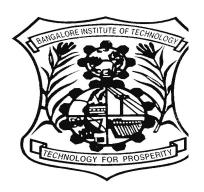
BANGALORE INSTITUTE OF TECHNOLOGY

K.R.ROAD, V.V. PURA, BENGALURU -560 004



Department of Computer Science and Engineering

BCSL404

Analysis and Design of Algorithms Laboratory Record

IV-Semester

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K. R. Road, V. V. Pura, Bengaluru- 560004

Department of Computer Science and Engineering

VISION:

To be a center of excellence in computer engineering education, empowering graduates as highly skilled professionals.

MISSION:

M1	To provide a platform for effective learning with emphasis on technical excellence.
M2	To train the students to meet current industrial standards and adapt to emerging technologies.
M3	To instill the drive for higher learning and research initiatives.
M4	To inculcate the qualities of leadership and Entrepreneurship.

PROGRAM EDUCATIONAL OBJECTIVES (PEO)

PEO-1	Graduates will apply fundamental and advanced concepts of Computer Science and Engineering for solving real world problems.					
PEO-2	Graduates will build successful professional careers in various sectors to facilitate societal needs.					
PEO-3	Graduates will have the ability to strengthen the level of expertise through higher studies and research.					
PEO-4	Graduates will adhere to professional ethics and exhibit leadership qualities to become an Entrepreneur.					

PROGRAM SPECIFIC OUTCOMES (PSOs)

PSO-1	The graduates of the program will have the ability to build software products by applying theoretical concepts and programming skills.
PSO-2	The graduates of the program will have the ability to pursue higher studies and Research in the modern computing era.

PROGRAM OUTCOMES (POs)

Engineering Graduates will be able to:

- 1. **Engineering knowledge**: Apply the knowledge of mathematics, science, engineering fundamentals, and an engineering specialization to the solution of complex engineering problems.
- 2. **Problem analysis**: Identify, formulate, review research literature, and analyze complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.
- 3. **Design/development of solutions**: Design solutions for complex engineering problems and design system components or processes that meet the specified needs with appropriate consideration for the public health and safety, and the cultural, societal, and environmental considerations.
- 4. **Conduct investigations of complex problems**: Use research-based knowledge and research methods including design of experiments, analysis and interpretation of data, and synthesis of the information to provide valid conclusions.
- 5. **Modern tool usage**: Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools including prediction and modeling to complex engineering activities with an understanding of the limitations.
- 6. **The engineer and society**: Apply reasoning informed by the contextual knowledge to assess societal, health, safety, legal and cultural issues and the consequent responsibilities relevant to the professional engineering practice.
- 7. **Environment and sustainability**: Understand the impact of the professional engineering solutions in societal and environmental contexts, and demonstrate the knowledge of, and need for sustainable development.
- 8. **Ethics**: Apply ethical principles and commit to professional ethics and responsibilities and norms of the engineering practice.
- 9. **Individual and team work**: Function effectively as an individual, and as a member or leader in diverse teams, and in multidisciplinary settings.
- 10. **Communication**: Communicate effectively on complex engineering activities with the engineering community and with society at large, such as, being able to comprehend and write effective reports and design documentation, make effective presentations, and give and receive clear instructions.
- 11. **Project management and finance**: Demonstrate knowledge and understanding of the engineering and management principles and apply these to one's own work, as a member and leader in a team, to manage projects and in multidisciplinary environments.
- 12. **Life-long learning**: Recognize the need for, and have the preparation and ability to engage in independent and life-long learning in the broadest context of technological change.

Bangalore Institute of Technology

K R Road, VV pura, Bengaluru 560004

Department of Computer Science and Engineering

Course Learning Objectives (CLO):

This course will enable students to:

- 1. Explain the methods of analyzing the algorithms and to analyze performance of algorithms.
- 2. State algorithm's efficiencies using asymptotic notations.
- 3. Solve problems using algorithm design methods such as the brute force method, greedy method, divide and conquer, decrease and conquer, transform and conquer, dynamic programming, backtracking and branch and bound.
- 4. Choose the appropriate data structure and algorithm design method for a specified application.
- 5. Introduce P and NP classes.

Course Outcomes (COs):

At the end of the course the student will be able to:

- 1. Develop programs to solve computational problems using suitable algorithm design strategy.
- 2. Compare algorithm design strategies by developing equivalent programs and observing running times for analysis (Empirical).
- 3. Make use of suitable integrated development tools to develop programs.
- 4. Demonstrate and present the development of program, its execution and running time(s) and record the results/inferences.

	PO 1	PO 2	PO 3	PO 4	PO 5	PO 6	PO 7	PO 8	PO 9	PO 10	PO 11	PO 12	PSO1	PSO2
CO1	3	2	3		1								2	
CO2	3	3	3		1	2				2		1	2	
CO3	2				3						1	1	2	
CO4			2							2			1	

Sl No.	Name of Experiment
1	Design and implement C/C++ Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Kruskal's algorithm.
2	Design and implement C/C++ Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Prim's algorithm.
3	a. Design and implement C/C++ Program to solve All-Pairs Shortest Paths problem using Floyd's algorithm.
	b. b. Design and implement C/C++ Program to find the transitive closure using Warshal's algorithm.
4	Design and implement C/C++ Program to find shortest paths from a given vertex in a weighted connected graph to other vertices using Dijkstra's algorithm
5	Design and implement C/C++ Program to obtain the Topological ordering of vertices in a given digraph.
6	Design and implement C/C++ Program to solve 0/1 Knapsack problem using Dynamic Programming method.
7	Design and implement C/C++ Program to solve discrete Knapsack and continuous Knapsack problems using greedy approximation method
8	Design and implement $C/C++$ Program to find a subset of a given set $S = \{sl, s2,,sn\}$ of n positive integers whose sum is equal to a given positive integer d
9	Design and implement C/C++ Program to sort a given set of n integer elements using Selection Sort method and compute its time complexity. Run the program for varied values of n> 5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.
10	Design and implement C/C++ Program to sort a given set of n integer elements using Quick Sort method and compute its time complexity. Run the program for varied values of n> 5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.
11	Design and implement C/C++ Program to sort a given set of n integer elements using Merge Sort method and compute its time complexity. Run the program for varied values of n> 5000, and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator
12	Design and implement C/C++ Program for N Queen's problem using Backtracking

