

# CS2610: Computer Organization and Architecture

## Lab Report 3

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### Objective:

*To analyse the efficiency of array multiplier and Multiplier using Booth's algorithm in terms of processing delay and power consumption.*

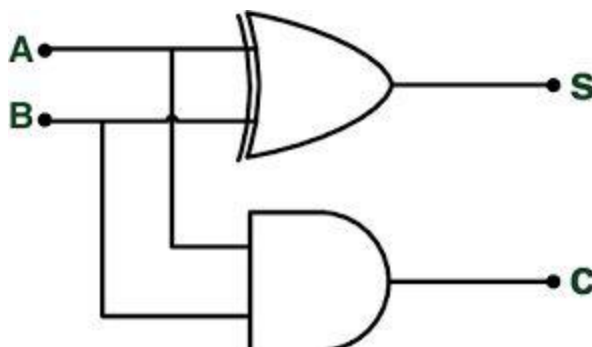
### Problem:

*To implement the below mentioned scenarios in Verilog:*

- a.) Implement an array multiplier that would multiply two unsigned numbers  $A$  and  $B$ .*
- b.) Implement Booth's algorithm that would multiply two 2's complement numbers  $A$  and  $B$ .*

### Theory and Approach:

***Half Adder:** A half adder is a type of adder, an electronic circuit that performs the addition of numbers. The half adder is able to add two single binary digits and provide the output plus a carry value. It has two inputs, called  $A$  and  $B$ , and two outputs  $S$  (sum) and  $C$  (carry).*

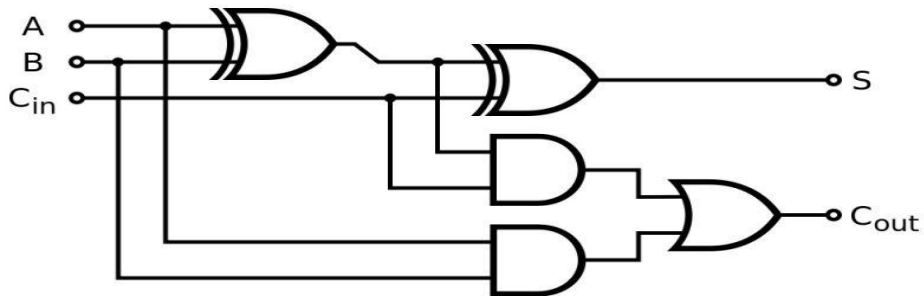


*Figure: Half Adder*

## *Code for half adder*

```
module half_adder(input a,b, output c,s);  
  
    assign s = a^b;  
    assign c = a&b;  
  
endmodule
```

***Full Adder:** A full adder is a logical circuit that performs an addition operation on three one-bit binary numbers. The full adder produces a sum of the two inputs and carry value. It gives a two bit output. First bit as sum and the second bit is the carry bit. Both these represents the sum of the given three one bit inputs.*



*Figure 1: Full Adder Circuit*

## *Code for full adder*

```
module full_adder(input a, input b, input c, output s, output C);  
  
    assign s = a^b^c;  
    assign C = ((a^b)&c) | (a&b);  
  
endmodule
```

## 1.) Array multiplier of two unsigned numbers

Array multiplier is well known due to its regular structure. Multiplier circuit is based on add and shift algorithm. Each partial product is generated by the multiplication of the multiplicand with one multiplier bit. The partial product are shifted according to their bit orders and then added.

We have used full adders and half adders to implement the array multiplier in the way as given in the below image.

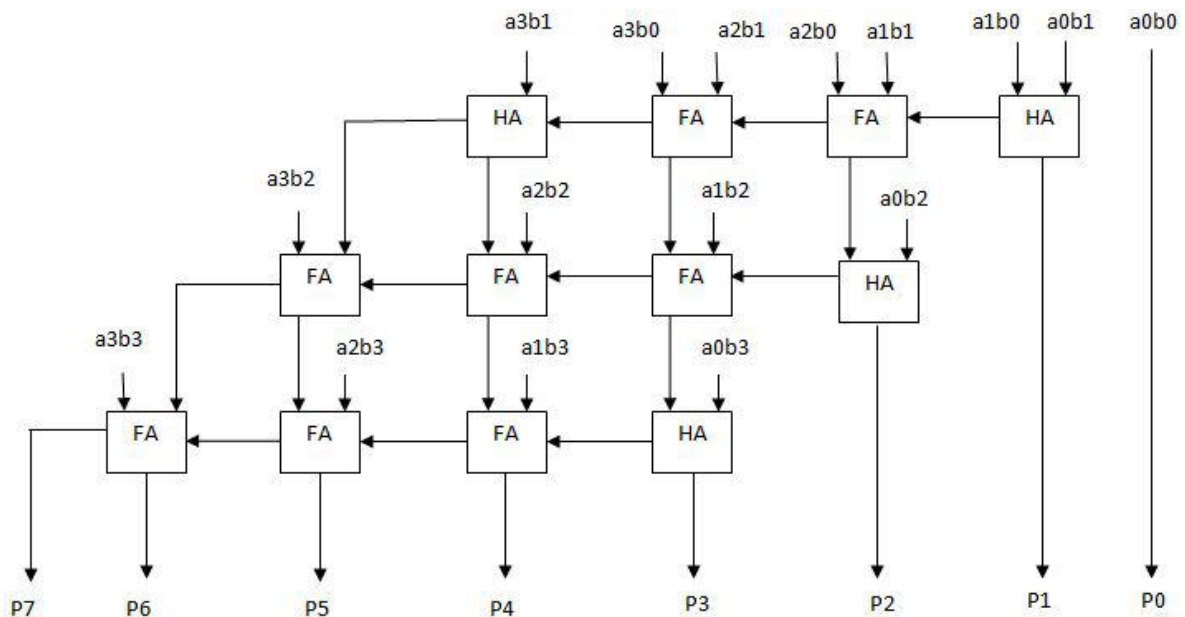


Figure 2: 4 Array Multiplier

## Code for Array Multiplier

```
module main(input [3:0]a,b,
            output [7:0] out);

    wire c1,c2,c3,c4,c5,c6,c7,c8,c9,c10,c11;
    wire s1,s2,s3,s4,s5,s6;
```

