



**A COMPILER DESIGN LAB Project**

Submitted to

**Institute Of Engineering & Management, Kolkata**

Towards fulfilment for The

Bachelor of Technology

in

**INFORMATION TECHNOLOGY**

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## CERTIFICATE OF APPROVAL

This is to certify that the work embodied in this project entitled **Compiler Design Lab** submitted by our group to the Department of Information Technology, have been carried out under my direct supervision and guidance. The project work has been prepared as per the regulations of Institute of Engineering and Management and I strongly recommend that this project work be accepted in fulfillment of the minor Compiler Design Lab project.

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## **Experiment No. :- 1**

### **Aim:-**

To identify whether a given line is Comment Line or not.

### **Theory :-**

The goal of this experiment is to identify whether a given line is a comment line or not.

### **COMMENT LINE**

In programming, comments are hints that a programmer can add to make their code easier to read and understand.

### **TYPES OF COMMENT LINE**

Types of comments in programs:

Single Line Comment: Comments preceded by a Double Slash ('//')

Multi-line Comment: Comments starting with ('/\*') and ending with ('\*/').

input:-

```
#include <stdio.h>
int main() {
    // print Hello World to the screen
    printf("Hello World");
    return 0;
}
```

output:-

Hello World

## 1. Single-line Comments in C

In C, a single line comment starts with `//`. It starts and ends in the same line. For example,

```
#include <stdio.h>
int main() {

    // create integer variable
    int age = 25;

    // print the age variable
    printf("Age: %d", age);

    return 0;
}
```

output:-

Age: 25

## 2. Multi-line Comments in C

In C programming, there is another type of comment that allows us to comment on multiple lines at once, they are multi-line comments. To write multi-line comments, we use the `/*...*/` symbol. For example,

```
#include <stdio.h>
int main() {

    // create integer variable
    int age = 25;

    // print the age variable
    printf("Age: %d", age);

    return 0;
}
```

output:-

Age : 25

## **Procedure:-**

The simulator tab would:

- Please Enter Any Line in Text-Box Area.
- Enter Submit .
- It will Check if 1st position and 2nd position of a String is equal to "/" then "IT IS COMMENT LINE".
- If above is not True then Check if 1st position and last position of String is "/" and 2nd position and 2nd-last position of string is "\*", then "IT IS COMMENT LINE".
- If both are not True then "NOT A COMMENT LINE".
- Press clear button for clear Text Box.

## **Simulation:-**

The screenshot shows a web application interface for a simulation. At the top, the title is "Simulation 1:- To Check Given String Line is Comment Line or Not". Below the title, there is a label "Enter any line :". Underneath the label is a text input field containing the text "//hello". Below the input field, there are two buttons: a green "submit" button and a red "clear" button. To the right of these buttons, the text "IT IS COMMENT LINE" is displayed in a light gray box.

## **References:-**

- <https://www.javatpoint.com/comments-in-c>

## **Experiment No. :- 2**

### **Aim:-**

Design a lexical analyzer for given language and the lexical analyzer should ignore redundant spaces, tabs and new lines. It should also ignore comments, the length of the user input string can be user specific. From this particular statement count total number of individual token.

### **Theory:-**

- Lexical analysis is the very first phase in the compiler design.
- Lexemes and Tokens are the sequence of characters that are included in the source program according to the matching pattern of a token.
- Lexical analyzer is implemented to scan the entire source code of the program.
- Lexical analyzer helps to identify token into the symbol table .
- Removes one character from the remaining input is useful Error recovery method.
- Lexical Analyser scan the input program while parser perform syntax analysis.
- It eases the process of lexical analysis and the syntax analysis by eliminating unwanted tokens

### **Basic Terminologies:-**

#### **What's a lexeme?**

A lexeme is a sequence of characters that are included in the source program according to the matching pattern of a token. It is nothing but an instance of a token.