Subject Code: BCSL404 CIE Marks: 50

Hours/Week: 0:0:2 Total Hours: 24

EVALUATION CRITERIA

Program Write-up and Execution rubrics (Max: 07 Marks)

		Good	Average
a.	Understanding of problem and approach to solve.	Demonstrate good knowledge of design concepts of given problem/algorithm and implementation in Java.	Moderate understanding of the algorithm and its implementation in Java. (2)
	(3 Marks)	(3)	
	Execution and Testing	Program hands all possible	Average conditions are
b.	(2 Marks)	conditions while execution with proper results.	defined and verified during execution.
		(2)	(1)
	Results and	Meticulous documentation	Acceptable documentation
	Documentation	of results obtained in data	shown in data sheets and
c.	(2 Marks)	sheets and manual.(2)	manual.(1)

Viva Rubrics (Max: 03 Marks)

		Good	Average
	Conceptual Understanding (2 Marks)	Explains different steps of algorithm/problem and related concepts involved.(2)	Adequately provides explanation on algorithm/problem.(1)
a.			
	Formulation of program/Program Explanation.	Different strategies and concepts to derive solutions to specified program or	Intuitive use of strategies and conceptsto derive solutions to specified program.
b.	(1 Marks)	Tracing of the program. (1)	(0.5)

Subject Code: BCSL404 CIE Marks : 50

Hours/Week: 0:0:2 Total Hours: 24

Schedule of Experiments

Sl. No	Name of Experiment	To be Completed	
1	Sample programs	Week1	
2	Selection sort	Week2	
3	Quick Sort with time complexity	Week3	
4	Merge sort with time complexity.	Week4	
5	Topological ordering	Week5	
6	Prim's , Kruskal's Algorithm to find MST	Week6	
7	Dijkstra's Algorithm to find shortest path to other vertices	Week7	
8	Lab Test -1	Week8	
9	Knapsack Using Greedy Method	Week9	
10	0/1 Knapsack Using Dynamic programing	Week10	
11	Floyd's Algorithm to implement All Pair Shortest Path Problem and Warshal's Algorithm to compute Transitive Closure	Week11	
12	Sum of Subset Problem, N-Queens Problem	Week12	
13	Final Lab Test	Week13	

Subject Code: BCSL404 CIE Marks : 50

Hours/Week: 0:0:2 Total Hours: 24

PROCEDURE FOR PROGRAM EXECUTION

ENVIRONMENT: Eclipse tool on Linux platform

- 1. Click on Applications \rightarrow Programming \rightarrow Eclipse
- 2. New \rightarrow Project \rightarrow C Project
- 3. Project name \rightarrow Enter USN
- 4. Project type \rightarrow Empty Project
- 5. Toolchains \rightarrow Linux GCC
- 6. Click on next \rightarrow
 - i. Select configurations \rightarrow select both Debug and Release or click on select all option on the right end
- 7. Finish

1. Design and implement C Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Kruskal's algorithm.

Aim: To apply Kruskal's Algorithm for computing the minimum spanning tree is directly based on the generic MST algorithm. It builds the MST in forest.

Definition: Kruskal's algorithm is an algorithm in graph theory that finds a minimum spanning tree for a connected weighted graph. This mean it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. If the graph is not connected, then it finds a minimum spanning forest .It is an example of a greedy algorithm.

Efficiency: With an efficient sorting algorithm, the time efficiency of Kruskal's algorithm will be in O(E log E)

Algorithm

Start with an empty set A, and select at every stage the shortest edge that has not been chosen or rejected, regardless of where this edge is situated in graph.

- Initially, each vertex is in its own tree in forest.
- Then, algorithm consider each edge in turn, order by increasing weight.
- If an edge (u, v) connects two different trees, then (u, v) is added to the set of edges of the MST, and two trees connected by an edge (u, v) are merged into a single tree.
- On the other hand, if an edge (u, v) connects two vertices in the same tree, then edge (u, v) is discarded.

Kruskals algorithm can be implemented using **disjoint set** data structure or **priority queue** data structure.

MST KRUSKAL (G, w)

- 1. $A \leftarrow \{\}$ // A will ultimately contains the edges of the MST
- 2. for each vertex v in V[G]
- 3. do Make Set (v)
- 4. Sort edge of E by nondecreasing weights w
- 5. for each edge (u, v) in E
- 6. do if FIND SET $(u) \neq$ FIND SET (v)
- 7. then $A = AU\{(u, v)\}$
- 8. UNION (u, v)
- 9. Return A

```
#include<stdio.h>
int i,j,k,a,b,u,v,n,ne=1;
int min,mincost=0,cost[9][9],parent[9];
int find(int);
int uni(int,int);
void main()
printf("\n\t Implementation of Kruskal's
algorithm\n"); printf("\nEnter the no. of vertices:");
scanf("%d",&n);
printf("\nEnter the cost adjacency matrix:\n");
for(i=1;i \le n;i++)
{
   for(j=1;j<=n;j++)
           scanf("%d",&cost[i][j]);
           if(cost[i][j]==0)
                   cost[i][j]=999;
    }
}
printf("The edges of Minimum Cost Spanning Tree are\n");
while(ne < n)
    for(i=1,min=999;i<=n;i++)
{
           for(j=1;j \le n;j++)
           {
                   if(cost[i][j] < min)
   min=cost[i][j];
   a=u=i;
                          b=v=j;
           }
u=find(u);
v = find(v);
```

```
if(uni(u,v))
 {
            printf("%d edge (%d,%d) =%d\n",ne++,a,b,min);
            mincost +=min;
     }
     cost[a][b]=cost[b][a]=999;
   printf("\n\tMinimum cost = %d\n",mincost);
}
   int find(int i)
            while(parent[i])
                    i=parent[i];
             return i;
   }
   int uni(int i,int j)
   {
            if(i!=j)
             {
                    parent[j]=i;
                    return 1;
            return 0;}
```

```
Implementation of Kruskal's algorithm

Enter the no. of vertices:4

Enter the cost adjacency matrix:
0 2 1 4
2 0 3 1
1 3 0 2
4 1 2 0

The edges of Minimum Cost Spanning Tree are
1 edge (1,3) =1
2 edge (2,4) =1
3 edge (1,2) =2

Minimum cost = 4
```

