

Q1 : PING command :

- a) To specify the number of echo requests : `ping -c count <host_name>`
b) To set time interval btw 2 successive pings : `ping -i interval <host_name>`
* *min value of interval : 0.2 sec for non super-users*
* *default value of interval : 1 sec for all users*
c) To send packets without waiting for reply : `ping -l count <host_name>`
* *max value of count : 3 for non super-users*
d) To specify packet size : `ping -s size <host_name>`
* *Additional 8 bytes for ICMP header and 20 bytes for IP header*
* *For packet size = 32 bytes, total packet size = 32 + 28 = 60 bytes*

Q2 : PING - RTT experiment :

- * All readings were taken via <http://www.spfld.com/ping.html> (server location: New Jersey)

```
PING stackoverflow.com (151.101.65.69) 56(84) bytes of data.  
  
--- stackoverflow.com ping statistics ---  
30 packets transmitted, 30 received, 0% packet loss, time  
29016ms  
rtt min/avg/max/mdev = 4.754/4.887/5.305/0.111 ms
```

```
--- wikipedia.org ping statistics ---  
30 packets transmitted, 0 received, 100% packet loss, time 29001ms
```

PACKET LOSS :

There was no case of packet loss greater than 0% for packet size 64 bytes. However, packet losses might occur due to network congestion. Sometimes, 100% packet loss is also witnessed, as hosts block ICMP packets (low-priority) to reduce traffic.

CORRELATION B/W RTT AND PACKET SIZE :

- * From the chart below, the following observations can be made :
- For smaller packet sizes (less than 2048 bytes), RTT is more or less a stable value.
 - For packet size of 2048, only 2 out of 6 hosts received any packets. Hence, only imdb, amazon map the RTT corresponding to 2048 byte packet size.
- * These observations can be explained from the fact that MTU (Maximum Transmission Unit, i.e., the maximum number of octets that the network interface can handle) has a default value of 1500, Therefore, packets of size greater than 1500 are either rejected or broken into smaller packets resulting into a higher RTT.

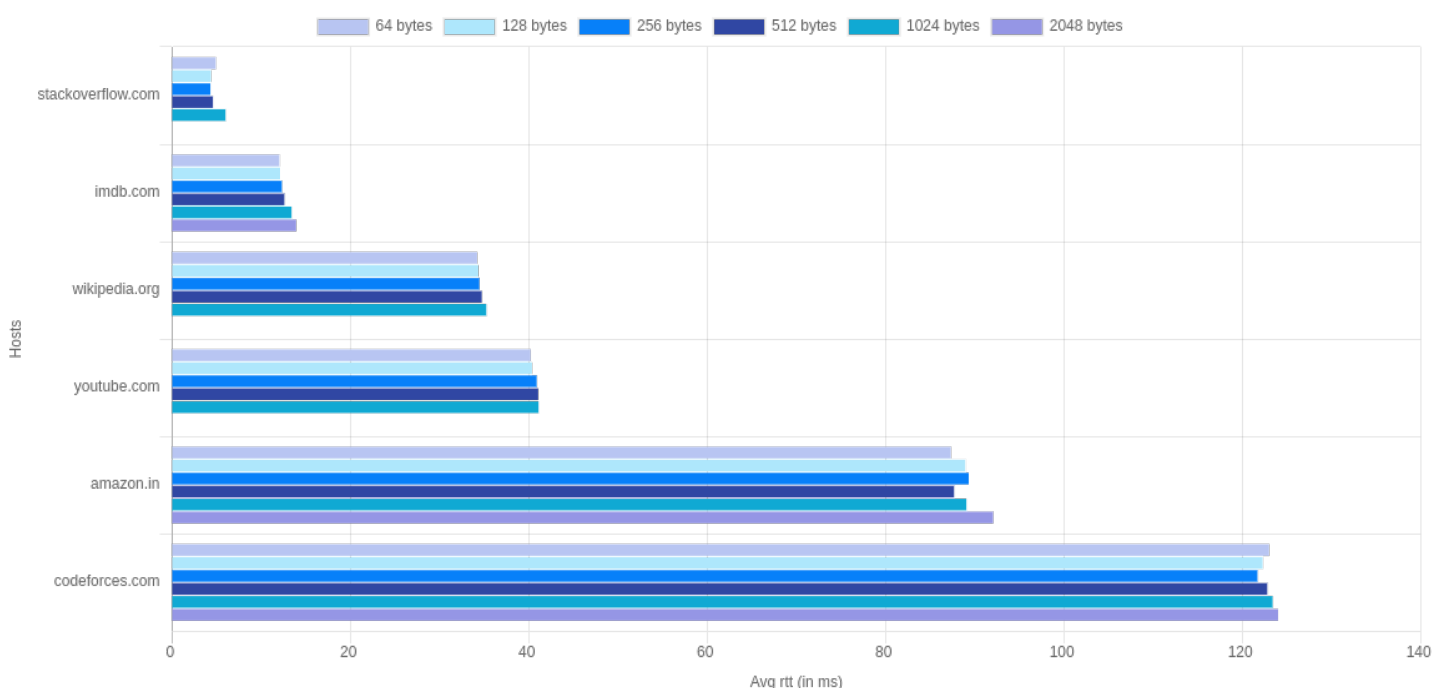


Illustration 1: CORRELATION BETWEEN RTT AND PACKET SIZE

STRONG CORRELATION B/W RTT AND GEOGRAPHICAL DISTANCE :

- * The *ROUND TRIP TIME (RTT)* is shows "*strong, positive linear relation*" with distance bcs :
 - a) Increase in propagation delay – This occurs bcs of increased distance
 - b) Increase in processing and queuing delay – This occurs bcs of more no. of hops
- * This can be seen in the table below. RTT increases with distance, especially in case of inter continental pings (Russia, Ireland from USA).
- * **DEVIATIONS** : Significant deviations from linear relation arise due to large-scale routing behaviours, like the noise introduced by the procedure.