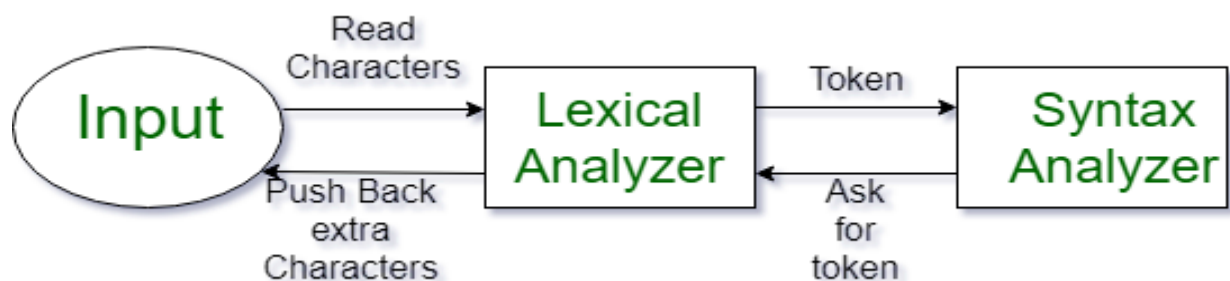
## What's a token?

Tokens in compiler design are the sequence of characters which represents a unit of information in the source program.

## Roles of the Lexical analyzer:-

Lexical analyzer performs below given tasks:

- Helps to identify token into the symbol table
- Removes white spaces and comments from the source program
- Correlates error messages with the source program
- Helps you to expands the macros if it is found in the source program
- Read input characters from the source program



## Procedure:-

1. Write some code in C programming language (NOTE:- Give Space every new word/character/operator).
2. Then click on given Submit button.
3. It will generate Tokens, identifiers and keywords list in the page.

## Simulation:-

**LEXICAL ANALYZER**

Follow the Following steps:

1. Write some code in C programming language.
2. The click on given four button to do that thing what written over there

```
#include
//helelo
int main ( )
{
char a ;
//showing how to write the simple
//code with space
return a ;
}
```

**TOKEN**

TOKEN	TYPE
int	KEYWORD
main	IDENTIFIER
(	Special Character
)	Special Character
char	KEYWORD

**NON\_TOKEN**

NON_TOKEN	DESCRIPTION
#include	IT's a pre processor
//helelo	IT's a comment
//showing how to write the simple	IT's a comment
//code with space	IT's a comment

## Reference:-

- <https://www.guru99.com/compiler-design-lexical-analysis.html>



## Experiment No. :- 3

### Aim:-

To recognize strings under 'a\*', 'a\*b+', abb.

### Theory:-

set of **productions** (P). The productions of a grammar specify the manner in which the terminals and non-terminals can be combined to form strings. Each production consists of a **non-terminal** called the left side of the production, an arrow, and a sequence of tokens and/or **on- terminals**, called the right side of the production.

a sequence of production rules, in order to get the input string. During parsing, we take two decisions for some sentential form of input:

- Deciding the non-terminal which is to be replaced.
- Deciding the production rule, by which, the non-terminal will be replaced.

### Procedure:-

- Enter Expression in Text-box.
- Click Button Clear for clear the text box and result box.
- Click Button submit and then .

Inside Algorithm will work:-

- Get String/Expression from user.
- Take a counter variable c and initialize 1.
- Check 1<sup>st</sup> position is not 'a' then break And print "NOT A VALID ENTRY" else next step.
- Now create a loop for check every character , if character is 'b' then c=0.
- Elseif character is 'a' and c==0 then print "NOT A VALID ENTRY" and break.
- And one more condition, if character is neither 'a' nor 'b' then print "NOT A VALID ENTRY".

- Outside loop we have to check whether loop reached last character or not if it reached then print "VALID ENTRY".

## **Simulation:-**

Simulation 1:- To Check Recognize string under 'a','a\*b+','abb'

Enter any line :

NOT A VALID ENTRY

## **References:-**

- [https://www.tutorialspoint.com/compiler\\_design/compiler\\_design\\_syntax\\_analysis.htm#:~:text=The%20productions%20of%20a%20grammar,right%20side%20of%20the%20production.](https://www.tutorialspoint.com/compiler_design/compiler_design_syntax_analysis.htm#:~:text=The%20productions%20of%20a%20grammar,right%20side%20of%20the%20production.)

## **Experiment No. :- 4**

### **Aim:-**

To check whether an given identifier is valid or not.