Input/Output:

```
Enter the n value:5
Enter the cost adjacency matrix:
0 10 15 9 999
10 0 999 17 15
15 999 0 20 999
9 17 20 0 18
999 15 999 18 0

The edges of minimum cost spanning tree are:
1 edge (1,4)
2 edge (1,2)
3 edge (1,3)
4 edge (2,5)
Minimum cost =49
```

2. Design and implement C Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Prim's algorithm.

Aim: To find minimum spanning tree of a given graph using prim's algorithm

Definition: Prim's is an algorithm that finds a minimum spanning tree for a connected weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all edges in the tree is minimized. Prim's algorithm is an example of a greedy algorithm.

Algorithm

```
MST_PRIM (G, w, v)
// Prim's algorithm for constructing a minimum spanning tree
// Input: A weighted connected graph G=(V,E)
// Output: The set of edges composing a minimum spanning tree of G
1. Q ← V[G]
2. for each u in Q do
3. key [u] ← ∞
4. key [r] ← 0
```

```
6. while queue is not empty do
7.  u ← EXTRACT_MIN (Q)
8. for each v in Adj[u] do
9. if v is in Q and w(u, v) < key [v]</li>
10. then π[v] ← w(u, v)
11. key [v] ← w(u, v)
```

Program:

5. $\pi[r] \leftarrow \text{NII}$

```
#include<stdio.h>
int visited[10]={0}, cost[10][10], min, mincost=0;
int i,j,ne=1, a, b, u, v;

int main()
{
    int num;
    printf("\n\t\t\Prim's Algorithm");
    printf("\n\nEnter the number of nodes= ");
    scanf("%d", &num);
    printf("\nEnter the adjacency matrix\n\n");

    for(i=1; i<=num; i++)
    {
}</pre>
```

```
for(j=1; j<=num; j++)
               {
                     scanf("%d", &cost[i][j]);
                     if(cost[i][j]==0)
                             cost[i][j]=999;
              }
       }
       visited[1]=1; while(ne < num)</pre>
              for(i=1,min=999;i<=num;i++)
              for(j=1;j<=num;j++)
                     if(cost[i][j]< min)
                     if(visited[i]!=0)
                             min=cost[i][j];
                             a=u=i;
                             b=v=j;
                      printf("\n Edge %d:(%d - %d) cost:%d",ne++,a,b,min);
                                    mincost=mincost+min;
                     visited[b]=1;
       cost[a][b]=cost[b][a]=999;
       printf("\n\n\n Minimun cost=%d",mincost);
       return 0; }
1)
```

```
Prim's Algorithm

Enter the number of nodes= 4

Enter the adjacency matrix

0 2 1 4
2 0 3 1
1 3 0 2
4 1 2 0

Edge 1:(1 - 3) cost:1
Edge 2:(1 - 2) cost:2
Edge 3:(2 - 4) cost:1
```

2) Input/output:

Enter n value:3

Enter the graph data:

0 10 1

1006

160

Enter the souce node:1

 $1 -> 3 \cos t = 1$

 $3 -> 2 \cos t = 7$

minimum Cost=7

3) a. Design and implement C Program to solve All-Pairs Shortest Paths problem using Floyd's algorithm.

Definition: The **Floyd algorithm** is a graph analysis algorithm for finding shortest paths in a weighted graph (with positive or negative edge weights). A single execution of the algorithm will find the lengths (summed weights) of the shortest paths between *all* pairs of vertices though it does not return details of the paths themselves. The algorithm is an example of dynamic programming.

Algorithm:

Floyd's Algorithm

Accept no .of vertices

Call graph function to read weighted graph // w(i,j)

```
Set D[] <- weighted graph matrix // get D\{d(i,j)\} for k=0
// If there is a cycle in graph, abort. How to find?
        Repeat for k = 1 to n
                Repeat for i = 1 to n
                   Repeat for j = 1 to n
                        D[i,j] = min \{D[i,j], D[i,k] + D[k,j]\}
Print D
#include<stdio.h>
#define INF 999
int min(int a,int b)
{
        return(a<b)?a:b;
}
void floyd(int p[][10],int n)
{
        int i,j,k;
        for(k=1;k \le n;k++)
        for(i=1;i \le n;i++)
        for(j=1;j \le n;j++)
        p[i][j]=min(p[i][j],p[i][k]+p[k][j]);
}
void main()
int a[10][10],n,i,j;
        printf("\nEnter the n value:");
        scanf("%d",&n);
        printf("\nEnter the graph data:\n");
        for(i=1;i \le n;i++)
        for(j=1;j \le n;j++)
        scanf("\%d",\&a[i][j]);
        floyd(a,n);
        printf("\nShortest path matrix\n");
        for(i=1;i \le n;i++)
                for(j=1;j \le n;j++)
                printf("%d ",a[i][j]);
                printf("\n");
        }
```

Output:

```
the
       number of vertices
  the
       cost adjacency matrix
13
  1
  2
Θ
Pairs
      Shortest
                 Path
1
  3
  1
3
0
  2
```

3 b). Design and implement C Program to find the transitive closure using Warshal's algorithm.

