

SWPP 2023

Compiler Optimization Ideas

Team 1 - 박종한 배준익 이경훈 현재익

A dark blue diagonal gradient bar that starts from the bottom left and extends towards the top right, covering the lower half of the slide.

Introducing our Team

```
struct Team1 {  
    member 박종한; // foxisobese  
    member 배준익; // heatz123  
    member 이경훈; // pzmaus  
    member 현재익; // SyphonArch  
}
```

SNUCSE 19 // 3학년

아이디어 1

Lazy Evaluation With SIMD

`/* Cluster integer additions into SIMD style operations */`

덧셈 연산은 곱셈 연산에 비해 **cost**가 매우 크다.

이 **cost**는 **int**의 **bit** 너비와 무관하다.

Condition:

- 반복문 등에서 순서가 중요치 않은, 상대적으로 작은 **int**형의 덧셈이 여러 번 일어나는 경우

Optimization:

- 연산을 뒤로 미룸과 동시에 하나의 **register**로 모아서, 한번의 덧셈으로 처리한다.

Integer Sum	<code><reg> = sum <val1> ... <val18> <bw></code> <code><bw> := 1 8 16 32 64</code>	10
-------------	---	----

example

```
3    %c1 = add i4 %a1, %b1
4    %c2 = add i4 %a2, %b2
5    ...
6    %c16 = add i4 %a16, %b16
7    //store c1 ~ c16
8
9
10   %ptr1 = bitcast [16 * i4]* %a to i64*
11   %ptr2 = bitcast [16 * i4]* %b to i64*
12   %val1 = load i64, i64& %ptr1
13   %val2 = load i64, i64& %ptr2
14   %c = add i64 %val1, %val2
15   %ans = bitcast i64 %c to [16 * i4]*
16   //store
```


example

```
1  define i32 @main() {
2      ; Initialize variables
3      %i = add i32 1, 0
4      %sum = add i32 0, 0
5
6      ; Loop from 1 to 10
7      br label %loop_cond
8
9  loop_cond:
10     %cmp = icmp sle i32 %i, 10
11     br i1 %cmp, label %loop_body, label %loop_exit
12
13  loop_body:
14     %sum_temp = add i32 %sum, %i
15     %i_inc = add i32 %i, 1
16     br label %loop_cond
17
18  loop_exit:
19     ; Print the sum
20     %format_str = c"Sum: %d\n\00"
21     call i32 (@i8*, ...) @printf(i8* %format_str, i32 %sum)
22
23     ret i32 0
24 }
25
```

```
1  define i32 @main() {
2      ; Initialize variables
3      %i = add i32 1, 0
4      %sum = add i32 0, 0
5
6      ; Loop from 1 to 10
7      br label %loop_cond
8
9  loop_cond:
10     %cmp = icmp sgt i32 %i, 10
11     br i1 %cmp, label %loop_exit, label %loop_body
12
13  loop_body:
14     %sum_temp = add i32 %sum, %i
15     %i_inc = add i32 %i, 1
16     br label %loop_cond
17
18  loop_exit:
19     ; Print the sum
20     %format_str = c"Sum: %d\n\00"
21     call i32 (@i8*, ...) @printf(i8* %format_str, i32 %sum)
22
23     ret i32 0
24 }
25
```

아이디어 3

Registers over Stack over Heap

`/* Heap memory is extra expensive! */`

Heap의 cost는 바이트당 1024이며, load/store cost가 추가로 있다.

Stack은 cost가 없으며, load/store cost가 있다.

Register는 cost도 없고, load/store overhead도 없다!

스펙에 의해 33개의 general purpose register가 있으므로, 가능하다면 이를 최대한 활용한다.

Memory usage cost.

