

Data Structures and Algorithms

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Syllabus

Basic Concepts of Data Structures

Definitions; Data Abstraction; Performance Analysis - Time & Space Complexity, Asymptotic Notations; Polynomial representation using Arrays, Sparse matrix (Tuple representation); Stacks and Queues - Stacks, Multi-Stacks, Queues, Circular Queues, Double Ended Queues; Evaluation of Expressions- Infix to Postfix, Evaluating Postfix Expressions.

1 Definitions

- **Data Structures:** ways of organizing and storing data in a computer so that it can be accessed and modified efficiently. Types:
 1. Linear: Arrays, Linked Lists, Stacks, Queues
 2. Non-linear: Trees, Graphs
- **Data Abstraction:** concept of hiding the internal details of how data is stored or maintained and only showing the essential features or operations that can be performed on the data.

2 Performance Analysis

1. **Time Complexity**
2. **Space Complexity**

3 Stack

It follows FILO (First In Last Out) scheme

- pop - Removes from top
- push - Adds to top
- peek/top - See topmost element

3.1 Implementation of Stack

```
1  #include <stdio.h>
2  #define MAX 2
3
4  int stack[MAX];
5  int top = -1;
6
7  void push(int a) {
8      if (top + 1 >= MAX) {
9          printf("Stack Overflow\n");
10     } else {
11         stack[++top] = a;
12     }
13 }
14
15 int pop() {
16     if (top == -1) {
17         printf("Stack underflow \n");
18         return -1;
19     } else {
20         return stack[top--];
21     }
22 }
23
24 void display() {
25     if (top == -1) {
26         printf("Stack is empty!\n");
27     } else {
28         for (int i = top; i > -1; i--) {
29             printf("%d ", stack[i]);
30         }
31         printf("\n");
32     }
33 }
34
35 int main() {
36     pop();
37
38     push(5);
39     push(8);
40
41     display();
42
43     pop();
44
45     display();
46
47     return 0;
48 }
```

4 Queue

It follow FIFO (First In First Out) scheme

- pop - Removes from front
- push - Adds to rear
- peek/top - See frontmost element

4.1 Implementation of Queue

```
1 #include <stdio.h>
2 #define MAX 3
3
4 int queue[MAX];
5 int rear = -1, front = -1;
6
7 void enqueue(int a) {
8     if (rear + 1 >= MAX) {
9         printf("Queue Overflow\n");
10    } else {
11        if (front == -1) front = 0;
12        queue[++rear] = a;
13    }
14 }
15
16 int dequeue() {
17     if (front > rear) {
18         printf("Queue underflow \n");
19         return -1;
20     } else {
21         return queue[front++];
22     }
23 }
24
25 void display() {
26     if (rear == -1 || front > rear) {
27         printf("Queue is empty!\n");
28     } else {
29         for (int i = rear; i >= front; i--) {
30             printf("%d ", queue[i]);
31         }
32         printf("\n");
33     }
34 }
```

5 Multi-Stacks

2 or more stacks in a single array.

5.1 2 stacks in one array

Stack 1 grows from left to right. Stack2 grows from right to left. **Condition:**

- To push into stack1:

$$top1 + 1 < top2$$

- To push into stack2:

$$top2 - 1 > top1$$

6 Addition of sparse polynomial

All the polynomials are stored inside an array of structures:

```

1 // Structure to represent a term
2 typedef struct {
3     int coeff;
4     int expo;
5 } Term;
6
7 Term polynomial[] = {{2, 3},{4, 0}} // 2x^3 + 4
```

Listing 3: Sparse Polynomial Addition Outline

6.1 Logic when adding 2 polynomials

let it be poly1 (with i as indexing), poly2 (with j as indexing) & result (with k as indexing)

- If $\text{poly1}[i].\text{exp} == \text{poly2}[j].\text{exp}$: add coefficients
- If $\text{poly1}[i].\text{exp}$ is greater than $\text{poly2}[j].\text{exp}$: Copy over $\text{poly1}[i]$
- If $\text{poly1}[i].\text{exp}$ is less than $\text{poly2}[j].\text{exp}$: Copy over $\text{poly2}[j]$

7 Sparse Matrix

Sparse matrix is a matrix with most of its elements are zero. Eg:

15	0	0	0	91	0
0	11	0	0	0	0
0	3	0	0	0	28
22	0	-6	0	0	0
0	0	0	0	0	0
-15	0	0	0	0	0

7.1 Finding transpose of sparse matrix

We use an array of structs to save values of non-zero values of matrix.

```
1 typedef struct {
2     int row;
3     int col;
4     int value;
5 } Term;
```

Listing 4: Representing sparse matrix

The first row of the array would be the **metadata** of the matrix: Number of rows, Number of columns and number of non-zero elements.

One such array for representation might look like this:

```
1 Term a[] = {
2     {6, 6, 8}, // metadata: 6 rows, 6 cols, 8 non-zero
3     values
4     {0, 0, 15},
5     {3, 0, 22},
6     {5, 0, -15},
7     {1, 1, 11},
8     {2, 1, 3},
9     {3, 2, -6},
10    {0, 4, 91},
11    {2, 5, 28}
12 };
```

7.1.1 Program for transpose

```
1 void transpose(Term a[], Term b[]) {
2     int n = a[0].value;
3
4     b[0].row = a[0].col;
5     b[0].col = a[0].row;
6     b[0].value = n;
7 }
```

```

8      if (n > 0) {
9          int indexb = 1; // To keep track of values in b
10         for (int i = 0; i < a[0].col; i++) { // For sorting
11             for (int j = 1; j <= n; j++) { // For iteration
12                 if (a[j].col == i) {
13                     b[indexb].row = a[j].col;
14                     b[indexb].col = a[j].row;
15                     b[indexb].value = a[j].value;
16                     indexb++;
17                 }
18             }
19         }
20     }
21 }

```

Listing 5: Finding the transpose of a sparse matrix