## In21-S4-CS3513- Programming Languages

# **Programming Project 01**

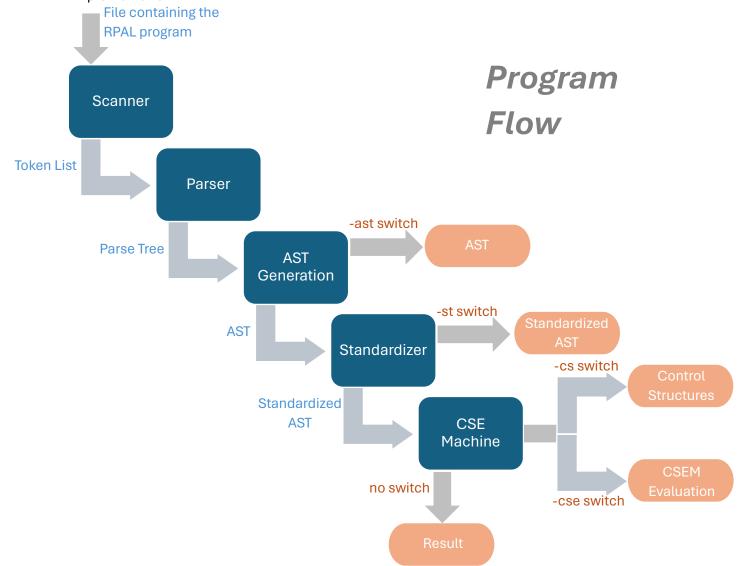
#### **GROUP 50**

W.A.R.T. FONSEKA – 210170G K.S. RANASINGHE – 210518H

#### **Project description**

The program should read an input file which contains a RPAL program. Then the lexical analyser tokenizes the program. The parser then builds a parse tree using the token list. Then we build a Abstract Syntax Tree using the parse tree. Then we standardize the AST. Finally, the CSE machine flattens the standardized AST and evaluate the program and output the result. The complete project is written in python.

As mentioned above there are multiple components in the program which we have implemented separately for ease of use. We added some additional functionality as improvements to program as well. Diagram given below gives a high-level definition for the implementation.



Next let's analyse the individual components. For this we will go through the individual python files as the components were initialized in separate files.

#### Lexical Analyser: Scanner.py

In this we have the tokenize(program) function which breaks the program into a token list. The different lexical rules are defined as shown below.

```
IDENTIFIER = '<IDENTIFIER>'
OPERATOR = '<OPERATOR>'
STRING = '<STRING>'
PUNCTUATION = ['(', ')', ';', ',']
DELETE = '<DELETE>'
SPECIAL = ['let', 'in', 'fn', 'where', 'aug', 'or', 'and', 'not', 'eq', 'ne', 'rec', 'within',
            'true', 'false', 'nil', 'dummy', 'gr', 'ge', 'ls', 'le']
# Regular expressions for lexicon rules
identifier_pattern = r'[a-zA-Z_][a-zA-Z0-9_]*'
integer_pattern = r'\d+'
operator_pattern = r'[+\-*<>&.@/:=[|$\#!%^_[\]{}"?]'
string_pattern = r"'(?:\\[tn\\';,() ]|\\['"?]|[a-zA-ZO-9_+\-*<>&.@/:=\[|$\#!%^_[\]{}'"])+'"
spaces_pattern = r'[ \t]+'
comment pattern = r'//.*
newline_pattern = r'\n'
# Special patterns given in the grammer
special_pattern = r'->'
special_pattern1 = r'\*\*'
```

Next for each combination of characters in the program we need to identify the different token types to assign. There is also a precedence for token types to avoid any identification issues as given below.

```
# Matching patterns
for pattern, token_type in [
    (identifier_pattern, IDENTIFIER),
    (integer_pattern, INTEGER),
    (comment_pattern, DELETE),
    (special_pattern, OPERATOR),
    (special_pattern1, OPERATOR),
    (special_pattern1, OPERATOR),
    (operator_pattern, OPERATOR),
    (string_pattern, STRING),
    (spaces_pattern, DELETE),
    (newline_pattern, DELETE),
];
# Matching pattern,

# Comment is given higher precedence than Operator for it to be deleted
# Precendence of special_pattern and special_pattern1 is higher than operator_pattern
# Same as above is done here. Otherwise the two characters will be read as 2 tokens

(string_pattern, DELETE),
    (newline_pattern, DELETE),

| Comment is given higher precedence than Operator for it to be deleted
# Precendence of special_pattern and special_pattern1 is higher than operator_pattern
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
# Same as above is done here. Otherwise the two characters will be read as 2 tokens
```