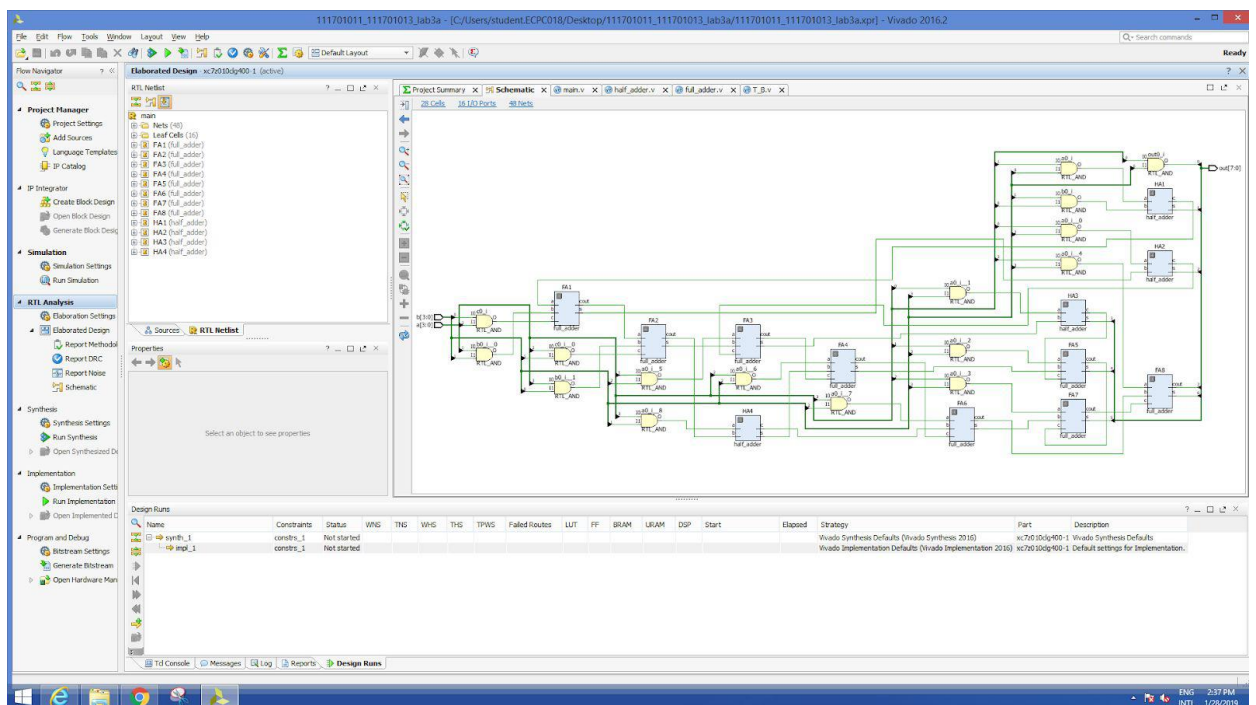
```

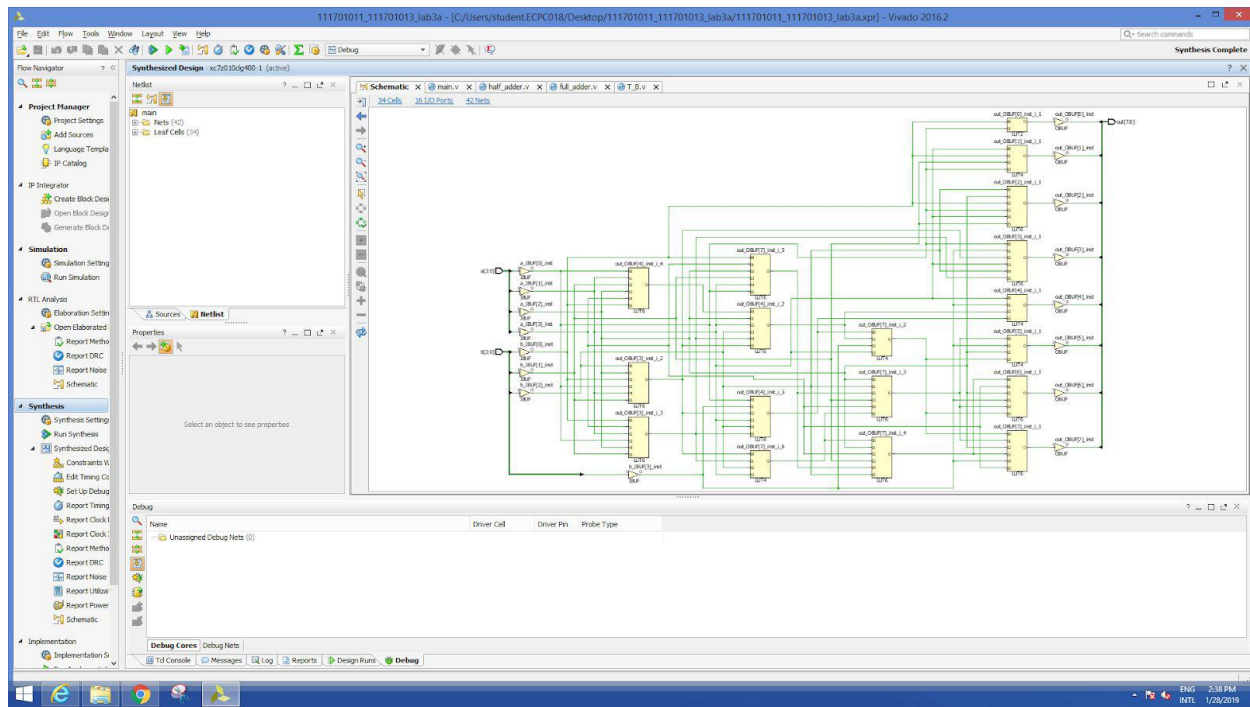
assign out[0] = a[0]&b[0];
half_adder HA1 (a[1]&b[0],b[1]&a[0],c1,out[1]);
full_adder FA1 (c1,a[1]&b[1],b[0]&a[2],c2,s1);
half_adder HA2 (a[0]&b[2],s1,c3,out[2]);
full_adder FA2 (c2,a[3]&b[0],a[2]&b[1],c4,s2);
full_adder FA3 (a[1]&b[2],s2,c3,c5,s3);
half_adder HA3 (a[0]&b[3],s3,c6,out[3]);
half_adder HA4 (a[3]&b[1],c4,c7,s4);
full_adder FA4 (a[2]&b[2],s4,c5,c8,s5);
full_adder FA5 (a[1]&b[3],s5,c6,c9,out[4]);
full_adder FA6 (a[3]&b[2],c7,c8,c10,s6);
full_adder FA7 (a[2]&b[3],s6,c9,c11,out[5]);
full_adder FA8 (a[3]&b[3],c10,c11,out[7],out[6]);

