

VERIFICATION DOCUMENT- ALU

CONTENTS	Page Number	
1.1 ALU introduction.	02	
1.2 Advantages of ALU.	02	
1.3 Disadvantages of ALU.	03	
1.4 Use cases of ALU.	03	
1.5 Project Overview of ALU.	04	
1.6 Design Features.	05	
1.7 Design Limitation.	06	
1.8 Design diagram with interface signals.	07	
2.1 Verification Architecture	07	
2.2 Verification Architecture for ALU	09	
2.3 flow chart of sv components	11	

CHAPTER 1 – DESIGN OVERVIEW

1.1 ALU introduction:

The Arithmetic Logic Unit (ALU) serves as the central processing component in any digital architecture. Imagine it as an advanced computational engine capable of executing both logical and arithmetic tasks. This implementation is notable for its adaptability it is designed as a parameterized module, allowing us to specify the bitwidth to match diverse application requirements.

What sets this ALU apart is its all-encompassing treatment of digital computation. Beyond standard operations like addition and subtraction, it supports advanced features such as bitwise rotations evaluations and incorporates mechanisms for validating inputs and detecting errors. With synchronous operation and reliable clock and reset integration, the design aligns perfectly with the needs of contemporary digital systems.

1.2 Advantages of ALU:

• Scalable Design:

You can easily scale to 8, 16, 32, 64-bit, or any-bit without changing the design. Saves time and avoids new bugs.

• Wide Range of Operations:

Supports 14 arithmetic and 14 logic operations — from simple add to complex rotate. So, fewer external components are needed.

• Smart Input Handling:

It uses INP_VALID to handle inputs arriving at different times. A timeout avoids system hangs if inputs are missing.

• Status Flags:

Each operation gives useful flags like carry, overflow, greater, less, equal, and error — helping the system make smart choices.

• Power Saving:

Clock enable (CE) lets you turn off the ALU when not needed — great for saving power in battery devices.

• Error Detection:

Catches invalid operations, especially in rotate cases where bad input patterns can cause problems.

1.3 Disadvantages of ALU:

• Timeout Isn't Flexible:

The ALU waits exactly 16 clock cycles for inputs — no more, no less. That might be too long for fast systems or too short for slower ones, but you can't adjust it.

• Commands Can Be Confusing:

The same command number does different things depending on the mode. For example, CMD = 0 means ADD in one mode and AND in another. Easy to mix up if you're not careful.

• One Error Signal for Everything:

If something goes wrong, you just get a generic error flag. It doesn't tell you what the problem is — was it a timeout, a wrong command, or a bad input? Makes troubleshooting harder.

• Too Big for Simple Jobs:

The ALU supports a lot of features, which is great, but it also uses more hardware. If your system only needs basic stuff, this design might be overkill.

• Rotate Operation Is Tricky:

Rotate left/right needs operand B to follow strict rules — like making sure bits [7:4] are zero. It's easy to mess this up in software, leading to errors.

1.4 Use cases of ALU:

• Arithmetic Operations

- Performs basic math like addition, subtraction, multiplication, and division
- Essential for tasks like updating counters, calculating addresses, or doing math in a program

• Logical Operations

• Executes AND, OR, XOR, NOT operations

• Used in comparing values, bit masking, and decision-making logic in CPUs and digital circuits

• Bit Shifting and Rotation

- Shifts bits left or right, rotates bits for specific applications
- Common in encryption, encoding/decoding, and data manipulation

Comparisons

- Compares two numbers to check greater than, less than, or equal
- Useful for if-else, loops, and conditional branching in software and hardware

Address Calculations

Used in memory operations to calculate next instruction or data address

• Checksum and CRC Computation

• Helps compute checksums and cyclic redundancy checks for error detection in communication systems

• Control Signal Generation

• Uses comparison results to generate control signals in FSMs (Finite State Machines) and controllers

• Overflow and Carry Detection

- Detects arithmetic overflow or carry during operations
- Important for signed/unsigned number handling and ensuring accuracy

