Quick Sort Algorithm Overview

Here's a complete and clean summary of the entire Quick Sort explanation video:

What is Quick Sort?

Quick Sort is a **Divide and Conquer** sorting algorithm that:

- Sorts an array in **ascending** or **descending** order.
- Works by choosing a pivot, placing it in its correct sorted position, and recursively sorting elements to the left and right.

Quick Sort Steps (High-Level Intuition)

1. Pick a Pivot

- Choose any element: first, last, middle, or random.
- In this tutorial: **first element** is chosen as the pivot.

2. Place Pivot in Its Correct Position

- All elements smaller than the pivot go to the left.
- All elements greater than the pivot go to the right.
- This divides the array into two subarrays (left and right of pivot).

3. Recursively Apply Quick Sort

- Sort the left subarray.
- Sort the right subarray.
- Stop when the subarray has ≤ 1 element (already sorted).

Key Terms

• Partition Index: Final position of the pivot where left < pivot < right.

- Divide and Conquer:
 - Divide the array using pivot (partitioning).
 - Conquer by recursively sorting subarrays.

How Partitioning Works (Two Pointer Approach)

- **1.** Initialize two pointers:
 - i at low (start of array)
 - j at high (end of array)
- **2.** While i <= j:
 - Move i right until arr[i] > pivot
 - Move j left until arr[j] < pivot
 - If i < j:swap arr[i] and arr[j]
 - If $i \ge j$: swap pivot with $arr[j] \rightarrow Now pivot is at correct place$
- 3. Return j as the partition index.

Pseudo Code

```
void quickSort(vector<int>& arr, int low, int high) {
  if (low < high) {
    int pIndex = partition(arr, low, high);
    quickSort(arr, low, pIndex - 1);  // Left part
    quickSort(arr, pIndex + 1, high);  // Right part
  }
}
int partition(vector<int>& arr, int low, int high) {
  int pivot = arr[low];
  int i = low + 1;
  int j = high;
```

```
while (i <= j) {
    while (i <= high && arr[i] <= pivot) i++;
    while (j >= low && arr[j] > pivot) j--;
    if (i < j) swap(arr[i], arr[j]);
}

swap(arr[low], arr[j]); // Place pivot at correct position
    return j;
}</pre>
```

Time and Space Complexity

- Time Complexity:
 - Best/Average Case: O(N log N)
 - Worst Case (pivot at ends or sorted array): 0(N^2)
- Space Complexity:
 - O(log N) for recursive calls (stack space)
 - No extra array (unlike Merge Sort)

Why Quick Sort?

- ✓ Faster than Merge Sort in practice (no extra space)
- Efficient for large datasets
- \triangle But worst-case is $O(N^2)$ if pivot selection is poor

◆ To Practice Further

As an **assignment**, try:

- Implementing Quick Sort in descending order by tweaking comparisons (< becomes</p>
- > and vice versa)

Let me know if you want a C++, Python, or Java version of this implementation!