HOST_NAME	IP_ADDRESS	LOCATION	RTT-3:15PM	RTT-7:30PM	RTT-10:30PM	AVG_RTT
stackoverflow.com	151.101.65.69	San Francisco, USA	4.887 ms	3.835 ms	4.818 ms	4.513 ms
imdb.com	52.94.225.248	Virginia, USA	12.019 ms	11.739 ms	11.674 ms	11.810 ms
wikipedia.org	208.80.154.224	New York, USA	34.197 ms	34.174 ms	33.510 ms	33.960 ms
youtube.com	64.233.177.91	California, USA	40.194 ms	40.540 ms	40.705 ms	40.479 ms
amazon.in	52.95.116.115	Dublin, Ireland	87.318 ms	86.243 ms	87.009 ms	86.856 ms
codeforces.com	81.27.240.126	Moscow, Russia	123.006 ms	121.192 ms	121.803 ms	122 ms

CORRELATION B/W RTT AND TIME OF THE DAY :

- * From the data in the above table, it can be observed that RTT varies with time of the day. This is because higher congestion leads to more ROUND TRIP TIME. Hence it can said that the network traffic is less at 7:30 PM IST as compared to 3:15 PM and 10:30 PM IST.

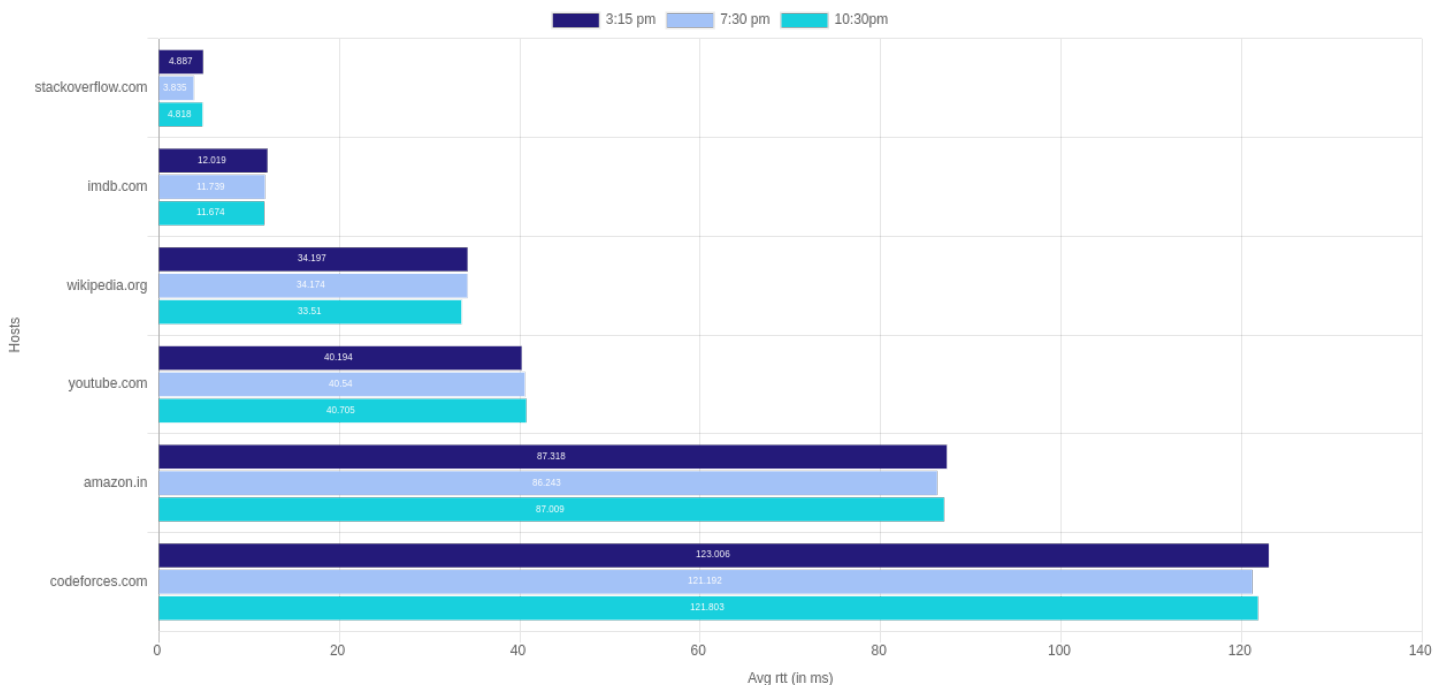


Illustration 2 : CORRELATION BTW RTT ,TIME AND DISTANCE

Q3 : PING -n and PING -p :

Command	Packets Sent	Packets Received	Packet Loss rate	Min Latency	Max Latency	Mean Latency	Median Latency
ping -n -c 1000 172.16.112.12	1,000	1,000	0%	0.246	16	0.595	0.5
ping -p ff00 -c 1000 172.16.112.12	1,000	999	0.1%	0.258	16.1	0.617	0.504

- * Both the commands are sending 1000 packets to the same host. Hence, they are very similar, except for two aspects which are as follows :
 - a) The **-n (numeric only)** option is expected to be **faster**, as it displays the raw ip addresses without looking for more human friendly host names (like google.com)
 - b) The **-p ff00 (pattern)** option fills the packet with **16 pad bits**. This is used for dianosing data-

dependent problems, like **synchronisation problem** of clocks due to only one transition from 1 to 0 in 1111111100000000 (#00).