```

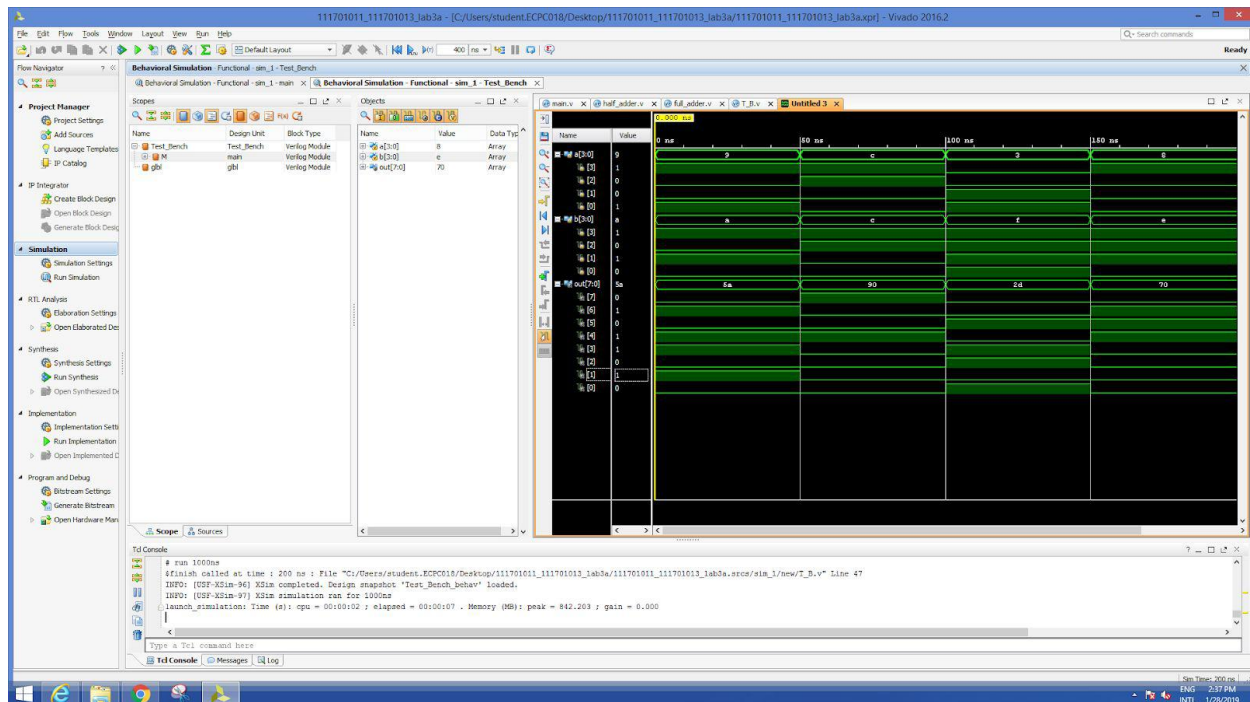
endmodule

## Elaborated Design

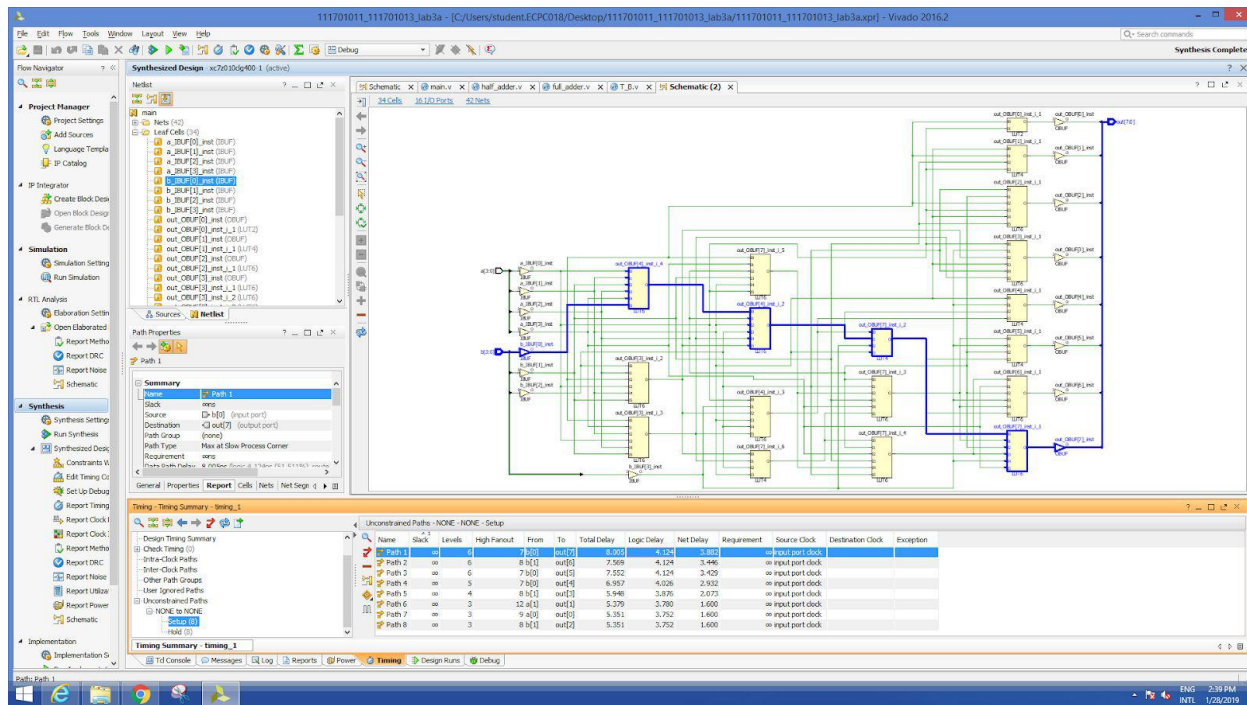




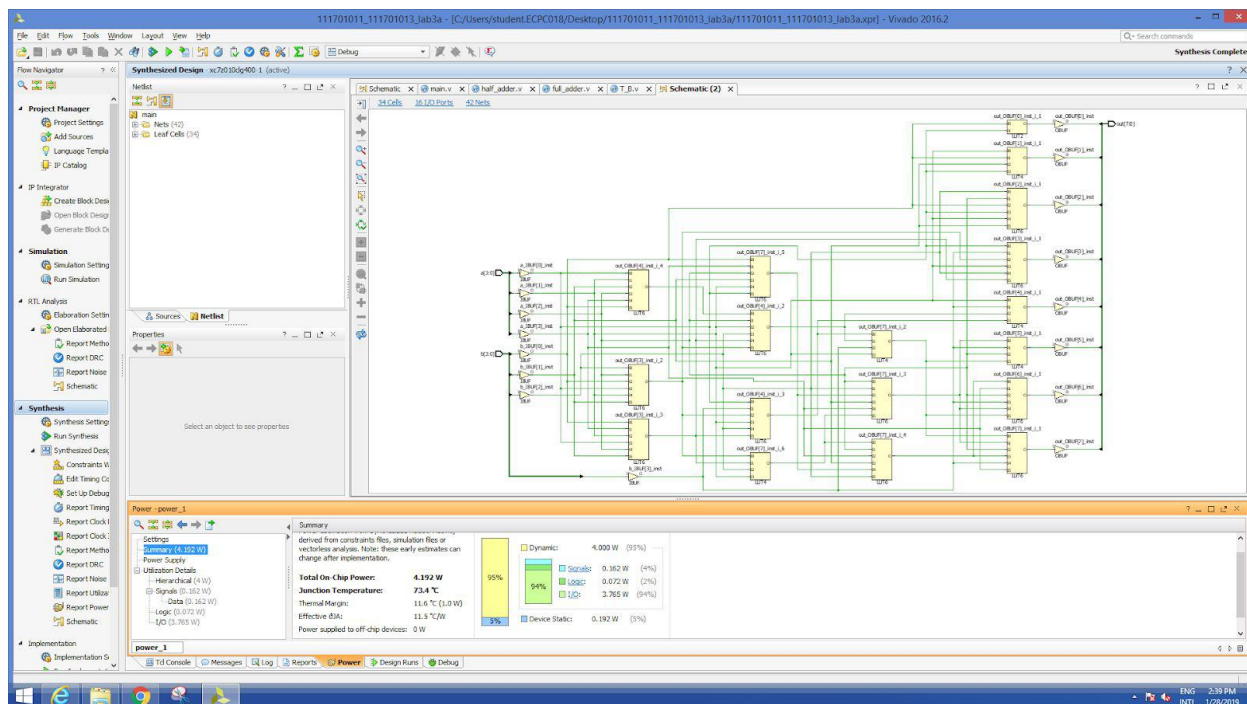
## Simulation Results



## Timing Results



## Schematic and On chip