

# CS5014 Foundations of Data Science & Machine Learning

## Quiz 04

May 10, 2021 — 8.30 - 9.30 AM

### Instructions

1. This is a **zoom proctored** exam. Please adjust your seating so that **your face, hands, answer book and the mobile phone that you will use to scan the sheets are always in the webcam view**. Do not leave your seat or talk to anyone during the exam.
2. Write your answer on plain paper with your **name and roll number on the first sheet**.
3. This is a **closed book** exam. Do not refer to any books, notes, the Internet or any other person during the exam.
4. You can take **maximum 5 minutes after the exam to scan** the sheets into a **single PDF** file and upload to Moodle. Submissions made after 9:35 AM will be evaluated only if there is a genuine reason for the delay.

### Questions

**Question 1.** Let  $H$  be the 3-dimensional lattice  $[n]^3$  with self loops at all the boundary vertices. All the non-loop edges have weight 1 while the self-loop at a boundary vertex  $v$  has weight  $6 - d_v$ , where  $d_v$  is the degree of  $v$ . Hence the total weight of all edges incident to any vertex is 6. Find the normalised conductance  $\phi(H)$  of  $H$ .

**Solution.** We will ignore rounding errors in this analysis throughout. If you choose  $S$  to be the left half vertices of the lattice, i.e.,  $S = [n/2] \times [n] \times [n]$ , we get  $E(S, \bar{S}) = n^2$ ,  $w(S) = 6|S| = 3n^3$ , and hence  $\phi(S, \bar{S}) = n^2/(3n^3) = 1/3n$ . Hence  $\phi(H) \leq 1/3n$ .

Showing that  $\phi(H) \geq 1/3n$  rigorously is a bit more tricky. One can do some hand waving and say that any  $S$  which maximises  $\phi(S, \bar{S})$  will be a rectangular box sticking to a corner of  $H$  and then analyse three cases based on how many dimensions of the box are full.

A more rigorous argument can be made as follows. Let  $S$  be a subset of  $V(H)$  which maximises  $\phi(S, \bar{S})$  and size *less than*  $n^3/2$ . Then there is no vertex  $v$  outside  $S$  such that it has a total edge weight of 3 or more into  $S \cup \{v\}$ . Otherwise  $S \cup \{v\}$  has a larger  $\phi$ -value than  $S$  since the weight of escaping edges did not increase while the weight of the set itself increased by one. Geometrically this means that there are no “concave corners” in  $S$ . The same argument can be extended to show that there is no pair of adjacent vertices  $u$  and  $v$  outside of  $S$  such that the total edge weight between  $\{u, v\}$  and  $S \cup \{u, v\}$  is at least 6. Otherwise  $S \cup \{u, v\}$  is better than  $S$ . This geometrically means that there are no “concave edges” in  $S$ . Putting it together we get that  $S$  is convex and being a subset of the lattice, it has to be a rectangular box. It is an easy argument that this rectangular box has to touch a corner of  $H$  if it is maximising  $\phi$ . Further, because of the self loops at the boundary of the lattice, a pair of boundary vertices  $\{u, v\}$  of the lattice next to an edge of  $S$  has a total edge weight of at least 6 between  $\{u, v\}$  and  $S \cup \{u, v\}$ . Hence  $S$  has size exactly  $n^3/2$ . Once the size of  $S$  is  $n^3/2$  we cannot do the fiddling around by adding more vertices. But we can still remove a vertex from a convex corner of  $S$  and add it to a concave corner, remove a line from a convex edge and add it to a concave edge etc. Thus we can always end up with a  $\phi$ -maximising  $S$  of size  $n/2$  which is the left half vertices of  $H$ . Thus our upper bound from the previous paragraph is tight.

**Question 2.** Let  $C_4$  be the cycle on 4-vertices  $\{0, 1, 2, 3\}$ . Find the hitting time  $h(0, 1)$ , for the standard random walk on  $C_4$ .

*Hint.* Search for symmetries.

*Warning.* It is recommended that you do this question at the end.

**Solution.** By symmetry  $h(i, i+1)$  should be the same for all  $i \in \{0, \dots, 3\}$ , where addition is modulo 4. Call this common value  $x$ . Symmetry again tells you that  $\forall i$ ,  $h(i, i-1) = x$ . Similarly  $h(i, i+2)$  should be the same for all  $i \in \{0, \dots, 3\}$ , where addition is modulo 4. Call this common value  $y$ . Now consider

$$\begin{aligned} h(0, 2) &= \frac{1}{2}(1 + h(1, 2)) + \frac{1}{2}(1 + h(3, 2)) \\ y &= \frac{1}{2}(1 + x) + \frac{1}{2}(1 + x) \\ y &= 1 + x \end{aligned}$$

$$\begin{aligned} h(1, 2) &= \frac{1}{2}1 + \frac{1}{2}(1 + h(0, 2)) \\ x &= \frac{1}{2}1 + \frac{1}{2}(1 + y) \\ x &= 1 + y/2 \end{aligned}$$

Solving the above two equations gives  $x = 3$  and  $y = 4$ . Hence our answer  $h(0, 1) = x = 3$ .

**Question 3.** Let  $A$  be a real symmetric square matrix. Find a relation between

1. Singular vectors and eigenvectors of  $A$ .
2. Singular values and eigenvalues of  $A$ .

Prove your answer.

**Solution.** Since  $A$  is real symmetric, it is always diagonalisable and has  $n$  eigenvalues  $\lambda_1, \dots, \lambda_n$  and  $n$  orthonormal eigenvectors  $e_1, \dots, e_n$  corresponding to them. Proving this is a typical linear algebra exercise:

For any vector  $x = x_1e_1 + \dots + x_ne_n$  in  $\mathbb{R}^n$   $y = Ax = \lambda_1x_1e_1 + \dots + \lambda_nx_ne_n$ . Hence one of the unit vectors which maximises  $\|Ax\|$  is  $e_1$ . That is,  $e_1$  qualifies to be a first singular vector of  $A$  with singular value  $\sigma_1 = |\lambda_1|$ . Any vector  $x$  orthogonal to  $e_1$  can be represented as  $x = x_2e_2 + \dots + x_ne_n$  and it is easy to see that one of the vectors maximising  $\|Ax\|$  from this set is  $e_2$ . Hence  $e_2$  qualifies to be a second singular vector with  $\sigma_2 = |\lambda_2|$ .

Hence  $\forall i \in [n]$ ,  $\sigma_i = |\lambda_i|$ . The set of eigenvectors  $e_1, \dots, e_n$  of  $A$  qualify to be singular vectors of  $A$  also. But remember that neither the set of eigenvectors nor the set of singular vectors is unique. Hence we can only say that any set of orthonormal eigenvectors of  $A$  will also qualify to be singular vectors of  $A$ .