Based on whatever the token type identified we add them to the token list. We handle spacing and new lines also here. Also, if any incorrect character combination would output an error.

```
pos += 1
elif char == '\n':
    tokens.append(('NEWLINE', '\n'))
    pos += 1
else:
    raise SyntaxError(f"Invalid character '{char}' at position {pos}")
else:
```

Finally, we have the lex(filename) which reads the file and input the program to the tokenize function and return the token list.

```
# Main function reads the input file and tokenize and return the token list
def lex(filename):
    try:
        with open(filename, 'r') as file:
            program = file.read()
    except FileNotFoundError:
        print(f"File '{filename}' not found.")
        return []

# Tokenize the program
tokens = tokenize(program)
```

#### Parser: Parser1.py

Here, we perform syntax analysis to catch any syntax errors and throw an error message if any are found. Our main function, parse(TokenList), takes the output token list from the scanner, calls the Read(Token) function for initialization itself, and invokes the E(), starting function of the recursive descent parser. And finally, it's returns the parser tree to next step.

```
# Main function parses the token list and returns the parsed tree
def parse(TokenList):
    global tokens
    tokens = TokenList
    Read("")
    E()
    return Parse_tree
```

The Read(Token) function is responsible for verifying that the upcoming token is syntactically correct and consuming it. If the token is syntactically incorrect, it tries to provide a comprehensive error message, detailing what the error is, what the expected token is, and where the error occurs.

The Build\_Tree(node\_type, num\_children) function is responsible for building the parser tree in postorder and appending it to a list.

```
def Build_Tree(node_type, num_children):
    # Check if the node type is an identifier, integer, or string
    if node_type in ["<IDENTIFIER>", "<INTEGER>", "<STRING>"]:
        # If the node type is an identifier, construct the corresponding token
        if node_type == "<IDENTIFIER>";
            Token = "<ID:" + Curr_Token + ">"
            # If the node type is an integer, convert the current token to string and construct the token
        if node_type == "<INTEGER>":
            Token = "<INT:" + str(Curr_Token) + ">"
            # If the node type is a string, construct the corresponding token
        if node_type == "<STRING>":
            Token = "<STR:" + Curr_Token + ">"
            # Append the constructed token to the Parse tree along with the number of children
            Parse_tree.append((Token, num_children))
        else:
            # If the node type is not an identifier, integer, or string, append it to the Parse tree along with the number of children
            Parse_tree.append((node_type, num_children))
```

We have separate functions for each grammar rule, and within these functions, we call the Read(Token) function and the Build\_Tree(node\_type, num\_children) function in appropriate places. Some example functions are as follows. Some of these functions have error-handling methods to manage syntax errors that cannot be handled by the Read(Token) function itself.

```
def E():
   if Next_Token_Type == "let":
       Read("let")
       D()
       Read("in")
       E()
       Build_Tree("let", 2)
   elif Next_Token_Type == "fn":
       Read("fn")
       Vb()
       while Next_Token_Type in ["<IDENTIFIER>", "("]:
           Vb()
       Read(".")
       E()
       Build_Tree("lambda", n + 1)
       Ew()
```

```
def R():
    Rn() # R -> Rn
    # R ->R Rn => 'gamma'
    while Next_Token_Type in [
        "<IDENTIFIER>",
        "<STRING>",
        "true",
        "false",
        "nil",
        "(",
        "dummy",
    ]:
    Rn()
    Build_Tree("gamma", 2)
```

```
def V1():
    # V1 -> '<IDENTIFIER>' list ',' => ','?
    if Next_Token_Type == "<IDENTIFIER>":
        Read("<IDENTIFIER>")
        Build_Tree("<IDENTIFIER>", 0)
        n = 0
        while Next_Token_Type == ",":
            Read(",")
            Read("<IDENTIFIER>")
            Build_Tree("<IDENTIFIER>", 0)
            n += 1
        if n > 0:
            Build_Tree(",", n + 1)
        else:
        # handle error
        print("Error:error occurs near ", Next_Token, " .IDENTIFIER expected.")
```