power Results





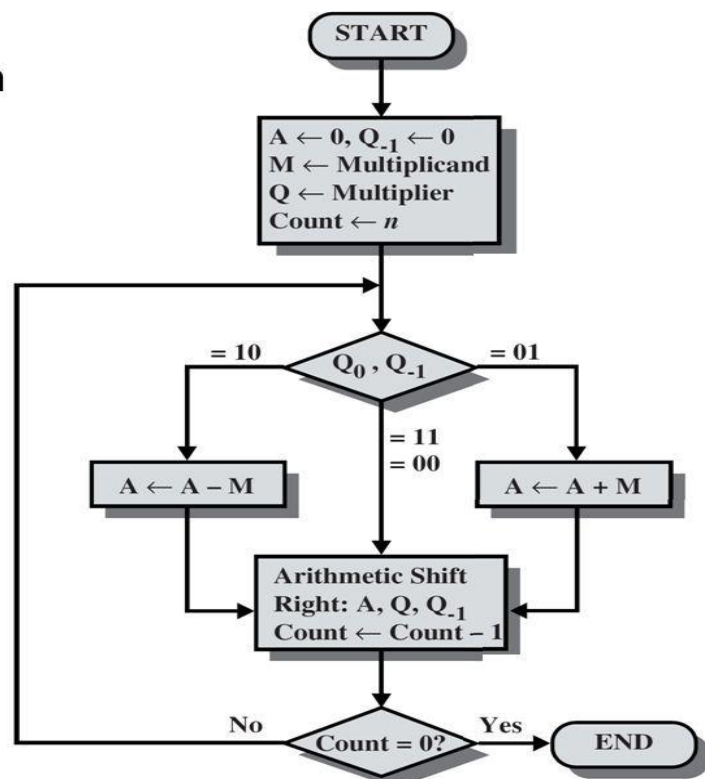
## 2.) Multiplier using Booth's Algorithm

Booth's multiplication algorithm is a multiplication algorithm that multiplies two signed binary numbers in two's complement notation.