• Executing CPU Instructions

• The ALU is the core part of the execution stage of a CPU — it carries out actual operations for instructions

1.5 Project Overview of ALU:

This project focuses on building a powerful, yet flexible ALU (Arithmetic Logic Unit) that can be used in a wide range of digital systems, from processors to embedded controllers. The key idea is parameterization — allowing the ALU to work with different bit-widths such as 16, 32, 64, and 128 bits, making it suitable for both low-end and high-performance applications. At its core, the ALU processes two operands (OPA and OPB) and operates in two distinct modes: arithmetic mode (MODE=1) and logical mode (MODE=0). Each mode supports 14 operations,

ranging from basic addition, subtraction, and multiplication to advanced bitwise logic and rotate instructions. This dual-mode structure efficiently utilizes a 4-bit command system while keeping arithmetic and logic functions clearly separated.

To support real-world system requirements, the design includes advanced control features. The INP_VALID signal helps manage inputs that may not arrive at the same time, a common issue in pipelined or asynchronous environments. A built-in timeout mechanism ensures that the ALU doesn't wait forever for missing inputs, improving system reliability. The ALU also offers comprehensive status reporting, including overflow detection, carry flags, and comparison results (greater than, less than, equal), enabling smarter decisions by system software. For example, a control unit or CPU can use these flags for branching or error handling.

Error detection is another strong point of this design. It has dedicated error management for invalid operations—especially for complex ones like variable bit rotations, where certain patterns in operand B can cause undefined behaviour. Instead of failing silently, the ALU flags these issues, making debugging and system integration much easier. Overall, the project aims to create a highly reusable, scalable, and robust ALU that balances performance, flexibility, and reliability—essential for modern digital design.

1.6 Design Features:

• Synchronous with Asynchronous Reset

Works on the rising edge of the clock and can reset instantly using an async reset — useful for reliable startup and emergency resets.

• Flexible Bit-Width

Uses parameters to set data width (16, 32, 64, 128 bits, etc.), so it can be easily adapted without changing the core code.

Smart Operand Control

A 2-bit INP_VALID signal shows if none, one, or both inputs are ready — great for handling inputs that come at different times.

Advanced Rotate Operations

Supports variable rotate left/right based on operand B. Includes error checks to catch bad input values.

• Rich Comparison Output

Gives all comparison results — greater (G), less (L), and equal (E) — at the same time for faster decisions in control logic.

• Built-in Overflow Detection

Automatically detects overflow in arithmetic operations to help avoid errors in calculations.

• Power Saving Option

Has a CE (clock enable) input so the ALU can be turned off when not needed — helpful for low-power systems.

1.7 Design Limitation:

• Fixed Timeout

The 16-cycle timeout is hardcoded. If a system needs a shorter or longer wait, the design must be changed.

Mode-Based Command Confusion

Same command code does different things in arithmetic and logical modes — this can lead to software mistakes.

• One Error Bit for All Issues

Only one ERR signal is used for all errors, so you can't tell if it was a timeout, invalid command, or input problem.

• Limited Rotate Range

Rotate operations only support up to 8 positions — may not be enough for larger bit-widths.

• Fixed Operand Priority

In case of timeout, the ALU always uses the latest operand — which may not match what some systems expect.

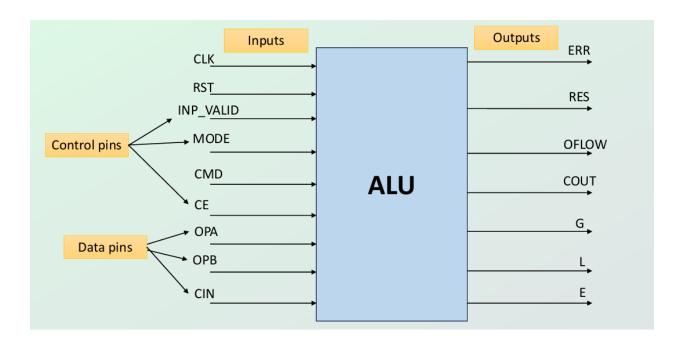
• No Pipelining

It executes one operation at a time with no pipeline support, which can slow down performance in fast systems.