## **AST Generation : AST.py**

In here we have the function <code>Build\_Preorder\_Tree(AST)</code> which takes the parse tree as the input and convert it to Abstract Syntax Tree. In here we are generating the AST structure that is given in the project document. We start by reversing input parse tree list and iterating the list while creating a new list with the AST structure as shown below.

```
for node in Postorder AST:
    num child = node[1]
   node val = node[0]
    # Add the node value with appropriate indentation to the preorder AST
    Preorder_AST[i] = "." * depth + node_val
   depth += 1
   # Insert None values for the children of the current node
    for j in range(num child):
        i += 1
        Preorder AST.insert(i, None)
    # Adjust depth and index if the current node has no children
    if num child == 0:
        depth -= 1
        i -= 1
        while Preorder_AST[i] != None and i > -1:
           depth -= 1
           i -= 1
return Preorder AST
```

This function returns a list containing the AST structure. Which we will be output using the Print\_AST(Preorder\_AST) if the -ast switch is provided.

```
# This function prints the AST using the parse tree
def Print_AST(Preorder_AST):
    for node in Preorder_AST:
        print(node)
```

## **AST Standardization: ST.py**

In the previous step(AST generation) we did not create a tree structure using the AST. In here we first create the tree structure using the build tree(ast list) function as given below.

```
class ASTNode:
    def __init__(self, type, parent=None):
        self.type = type
        self.children = []
        self.parent = parent

def add_child(self, child):
        child.parent = self
        self.children.append(child)

def remove_last_child(self):
        if self.children:
            return self.children.pop()
        else:
            return None
```

```
def build_tree(ast_list):
    root = ASTNode(ast_list[0])  # Assuming the first element is the root node
    stack = [root]  # Stack to keep track of parent nodes

for item in ast_list[1:]:
    level = item.count('.')
    node_type = item[level:]

    node = ASTNode(node_type)

# Pop nodes from stack until the proper parent is found
    while level <= len(stack) - 1:
        stack.pop()

parent = stack[-1]
    parent.add_child(node)  # Use the add_child method to set the parent of the node
    stack.append(node)

return root</pre>
```

Next we will iterate through this tree while standardizing the necessary nodes. This is carried out by the standardize(node) function. This function takes the root node of the AST as input. As given by the CSE machine rules we don't have to standardize every node. All the nodes we standardized are given below.

```
if node.type == 'let':
    # let gamma
# / \ / \
# = P -> lambda E
# / \ / \
# X E X P
equal_node = node.children[0]
e = equal_node.children[1]
equal_node.children[1] = node.children[1]
node.children[1] = e
equal_node.type = 'lambda'
node.type = 'gamma'
```

```
elif node.type == 'function form':
   equal node = ASTNode('=')
    equal node.children.append(node.children[0])
    equal node.children.append(ASTNode('lambda'))
    node.children.pop(0)
    last_child = node.remove_last_child()
    curr node = equal node.children[1]
    while node.children:
       curr_node.add_child(node.children[0])
       node.children.pop(0)
        if node.children:
           curr node.add child(ASTNode('lambda'))
           curr node.add child(last child)
        curr node = curr node.children[1]
    node.type = equal node.type
    node.children = equal node.children
```

```
elif node.type == 'where':

# where gamma
# / \ / \
# P = -> lambda E
# / \ / \
# X E X P
equal_node = node.children[1]
node.children[1] = node.children[0]
node.children[0] = equal_node
e = equal_node.children[1]
equal_node.children[1] = node.children[1]
node.children[1] = e
equal_node.type = 'lambda'
node.type = 'gamma'
```

Altogether we standardized 8 node types as 'let', 'where', 'fcn\_form', 'within', '@', 'lambda', 'and' and 'rec'. All the other node types will be handled by the cse machine as given by the rules.

There is also the print\_tree(root, indent=0) function which would output the standardized AST if the -st switch is provided.

```
def print_tree(root, indent=0):
    print("." * indent + root.type)
    for child in root.children:
        print_tree(child, indent + 1)
```

## **CSE Machine: ControlStructure.py, CSEM.py**

We are using two python scripts for the CSE machine process. ControlStructure.py generates the control structures for the standardized AST and the CSEM.py evaluates the control structures to generate the result.