Definition: The Warshall algorithm is a graph analysis algorithm used to determine the transitive closure of a directed graph or all paths in a directed graph by using the adjacency matrix.. The algorithm is an example of Dynamic programming.

Algorithm

```
//Input: Adjacency matrix of digraph
//Output: R, transitive closure of digraph
Accept no .of vertices
Call graph function to read directed graph
Set R[] <- digraph matrix // get R \{r(i,j)\} for k=0
Print digraph
Repeat for k = 1 to n
                Repeat for i = 1 to n
                    Repeat for j = 1 to n
                         R(i,j) = 1 if
\{r_{ij}^{(k-1)} = 1 \text{ OR }\}
                                 r_{ik}^{(k-1)} = 1 and r_{ki}^{(k-1)} = 1
Print R
Program b):
#include<stdio.h>
void warsh(int p[][10],int n)
{
        int i,j,k;
```

```
for(k=1;k \le n;k++)
       for(i=1;i \le n;i++)
       for(j=1;j \le n;j++)
       p[i][j]=p[i][j] \parallel p[i][k] \&\& p[k][j];
}
int main()
       int a[10][10],n,i,j;
       printf("\nEnter the n value:");
       scanf("%d",&n);
       printf("\nEnter the graph data:\n");
       for(i=1;i \le n;i++)
       for(j=1;j \le n;j++)
       scanf("%d",&a[i][j]);
       warsh(a,n);
       printf("\nResultant path matrix\n");
       for(i=1;i \le n;i++)
        {
               for(j=1;j \le n;j++)
               printf("%d ",a[i][j]);
               printf("\n");
       return 0;
}
Output:
Enter the n value:4
Enter the graph data:
0100
0001
0000
1010
Resultant path matrix
1111
1111
0000
1111
```

4) Design and implement C Program to find shortest paths from a given vertex in a weighted connected graph to other vertices using Dijkstra's algorithm.

Dijikstra's algorithm: For a given source vertex(node) in the graph, the algorithm finds the path with lowest cost between that vertex and every other vertex. It can also be used for finding cost of shortest paths from a single vertex to a single destination vertex by stopping the algorithm once the shortest path to the destination vertex has been determined.

Efficiency:

1) $\Theta(|V^2|)$ graph represented by weighted matrix and priority queue as unordered array

2)O(E log₂ V) graph represented by adjacency lists and priority queue as min-heap

```
Dijkstra(G,s)
```

{

```
// Dijkstra's algorithm for single-source shortest paths
   // Input : A weighted connected graph G=(V,E) with nonnegative weights and
              its vertex s
  // Output : The length d<sub>v</sub> of a shortest path from s to v and its penultimate vertex
               pv for every v in V.
   Initialise(Q)
                                         // Initialise vertex priority queue to empty
   for every vertex v in V do
     {
               d_v \leftarrow \infty;
               pv←null
               Insert(Q,v,dv
     //Initialise vertex priority in the priority queue
       d_s \leftarrow 0;
       Decrease(Q,s ds)
                                                   //Update priority of s with ds
       Vt←Ø
       for i \leftarrow 0 to |v| - 1 do
                                                //delete the minimum priority element
       u^* \leftarrow DeleteMin(Q)
       Vt \leftarrow Vt U \{u^*\}
       for every vertex u in V-V<sub>T</sub> that is adjacent to u* do
       if d_u^* + w(u^*, u) < d_u
       d_u \leftarrow d_u^* + w(u^*, u);
       p<sub>u</sub>←u*
       Decrease(Q,u,du)
#include<stdio.h>
#define INF 999
void dijkstra(int c[10][10],int n,int s,int d[10])
     int v[10],min,u,i,j;
     for(i=1;i \le n;i++)
         d[i]=c[s][i];
         v[i]=0;
     }
```

```
v[s]=1;
    for(i=1;i \le n;i++)
       min=INF;
       for(j=1;j \le n;j++)
       if(v[j]==0 \&\& d[j] < min)
               min=d[j];
               u=j;
v[u]=1;
       for(j=1;j \le n;j++)
       if(v[j]==0 && (d[u]+c[u][j]) < d[j])
       d[j]=d[u]+c[u][j];
    }
}
int main()
{
    int c[10][10],d[10],i,j,s,sum,n;
    printf("\nEnter n value:");
    scanf("%d",&n);
    printf("\nEnter the graph data:\n");
    for(i=1;i \le n;i++)
    for(j=1;j \le n;j++)
    scanf("%d",&c[i][j]);
    printf("\nEnter the souce node:");
    scanf("%d",&s);
    dijkstra(c,n,s,d);
    for(i=1;i \le n;i++)
    printf("\nShortest distance from %d to %d is %d",s,i,d[i]);
  return 0;
Input/Output
1) Enter n value:6
Enter the graph data:
0 15 10 999 45 999
999 0 15 999 20 999
20 999 0 20 999 999
999 10 999 0 35 999
999 999 999 30 0 999
999 999 999 4 999 0
```

Enter the souce node:2

```
Shortest distance from 2 to 1 is 35
Shortest distance from 2 to 2 is 0
Shortest distance from 2 to 3 is 15
Shortest distance from 2 to 4 is 35
Shortest distance from 2 to 5 is 20
Shortest distance from 2 to 6 is 999
```

5) Design and implement C Program to obtain the Topological ordering of vertices in a given digraph.

Aim: To find topological ordering of given graph

Definition: Topological ordering is that for every edge in the graph, the vertex where the edge starts is listed before the edge where the edge ends.