Wasted Resources

Even unused operations still take up hardware space due to the parameterized design — not ideal for small or resource-limited chips.

1.8 Design diagram with interface signals:



Signal Name	Type	Description
INP_VALID	Input	Input valid signal - indicates when input data is valid
MODE	Input	Mode selection signal - determines ALU operation mode
CMD	Input	Command signal - specifies the specific ALU operation

OPA	Input	Operand A - first arithmetic/logic operand
ОРВ	Input	Operand B - second arithmetic/logic operand
CIN	Input	Carry In - input carry for arithmetic operations

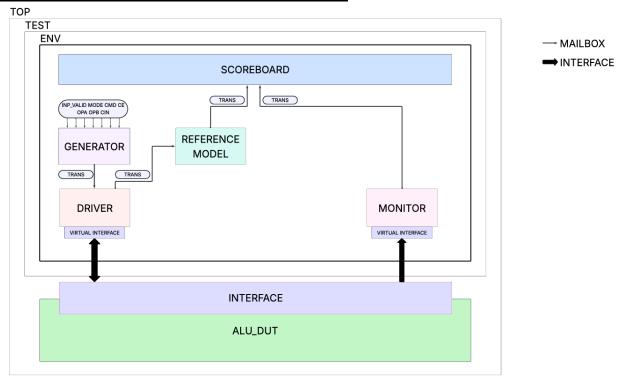
Output ports:

Signal Name	Туре	Description
ERR	Output	Error signal - indicates if an error occurred during operation
RES	Output	Result - the output result of the ALU operation
OFLOW	Output	Overflow - indicates arithmetic overflow condition
COUT	Output	Carry Out - output carry from arithmetic operations
G	Output	
L	Output	Less than - comparison result flag
E	Output	Equal - comparison result flag

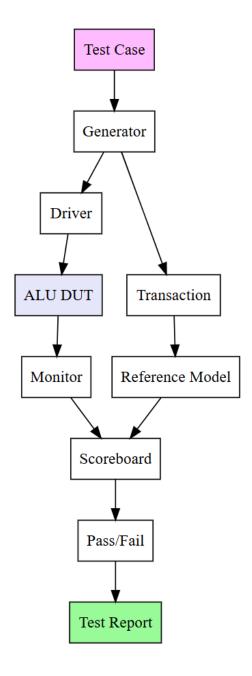
CHAPTER 2 - Verification Architecture

2.1 Verification Architecture:

2.2 Verification Architecture for ALU:



The figure shows general testbench architecture used for verifying digital designs. At the top module is the Design Under Verification (DUV). The test has the testbench environment that includes a transaction generator that creates and manages the input testcases for testing. These transactions are randomized within the generator and then transmitted to the driver, which applies them as signals to the DUV through a virtual interface. Outputs from the DUV are monitored by a monitor, which forwards this data to a scoreboard for checking correctness against expected results calculated by a reference model. The environment block encapsulates all these components, ensuring coordination among them, while the scoreboard overseas the entire verification process to validate the DUV.



The ALU verification project follows a structured testbench architecture to thoroughly validate the Arithmetic Logic Unit's functionality. The process begins with the Test Case, which defines specific operations to verify, such as addition or logical functions. The Generator creates input transactions (data and operations), which are passed to the Driver for conversion into signal-level inputs. These signals are applied to the ALU (DUT) through a Virtual Interface (v1), while a

Reference Model independently calculates expected results. The Monitor observes the ALU's outputs and converts them into readable data, which is then compared against the Reference Model's predictions by the Scoreboard to identify discrepancies.

Communication between components is managed via a Mailbox, ensuring synchronized data transfer, such as sending transactions from the Generator to both the Driver and Reference Model. The entire flow is designed to systematically validate the ALU's behavior, from stimulus generation to result checking. Finally, a Test Report summarizes the outcomes, highlighting any failures and confirming successful operations. This approach ensures comprehensive verification, catching design errors early and guaranteeing the ALU's reliability before deployment.