#### **Control Structure Generation: Control Structure.py**

The standardized tree created in the ST.py is first given as input to the function given below. preorder\_traversal(node, control\_stack, lambda\_stacks)

This function flattens the tree into a list containing the control structures. control\_stack contains the starting structure while the rest of structures are in lambda\_stacks. When flattening some nodes need to be handled explicitly as shown below.

#### Lambda node

```
# if we come across a lambda node we add an identifier(number) and create a seperate control
# structure for it's children

if node.type == 'lambda':
    node.type = 'lambda' + str(i)
    i += 1
    control_stack.append(node)
    lambda_stack = []
    preorder_traversal(node, lambda_stack, lambda_stacks)
    lambda_stacks.append(lambda_stack)
```

#### Beta node

```
# if we have a conditional we replace it with a beta symbol then add two delta nodes for
# the conditions and create seperate structures for those delta nodes. The naming system
#for both delta and lambda nodes act as a pointer system to their control structures.

elif node.type == '->':
    node.type = 'β'
    control_stack.append(ASTNode('delta' + str(j)))
    control_stack.append(ASTNode('delta' + str(j+1)))
    control_stack.append(node)
    preorder_traversal(node.children[0], control_stack, lambda_stacks)

lambda_stack = []
    k = j
        j +=2
    lambda_stack.append(ASTNode('delta' + str(k)))
    preorder_traversal(node.children[1], lambda_stack, lambda_stacks)

lambda_stack = []
    lambda_stack.append(ASTNode('delta' + str(k+1)))
    preorder_traversal(node.children[2], lambda_stack, lambda_stacks)
lambda_stacks.append(lambda_stack)
```

#### Tau node

#### Token type

```
# we have to remove token types from the names here as they are of no use here on.
else:
    if node.type.startswith('<ID:'):
        node.type = node.type[4:-1]
    elif node.type = node.type[5:-1]
    elif node.type = startswith('<STR:'):
        if node.type == "<STR:''>":
            node.type = "''"
        else:
            node.type = node.type[5:-1]
        control_stack.append(node)
        for child in node.children:
            preorder_traversal(child, control_stack, lambda_stacks)
```

The above function is called inside the main function <code>create\_control\_structure(ast\_list)</code>. This function first runs the flattening function and gets the lists containing the control structures and further modify them to get the exact nature of the required control structures. First we will remove the tree structure from all the elements for ease of use.

Next, we will use a dictionary to capture the relationship between the lambda nodes and their respective variables.

```
# this is a dictionary with the lambda node set as key and the variable set as value
# λ1[x] -> lambda_values[λ1] = x
lambda_values = {}

# Iterate over each lambda stack
for stack in l_stack:
    if stack[0].startswith('lambda'):
        lambda_values[stack[0]] = stack[1] # Store the lambda node and its value
```

Since the lambda nodes are stored in a separate dictionary, we can remove them, and the variables associated with them from the lists we are working on. Also, we can set the exact nature we want the lambda nodes to look in the control structure.

```
# we can remove lambda and the variable from the control structures now as it is stored
# in the previously created dictionary
for stack in l_stack:
    if stack[0].startswith('lambda'):
        stack.pop(0)
        stack.pop(0)

# this sets the way we want the lambda node to look like -> \( \lambda \)_1['x']
for index, value in enumerate(c_stack):
    if value in lambda_values:
        c_stack[index] = f"{value}['{lambda_values[value]}']"
```

Now we can build the dictionary with the control structures used in the CSE machine as all the necessary modifications have been made.

```
# now we will create the main dictionary for storing all the control structures
result_dict = {'\delta_0': c_stack.copy()}

for i, stack in enumerate(l_stack):
    result_dict[f'\delta_{i+1}'] = stack.copy()
```

Also made some changes to the Greek symbols used in the control structure for readability.

Finally, we had to handle the case of delta nodes. We need to point to the correct control structures when handling conditionals.