- * Due to these aspects, the following observations occur in the mapping of <frequency,latency> :
- The mean latency of -n command < the mean latency of -p command ... (1st aspect)
 - Since synch is data (pattern) dependent, there are higher chances of packet losses, which is observed as shown in the table above. ... (2nd aspect)

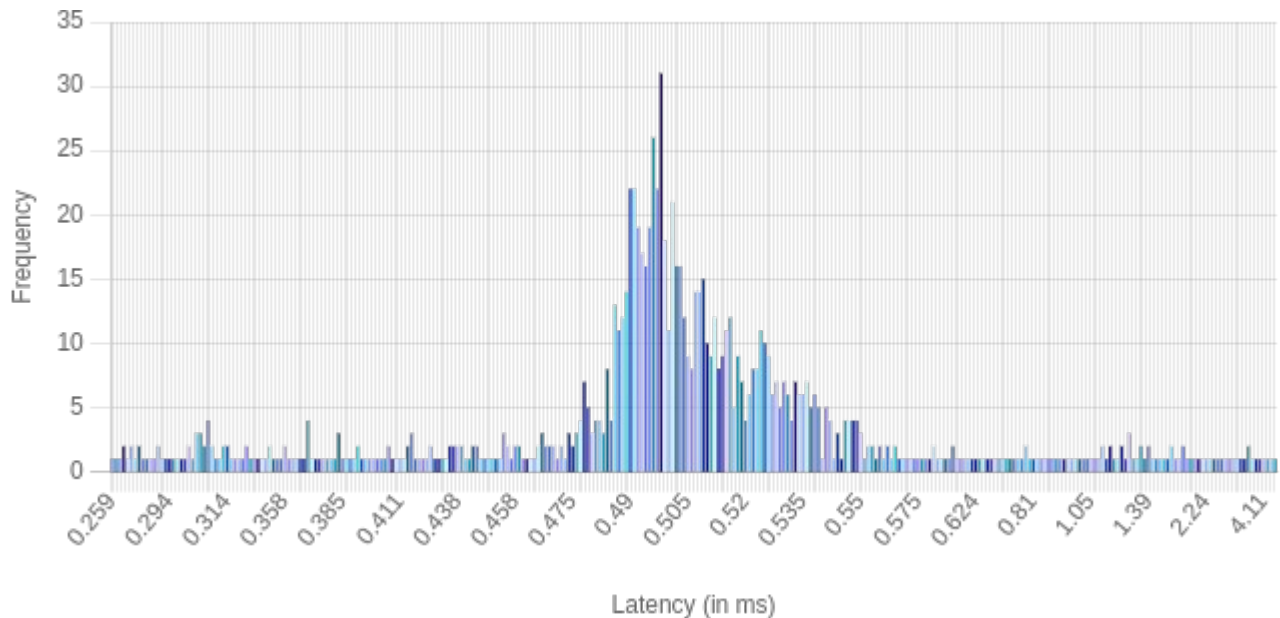


Illustration 3 : Normal Distribution for ping -n c 1000 172.16.112.12

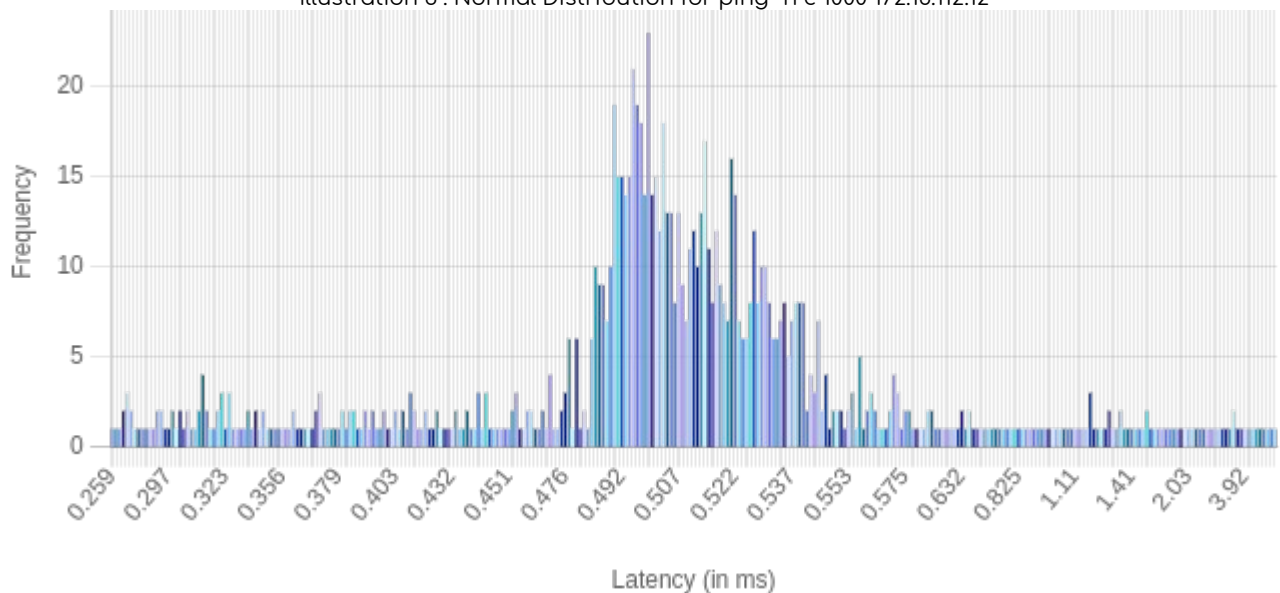


Illustration 4 : Normal Distribution for ping -p ff00 -c 1000 172.16.112.12

Q4 : IFCONFIG and ROUTE command :

```

rashl@rashls-asptre-e5-575g:~$ ifconfig
enp4s0f1: flags=4099<UP,BROADCAST,MULTICAST> mtu 1500
    ether a8:1e:84:55:ab:57 txqueuelen 1000 (Ethernet)
    RX packets 0 bytes 0 (0.0 B)
    RX errors 0 dropped 0 overruns 0 frame 0
    TX packets 0 bytes 0 (0.0 B)
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0

lo: flags=73<UP,LOOPBACK,RUNNING> mtu 65536
    inet 127.0.0.1 netmask 255.0.0.0
    inet6 ::1 prefixlen 128 scopeid 0x10<host>
    loop txqueuelen 1000 (Local Loopback)
    RX packets 402 bytes 34644 (34.6 KB)
    RX errors 0 dropped 0 overruns 0 frame 0
    TX packets 402 bytes 34644 (34.6 KB)
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0

wlp3s0: flags=4163<UP,BROADCAST,RUNNING,MULTICAST> mtu 1500
    inet 10.42.0.161 netmask 255.255.255.0 broadcast 10.42.0.255
    inet6 fe80::ed82:63a2:ba94:42a3 prefixlen 64 scopeid 0x20<link>
    ether 3c:a0:67:6c:14:2d txqueuelen 1000 (Ethernet)
    RX packets 12519 bytes 15344697 (15.3 MB)
    RX errors 0 dropped 0 overruns 0 frame 0
    TX packets 5747 bytes 1354683 (1.3 MB)
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0

```

a) I FCONFIG command :

- * Inter Face **CONFIG**uration
- * Shows/Configures the list of network interfaces that are "UP" (enabled)
- * Normal Invocation :
ifconfig interface (address (parameters))
- * It is used while booting to set up interfaces and while debugging and system tuning.

As observed from the image, it displays the <interface-name> followed by its associated flags, address info and the packet statistics. (RX : received and TX : transmitted)

```

enp4s0f1:
|   |   |
v   |   |
en  |   | -----> ethernet
v   |   |
p4  |   | -----> bus no (4)
v   |   |
s0f1 --> slot no (0f1)

```

INTERFACE NAMING :

<type-of-interface, bus no, slot no>

Similarly, wlp3s0 refers to the following :

```

wl  -   wireless
p3  -   bus no 3
s0  -   slot 3

```

Interface "lo" (as seen in ifconfig result on prev page) however represents **loopback**, which is the computer's ref to itself. It is a virtual network interface to communicate with itself for diagnostics, troubleshooting and to connect to servers running on the local machine.

FLAGS : Represent the current status of the interface, for eg : **<UP, BROADCAST, RUNNING, MULTICAST>**

