

NOTE: when using auxiliary data structures (e.g. stack and queue) can just publically inherit and add logic for the common interface methods

# Implementation

---

## Binary Search Tree

---

- many nullary public methods have a corresponding private method that takes a node (root to start) as its only argument
- constructors:
  - default -> root = nullptr
  - copy
  - move
- destructor:
  - call private clear() and then set root = nullptr
- traversal:
  - NOTE: each can potentially take a visit function (pointer, lambda, etc)
  - preorder (recursive):
    - nullary calls unary with one Node
    - if node != nullptr
      - visit node
      - preorder(node->left)
      - preorder(node->right)
  - inorder (recursive):
    - nullary calls unary with one Node
    - if node != nullptr
      - inorder(node->left)
      - visit node
      - inorder(node->right)
  - postorder (recursive):
    - nullary calls unary with one Node
    - if node != nullptr
      - postorder(node->left)
      - postorder(node->right)
      - visit node

- breadth first:
  - create queue
  - node = root
  - if node != nullptr
    - enqueue node
    - while !queue.empty
      - node = queue.dequeue
      - visit node
      - enqueue node-> left if node->left != nullptr
      - enqueue node-> right if node->right != nullptr
- iterative preorder:
  - create stack
  - node = root
  - if node != nullptr
    - push node
    - while !stack.empty
      - node = stack.pop
      - visit node
      - push node-> left if node->left != nullptr
      - push node-> right if node->right != nullptr
- iterative inorder:
  - NOTE: iterative inorder takes additional work
  - create stack
  - node = root
  - while node != nullptr
    - while node != nullptr
      - if node->right != nullptr
        - push node->right
      - push node
      - node = node->left
    - node = stack.pop
    - while !stack.empty
      - visit node
      - node = stack.pop
    - visit node
    - is !stack.empty
      - node = stack.pop
    - else node = nullptr
- iterative postorder:
  - NOTE: this is not just reordering preorder, visit has to follow pop in the while loop for the ordering to be any different
  - NOTE: can alternatively use two stacks to achieve the same results

- create stack
- current = root
- previous = root
- if node != nullptr
  - for ; current->left != nullptr; current = current->left
    - push current
  - while current->right != nullptr or current->right = previous
    - visit current
    - return if stack.empty
    - current = stack.pop
  - push current
  - current = current->right
- threaded inorder:
- morris inorder:
  - NOTE: modifies tree by making nodes with a left descendant the right child of the rightmost node in the node's left descendant
  - create node and tmp
  - while node != nullptr
    - if node->left == nullptr
      - visit node
      - node = node->right
    - else
      - tmp = node->left
      - while tmp->right != nullptr and tmp->right != node
        - tmp = tmp->right
      - if tmp->right == nullptr
        - tmp->right = node
        - node = node->left
      - else
        - visit node
        - tmp->right = nullptr
        - node = node->right
- search:
  - return Value\*
  - take Node and Key&
  - while node != nullptr
    - if key == node->key
      - return &node->value
    - else if key < node->key
      - node = node->left
    - else node = node->right

- return nullptr
- insert/put:
  - iterative:
    - take Key& key, Value& value
    - if root is nullptr
      - create new node with key, value
      - set as root
      - return
    - create current = root and previous = nullptr
    - while current != nullptr
      - previous = current
      - if key < current->key
        - current = current->left
      - else current = current->right
    - if key < previous->key
      - previous->left = new node with key, value
    - else
      - previous->right = new node with key, value
  - recursive:
    - two methods:
      - one takes key, value
      - one takes node, key, value
    - if node == nullptr
      - return new node with key, value
    - if key < node->key
      - node->left = recursively call(node->left, key, val)
    - else if key > node->key
      - node->right = recursively call(node->right, key, val)
    - else node->value = value
    - (set size of node if stored)
    - return
- get:
  - return Value\*
  - use search algorithm above and return nullptr for search miss
- delete:
  - delete by merging
  - find and delete by merging
  - delete by copying