There is also another function print\_dict(result\_dict) which outputs the control structures created by the above function.

```
# this function prints the dictionary containing the control structures

def print_dict(result_dict):
    for key, value in result_dict.items():
        print(key, end=' : ')
        for i in value:
            print(i, end=' ')
        print()
```

#### **Control Structure Generation : CSEM.py**

The dictionary containing control structures is evaluated in here by the function given here evaluate(control, stack, result\_dict, env, out=True)

We have to input both the control and stack used in the evaluation as well. The 'out' parameter is used to determine whether the evaluation process needs to be shown as output. If it is set to true, we output the evaluation of control and stack.

In this function we will be handling all the CSE rules as mentioned in the notes. Given below are some of the implementations for the rules.

```
# we will be evaluating the control until it becomes empty
while control:
    # RULE 1 - get the top element from the control
    val = control.pop()

# RULE 3 - handling stack incase of a gamma node
    if val == '\vec{y}':
        # we take the top value from the stack for evaluation
        rator = stack.pop(0)
        # if rator is a list we need to pop the top value(n) form stack again
        # and get the nth instance from the rator list and add to the stack
        if isinstance(rator, list):
            i = int(stack.pop(0))
            stack.insert(0,rator[i-1])
            continue
```

```
# RULE 12 - defining the functionality of Ystar node
elif rator == 'Ystar':
    var = stack.pop(0)
    eta = 'ŋ' + var[1:]
    stack.insert(0,eta)

# RULE 13 - defining the functionality of eta node
elif rator.startswith('ŋ'):
    stack.insert(0,rator)
    lamda = '\lambda' + rator[1:]
    stack.insert(0,lamda)
    control.append('\lambda')
    control.append('\lambda')
```

```
if rator.startswith('λ_'):
   num = rator[2]
   var = rator[5:-2]
   rep = stack.pop(0)
   var_list = []
   if ',' in var:
       if var.startswith('['):
           var = var[2:-2]
       var_list = var.split("', '")
       for i in range(0,len(var_list)):
           env[var_list[i]] = rep[i]
       env[var] = rep
       for i in range(0,len(control)):
           if control[i] == var:
               control[i] = rep
   control.append('e' + num)
   stack.insert(0, 'e' + num)
   for item in result_dict['\delta_' + num]:
       control.append(item)
```

```
# RULE 8 - handling conditional
elif val == '\beta':
    rand = stack.pop(0)
    if rand == 'true':
        control.pop()  # this holds the false case
        var = control.pop()  # this holds the true case
else:
        var = control.pop()  # this holds the false case
        control.pop()  # this holds the true case

# based on the case we choose we find the control structure relevant
# to it and add it to the control
for item in result_dict[var]:
        control.append(item)
```

We also had to handle the built-in functions of RPAL language. Finding the definitions of some of these functions were difficult as they had varying definitions in different sources. So, we used the implementation as shown by the RPAL interpreter given to us.

```
# We will handle the built in functions of RPAL here
elif rator == 'Order':
    val = stack.pop(0)
    if isinstance(val, list):
        num = len(val)
        stack.insert(0,str(num))

elif rator == 'Print':
    val = stack.pop(0)
    # the output format is not the same as what we have in the control structures.
    # they are hanlded here
    formatted_list = str(val).replace('[', '(').replace(']', ')').replace("'", '').replace("'", '')
    print(formatted_list)
```

```
elif rator == 'Isinteger':
    val = stack.pop(0)
    try:
        int(val)
        stack.insert(0,'true')
        stack.insert(0,'false')
elif rator == 'Istruthvalue':
    val = stack.pop(0)
        stack.insert(0,'true')
        stack.insert(0,'false')
elif rator == 'Isstring':
    val = stack.pop(0)
    try:
        int(val)
        stack.insert(0,'false')
    except ValueError:
        stack.insert(0,'true')
elif rator == 'Istuple':
    val = stack.pop(0)
    if isinstance(val, list):
        stack.insert(0,'true')
        stack.insert(0,'false')
elif rator == 'Isfunction':
    val = stack.pop(0)
    if var.startswith('λ'):
        stack.insert(0,'true')
        stack.insert(0,'false')
elif rator == 'Isdummy':
    val = stack.pop(0)
    if val == 'dummy':
        stack.insert(0,'true')
    else:
        stack.insert(0,'false')
```

```
# The definitions Stem and Stern are different in different sources.
# we applied the same functionality shown in the RPAL interpreter provided
# Stern removes the first element from a string
elif rator == 'Stern':
    val = stack.pop(0)
    val = val[2:]
    val = "'" + val
    stack.insert(0,val)

# Stem returns the first element of a string
elif rator == 'Stem':
    val = stack.pop(0)
    s = val[:2] + "'"
    stack.insert(0,s)

# Conc concatenates two strings
elif rator == 'Conc':
    val1 = stack.pop(0)
    val2 = stack.pop(0)
    control.pop()
    res = val1[:-1] + val2[1:]
    stack.insert(0,res)
```

```
# this is also a built in function of RPAL
# using this we can add elements to tuples
elif val.startswith('aug'):
    var = stack.pop(0)
    if isinstance(var, list):
        var.append(stack.pop(0))
        stack.insert(0,var)
    # if there is no tuple we create a tuple with only one element
    else:
        tup = [stack.pop(0)]
        stack.insert(0,tup)
```