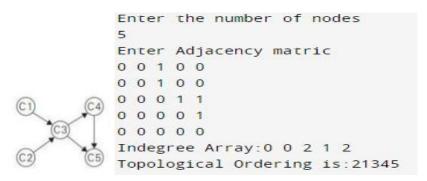
Algorithm:

- 1. Repeatedly identify in a remaining digraph a source which is a vertex with no incoming edges and delete it along with all edges outgoing from it
- 2. The order in which the vertices are deleted yields a solution to the topological sorting.

```
#include<stdio.h>
void topo sort(int a[20][20], int n)
{
        int t[10],vis[10],stack[10],i,j,indeg[10],top=0,ele,k=1;
        for(i=1;i \le n;i++)
        {
               t[i]=0;
               vis[i]=0;
               indeg[i]=0;
        for(i=1;i \le n;i++)
               for(j=1;j \le n;j++)
                        if(a[i][j]==1)
                                indeg[j]=indeg[j]+1;
                        }
                }
        printf("Indegree Array:");
        for(i=1;i \le n;i++)
               printf("%d ",indeg[i]);
```

```
for(i=1;i<=n;i++)
               if(indeg[i]==0)
                      stack[++top]=i;
                      vis[i]=1;
       while(top>0)
               ele=stack[top--];
               t[k++]=ele;
               for(j=1;j \le n;j++)
                      if(a[ele][j]==1 && vis[j]==0)
                              indeg[j]=indeg[j]-1;
                              if(indeg[j]==0)
                                     stack[++top]=j;
                                     vis[j]=1;
                              }
                      }
               }
       printf("\nTopological Ordering is:");
       for(i=1;i \le n;i++)
       printf("%d",t[i]);
int main()
{
       int n,a[20][20],i,j;
       printf("Enter the number of nodes\n");
       scanf("%d",&n);
       printf("Enter Adjacency matric\n");
       for(i=1;i<=n;i++)
               for(j=1;j<=n;j++)
                      scanf("%d",&a[i][j]);
       topo_sort(a,n);
```

Output:



6. Design and implement C Program to solve 0/1 Knapsack problem using Dynamic Programming method.

Aim: To implement 0/1 Knapsack problem using Dynamic programming

Definition: using **Dynamic programming**

It gives us a way to design custom algorithms which systematically search all possibilities (thus guaranteeing correctness) while storing results to avoid recomputing (thus providing efficiency).

We are given a set of n items from which we are to select some number of items to be carried in a knapsack(BAG). Each item has both a *weight* and a *profit*. The objective is to choose the set of items that fits in the knapsack and maximizes the profit. Given a knapsack with maximum capacity W, and a set S consisting of n items, Each item i has some weight wi and benefit value bi (all wi, bi and W are integer values) Problem: How to pack the knapsack to achieve maximum total value of packed items?

ALGORITHM

```
\label{eq:continuous_series} \begin{subarray}{ll} \begin{subarray}{ll}
```

PROGRAM:

```
#include<stdio.h>
int w[10],p[10],n;
int max(int a,int b)
{
       return a>b?a:b;
int knap(int i,int m)
       if(i==n) return w[i]>m?0:p[i];
       if(w[i] > m) return knap(i+1,m);
       return \max(\text{knap}(i+1,m),\text{knap}(i+1,m-w[i])+p[i]);
}
int main()
{
       int m,i,max profit;
       printf("\nEnter the no. of objects:");
       scanf("%d",&n);
       printf("\nEnter the knapsack capacity:");
       scanf("%d",&m);
       printf("\nEnter profit followed by weight:\n");
       for(i=1;i \le n;i++)
       scanf("%d %d",&p[i],&w[i]);
       max_profit=knap(1,m);
       printf("\nMax profit=%d",max_profit);
       return 0;
}
Input/Output:
Enter the no. of objects:4
Enter the knapsack capacity:6
Enter profit followed by weight:
78 2
45 3
92 4
71 5
```

7. Design and implement C Program to solve discrete Knapsack and continuous Knapsack problems using greedy approximation method.

This program first calculates the profit-to-weight ratio for each item, then sorts the items based on this ratio in non-increasing order. It then fills the knapsack greedily by selecting items with the highest ratio until the knapsack is full. If there's space left in the knapsack after selecting whole items, it adds fractional parts of the next item. Finally, it prints the optimal solution and the solution vector.

Here's a simplified version of the C program to solve discrete Knapsack and continuous Knapsack problems using the greedy approximation method:

```
#include<stdio.h>
void knapsack(int n, float weight[], float profit[], float capacity)
  float x[20], tp = 0;
 int i, j, u;
  u = capacity;
  for (i = 0; i < n; i++)
    x[i] = 0.0;
  for (i = 0; i < n; i++)
    if (weight[i] > u)
      break;
    else {
      x[i] = 1.0;
      tp = tp + profit[i];
      u = u - weight[i];
  }
  if (i \le n)
    x[i] = u / weight[i];
  tp = tp + (x[i] * profit[i]);
  printf("\nThe result vector is:- ");
  for (i = 0; i < n; i++)
    printf("%f\t", x[i]);
```

```
printf("\nMaximum profit is:- %f", tp);
}
int main()
  float weight[20], profit[20], capacity;
 int num, i, j;
 float ratio[20], temp;
 printf("\nEnter the no. of objects:- ");
 scanf("%d", &num);
 printf("\nEnter the wts and profits of each object:- ");
 for (i = 0; i < num; i++) {
    scanf("%f %f", &weight[i], &profit[i]);
  }
 printf("\nEnter the capacity of knapsack:- ");
  scanf("%f", &capacity);
 for (i = 0; i < num; i++) {
   ratio[i] = profit[i] / weight[i];
  }
 for (i = 0; i < num; i++)
    for (j = i + 1; j < num; j++) {
     if (ratio[i] < ratio[j]) {
       temp = ratio[j];
       ratio[j] = ratio[i];
        ratio[i] = temp;
        temp = weight[j];
        weight[i] = weight[i];
        weight[i] = temp;
        temp = profit[j];
       profit[j] = profit[i];
       profit[i] = temp;
     }
   }
  }
```

```
knapsack(num, weight, profit, capacity);
return(0);
}

Input/Output:
1) Enter the no. of objects:- 4
Enter the wts and profits of each object:- 56 23 78 45 98 76 78 78
Enter the capacity of knapsack:- 100
The result vector is:-: 1.00000 0.000000 0.000000 0.000000
Maximum profit is:- 78.0

2) Enter the no. of objects:- 3
Enter the wts and profits of each object:- 20 30 25 40 10 35
Enter the capacity of knapsack:- 40
The result vector is:-: 1.00000 1.000000 0.25
Maximum profit is:- 82.5
```

8. Design and implement C/C++ Program to find a subset of a given set $S = \{sl, s2,....,sn\}$ of n positive integers whose sum is equal to a given positive integer d.