Booth's algorithm examines adjacent pairs of bits of the ' $N$ '-bit multiplier  $Y$  in signed two's complement representation, including an implicit bit below the least significant bit,  $y_{-1} = 0$ . For each bit  $y_i$ , for  $i$  running from 0 to  $N - 1$ , the bits  $y_i$  and  $y_{i-1}$  are considered. Where these two bits are equal, the product accumulator  $P$  is left unchanged. Where  $y_i = 0$  and  $y_{i-1} = 1$ , the multiplicand times  $2^i$  is added to  $P$ ; and where  $y_i = 1$  and  $y_{i-1} = 0$ , the multiplicand times  $2^i$  is subtracted from  $P$ . The final value of  $P$  is the signed product.

The representations of the multiplicand and product are both also in two's complement representation, like the multiplier. As stated here, the order of the steps is not determined. Typically, it proceeds from *LSB* to *MSB*, starting at  $i = 0$ ; the multiplication by  $2^i$  is then typically replaced by incremental shifting of the  $P$  accumulator to the right between steps; low bits can be shifted out, and subsequent additions and subtractions can then be done just on the highest  $N$  bits of  $P$ .

### Booth's Algorithm Flowchart



## 2's Comp Multiplication Booth's Algorithm

<pre>       0111     x0011  (0)     11111001 1-0     0000000 1-1     000111 0-1     00010101 (-21)           </pre> <p>(a) <math>(7) \times (3) = (21)</math></p>	<pre>       0111     x1101  (0)     11111001 1-0     0000111 0-1     111001 1-0     11101011 (-21)           </pre> <p>(b) <math>(7) \times (-3) = (-21)</math></p>
<pre>       1001     x0011  (0)     00000111 1-0     0000000 1-1     111001 0-1     11101011 (-21)           </pre> <p>(c) <math>(-7) \times (3) = (-21)</math></p>	<pre>       1001     x1101  (0)     00000111 1-0     1111001 0-1     000111 1-0     00010101 (21)           </pre> <p>(d) <math>(-7) \times (-3) = (21)</math></p>

**Figure 9.14 Examples Using Booth's Algorithm**

### *Code for Booth's Multiplier*

```

module main( input [3:0] M,input [3:0] Q,output [7:0] O);

  wire [7:0] A, B;
  reg [7:0] temp, A1, X, Y, P, C1;

  adder a1(X, Y, A);
  adder a2(A, P, B);
  adder a3(B, temp, O);

  always@*
  begin
    if(M[3] == 0)
      begin
        A1 = 8'b00000000;
      end
    else
      begin
        A1 = 8'b11110000;
      end
    if(Q[0] == 1)

```



```

begin
    A1[0] = M[0];
    A1[1] = M[1];
    A1[2] = M[2];
    A1[3] = M[3];

    C1[0] = ~A1[0];
    C1[1] = ~A1[1];
    C1[2] = ~A1[2];
    C1[3] = ~A1[3];
    C1[4] = ~A1[4];
    C1[5] = ~A1[5];
    C1[6] = ~A1[6];
    C1[7] = ~A1[7];
    C1 = C1 + 1;
    X = C1;
end
else
begin
    X = 8'b00000000;
end
if(Q[0] == 1 && Q[1] == 0)
begin
    A1[1] = M[0];
    A1[2] = M[1];
    A1[3] = M[2];
    A1[4] = M[3];
    Y = A1;
end
else if(Q[0] == 0 && Q[1] == 1)
begin
    A1[1] = M[0];
    A1[2] = M[1];
    A1[3] = M[2];
    A1[4] = M[3];
    C1[0] = ~A1[0];
    C1[1] = ~A1[1];
    C1[2] = ~A1[2];
    C1[3] = ~A1[3];
    C1[4] = ~A1[4];
    C1[5] = ~A1[5];
    C1[6] = ~A1[6];
    C1[7] = ~A1[7];

```

```

        C1 = C1 + 1;

    Y = C1;
end
else
begin
Y = 8'b00000000;
end

if(Q[1] == 1 && Q[2] == 0)
begin
    A1[0] = 1'b0;
    A1[1] = 1'b0;
    A1[2] = M[0];
    A1[3] = M[1];
    A1[4] = M[2];
    A1[5] = M[3];
    P = A1;
end
else if(Q[1] == 0 && Q[2] == 1)
begin
    A1[0] = 1'b0;
    A1[1] = 1'b0;
    A1[2] = M[0];
    A1[3] = M[1];
    A1[4] = M[2];
    A1[5] = M[3];
    C1[0] = ~A1[0];
    C1[1] = ~A1[1];
    C1[2] = ~A1[2];
    C1[3] = ~A1[3];
    C1[4] = ~A1[4];
    C1[5] = ~A1[5];
    C1[6] = ~A1[6];
    C1[7] = ~A1[7];
    C1 = C1 + 1;

    P = C1;
end
else
begin
P = 8'b00000000;
end

```

```

if(Q[2] == 1 && Q[3] == 0)
begin
    A1[0] = 1'b0;
    A1[1] = 1'b0;
    A1[2] = 1'b0;
    A1[3] = M[0];
    A1[4] = M[1];
    A1[5] = M[2];
    A1[6] = M[3];
    temp = A1;
end
else if(Q[2] == 0 && Q[3] == 1)
begin
    A1[0] = 1'b0;
    A1[1] = 1'b0;
    A1[2] = 1'b0;
    A1[3] = M[0];
    A1[4] = M[1];
    A1[5] = M[2];
    A1[6] = M[3];
    C1[0] = ~A1[0];
    C1[1] = ~A1[1];
    C1[2] = ~A1[2];
    C1[3] = ~A1[3];
    C1[4] = ~A1[4];
    C1[5] = ~A1[5];
    C1[6] = ~A1[6];
    C1[7] = ~A1[7];
    C1 = C1 + 1;

    temp = C1;
end
else
begin
    temp = 8'b00000000;
end
end