- a) UP : The interface is accessible to the IP layer, hence it has been assigned an address and routing table.
- b) BROADCAST : This indicates that the interface supports broadcasting.
- c) RUNNING : This indicates that the network driver has been loaded and has initialized the interface.
- d) MULTICAST : MULTICAST is like BROADCAST except that instead of automatically including everybody, the only people who receive packets sent to a multicast address are those programmed to listen to it.

ADDRESS INFO. :

a) inet : Same as **IPv4***

b) inet6 : Same as **IPv6****

```

inet 10.150.37.172 netmask 255.255.248.0 broadcast 10.150.39.255
inet6 fe80::1c9c:a06b:4ec0:504 prefixlen 64 scopeid 0x20<link>
ether 3c:a0:67:6c:14:2d txqueuelen 1000 (Ethernet)

```

**IPv4 : 32 bit numeric IP address*

***IPv6 : 128 bit alphanumeric IP address*

c) netmask : Network Mask for the associated IP address

d) broadcast : 32 bit broadcast address

e) MAC addr (H/w addr) : Address assigned to the LAN card. This addr uniquely defines the device

f) scopeID : This defines the scope of the IP address. For eg, in the img above, IPv6 addr is local.

g) prefixLen : No of bits in the IP address that are to be used as the subnet mask

h) TXqueueLen : Limits the number of packets in the transmission queue in the interface's device driver

i) MTU : Max number of octets the interface is able to handle in one transaction

b) IFCONFIG instructions :

- a) Ifconfig -a : shows all network interfaces including the DOWN ones
- b) ifconfig <interface-name> **UP/DOWN** : enables/disables the mentioned interface
- c) ifconfig <interface-name> **mtu count** : changes MTU value for the mentioned interface
- d) ifconfig <interface-name> **netmask addr** : assigns given netmask addr to the given interface
- ifconfig <interface-name> **broadcast addr** : assigns given broadcast addr to the given interface
- ifconfig <interface-name> **addr** : assigns given addr as IP to the given interface
- e) ifconfig <interface-name> **promisc** : enables promiscuous mode
- ifconfig <interface-name> **-promisc** : disables promiscuous mode

in **promiscuous** mode, the driver does not check whether the packet is meant/not meant for itself and simply **accepts all packets

c) ROUTE command :

Normal invocation : route (-f) (-p) (command (destination) (mask subnetmask) (gateway) (metric))

```
rashi@rashits-aspire-e5-575g:~$ route
```

Kernel IP routing table

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
default	_gateway	0.0.0.0	UG	600	0	0	wlp3s0
10.150.32.0	0.0.0.0	255.255.248.0	U	600	0	0	wlp3s0
link-local	0.0.0.0	255.255.0.0	U	1000	0	0	wlp3s0

Route is used to show, manipulate kernel's IP routing table

Primary use : To set up static routes to specific hosts/networks via interface after configuring with ifconfig

IMPORTANT KEYWORDS :