AIM: An instance of the Subset Sum problem is a pair (S, t), where $S = \{x1, x2, ..., xn\}$ is a set of positive integers and t (the target) is a positive integer. The decision problem asks for a subset of S whose sum is as large as possible, but not larger than t.

Algorithm:

```
SumOfSub (s, k, r) 

//Values of x[j], 1 \le j \le k, have been determined 

//Node creation at level k taking place: also call for creation at level K+1 if possible 

//s = sum of 1 to k-1 elements and r is sum of k to n elements 

//generating left child that means including k in solution 

Set x[k] = 1 

If (s + s[k] = d) then subset found, print solution 

If (s + s[k] + s[k+1] \le d) then SumOfSum (s + s[k], k+1, r - s[k]) 

//Generate right child i.e. element k absent 

If (s + r - s[k] \ge d) AND (s + s[k+1]) \le d 

THEN \{x[k] = 0; 

SumOfSub(s, k+1, r - s[k])
```

Program

#include<stdio.h>

#define MAX 10

```
int s[MAX],x[MAX],d;
void sumofsub(int p,int k,int r)
{
       int i;
       x[k]=1;
       if((p+s[k])==d)
       {
               for(i=1;i \le k;i++)
               if(x[i]==1)
                      printf("%d ",s[i]);
               printf("\n");
       }
       else
       if(p+s[k]+s[k+1] \le d)
               sumofsub(p+s[k],k+1,r-s[k]);
       if((p+r-s[k]>=d) && (p+s[k+1]<=d))
       {
               x[k]=0;
               sumofsub(p,k+1,r-s[k]);
}
}
int main()
{
       int i,n,sum=0;
       printf("\nEnter the n value:");
       scanf("%d",&n);
       printf("\nEnter the set in increasing order:");
```

```
for(i=1;i \le n;i++)
               scanf("%d",&s[i]);
       printf("\nEnter the max subset value:");
       scanf("%d",&d);
       for(i=1;i \le n;i++)
               sum=sum+s[i];
       if(sum<d \parallel s[1]>d)
               printf("\nNo subset possible");
       else
       sumofsub(0,1,sum);
       return 0;
}
Input/output:
Enter the n value:9
Enter the set in increasing order:1 2 3 4 5 6 7 8 9
Enter the max subset value:9
126
1 3 5
18
234
27
36
4 5
9
```

9) Design and implement C/C++ Program to sort a given set of n integer elements using Selection Sort method and compute its time complexity. Run the program for varied values of n> 5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.

Aim: Sort a given set of elements using Selection sort and determine the time required to sort elements. Repeat the experiment for different values of n, the number of elements in the list to be sorted and plot a graph of the time taken versus n.

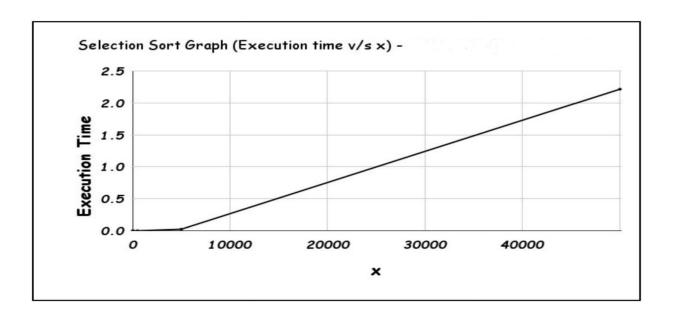
Definition: selection sort is a sorting routine that scans a list of items repeatedly and, on each pass, selects the item with the lowest value and places it in its final position. It is based on brute force approach. Sequential search is a $\Theta(n^2)$ algorithm on all inputs.

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
// Function to perform selection sort
voidselectionSort(intarr[], int n)
inti, j, min idx, temp;
  // One by one move boundary of unsorted subarray
for (i = 0; i \le n-2; i++)
     // Find the minimum element in unsorted array
min idx = i;
for (j = i+1; j \le n-1; j++)
if (arr[j] <arr[min_idx])</pre>
min idx = j;
     // Swap the found minimum element with the first element
temp=arr[i];
arr[i]=arr[min_idx];
arr[min idx]=temp;
}
int main()
int n;// Number of elements
printf("enter the value of n\n");
scanf("%d",&n);
intarr[n];
srand(time(NULL)); // Seed for random number generation
  // Generate random numbers and fill the array
printf("Original Array: ");
```

```
for (inti = 0; i< n; i++)
{
arr[i] = rand() % 1000; // Generating random numbers between 0 and 999
printf("%d ", arr[i]);
  }
printf("\n");
  // Perform selection sort
clock tstart time = clock();
selectionSort(arr, n);
clock tend time = clock();
doubletotal time = (double)(end time - start time) / CLOCKS PER SEC;
  // Display sorted array
printf("Sorted Array: ");
for (inti = 0; i < n; i++)
printf("%d ", arr[i]);
printf("\n");
  // Print time complexity
printf("Time Complexity: O(n^2)\n");
printf("Execution Time: %f seconds\n", total time);
return 0;
}
```

Output:

Note down the different values of n(eg: 5000,6000,7000,8000,9000,10000) and required time for all values of n respectively. Plot the graph for the same.



10) Design and implement C/C++ Program to sort a given set of n integer elements using Quick Sort method and compute its time complexity. Run the program for varied values of n> 5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.

Aim:

The aim of this program is to sort 'n' randomly generated elements using Quick sort and Plotting the graph of the time taken to sort n elements versus n.

<u>Definition:</u> Quicksort is based on the Divide and conquer approach. Quick sort divides array according to their value. Partition is the situation where all the elements before some position s are smaller than or equal to A[s] and all the elements after position s are greater than or equal to A[s].

Efficiency: Cbest(n) Θ (nlog2n),Cworst(n) Θ (n²),Cavg(n) Θ 1.38nlog2n