```

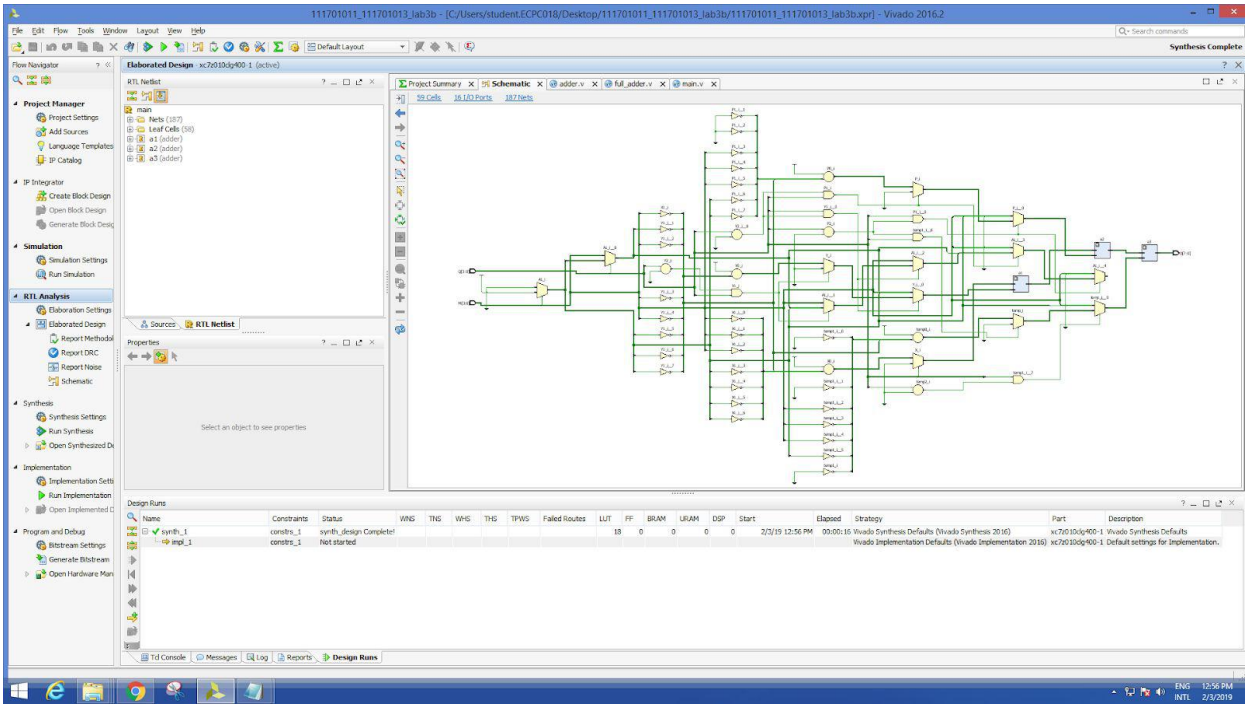
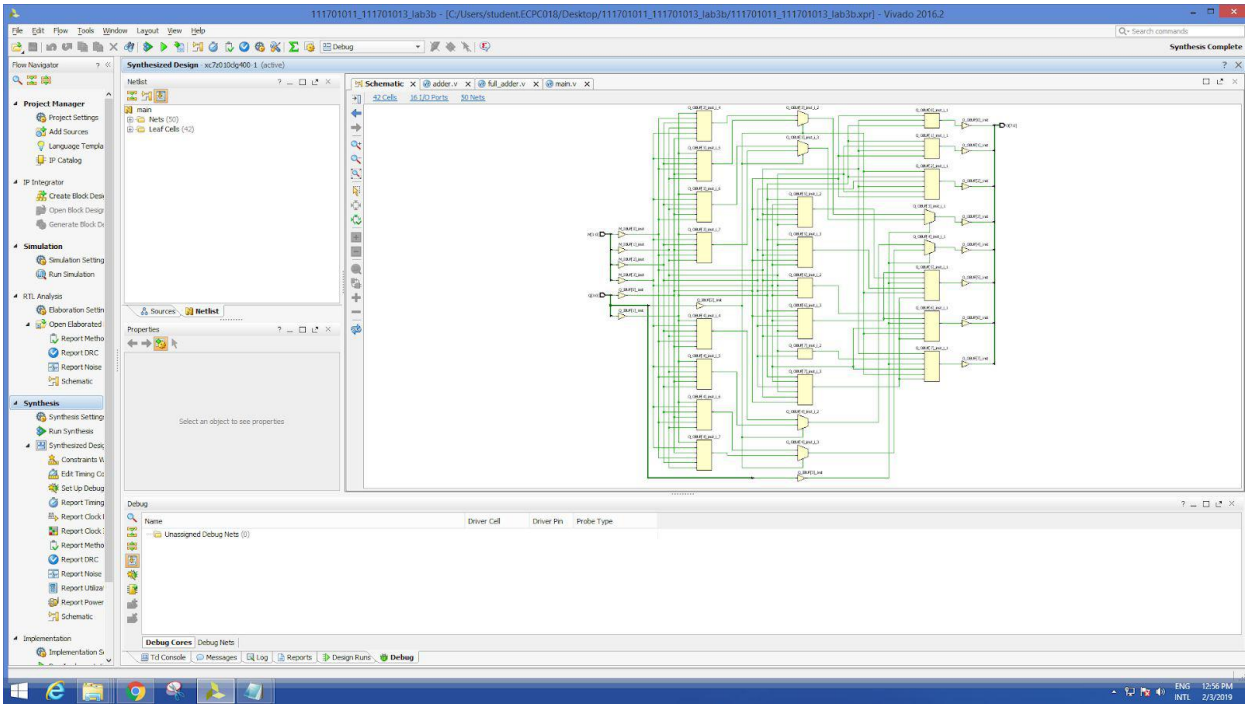
Endmodule