- The memory usage cost is 1024 times the maximum heap-allocated byte size at any moment.
- For example, the memory usage cost of

```
r1 = malloc 8
free r1
r2 = malloc 8
free r2
```

is $1024 * 8 = 8192$, because the maximum memory usage is 8 bytes.

example

```
6  define i32 @main() {
7      ; Allocate a block of 16 bytes on the heap and get a pointer to it
8      %ptr = call i8* @malloc(i64 16)
9
10     ; Cast the pointer to i32* so we can store integers in it
11     %array_ptr = bitcast i8* %ptr to i32*
12
13     ; Store integers in the four 4-byte chunks
14     store i32 1, i32* %array_ptr
15     store i32 2, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 1)
16     store i32 3, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 2)
17     store i32 4, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 3)
18
19     ; Add up the integers in the array
20     %sum1 = load i32, i32* %array_ptr
21     %sum2 = load i32, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 1)
22     %sum3 = load i32, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 2)
23     %sum4 = load i32, i32* getelementptr inbounds (i32, i32* %array_ptr, i64 3)
24     %sum = add i32 %sum1, %sum2
25     %sum = add i32 %sum, %sum3
26     %sum = add i32 %sum, %sum4
27
28     ; Print the sum
29     %format_str = c"Sum: %d\n\00"
30     call i32 (i8*, ...) @printf(i8* %format_str, i32 %sum)
31
32     ; Free the allocated memory
33     call void @free(i8* %ptr)
34
35     ret i32 0
36 }
37
```

```
1  define i32 @main() {
2      ; Store integers in registers
3      %num1 = alloca i32
4      %num2 = alloca i32
5      %num3 = alloca i32
6      %num4 = alloca i32
7      store i32 1, i32* %num1
8      store i32 2, i32* %num2
9      store i32 3, i32* %num3
10     store i32 4, i32* %num4
11
```

```
12     ; Add up the integers
13     %sum1 = load i32, i32* %num1
14     %sum2 = load i32, i32* %num2
15     %sum3 = load i32, i32* %num3
16     %sum4 = load i32, i32* %num4
17     %sum = add i32 %sum1, %sum2
18     %sum = add i32 %sum, %sum3
19     %sum = add i32 %sum, %sum4
20
```

```
21     ; Print the sum
22     %format_str = c"Sum: %d\n\00"
23     call i32 (i8*, ...) @printf(i8* %format_str, i32 %sum)
24
25     ret i32 0
26 }
27
```

```
1  define i32 @main() {
2      ; Store integers in registers
3      %num1 = add i32 1, 0
4      %num2 = add i32 2, 0
5      %num3 = add i32 3, 0
6      %num4 = add i32 4, 0
7
8      ; Add up the integers
9      %sum = add i32 %num1, %num2
10     %sum = add i32 %sum, %num3
11     %sum = add i32 %sum, %num4
12
13     ; Print the sum
14     %format_str = c"Sum: %d\n\00"
15     call i32 (i8*, ...) @printf(i8* %format_str, i32 %sum)
16
17     ret i32 0
18 }
19
```


아이디어 4

Prefer mul/div over shifting

`/* Never use shifts */`

Shifting보다는 $*2^n$, $/2^n$ 을 이용

Shifting cost = 4

Mul / div cost = 1

Const 2^n 을 곱하거나 나누는 방식이 더 효율적

Kind	Name	Cost
Integer Multiplication/Division	<code><reg> = udiv <val1> <val2> <bw></code> <code><reg> = sdiv <val1> <val2> <bw></code> <code><reg> = urem <val1> <val2> <bw></code> <code><reg> = srem <val1> <val2> <bw></code> <code><reg> = mul <val1> <val2> <bw></code> <code><bw> := 1 8 16 32 64</code>	1
Integer Shift/Logical Operations - shl: shift-left - lshr: logical shift-right - ashr: arithmetic shift-right	<code><reg> = shl <val1> <val2> <bw></code> <code><reg> = lshr <val1> <val2> <bw></code> <code><reg> = ashr <val1> <val2> <bw></code> <code><reg> = and <val1> <val2> <bw></code> <code><reg> = or <val1> <val2> <bw></code> <code><reg> = xor <val1> <val2> <bw></code> <code><bw> := 1 8 16 32 64</code>	4

example

```
1
2  %after = shl i32 %before, 2
3  %after = shr i32 %before, 2
4
```

```
1
2  %after = mul i32 %before, 4
3  %after = udiv i32 %before, 4
4
```

아이디어 5

Aload when possible

/* Load is expensive, aload is cheap */

착안점

Load의 cost는 20/30, async load의 cost는 24/34이지만, waiting을 최소화하면 cost를 1로 줄일 수 있다.

