list is $[a_1, a_2, a_3, \dots, a_n]$, then the list $[|a_1 - a_2|, |a_2 - a_3|, |a_3 - a_4|, \dots, |a_{n-1} - a_n|]$ has exactly k distinct integers.

If there are multiple answers, print any of them.

Example 1:

Input: $n = 3, k = 1$

Output: $[1, 2, 3]$

Explanation: The $[1, 2, 3]$ has three different positive integers ranging from 1 to 3 , and the $[1, 1]$ has exactly 1 distinct integer: 1 .

Example 2:

Input: $n = 3, k = 2$

Output: $[1, 3, 2]$

Explanation: The $[1, 3, 2]$ has three different positive integers ranging from 1 to 3 , and the $[2, 1]$ has exactly 2 distinct integers: 1 and 2 .

Note:

1. The n and k are in the range $1 \leq n, k \leq 10^4$.

Solution 1

if you have n number, the maximum k can be $n - 1$;

if n is 9, max k is 8.

This can be done by picking numbers interleavingly from head and tail,

```
// start from i = 1, j = n;  
// i++, j--, i++, j--, i++, j--
```

```
1  2  3  4  5  
  9  8  7  6  
out: 1 9 2 8 3 7 6 4 5  
dif: 8 7 6 5 4 3 2 1
```

Above is a case where k is exactly $n - 1$

When k is less than that, simply lay out the rest (i, j) in incremental order (all diff is 1). Say if k is 5:

```
      i++ j-- i++ j--  i++ i++ i++ ...  
out: 1  9  2  8  3  4  5  6  7  
dif:  8  7  6  5  1  1  1  1
```

C++

```
class Solution {  
public:  
    vector<int> constructArray(int n, int k) {  
        vector<int> res;  
        for (int i = 1, j = n; i <= j; ) {  
            if (k > 1) {  
                res.push_back(k-- % 2 ? i++ : j--);  
            }  
            else {  
                res.push_back(k % 2 ? i++ : j--);  
            }  
        }  
        return res;  
    }  
};
```

C++ Compact

```
class Solution {  
public:  
    vector<int> constructArray(int n, int k) {  
        vector<int> res;  
        for (int i = 1, j = n; i <= j; )  
            res.push_back(k > 1 ? (k-- % 2 ? i++ : j--) : (k % 2 ? i++ : j--));  
        return res;  
    }  
};
```

Java

```
class Solution {
    public int[] constructArray(int n, int k) {
        int[] res = new int[n];
        for (int i = 0, l = 1, r = n; l <= r; i++)
            res[i] = k > 1 ? (k-- % 2 != 0 ? l++ : r--) : (k % 2 != 0 ? l++ : r--);
        return res;
    }
}
```

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Solution 2

The requirement of k distinct distance can be achieved from $1, 2, \dots, k+1$ ($\leq n$), by the following strategy:

$1, k+1, 2, k, 3, k-1 \dots$;

The distance **of this** sequence **is** $k, k-1, k-2, \dots, 2, 1$

Then append the remaining numbers to the list.

```
class Solution {
public:
    vector<int> constructArray(int n, int k) {
        int l = 1, r = k+1;
        vector<int> ans;
        while (l <= r) {
            ans.push_back(l++);
            if (l <= r) ans.push_back(r--);
        }
        for (int i = k+2; i <= n; i++)
            ans.push_back(i);
        return ans;
    }
};
```

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Solution 3

When $k = n-1$, a valid construction is $[1, n, 2, n-1, 3, n-2, \dots]$. One way to see this is, we need to have a difference of $n-1$, which means we need 1 and n adjacent; then, we need a difference of $n-2$, etc.

This leads to the following idea: we will put $[1, 2, \dots, n-k-1]$ first, and then we have $N = k+1$ adjacent numbers left, of which we want k different differences. This is just the answer above translated by $n-k-1$: we'll put $[n-k, n, n-k+1, n-1, \dots]$ after.

```
def constructArray(self, n, k):
    ans = range(1, n - k)
    for d in xrange(k+1):
        if d % 2 == 0:
            ans.append(n-k + d/2)
        else:
            ans.append(n - d/2)

    return ans
```

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