```
Algorithm: Quicksort
(A[1...r])

//Sorts a sub array by quick
sort

//Input:Asub arrayA[1..r]ofA[0..n-1] ,defined by its left and right indices 1

//andr
```

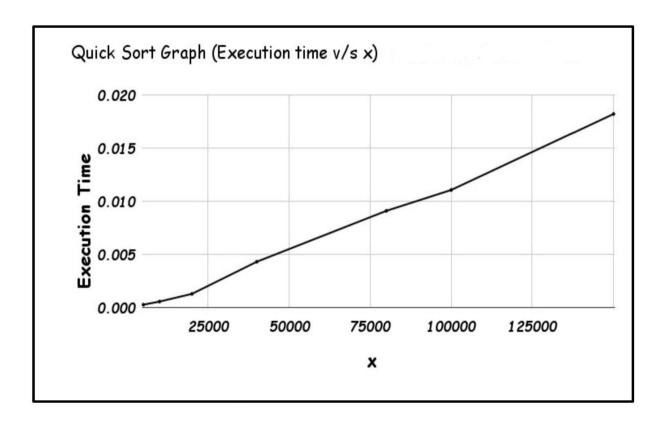
```
//Output:The subarrayA[1..r] sorted in nondecreasing
         order
  ifl < r
           s=Partition(A[l..r])//sisasplitpositio
           n Quick sort (A[1 ...s-1])
           Quicksort (A[s+1...r])
ALGORITHM Partition (A[1...r])
//Partitionasub arraybyusingits firstelement asa pivot
//Input :Asubarray A[1...r]ofA[0...n-1]definedby itsleft and right indices and// r(1
<r)
//Output:ApartitionofA[1...r], with the split position returned as this function's value p=A[1]
i=1; j=r+1; repeat
  delay(500);
  repeati=i+1untilA[i]>
  =p repeat j=j-1 until
  A[J] \le p
                    Swap
  (A[i],A[j])
until I>=j
Swap(A[i],A[j])//UndolastSwapwheni>=j
Return j
Program:
#include <stdio.h>
#include<stdlib.h>
#include <time.h>
//Functionto partitionthearray
intpartition(inta[],intlow,inthigh){ in
  t pivot = a[low]; // Pivot element
  inti=low,temp;//Indexofsmallerelement int
  j=high+1;// Index of largest element
  while(i \le j)
  {
```

```
doi++;while(pivot>=a[i]
     ); do j--
     ;while(pivot<a[j]);</pre>
       if(
        {
          temp=a[i]; a[i]=a[j]; a[j]=temp;
       }
}
  temp=a[low];
  a[low]=a[j];
  a[j]=temp;
  return j;
}
//Functiontoperform Quick Sort
void quickSort(inta[],intlow,int high)
{
       intk;
       if(low>high) return;
     //Partitioning index
     k=partition(a,low, high);
     //Separatelysortelements after partition
```

```
quickSort(a, low, k - 1);// left part of pivot element
     quickSort(a,k+1,high);//rightpartofpivotelement
}
  intmain()
  srand(time(0));//Seed for random number generation
  intn;
  printf("Enter the number of elements:");
   scanf("%d", &n);
  inta[n];
  printf("Randomly generated array:\n");
  for (int i = 0; i < n; i++) {
     a[i]=rand()%1000;//Generatingrandomnumbersbetween0and999
     printf("%d ", a[i]);
  a[n]=9999;
  clock tstart=clock();
  quickSort(a,0,n-1); clock t end=clock();
  printf("\nSortedarray:\n");
  for (int i = 0; i < n; i++) {
     printf("%d",a[i]);
  }
  printf("\n");
  printf("Total
                   time taken
                                to sort %d
                                                   elements
                   %lf\n",n,((double)(end-start)/CLOCKS PER SEC));
  return 0;
}
```

Output:

Note down the different values of n (eg: 5000,6000,7000,8000,9000,10000) and required time for all values of n respectively. Plot the graph for the same.



11. Design and implement C Program to sort a given set of n integer elements using Merge Sort method and compute its time complexity. Run the program for varied values of n> 5000, and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.

Aim:

The aim of this program is to sort 'n' randomly generated elements using Merge sort and Plotting the graph of the time taken to sort n elements versus n.

Definition: Merge sort is a sort algorithm based on divide and conquer technique. It divides the array element based on the position in the array. The concept is that we first break the list into two smaller lists of roughly the same size, and then use merge sort recursively on the sub problems, until they cannot subdivide anymore. Then, we can merge by stepping through the lists in linear time.

Its time efficiency is $\Theta(n \log n)$.

Algorithm: Merge sort (A[0...n-1]