***Adder:** Adder is 8 bit adder of two 8-bit given numbers. It gives only Sum and carry out (last bit ) is neglected.*

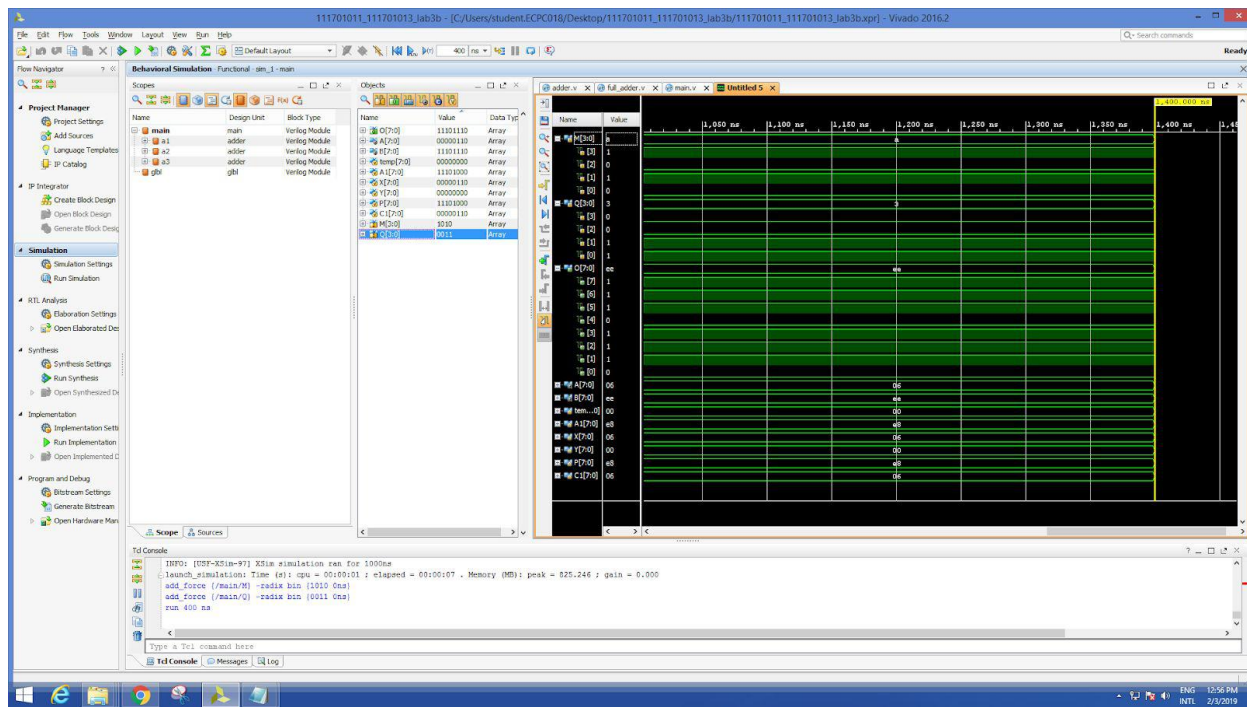
### ***Code for adder***

```
module adder(  
    input [7:0] A,  
    input [7:0] B,  
    output [7:0] S  
);  
  
    wire c1, c2, c3, c4, c5, c6, c7, c8;  
  
    fullAdder fa1(A[0], B[0], 0, S[0], c1);  
    fullAdder fa2(A[1], B[1], c1, S[1], c2);  
    fullAdder fa3(A[2], B[2], c2, S[2], c3);  
    fullAdder fa4(A[3], B[3], c3, S[3], c4);  
    fullAdder fa5(A[4], B[4], c4, S[4], c5);  
    fullAdder fa6(A[5], B[5], c5, S[5], c6);  
    fullAdder fa7(A[6], B[6], c6, S[6], c7);  
    fullAdder fa8(A[7], B[7], c7, S[7], c8);  
endmodule
```