- a) Destination : The destination network or destination host.
- b) Gateway : The gateway* address or 0.0.0.0, if none set.
- c) Genmask : The netmask** for the destination net
- d) Metric*** : The metric option assigns an integer cost metric (distance to target in hops)
- e) Ref : No of references to the route
- f) Use : The count of lookups for a route
- g) Iface : Interface to which packets for this route will be sent
- h) Flags : **U** (route is up), **H** (target is a host), **G** (use gateway), **C** (cache entry), **I** (reject route), **R** (reinstate route for dynamic routing), **D** (dynamically installed), **M** (modified)

* Gateways regulate traffic btw two dissimilar networks, while routers regulate traffic btw similar ones

** 255.255.255.255 for a host destination and '0.0.0.0' for the default route.

*** Range - [1-9999], used to calculate the fastest, most reliable, and least expensive routes.

d) ROUTE instructions :

```
rashi@rashis-aspire-e5-575g:~$ route -n
Kernel IP routing table
Destination      Gateway         Genmask        Flags Metric Ref    Use Iface
0.0.0.0          10.42.0.1      0.0.0.0        UG    20600  0      0 wlp3s0
10.42.0.0        0.0.0.0        255.255.255.0  U     600    0      0 wlp3s0
```

```
rashi@rashis-aspire-e5-575g:~$ route -e
Kernel IP routing table
Destination      Gateway         Genmask        Flags  MSS Window  irtt Iface
default          strix-15        0.0.0.0        UG      0  0      0 wlp3s0
10.42.0.0        0.0.0.0        255.255.255.0  U      0  0      0 wlp3s0
```

```
root@rashis-aspire-e5-575g:/home/rashi# route add -host 172.16.112.3 reject
root@rashis-aspire-e5-575g:/home/rashi# route -n
Kernel IP routing table
Destination      Gateway         Genmask        Flags Metric Ref    Use Iface
0.0.0.0          10.42.0.1      0.0.0.0        UG    20600  0      0 wlp3s0
10.42.0.0        0.0.0.0        255.255.255.0  U     600    0      0 wlp3s0
172.16.112.3     -              255.255.255.255 !H     0     -      0 -
```

route -n

Show numerical addresses instead of trying to find symbolic host names

route -e

Use netstat(8) - format for displaying the routing table.

route add -host <addr> reject

Rejecting routing to a particular host

**Notice !H in flags*

```
route add -net 127.0.0.0 netmask 255.0.0.0 metric 1024 dev lo
```

Adds the normal loopback entry, using netmask 255.0.0.0 and associated with the "lo" device

```
root@rashis-aspire-e5-575g:/home/rashi# route add -net 127.0.0.0 netmask 255.0.0.0 metric 1024 dev lo
root@rashis-aspire-e5-575g:/home/rashi# route -n
Kernel IP routing table
Destination      Gateway         Genmask        Flags Metric Ref    Use Iface
0.0.0.0          10.42.0.1      0.0.0.0        UG     600    0      0 wlp3s0
10.42.0.0        0.0.0.0        255.255.255.0  U     600    0      0 wlp3s0
127.0.0.0        0.0.0.0        255.0.0.0      U    1024    0      0 lo
```

Q5 : NETSTAT command :

a) **USE :** Netstat is a handy cross-platform (Linux, Windows etc.), which is mainly used for troubleshooting and debugging. It is used to :

- 1) Display incoming and outgoing network connections
- 2) Display routing tables
- 3) Display number of network interfaces
- 4) Display network protocol statistics

To show all established TCP connections :

```
netstat -at | grep
"ESTABLISHED"
```

```
rashi@rashis-aspire-e5-575g:~$ netstat -at | grep "ESTABLISHED"
tcp        0      0 rashis-aspire-e5-:57958 maa05s06-in-f3.1e:https ESTABLISHED
tcp        0      0 rashis-aspire-e5-:35362 maa03s31-in-f14.1:https ESTABLISHED
tcp        0      0 rashis-aspire-e5-:45662 maa03s29-in-f10.1:https ESTABLISHED
tcp        0      0 rashis-aspire-e5-:57388 edge-star-shv-02-:https ESTABLISHED
tcp        0      0 rashis-aspire-e5-:49014 maa05s04-in-f3.1e:https ESTABLISHED
tcp        0      0 rashis-aspire-e5-:40834 172.217.194.155:https ESTABLISHED
```

Description of different fields (in order):

Proto : The protocol (tcp, udp, udpl, raw) used by the socket.

Recv-Q : The # of bytes not copied (waiting to be sent) by the user program connected to socket.

Send-Q : The # of bytes not acknowledged (waiting to be read) by the remote host.

Local Address : <Address, port number> of the local end of the socket.

If the --numeric option is not specified,
the socket address --> canonical hostname
the port number i--> corresponding service name.

Foreign Address : <Address, port number> of the remote end of the socket.

State : This field refers to the state of the socket. It is left **blank** in **UDP/UDPLite** connections.

Some values are :

– **ESTABLISHED** (the socket has an established connection),

TIME_WAIT (The socket is waiting after connection is closed by remote machine to handle packets still in the network),

LISTEN (the socket is listening for incoming connections), etc.

c) netstat -r :

'MSS' : Maximum Segment

Size(size of largest datagram constructed for transmission),

'Window' : default window size(maximum amount of data the system will accept in one burst from a remote host),

'Irtt' indicates Initial Round Trip Time for TCP connections over this route.