Finally, we have the function print\_control\_stack(control, stack) which output the current state control and the stack. This function is called inside the evaluate function for each iteration to capture the whole evaluation process.

```
# this method prints the control and the stack in operation as mentioned in the lecture notes.

def print_control_stack(control, stack):
    # Get the width of the terminal window
    terminal_width, _ = shutil.get_terminal_size()

# Calculate the maximum length of the content (to determine the padding)
    max_length = max(len(' '.join(control)), len(' '.join(str(item) for item in stack)))

# Calculate the total length of the concatenated strings
    total_length = max_length * 2 + 1  # Total length includes the space between the lists

# Calculate the number of spaces needed to fill the remaining width
    num_spaces = int(max(terminal_width - total_length, 0) * 0.5)  # Ensure it's non-negative

# Construct the formatted string with the two lists separated by spaces
    formatted_string = ' '.join(control).ljust(max_length) + ' ' * num_spaces
    + '|' + ' ' * num_spaces + ' '.join(str(item) for item in stack).rjust(max_length)

print(formatted_string)
    print(terminal_width * '-')
```

# Main Program: myrpal.py

This is the base program which executes all the other programs. We first import all the necessary functions from the respective files.

```
# Importing lex function from scanner.py
from scanner import lex

# Importing parse function from parser1.py
# Had to name this file parser1.py because there exists a system file named parser.py
from parser1 import parse

# Importing Print_AST function from AST.py
from AST import Print_AST, Build_Preorder_Tree

# Importing standardisation function from ST.py
from ST import standardize, build_tree, print_tree

# Importing create_control_structure function from ControlStructure.py
from ControlStructure import create_control_structure, print_dict

# Importing create_control_structure function from ControlStructure.py
from CSEM import evaluate
```

Next we will simply execute these functions one by one according to the flow that was stated earlier

```
def main():
    try:
        if len(sys.argv) not in [2,3]:
            print("Usage: python scanner.py <filename> [-ast]")
        # Read filename from command-line argument
        filename = sys.argv[1]
        tokens = lex(filename)
        # Call the parse function to generate the Parse tree
        Parse tree = parse(tokens)
        # Next we will build the AST using the parse tree
        AST = Build_Preorder_Tree(Parse_tree)
        # We will standardize the generated AST using this function
        ST = standardize(build_tree(AST))
        # This will generate the control structures for CSE machine evaluation
        CS = create_control_structure(AST)
        control = []
        stack = []
        control.append('€0')
        stack.append('00')
        for val in CS['\delta_0']:
            control.append(val)
        env = \{\}
        if len(sys.argv) == 3:
            # if -ast switch is given we will print the AST
            if sys.argv[2] == "-ast":
                Print_AST(AST)
            if sys.argv[2] == "-st":
                print_tree(ST)
            # if -cs switch is given we will print the control structures for the
program
            if sys.argv[2] == "-cs":
                print_dict(CS)
            # if -cse switch is given we will print the CSE machine evaluation of
the control structures
            if sys.argv[2] == "-cse":
                evaluate(control,stack, CS, env)
        # if no switch is given we will just output the result of the program
        else:
            evaluate(control,stack, CS, env, False)
    except:
        print()
```