

Design and Analysis of Algorithms

Lecture - 11

Success is always inevitable with Hard Work and Perseverance

N. Ravitha Rajalakshmi

Learning Objective

• Discuss D&C strategy for classical problems

Sorting

Sorting

Input: An array A with n elements.

Output: Permutation of Array A where elements are arranged in non-decreasing order.

Initial Array	130	10	40	8	20	200
Sorted Array	8	10	20	40	130	200

Merge Sort Array Elements

130 10 40 8 20 200

Split the array into two halves

130 10 40

8 20 200

Sort them independently

10 40 130

8 20 200

Combine

8 10 20 40 130 200

Pause & Think

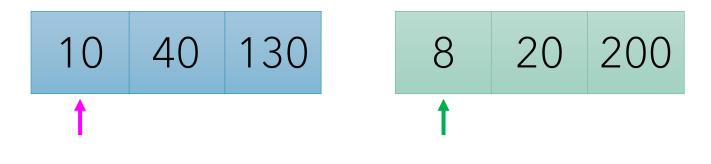
Can we split the array and do sorting in parallel?
 Yes

Will it affect the final answer?
 No

Is there a mechanism for combining the independent solutions?
 To be determined

Combine

• Given two sorted arrays, how to combine it such that resultant array is sorted?





Function Merge(A, low, mid, high)

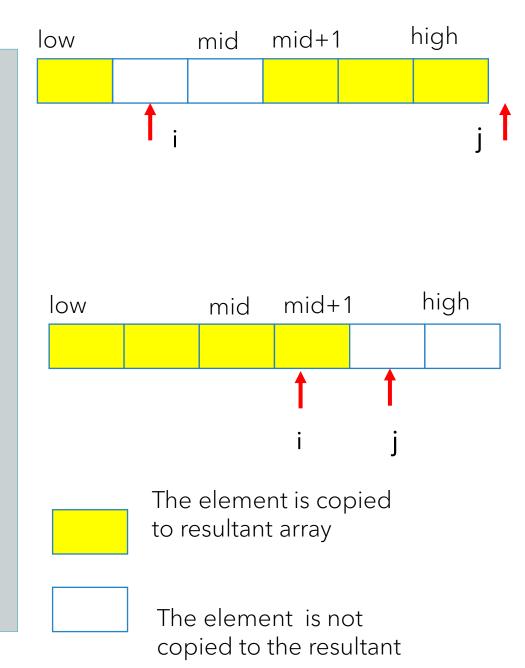
```
# First subarray contains elements from low to mid
# Second subarray contains elements from mid + 1 to high
       i <-low, k<-mid+1 n<-low B[low.... high]
       while(i \le mid \&\& k \le high){
              if(A[i] < = A[k])
                     B[n] \leftarrow A[i]
                     i++
                     n++
              else{
                     B[n] < -A[k]
                     k++
                     n++
```

Pause & Think

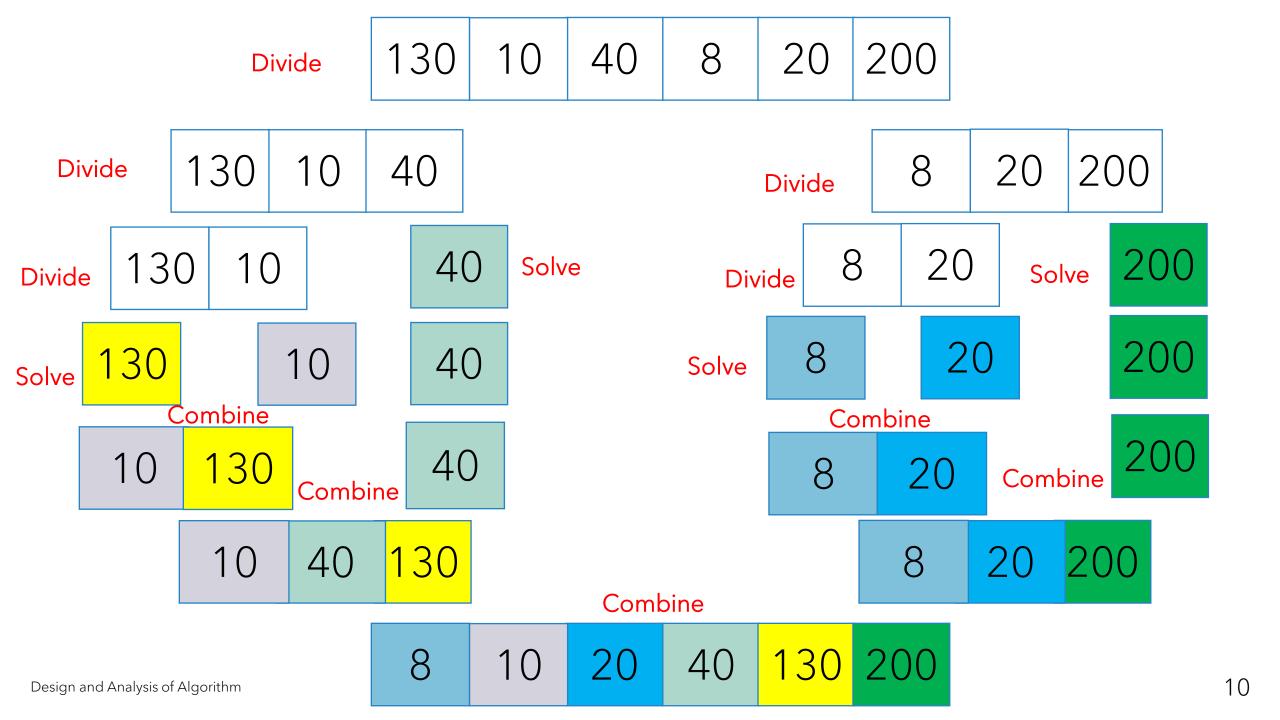
Is the merge algorithm complete?
 No

After the process, Are all the elements copied to resultant array?
 No , 200 is not yet copied to the resultant array

```
if(i \le mid)
      for(int m=i; m \le mid; m++){
             B[n] < -A[m]
              n++
if(k \le high){
      for(int m=k; m \le high; m++){
             B[n] < -A[m]
              n++
Copy the array elements from B to A
```



array



Function MergeSort(A, low, high)

```
# If the array contains more than one element
if (low <=high) {
      # Divide the array into two subarrays
      mid = (low + high)/2
      MergeSort(A, low, mid)
      MergeSort(A, mid+1, high)
      #Combine the solutions
      Merge(A, low, mid, high)
```

Analyzing Time Complexity

- Recurrence Relation
- Input size : n
- Basic Operation : Comparison
- Time Complexity : T(n)

$$T(1) = 0$$

$$T(n) = 2T\left(\frac{n}{2}\right) + n \text{ (Cost of Merge)}$$

• Using Master's theorem , $T(n) = O(n \log n)$

Pause & Think

Whether the algorithm is Stable?

Yes Order of elements with equal values is retained after sorting

• Whether the algorithm is in-place?

No

Extra memory equal to the size of the array is used in merge

Summary

How D&C is applied for Merge Sort

Thank You Happ Learning

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