Please refer Q4 part d for the remaining fields

```
Kernel IP routing table
Destination      Gateway         Genmask        Flags  MSS Window  irtt Iface
default          _gateway        0.0.0.0        UG      0  0      0 wlp3s0
10.150.32.0      0.0.0.0        255.255.248.0  U      0  0      0 wlp3s0
link-local       0.0.0.0        255.255.0.0    U      0  0      0 wlp3s0
```


d) netstat -i : My machine has 3 interfaces enp4s0f1, lo,wlp3s0. Please refer Interface naming in Q4

```
rashi@rashis-aspire-e5-575g:~$ netstat -i
Kernel Interface table
Iface      MTU      RX-OK RX-ERR RX-DRP RX-OVR      TX-OK TX-ERR TX-DRP TX-OVR Flg
enp4s0f1   1500      0      0      0 0          0      0      0      0 BMU
lo         65536    3676      0      0 0        3676      0      0      0 LRU
wlp3s0     1500   1672981      0      9 0       131134      0      0      0 BMRU
```

e) netstat -su : To show statistics of UDP connections

f) Loopback interface :

This is the computer’s ref to itself. It is a virtual network interface to communicate with itself for diagnostics, troubleshooting and to connect to servers running on the local machine.

The range is 1 – 127.0.0.0/8
Local Host : 127.0.0.1

```
rashi@rashis-aspire-e5-575g:~$ netstat -su
IcmpMsg:
  InType0: 13
  InType3: 49
  InType11: 404
  OutType3: 55
  OutType8: 61
Udp:
  6651 packets received
  30 packets to unknown port received
  0 packet receive errors
  8834 packets sent
  0 receive buffer errors
  0 send buffer errors
  IgnoredMulti: 324197
UdpLite:
IpExt:
  InMcastPkts: 127
  OutMcastPkts: 877
  InBcastPkts: 324492
  OutBcastPkts: 13
  InOctets: 198381975
  OutOctets: 14760039
  InMcastOctets: 13540
  OutMcastOctets: 162472
  InBcastOctets: 34281154
  OutBcastOctets: 795
  InNoECTPkts: 527574
```

Q6 : TRACEROUTE command : Prints the route that a packet takes to reach the host.

Useful when you want to know about the route and about all the **hops that a packet takes**. Seeing the traceroute information can help you determine why your connections to a given server might be poor and can help you **identify problems**. It also shows you how systems are connected to each other, letting you see how your ISP connects to the Internet as well as **how the target system is connected**.

a) Hop counts : Below are the tables showing hop counts for each host

HOSTS	7:30 am	7:30 pm	11:30 pm
youtube.com	11	14	13
wikipedia.org	14**	13	14
amazon.in	17**	14**	15**
imdb.com	17**	20**	28**
stackoverflow.com	13**	19**	25**
codeforces.com	12**	15**	15**

Since only 2 out of 6 hosts from question 2 were reached using trace command, below is a table with some extra hosts who were traced by traceroute

HOSTS	7:30 am	7:30 pm	11:30 pm
iitg.ac.in	2	2	2
google.com	11	11	11
twitter.com	11	10	10

*represents hosts that were not reached in 64 hops

a) Common hops : Just like the images below, the top 7 hops were observed to be common in all the traceroutes, like twitter.com, amazon.in, etc.

```
rashi@rashis-aspire-e5-575g:~$ traceroute youtube.com
traceroute to youtube.com (172.217.166.110), 64 hops max
 1  10.150.32.1  4.337ms  3.192ms  3.929ms
 2  192.168.193.1  1.715ms  1.731ms  1.671ms
 3  14.139.196.17  3.332ms  1.868ms  1.867ms
 4  10.119.254.241  2.315ms  2.170ms  1.979ms
 5  10.177.31.1  39.552ms  38.956ms  38.974ms
 6  10.255.238.205  40.711ms  39.061ms  38.981ms
 7  10.119.73.122  39.736ms  40.578ms  39.620ms
 8  72.14.213.20  43.572ms  43.329ms  43.355ms
 9  74.125.242.145  54.493ms  54.121ms  54.086ms
10  74.125.252.215  53.827ms  53.565ms  53.523ms
11  172.217.166.110  52.642ms  52.539ms  52.470ms
```

```
rashi@rashis-aspire-e5-575g:~$ traceroute google.com
traceroute to google.com (172.217.31.206), 64 hops max
 1  10.150.32.1  5.984ms  1.754ms  1.716ms
 2  192.168.193.1  1.709ms  1.601ms  1.693ms
 3  14.139.196.17  2.033ms  1.893ms  1.730ms
 4  10.119.254.241  19.260ms  2.193ms  2.051ms
 5  10.177.31.21  39.340ms  39.094ms  39.440ms
 6  10.255.238.205  39.863ms  39.596ms  39.490ms
 7  10.119.73.122  39.762ms  39.581ms  39.500ms
 8  72.14.195.128  43.459ms  43.326ms  46.662ms
 9  108.170.253.113  53.036ms  52.762ms  53.058ms
10  74.125.253.17  54.976ms  54.714ms  68.118ms
11  172.217.31.206  53.805ms  53.617ms  53.662ms
```

Here the first common entry is 10.150.32.1, which is my device. The other common hops occur because the packets pass through the same routers (for instance in the example above both servers lie in USA, hence first the packets are routed to reach USA, before diverging to their respective location-based gateways.)

b) Route changes based on time of day : Since the internet commonly follows packet switching, the data is sent in packets which are routed by layer 2, layer 3 routers. This routing is affected by the load balancing techniques to reduce the effect of congestion in the link, due to which packets might go through different paths to reach the same destination if one path is congested.