Kind	Syntax	Base Cost
Load	<code><reg> = load <size> <ptr></code> <code><size> := 1 2 4 8</code>	Stack area: 20 Heap area: 30 Cost reduced by 90% inside oracle
Store	<code>store <size> <val> <ptr></code> <code><size> := 1 2 4 8</code>	Stack area: 20 Heap area: 30 Cost reduced by 90% inside oracle
Async Load	<code><reg> = aload <size> <ptr></code> <code><size> := 1 2 4 8</code>	Stack area: 1 Heap area: 1 Cost to resolve Stack area: 24 Heap area: 34 Cannot use inside oracle

example

예시 코드

```
r2 = aload 4 r1
r3 = icmp eq r2 5
r12 = add r10 r11 32
r13 = add r11 r12 32
r11 = sub r11 r10 32
r11 = mul r10 r11 32
r10 = incr r10
br r3 true_bb false_bb
```

Async load의 cost: 24 -> 24-5-5-5-1-1

```
r2 = aload 4 r1
r12 = add r10 r11 32
r13 = add r11 r12 32
r11 = sub r11 r10 32
r11 = mul r10 r11 32
r10 = incr r10
r3 = icmp eq r2 5 ; deferred
br r3 true_bb false_bb
```

아이디어 6

Use the Oracle

`/* The Oracle makes load/stores cheap */`

Oracle 함수의 특징은

- load/store가 저렴하며
- (함수이므로) 리턴값이 하나인 것이다

Condition

- 여러번의 load/store 끝에 결과값이 store되거나, 하나의 i64값으로 줄어드는 경우

Optimization

- Oracle 내부로 로직을 이동시킨다. (Function inlining의 반대 느낌)

필요 시, oracle들을 chaining하여, store/load를 통해 workload를 이어받을 수 있다.

(1-1) Oracle Function

Syntax:

```
start oracle <Narg>:  
... (basic blocks)  
end oracle
```

- A function named oracle is treated specially.
- Unlike other functions, call oracle always costs 40 regardless of the number of arguments
- In this function, all load / store cost only 10% of its original cost (reduced by 90%)
- The interpreter will crash if call is used inside oracle
- The interpreter will crash if aload is used inside oracle
- The compiler will crash if oracle contains more than 50 LLVM IR instructions (excluding basic block labels)

example

```
1  define i32 @sum4(i32* %ptr1, i32* %ptr2, i32* %ptr3, i32* %ptr4) {
2      ; Load values from pointers
3      %val1 = load i32, i32* %ptr1
4      %val2 = load i32, i32* %ptr2
5      %val3 = load i32, i32* %ptr3
6      %val4 = load i32, i32* %ptr4
7
8      ; Calculate sum
9      %sum = add i32 %val1, %val2
10     %sum = add i32 %sum, %val3
11     %sum = add i32 %sum, %val4
12
13     ; Return sum
14     ret i32 %sum
15 }
16
```

```
1  define i32 @sum4(i32* %ptr1, i32* %ptr2, i32* %ptr3, i32* %ptr4) {
2
3      ; Chain oracles together
4      %result = call i64 @oracle(...) ; save temp results to memory
5      %result = call i64 @oracle(...) ; load from memory, and continue
6      %result = call i64 @oracle(...) ; and so on
7      %result = call i64 @oracle(...)
8
9      ; Return sum
10     ret i32 %result
11 }
12
13
```

아이디어 7

Optimize in code level

/* Basic optimizations: CSE, DCE */

Common Subexpression Elimination

이전에 연산했던 **expression** 이 있다면 재사용한다.

Dead Code Elimination

Reachable 하지 않은 영역을 제거한다.

example

```
4  define i32 @example_function(i32 %a, i32 %b) {
5      %c = add i32 %a, %b          ; Addition operation
6      %d = mul i32 %c, 3           ; Multiplication operation
7      %e = add i32 %c, %d          ; Redundant addition operation
8      %f = mul i32 %e, %b          ; Unused multiplication operation
9      ret i32 %f                   ; Return value
10 }
11
12 define i32 @example_function(i32 %a, i32 %b) {
13     %c = add i32 %a, %b          ; Addition operation
14     %d = mul i32 %c, 3           ; Multiplication operation
15     %f = mul i32 %c, %b          ; Optimized multiplication operation
16     ret i32 %f                   ; Return value
17 }
```


Other ideas...

Loop unrolling

Function inlining

And other general optimization ideas are all possible

<https://github.com/llvm/llvm-project/tree/main/llvm/lib/Transforms>