### **Theory:-**

#### **Identifier:-**

C identifiers represent the name in the C program, for example, variables, functions, arrays, structures, unions, labels, etc. An identifier can be composed of letters such as uppercase, lowercase letters, underscore, digits, but the starting letter should be either an alphabet or an underscore.

#### **Rules for constructing C identifiers**

- The first character of an identifier should be either an alphabet or an underscore, and then it can be followed by any of the character, digit, or underscore.
- It should not begin with any numerical digit.
- In identifiers, both uppercase and lowercase letters are distinct. Therefore, we can say that identifiers are case sensitive.
- Commas or blank spaces cannot be specified within an identifier.
- Keywords cannot be represented as an identifier.
- The length of the identifiers should not be more than 31 characters.
- Identifiers should be written in such a way that it is meaningful, short, and easy to read

### **Procedure:-**

1. Read the given input string from used by text-box.
2. Then Click on “Click” button .
3. It will Check the initial character of the string is numerical or any special character except ‘\_’ then print it is not a valid identifier.

4. Otherwise print it as valid identifier if remaining characters of string doesn't contains any special characters except '\_'.

## **Simulation:-**

Simulation :- TO CHECK GIVEN STRING IS VALID IDENTIFIER OR NOT

Now put the String in the given under filed

**IDENTIFIER**

*RESULT*

## **Reference:-**

- <https://www.javatpoint.com/c-identifiers#:~:text=C%20identifiers%20represent%20the%20name,an%20alphabet%20or%20an%20underscore.>

## **Experiment No. :- 5**

### **Aim:-**

To implementing Shift Reduce Parsing

### **Theory:-**

**Shift Reduce parser** attempts for the construction of parse in a similar manner as done in bottom-up parsing i.e. the parse tree is constructed from leaves(bottom) to the root(up). A more general form of the shift-reduce parser is the LR parser.

This parser requires some data structures i.e.

- An input buffer for storing the input string.
- A stack for storing and accessing the production rules

### **Basic Operations –**

- **Shift:** This involves moving symbols from the input buffer onto the stack.
- **Reduce:** If the handle appears on top of the stack then, its reduction by using appropriate production rule is done i.e. RHS of a production rule is popped out of a stack and LHS of a production rule is pushed onto the stack.
- **Accept:** If only the start symbol is present in the stack and the input buffer is empty then, the parsing action is called accept. When accepted action is obtained, it means successful parsing is done.
- **Error:** This is the situation in which the parser can neither perform shift action nor reduce action and not even accept action.

### **Procedure:-**

1. Put 1 or 2 or 3 any in very first text box.
2. Now when we put you are able to see the production in the another text area there.
3. Last just select given string as there you can see the dropdown
4. And click to submit
5. You able to see full way of **Shift Reduce Parse**.

## Simulation:-

Downloads x To\_STReck x +

File | D:/B.Tech/Compiler\_LA8/exp2.html

Apps classroom.google.c... login (67) Feed | Linkedln Free Coding Comp... 50+ tree questions... Binary Tree: Intervie... Inbox (2) - sumitksa... Problems - LeetCode Self learning >>

### SHIFT REDUCING PARSER

Put 1 or 2 or 3

1

Click

GRAMMER IS

S -> S + S

S -> S - S

S -> S / S

S -> S \* S

S -> id

CLEAR

Select the Input string from below menu:- a+b-c\*d/e

SUBMIT

Stack	Input Buffer	Parsing Action
S	343\$	Shift
S3	43\$	Shift
S34	3\$	Reduce By E->4
SSE	3\$	Shift
SSE3	\$	Reduce E->3E3
SE	\$	Accept
SS	-c*d/e\$	Shift
SS-	c*d/e\$	Shift
SS-c	*d/e\$	Reduce S->id
SS-S	*d/e\$	Reduce S->S-S
SS	*d/e\$	Shift
SS*	d/e\$	Shift
SS*d	/e\$	Reduce S->id
SS*S	/e\$	Reduce S->S*S
SS	/e\$	Shift
SS/	e\$	Shift
SS/e	\$	Reduce S->id
SS/S	\$	Reduce S->S/S
SS	\$	Accept

## References:-

- <https://www.geeksforgeeks.org/shift-reduce-parser-compiler/>