```
// Sorts array A[0..n-1] by Recursive merge sort
     // Input : An array A[0..n-1] elements
    // Output : Array A[0..n-1] sorted in non decreasing order
      If n > 1
        Copy A[0...(n/2)-1] to B[0...(n/2)-1]
        Copy A[0...(n/2)-1] to C[0...(n/2)-1]
        Mergesort (B[0...(n/2)-1])
        Mergesort (C[0...(n/2)-1])
        Merge(B,C,A)
   ALGORITHM Merge (B[0...p-1], C[0...q-1], A[0....p+q-1])
        // merges two sorted arrays into one sorted array
      // Input : Arrays B[0..p-1] and C[0...q-1] both sorted
      // Output : Sorted array A[0.... p+q-1] of the elements of B and C
         i = 0;
         j = 0;
         k=0;
        While i < p and j < q do
          if B[i] \leftarrow C[i]
           A[k]=B[i]; i=i+1;
           else
               A[k] = C[j]; j=j+1
          k=k+1;
         if i = p
             Copy C[j..q-1] to A[k...p+q-1]
          else
             Copy B[i ... p-1] to A[k ... p+q-1]
Program:
#include<stdlib.h>
#include<stdio.h>
#include<time.h>
void merge(int arr[], int l, int m, int r)
       int i, j, k;
       int n1 = m - 1 + 1;
       int n2 = r - m;
       // Create temp arrays
       int L[n1], R[n2];
```

{

```
// Copy data to temp array
       for (i = 0; i < n1; i++)
               L[i] = arr[1+i];
       for (j = 0; j < n2; j++)
               R[j] = arr[m + 1 + j];
       // Merge the temp arrays
       i = 0;
       j = 0;
       k = 1;
       while (i < n1 \&\& j < n2)
       {
               if (L[i] \leq R[j])
                       arr[k] = L[i];
                       i++;
               else
                       arr[k] = R[j];
                       j++;
               k++;
       // Copy the remaining elements of L[]
       while (i \le n1)
               arr[k] = L[i];
               i++;
               k++;
       }
// Copy the remaining elements of R[]
       while (j < n2)
               arr[k] = R[j];
               j++;
               k++;
       }
}
```

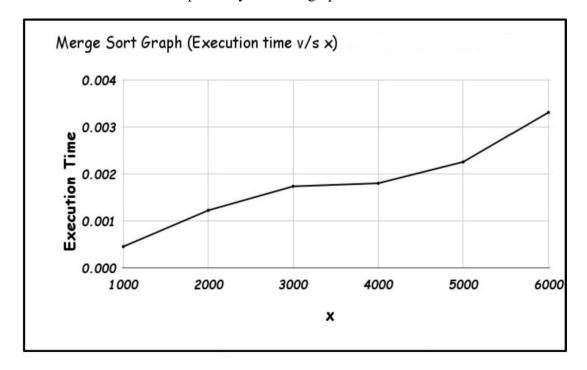
```
void mergeSort(int arr[], int l, int r)
       if (1 \le r)
        {
               // Finding mid element
               int m = 1+(r-1)/2;
               // Recursively sorting both the halves
               mergeSort(arr, 1, m);
               mergeSort(arr, m+1, r);
               merge(arr, l, m,r);
        }
}
void main()
{
       int n;
       printf("enter the value of n\n");
       scanf("%d",&n);
       int arr[n],i;
       srand(time(NULL));
       printf("the original array");
       for(i=0;i<n;i++)
        {
               arr[i]=rand()%1000;
               printf("%d ",arr[i]);
        }
       printf("\n");
       clock_t start=clock();
       mergeSort(arr, 0,n);
       clock t end=clock();
       printf("Sorted array:");
       for(i=0;i<n;i++)
        {
```

```
printf("%d ",arr[i]);
}
printf("\n");

printf("Total time taken to sort %d elements is %lf\n",n,((double)(end-start)/CLOCKS_PER_SEC));
}
```

Output:

Note down the different values of n (eg: 5000,6000,7000,8000,9000,10000) and required time for all values of n respectively. Plot the graph for the same.



12) Design and implement C/C++ Program for N Queen's problem using Backtracking.

Aim: To implement N Queens Problem using Back Tracking.

Definition:

The object is to place queens on a chess board in such a way as no queen can capture another one in a single

move.

Recall that a queen can move horizontal, vertical, or diagonally an infinite distance. This implies that no two queens can be on the same row, col, or diagonal We usually want to know how many different placements there are **Using Backtracking Techniques.**

Algorithm for new queen be placed	All solutions to the n-queens problem
Algorithm Place(k,i)	Algorithm NQueens(k, n)
//Return true if a queen can be placed in kth	// its prints all possible placements of n-
row & ith column	queens on an n×n chessboard.
//Other wise return false	{
(for i:=1 to n do{
for j:=1 to k-1 do	if Place(k,i) then
if(x[j]=i or Abs(x[j]-i)=Abs(j-k)))	1
then return false	X[k]:=I;
return true	if(k==n) then write (x[1:n]);
}	else NQueens(k+1, n);
60	1
	l în

Program

```
printf("Solution %d:\n\n",++count);
// No queen has placed initially
   for(int i=1;i<=n;i++)
     for(int j=1; j <=n; j++)
       c[i][j]='X';
// Place the queens on the chess board
   for(int i=1;i<=n;i++)
    c[i][x[i]]='Q';
// Place where the queens have been placed on the chess board
    for(int i=1; i \le n; i++)
     {
      for(int j=1; j <=n; j++)
       printf("%c",c[i][j]);
      printf("\n");
// Function to check whether the queens can be placed successfully or not
```

```
int place(int k,int xk)
{
  for(int i=1;i \le k-1;i++)
     //check whether two queens attach vertically or diagonally
    if((x[i]==xk)||(abs(i-k)==abs(x[i]-xk)))
     return False;//Queen cannot be placed in the kth column
     }
     return True;//kth queen can be successfully placed
}
void nqueen(int k)
  int j;
  for(int j=1;j<=n;j++)
    if(place(k,j))//Try to place in various columns
      x[k]=j;//Queen can be placed
      if(k==n)
      printsolution(),printf("\n");
      else
      nqueen(k+1);//Place the next queen in next row
```

```
void main()
  printf("enter the number of queens\n");
  scanf("%d",&n);
 nqueen(1);
 }
Input/Output:
Enter the no. of queens: 4
Solution 1
XQXX
XXXQ
QXXX
XXQX
Solution 2
XXQX
QXXX
XXXQ
XQXX
```