traceroute to imdb.com (52.94.237.74), 64 hops max					traceroute to imdb.com (52.94.237.74), 64 hops max				
1	10.42.0.1	1.772ms	1.622ms	1.603ms	2	10.12.0.254	7.808ms	4.099ms	4.375ms
2	10.12.0.254	4.651ms	4.131ms	5.786ms	3	172.17.0.50	2.449ms	3.130ms	2.429ms
3	172.17.0.50	2.403ms	2.339ms	2.518ms	4	172.17.0.1	5.611ms	1.898ms	3.714ms
4	172.17.0.1	2.110ms	2.225ms	3.283ms	5	192.168.193.1	1.781ms	1.760ms	1.797ms
5	192.168.193.1	2.025ms	1.834ms	2.986ms	6	14.139.196.17	1.914ms	1.903ms	4.471ms
6	14.139.196.17	2.032ms	1.832ms	1.946ms	7	10.119.254.241	2.134ms	1.876ms	1.950ms
7	10.119.254.241	2.232ms	2.178ms	2.882ms	8	10.177.31.1	39.317ms	39.115ms	41.867ms
8	10.177.31.1	41.875ms	39.231ms	39.253ms	9	10.255.238.205	39.461ms	39.232ms	39.164ms
9	10.255.238.205	39.438ms	40.658ms	39.034ms	10	10.119.73.122	39.738ms	43.130ms	41.851ms
10	10.119.73.122	40.081ms	39.915ms	41.579ms	11	125.22.85.29	49.149ms	50.191ms	49.057ms
11	* 115.248.104.230	50.882ms	*		12	116.119.49.240	86.751ms	86.907ms	94.537ms
12	80.81.65.97	54.563ms	53.405ms	53.552ms	13	120.29.215.241	80.273ms	80.516ms	80.551ms
13	62.216.135.230	324.193ms	409.357ms	409.332ms	14	180.87.96.22	426.719ms	409.731ms	408.807ms
14	85.95.26.109	409.221ms	254.799ms	251.437ms	15	180.87.12.1	381.163ms	379.311ms	277.131ms
15	85.95.25.5	255.481ms	252.521ms	257.402ms	16	180.87.67.33	149.812ms	146.685ms	147.649ms
16	85.95.26.233	193.872ms	191.163ms	191.377ms	17	120.29.217.07	303.092ms	193.293ms	232.012ms
17	85.95.26.41	404.276ms	409.421ms	254.714ms					

As seen above, the intermediate hops for imdb.com are different for 7:30pm (left) and 11:30pm (right)

c) Traceroute unable to find any route :

traceroute to imdb.com (52.94.237.74), 64 hops max				
2	10.12.0.254	7.808ms	4.099ms	4.375ms
3	172.17.0.50	2.449ms	3.130ms	2.429ms
4	172.17.0.1	5.611ms	1.898ms	3.714ms
24	52.93.249.165	395.669ms	408.794ms	287.679ms
25	52.95.62.115	327.258ms	272.473ms	311.444ms
26	* * *			
27	52.93.129.130	409.937ms	288.277ms	325.133ms
28	54.239.42.182	284.163ms	331.732ms	408.772ms
29	* * *			
30	* * *			
31	* * *			

As shown in the image on the left, traceroute may be unable to find a route to some hosts because traceroute follows ICMP protocol which shares data via low priority packets. Due to this, the packets are prone to being dropped by networks in case of heavy congestion. Furthermore, these packets are also prone to being blocked by firewalls. Consecutive rows of * * * (asterisk) symbolize that packets are not acknowledged by the router in the given timelimit.

d) To find route when ping fails : Yes, it is possible to find the route to certain hosts which do not respond to ping. While ping works with simple ICMP, there is another tool known as **Tracert** which works by **targeting the final hop**, but limiting the TTL and waiting for a time exceeded message, and then increasing it by one for the next iteration. Therefore, the response it gets is not an ICMP echo reply to the ICMP echo request from the host along the way, but a time exceeded message from that host. Thus, by increasing the TTL based on the last router's response, the packet is saved from being dropped and can reach the host.

Q7 : ARP command : arp (-v) (-i if) (-H type) -a (hostname)

ARP stands for **Address Resolution Protocol**. The primary function of this protocol is to resolve the IP address of a system to its mac address. ARP command shows the MAC addresses resolved from the IP address (please note : this resolution is not done by ARP command, it simply shows the ARP cache containing the resolved pairs)

a) Show complete ARP table : To show complete table, do not write any hostname after arp

arp: the image shows the complete table in default format of ARP

rashit@rashis-aspire-e5-575g:~\$ arp				
Address	HWtype	HWaddress	Flags	Mask
gateway	ether	00:25:b4:d9:f7:c0	C	
10.150.39.149	ether	8c:85:90:5b:d7:84	C	

arp -a : the image on the right shows the complete table in BSD style output format

rashit@rashis-aspire-e5-575g:~\$ arp -a				
gateway (10.150.32.1)	at	00:25:b4:d9:f7:c0	[ether]	on wlp3s0
? (10.150.39.149)	at	8c:85:90:5b:d7:84	[ether]	on wlp3s0

IMPORTANT KEYWORDS :

- a) Address : This is the IPv4 address of the dest which has been resolved to MAC address
- b) HWType : This represents the type of hardware, which in the image above is "ethernet"
- c) HWAddress : This is the MAC address of the dest which has been resolved from the IP address
- d) Flags : The flags convey the status of the address, check flag types from Q4 ROUTE FLAGS
- f) Interface : This is the name of the interface, check the convention from Q4 INTERFACE NAMING

b) Modifying ARP table : To modify the ARP table, the user must have super-user privileges