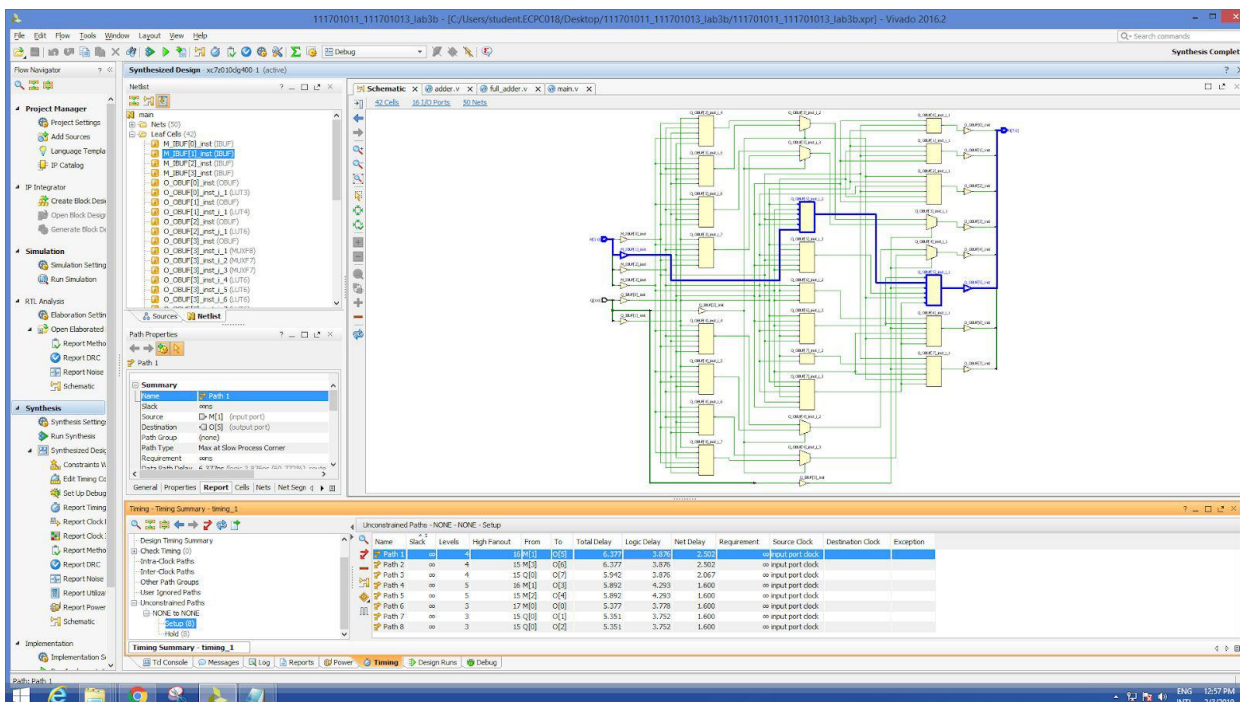
Elaborated design



# Simulation Results

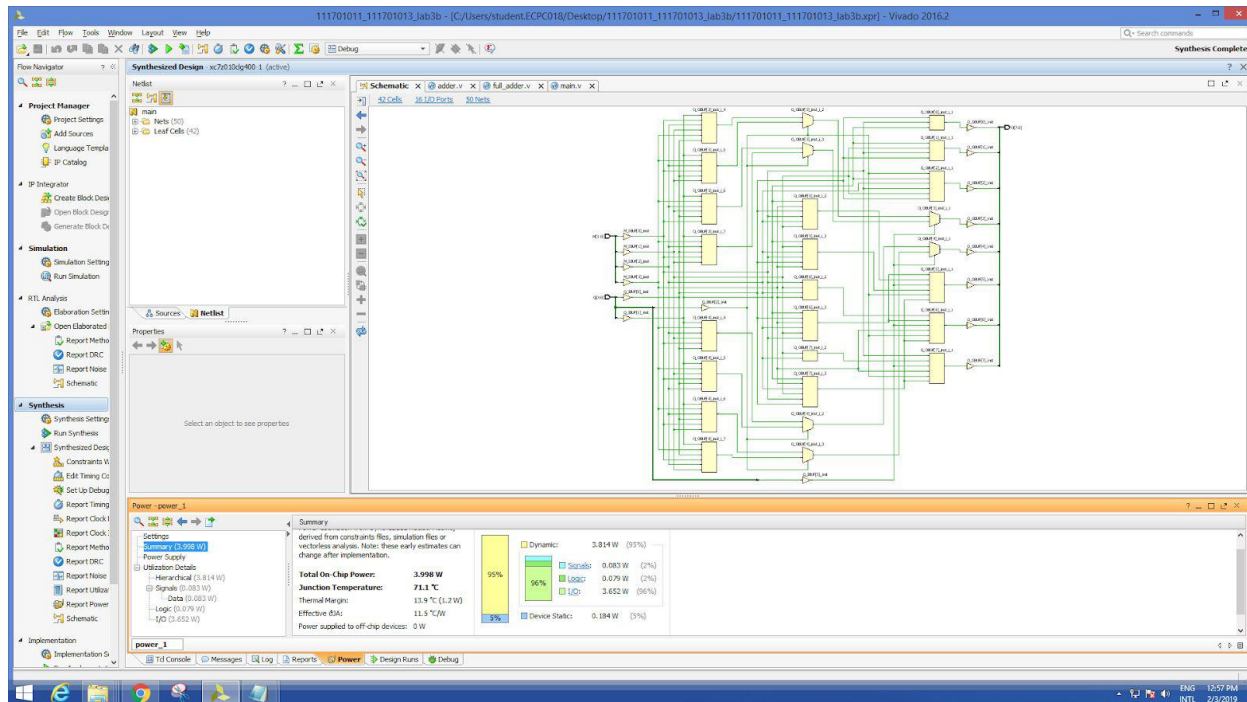


# Timing Results





## Schematic and Power Results



## Observations:

*Implement an array multiplier that would multiply two unsigned numbers  $A$  and  $B$ :*

*Dynamic Power used :  $4.000W$*

*Static Power used :  $0.192W$*

*Total On Chip Power used :  $4.192W$*

*Static Power is 5% of the total power while dynamic power is 95%.*

*In dynamic power,*

*94% is used for I/O operations =  $3.765W$*

*2% is used in logic =  $0.072W$*

*4% is used in signals =  $0.162W$*

*Junction Temperature was  $73.4^{\circ}C$*

*Implement Booth's algorithm that would multiply two 2's complement numbers  $A$  and  $B$ :*

*Dynamic Power used : 3.814W*

*Static Power used : 0.184W*

*Total On Chip Power used : 3.998W*

*Static Power is 5% of the total power while dynamic power is 95%.*

*In dynamic power,*

*96% is used for I/O operations = 3.652W*

*2% is used in logic = 0.079W*

*2% is used in signals = 0.083W*

*Junction Temperature was 71.1°C*

- 1.) In the first question, power consumption was more as compared to Multiplier circuit developed in second question .*
- 2.) Junction Temperature for the first question is also higher than the Junction temperatures in the second question.*
- 3.) In both the cases, dynamic power percentages are the same.*
- 4.) Most of the dynamic power is consumed by the I/O.*
- 5.) Rest of the dynamic power is used in logic and signals which is very less than input output dynamic power.*
- 6.) More buffers and logic was there in the first logic as compared to the second one.*

## **Results:**

- 1.) Power consumption is less when there are less number of components are used at the same time.*
- 2.) Power consumption reduces by a great factor when the number of components working parallelly is decreased.*
- 3.) But the time delay increases as the parallel logic is transformed into the pipeline logic.*
- 4.) But the static power remains constant in both the logic.*
- 5.) Time delay is lesser in second program than in the first program.*
- 6.) Booth's algorithm is more efficient as compared to the array multiplier logic as it takes lesser power consumption as well as lesser total delay.*
- 7.) Booth's Algorithm involve more dynamic processing as compared to the Array*

*multiplier algorithm.*

*8.) It takes 8 Full adder and 4 Half adder to implement  $4 \times 4$  bit array multiplier.*

## **Conclusions:**

- 1.) More power is consumed in the parallel programming as compared to sequential (pipeline) programming.*
- 2.) Static power consumed remains same as it is consumed when there is no activity on the circuit so it is independent on the circuit.*
- 3.) Processing delay is more in case of pipeline fashion while it is less in parallel fashion.*