- balance:
  - balance factor (either computed or stored):
    - one book says balance factor should be -1, 0, 1
    - return height of node->left - height of node->right
  - with sorted array (recursively):
    - take as arguments: vector data(sorted), first, last
    - if first <= last
      - middle = (first + last) / 2
      - insert/put middle into tree
      - balance(data, first, middle)
      - balance(data, middle + 1, last)
  - DSW algorithm:
    - TODO: implement this and review
    - also see Day-Stout-Warren on wiki for another approach
    - NOTE: creates a backbone (a linked list from root to right)
    - create backbone:
      - take root as argument
      - grandparent = nullptr
      - parent = root
      - left child
      - while parent != nullptr
        - if parent->left != nullptr
          - grandparent = rotate right(parent)
          - parent = left child
        - else
          - grandparent = parent
          - parent = parent->right
    - make rotations:
      - take an integer n as argument
      - parent = root
      - for n > 0; --n
        - rotate left(parent)
    - dsw balance:
      - if root != nullptr
        - create backbone
        - for tmp = root, n = 0; tmp != nullptr; tmp = tmp->right
          - ++n
        - $m = (2^{\lg(n + 1)} - 1)$

- make rotations (n - m)
  - while m > 1
    - m = m / 2
    - make rotations (m)
- for AVL tree:
  - if balance factor of node < -1
    - if balance factor node->right > 0
      - node->right = rotate right node->right
    - node = rotate left node
  - else if balance factor of node > 1
    - if balance factor node->left < 0
      - node->left = rotate right node->left
    - node = rotate right node
  - return node
- rotate right:
  - option 1:
    - take as arguments grandparent, parent, child
    - grandparent, parent != nullptr
      - parent->left = child-right
      - child->right = parent
      - set grandparent's parent pointer to child
  - option 2:
    - take one node pointer reference as argument (parent)
    - if parent != nullptr
      - child = parent->left
      - parent->left = child->right
      - child->right = parent
      - parent = child
    - (moves E and its left subtree up one level and S and its right subtree down one level)
- general:
  - contains
  - is empty
  - size
  - height:
    - if node == nullptr
      - return 0
    - return 1 + max of height of node->left and height of node->right

- symbol table:
  - keys
  - values
  - min
  - max
  - floor(Key)
  - ceiling(Key)
  - rank(Key)
  - select(int)
  - delete min
  - delete max
  - size(lo, hi)

## Threading Trees

---

- nodes hold pointers to predecessors and successors
- overload existing pointers:
  - left is either:
    - left child
    - predecessor
  - right is either:
    - right child
    - successor
- successor = top of tree after a leaf
- add a data member to indicate if the right points to child or successor

## Iterators (in C++)

---

- need a parent link in the node to just store a pointer
- useful to use a dummy root node to implement end
- with no parent link, can construct a linear structure (stack, queue, vector, etc) to store node pointers in a specific order but this will easily be invalidated
- can also always store parent and current ptr in iterator and set parent to nullptr for root (doesn't work, cannot move back up still)

## Handling misses in APIs that take keys as an argument

---

- C++ -> first option is to choose is throw `std::out_of_range` or related
- Java implementation allows for null, this requires a pointer type in C++
- the misses return null (Java)

- the C++ std library returns an iterator that references an element or is equivalent to the end() iterator
- Herb Sutter says for DAG/tree that hands out strong references to data work with shared\_ptr (since keys can be data/large data also, makes sense to use them here as well)

## Value semantics

---

- in general, do not make as much effort to avoid value semantics
- libraries use references for speed but value semantics simplifies APIs and can be more understandable

## Options

---

- APIs use references
- use pair to store key & value
- use smart pointers to store key & value
- add parent pointers to nodes
- to solve the issue of the usage of null in Java API's, instead of storing as a pointer, can just return the address of the data as a T\* (Non\_owning\_raw\_pointer)
- use weak\_ptr to return references to shared\_ptr