- To **add** or **change** an existing entry, follow `arp -s <IP_address> <H/W Address> ...1)`
- To **delete** an existing entry, follow `arp -d <IP_address> <H/W Address>`

```

root@rashis-aspire-e5-575g:/home/rashi# arp -s 10.150.39.148 ff:ff:ff:ff:ff:00
root@rashis-aspire-e5-575g:/home/rashi# arp -s 10.150.39.18 ff:ff:ff:ff:ff:00
root@rashis-aspire-e5-575g:/home/rashi# arp -s 10.150.39.16 ff:ff:ff:ff:ff:ff
root@rashis-aspire-e5-575g:/home/rashi# arp -s 10.150.39.1 ff:ff:ff:ff:ff:ff
root@rashis-aspire-e5-575g:/home/rashi# arp

```

Address	HWtype	HWaddress	Flags	Mask	Iface
_gateway	ether	00:25:b4:d9:f7:c0	C		wlp3s0
10.150.39.18	ether	ff:ff:ff:ff:ff:00	CM		wlp3s0
10.150.39.1	ether	ff:ff:ff:ff:ff:ff	CM		wlp3s0
10.150.39.149	ether	ff:ff:ff:ff:ff:00	CM		wlp3s0
10.150.39.148	ether	ff:ff:ff:ff:ff:00	CM		wlp3s0
10.150.39.16	ether	ff:ff:ff:ff:ff:ff	CM		wlp3s0

arp -s <IP> <H/W Address>

The image on the left shows the four hosts that were added by command 1)

c) ARP table timeout : The exact time when an entry would be removed can not be predicted as it depends on a number of parameter :

- gc_stale time : If an entry hasn't been used for these many sec, then it is marked stale (eligible for removal during garbage collection)
- gc_interval : This is the time after which the garbage collector runs periodically

It is important to note that permanent entries are not deleted, but normal entries after gc_stale time are make stale (eligible for deletion) and then wait for the next gc_interval to arrive to be deleted in garbage collection process.

ALGO TO PREDICT TIMEOUT :

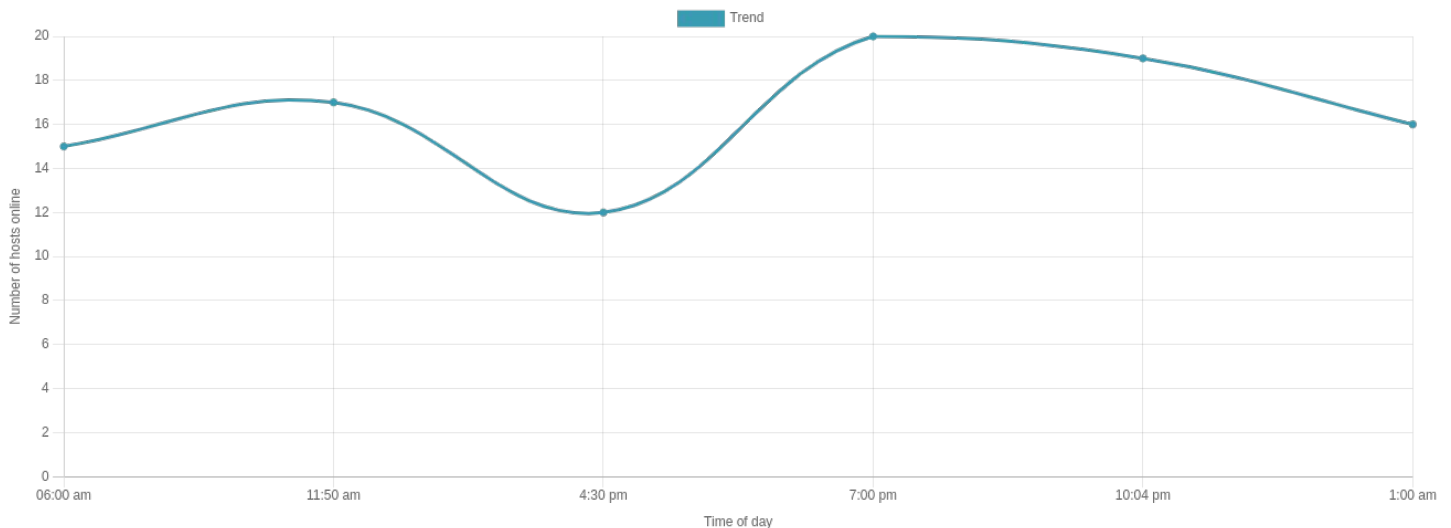
A method similar to poling can be used. For a newly added dummy entry, monitor the time after fixed intervals to see when the entry is deleted. A trial and error method to discover the timeout value is to add a temporary entry in the ARP table and keep on checking its presence after fixed intervals of time (say 5 seconds). Smaller the poling time, better would be the approximation.

Alternatively, one can use binary search for finding the cache time, for e.g. – Add a temporary entry in ARP and check after 5000ms. If then entry has been deleted, then add the entry again and check after 2500ms. The more no of iterations, the better will be the approximation.

d) Two IP address mapped to same Ethernet Addr : The scenario where two IP's can map to same Ethernet Address is when a router or a gateway connects two or more subnet ranges. When communicating with machines on the same subnet range, MAC address is used for directing the packages. In the ARP Table, the IP's of the devices which are connected in the other subnet range have the ethernet address/MAC address as that of the Router or Gateway which connects the two subnet ranges. ARP table is referred to convert these IP addresses to the MAC address and packets are sent to it(router/gateway). The router then uses it's routing table and sends the packet further to the correct device.

Q7 : ARP command : nmap -n -sP <subnet>/22

Subnet used : 10.16.0.254



OBSERVATION : As observed from the image above, the number of hosts are very low during the morning and class-hours. The number of hosts online increases from 6:00 am till 11:50 am. Then there is a sharp decrease and reaches a minimum of 12 at 4:30pm (Class hours). After class hours there is an increase and reaches a maximum at around 7:00pm and then gradually decreases again. These observations clearly state when the computers are switched ON or OFF in my LAN.